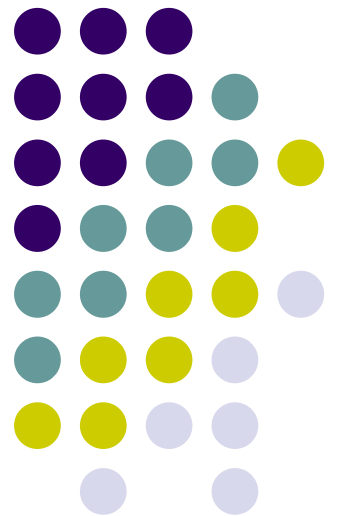


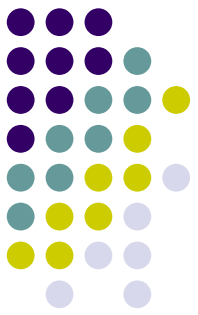
Computer Organization: A Programmer's Perspective

Machine-Level Programming (2: Conditions and Branches)

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Condition Codes



Single Bit Registers

CF Carry Flag (unsigned)

ZF Zero Flag

SF Sign Flag (for signed)

OF Overflow Flag (for signed)

Implicitly Set By Arithmetic Operations

`addl Src, Dest` same as C analog: `t = a + b`

CF set if carry out from most significant bit

Used to detect unsigned overflow

ZF set if `t == 0`

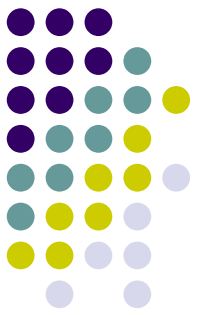
SF set if `t < 0`

OF set if two's complement overflow

`(a > 0 && b > 0 && t < 0) || (a < 0 && b < 0 && t >= 0)`

Not Set by `leal` instruction

Setting Condition Codes (cont.)



Explicit Setting by Compare Instruction

`cmpl Src2,Src1`

`cmpl b, a` like computing $a-b$ without setting destination

CF set if carry out from most significant bit

Used for unsigned comparisons

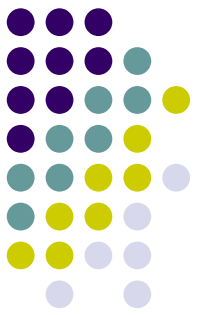
ZF set if $a == b$

SF set if $(a-b) < 0$

OF set if two's complement overflow

$(a > 0 \ \&\& \ b < 0 \ \&\& \ (a-b) < 0)$
 $|| \ (a < 0 \ \&\& \ b > 0 \ \&\& \ (a-b) > 0)$

Setting Condition Codes (cont.)



Explicit Setting by Test instruction

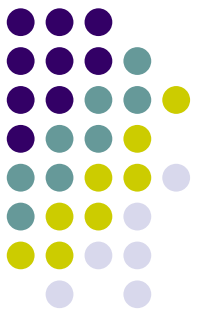
`testl Src2,Src1` like `Src1&Src2` in C

- Sets condition codes based on value of *Src1* & *Src2*
- Useful to have one of the operands be a mask

ZF set when `a&b == 0`

SF set when `a&b < 0`

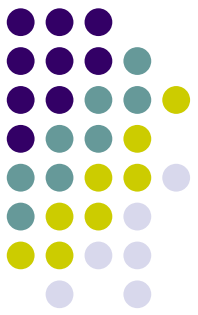
Reading Condition Codes



SetX Instructions

- Set lower byte of destination, no change to upper 7 bytes
- Based on combinations of condition codes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	\simZF	Not Equal / Not Zero
sets	SF	Negative
setns	\simSF	Nonnegative
setg	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed)
setge	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
setl	$(SF \wedge OF)$	Less (Signed)
setle	$(SF \wedge OF) \mid ZF$	Less or Equal (Signed)
seta	$\sim CF \ \& \ \sim ZF$	Above (unsigned)
setb	CF	Below (unsigned)



Reading Condition Codes (Cont.)

SetX Instructions (32 bits)

Set single byte based on combinations of condition codes

One of 8 addressable byte registers

Embedded within first 4 integer registers

Does not alter remaining 3 bytes

Typically use `movzbl` to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

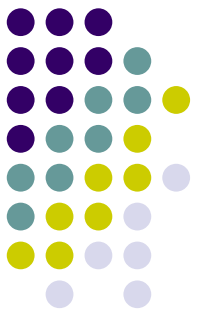
Body

```
movl 12(%ebp),%eax    # eax = y
cmpl %eax,8(%ebp)     # Compare x : y
setg %al              # al = x > y
movzbl %al,%eax       # Zero rest of %eax
```

%eax	%ah	%al
%edx	%dh	%dl
%ecx	%ch	%cl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

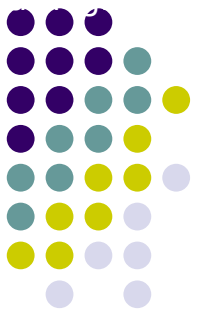
Note
inverted
ordering!

Jumping



Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	\simZF	Not Equal / Not Zero
js	SF	Negative
jns	\simSF	Nonnegative
jg	\sim (SF\wedgeOF) & \simZF	Greater (Signed)
jge	\sim (SF\wedgeOF)	Greater or Equal (Signed)
jl	(SF\wedgeOF)	Less (Signed)
jle	(SF\wedgeOF) ZF	Less or Equal (Signed)
ja	\simCF & \simZF	Above (unsigned)
jb	CF	Below (unsigned)



Conditional Branch Example (Old Style, 64 bits)

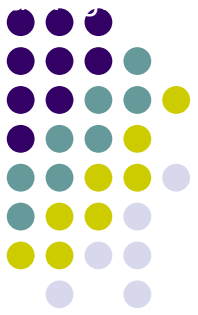
■ Generation

> gcc -Og -S -fno-if-conversion control.c

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi, %rdi    # x:y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:       # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value



Conditional Move

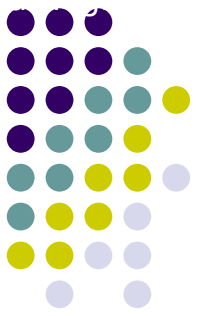
```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

absdiff:

```
movq    %rdi, %rax    # x
subq    %rsi, %rax    # result = x-y
movq    %rsi, %rdx
subq    %rdi, %rdx    # eval = y-x
cmpq    %rsi, %rdi    # x:y
cmovle  %rdx, %rax    # if <=, result = eval
ret
```

Bad Cases for Conditional Move



Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

- Both values get computed
- Only makes sense when computations are very simple

Risky Computations

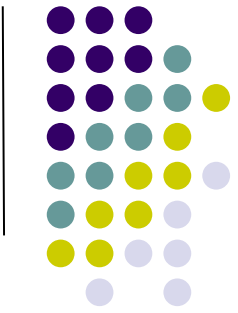
```
val = p ? *p : 0;
```

- Both values get computed
- May have undesirable effects

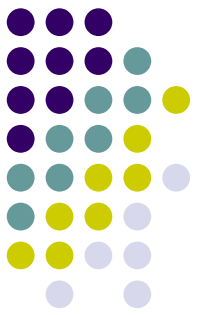
Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- Must be side-effect free



“Do-While” Loop Example



C Code

```
int fact_do
(int x)
{
    int result = 1;
    do {
        result *= x;
        x = x-1;
    } while (x > 1);
    return result;
}
```

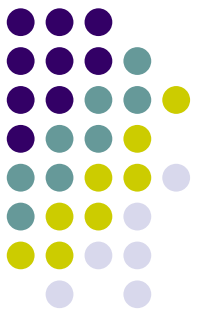
Goto Version

```
int fact_goto(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

Use backward branch to continue looping

Only take branch when “while” condition holds

“Do-While” Loop Compilation



Goto Version

```
int fact_goto
(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

32 bit Assembly

```
_fact_goto:
    pushl %ebp           # Setup
    movl %esp,%ebp      # Setup
    movl $1,%eax         # eax = 1
    movl 8(%ebp),%edx    # edx = x

L11:
    imull %edx,%eax      # result *= x
    decl %edx            # x--
    cmpl $1,%edx         # Compare x : 1
    jg L11               # if > goto loop

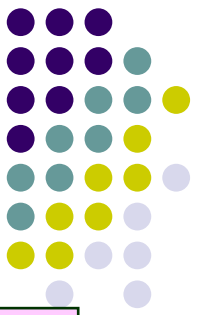
    movl %ebp,%esp      # Finish
    popl %ebp           # Finish
    ret                 # Finish
```

Registers

%edx x

%eax result

“While” Loop #1 (bad)



C Code

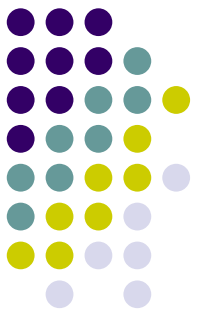
```
int fact_while
(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

First Goto Version

```
int fact_while_goto
(int x)
{
    int result = 1;
loop:
    if (!(x > 1))
        goto done;
    result *= x;
    x = x-1;
    goto loop;
done:
    return result;
}
```

Must jump out of loop if test fails: inefficient!

“While” Loop #2 (good)



C Code

```
int fact_while(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

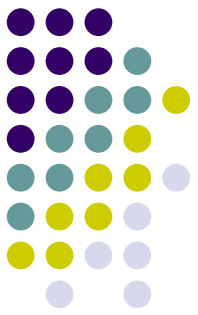
Second Goto Version

```
int fact_while_goto2
(int x)
{
    int result = 1;
    if (!(x > 1))
        goto done;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
done:
    return result;
}
```

Guards loop entry with extra test

Uses same inner loop as do-while version: efficient!

General “While” Translation



C Code

```
while (Test)  
    Body
```



Do-While Version

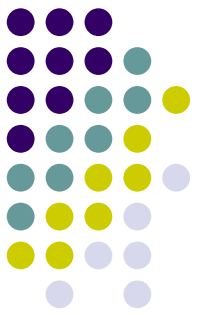
```
if (!Test)  
    goto done;  
do  
    Body  
    while(Test) ;  
done:
```



Goto Version

```
if (!Test)  
    goto done;  
loop:  
    Body  
    if (Test)  
        goto loop;  
done:
```


“For” Loop Example



General Form

```
int result;
for (result = 1;
    p != 0;
    p = p>>1) {
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

```
for (Init; Test; Update )
    Body
```

Init

```
result = 1
```

Test

```
p != 0
```

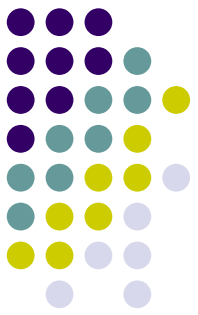
Update

```
p = p >> 1
```

Body

```
{
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

“For” → “While”



For Version

```
for (Init; Test; Update)  
    Body
```

While Version

```
Init;  
while (Test) {  
    Body  
    Update ;  
}
```

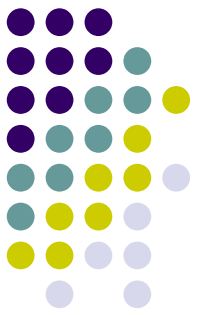
Do-While Version

```
Init;  
if (!Test)  
    goto done;  
do {  
    Body  
    Update ;  
} while (Test)  
done:
```

Goto Version

```
Init;  
if (!Test)  
    goto done;  
loop:  
    Body  
    Update ;  
    if (Test)  
        goto loop;  
done:
```

Switch Statements



```
typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
  op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
        case ADD :
            return '+';
        case MULT:
            return '*';
        case MINUS:
            return '-';
        case DIV:
            return '/';
        case MOD:
            return '%';
        case BAD:
            return '?';
    }
}
```

Implementation Options

Series of conditionals

Good if only few cases

Jump Table

Lookup branch target

Avoids conditionals

Possible when cases are
small integer consts

GCC: Picks one method

Bug in example code

No default given

Jump Table Structure



Switch Form

```
switch(op) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
    . . .  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

jtab:

Targ0
Targ1
Targ2
⋮
Targn-1

Jump Targets

Targ0:

Code Block
0

Targ1:

Code Block
1

Targ2:

Code Block
2

⋮

Targn-1:

Code Block
n-1

Approx. Translation (not valid C)

```
target = JTab[op];  
goto *target;
```

Switch Statement Example



Branching Possibilities

```
typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
  op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
        . . .
    }
}
```

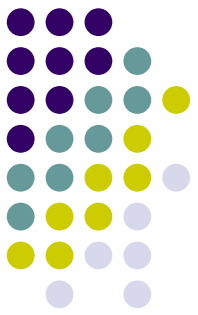
Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Setup:

```
unparse_symbol:
    pushl %ebp                # Setup
    movl %esp,%ebp           # Setup
    movl 8(%ebp),%eax         # eax = op
    cmpl $5,%eax              # Compare op : 5
    ja .L49                   # If > goto done
    jmp *.L57(,%eax,4)        # goto Table[op]
```

Assembly Setup Explanation



Symbolic Labels

Labels of form `.LXX` translated into addresses by assembler

Table Structure

Each target requires 4 bytes

Base address at `.L57`

Jumping

```
jmp .L49
```

Jump target is denoted by label `.L49`

```
jmp *.L57(, %eax, 4)
```

Start of jump table denoted by label `.L57`

Register `%eax` holds `op`

Must scale by factor of 4 to get offset into table

Fetch target from effective Address $.L57 + op * 4$

Jump Table

Table Contents

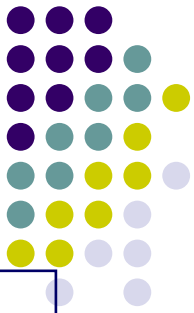
```
.section .rodata
    .align 4
.L57:
    .long .L51 #Op = 0
    .long .L52 #Op = 1
    .long .L53 #Op = 2
    .long .L54 #Op = 3
    .long .L55 #Op = 4
    .long .L56 #Op = 5
```

Enumerated Values

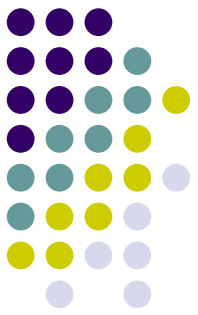
ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Targets & Completion

```
.L51:
    movl $43,%eax # '+'
    jmp .L49
.L52:
    movl $42,%eax # '*'
    jmp .L49
.L53:
    movl $45,%eax # '-'
    jmp .L49
.L54:
    movl $47,%eax # '/'
    jmp .L49
.L55:
    movl $37,%eax # '%'
    jmp .L49
.L56:
    movl $63,%eax # '?'
    # Fall Through to .L49
```



Switch Statement Completion



<code>.L49:</code>	<code># Done:</code>
<code>movl %ebp,%esp</code>	<code># Finish</code>
<code>popl %ebp</code>	<code># Finish</code>
<code>ret</code>	<code># Finish</code>

Puzzle

What value returned when `op` is invalid?

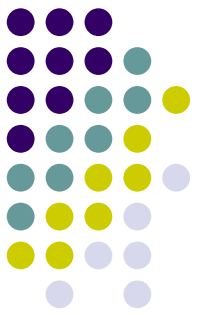
Answer

Register `%eax` set to `op` at beginning of procedure

This becomes the returned value

Advantage of Jump Table:
k-way branch in $O(1)$

Object Code



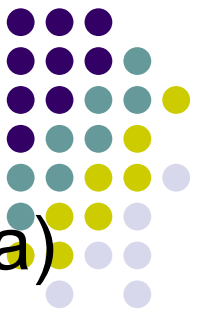
Setup

Label `.L49` becomes address `0x804875c`

Label `.L57` becomes address `0x8048bc0`

```
08048718 <unparse_symbol>:
8048718: 55                pushl    %ebp
8048719: 89 e5            movl    %esp, %ebp
804871b: 8b 45 08         movl    0x8(%ebp), %eax
804871e: 83 f8 05         cmpl    $0x5, %eax
8048721: 77 39            ja      804875c <unparse_symbol+0x44>
8048723: ff 24 85 c0 8b   jmp     *0x8048bc0(, %eax, 4)
```

Extracting Jump Table from Binary



Jump Table Stored in Read Only Data Segment (.rodata)

Various fixed values needed by your code

Can examine with objdump

```
objdump code-examples -s --section=.rodata
```

Show everything in indicated segment.

Difficult to read

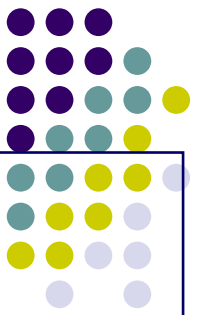
Jump table entries shown with reversed byte ordering

Contents of section .rodata:

```
8048bc0 30870408 37870408 40870408 47870408 0...7...@...G...
8048bd0 50870408 57870408 46616374 28256429 P...W...Fact(%d)
8048be0 203d2025 6c640a00 43686172 203d2025   = %ld..Char = %
...
```

E.g., 30870408 really means 0x08048730

Disassembled Targets



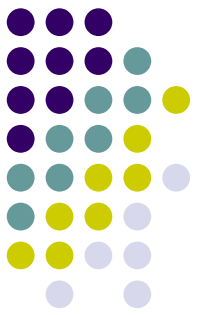
```
8048730:b8 2b 00 00 00 movl $0x2b,%eax
8048735:eb 25 jmp 804875c <unparse_symbol+0x44>
8048737:b8 2a 00 00 00 movl $0x2a,%eax
804873c:eb 1e jmp 804875c <unparse_symbol+0x44>
804873e:89 f6 movl %esi,%esi
8048740:b8 2d 00 00 00 movl $0x2d,%eax
8048745:eb 15 jmp 804875c <unparse_symbol+0x44>
8048747:b8 2f 00 00 00 movl $0x2f,%eax
804874c:eb 0e jmp 804875c <unparse_symbol+0x44>
804874e:89 f6 movl %esi,%esi
8048750:b8 25 00 00 00 movl $0x25,%eax
8048755:eb 05 jmp 804875c <unparse_symbol+0x44>
8048757:b8 3f 00 00 00 movl $0x3f,%eax
```

`movl %esi,%esi` does nothing

Inserted to align instructions for better cache performance

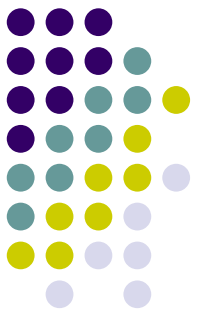
Done automatically by compiler, (remember the -O flag?)

Matching Disassembled Targets



Entry

0x08048730	8048730: b8 2b 00 00 00	movl
0x08048737	8048735: eb 25	jmp
0x08048740	8048737: b8 2a 00 00 00	movl
0x08048747	804873c: eb 1e	jmp
0x08048750	804873e: 89 f6	movl
0x08048757	8048740: b8 2d 00 00 00	movl
	8048745: eb 15	jmp
	8048747: b8 2f 00 00 00	movl
	804874c: eb 0e	jmp
	804874e: 89 f6	movl
	8048750: b8 25 00 00 00	movl
	8048755: eb 05	jmp
	8048757: b8 3f 00 00 00	movl



שאלות?