

Assembly

תרגול 5
תכנות באסמבלי, המשך

Condition Codes

- Single bit registers
 - ZF – zero flag
 - SF – sign flag
 - CF – carry flag
 - OF – overflow flag
- Relevant only for the most recent operation
- `leaq` does not alter any condition code
- In logical operations, CF and OF are set to 0
- For shift operations, OF is set to 0, CF is the last shifted out bit

ZF – Zero Flag

- Set if a result is zero

- Example:

- $\text{Addb } 0,0 \rightarrow ZF = 1$

SF – Sign Flag

- Set if a result is negative (MSB of the result = 1)■

- Example:

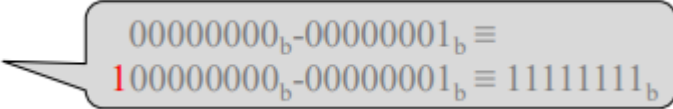
□ $\text{subb } 00000000_b, 00000001_b \rightarrow 11111111_b \rightarrow SF = 1$

CF – Carry Flag

- Used for unsigned arithmetic
- CF gets 1 if result is larger than capacity of the target operand
- Examples:

$$\square 11111111_b + 00000001_b = 1 \underbrace{00000000}_{0*8}_b = 0 \rightarrow CF = 1$$

$$\square 00000000_b - 00000001_b = \underbrace{11111111}_{1*8}_b \rightarrow CF = 1$$

unsigned arithmetic  (decimal) = 255 (decimal) → ~~Get the wrong answer~~

signed arithmetic → $11111111_b = -1(\text{decimal}) \rightarrow$
ignore CF flag

OF – Overflow Flag

- Used for signed arithmetic
- sign-bit-off operands + sign-bit-on result and vice versa → OF gets '1'

- Examples:

$$\square 01111111_b + 01000000_b = 1 \underbrace{00000000}_{0*7}_b \rightarrow OF = 1$$

$$\square 10000000_b + 10000000_b = 1 \underbrace{00000000}_{1*8}_b = 0 \rightarrow OF = 1$$

- Otherwise OF gets '0'

Usage examples

- Check equality of two values:
 - `subl %ecx, %edx` and `ZF = 1`

But then, the subtraction result we be saved in `%edx`

Setting Condition Codes

Instruction		Based on	Description
<code>cmpb</code>	S_2, S_1	$S_1 - S_2$	Compare bytes
<code>testb</code>	S_2, S_1	$S_1 \& S_2$	Test byte
<code>cmpw</code>	S_2, S_1	$S_1 - S_2$	Compare words
<code>testw</code>	S_2, S_1	$S_1 \& S_2$	Test word
<code>cmpd</code>	S_2, S_1	$S_1 - S_2$	Compare double words
<code>testd</code>	S_2, S_1	$S_1 \& S_2$	Test double word

- Similarly, `cmpq` and `testq`
- Incase we do not need to save the result

Test usage examples

- Check if a number is '0':
 - `testl %ecx, %ecx` and check if `ZF = 1`

Setting Condition Codes (cont.)

Explicit Setting by Test instruction

`testl Src2,Src1`

- Sets condition codes based on value of *Src1* & *Src2*
 - Useful to have one of the operands be a mask
- `testl b, a` like computing `a&b` without setting destination
- ZF set when `a&b == 0`
- SF set when `a&b < 0`

Accessing the Condition Codes

Instruction	Synonym	Effect	Set Condition
sete <i>D</i>	setz	$D \leftarrow ZF$	Equal / Zero
setne <i>D</i>	setnz	$D \leftarrow \sim ZF$	Not Equal / Not Zero
sets <i>D</i>		$D \leftarrow SF$	Negative
setns <i>D</i>		$D \leftarrow \sim SF$	Nonnegative
setg <i>D</i>	setnle	$D \leftarrow \sim (SF \wedge OF) \& \sim ZF$	Greater (Signed >)
setge <i>D</i>	setnl	$D \leftarrow \sim (SF \wedge OF)$	Greater or Equal (Signed >=)
setl <i>D</i>	setnge	$D \leftarrow SF \wedge OF$	Less (Signed <)
setle <i>D</i>	setng	$D \leftarrow (SF \wedge OF) \mid ZF$	Less or Equal (Signed <=)
seta <i>D</i>	setnbe	$D \leftarrow \sim CF \& \sim ZF$	Above (Unsigned >)
setae <i>D</i>	setnb	$D \leftarrow \sim CF$	Above or Equal (Unsigned >=)
setb <i>D</i>	setnae	$D \leftarrow CF$	Below (Unsigned <)
setbe <i>D</i>	setna	$D \leftarrow CF \& \sim ZF$	Below or Equal (Unsigned <=)

Figure 3.9: The set Instructions. Each instruction sets a single byte to 0 or 1 based on some combination of the condition codes. Some instructions have “synonyms,” i.e., alternate names for the same machine instruction.

Why does it work?

Let's take `cmpl b, a` for example

- If there is no overflow and $a \geq b$ → $SF=0, OF=0$
- If there is no overflow and $a < b$ → $SF=1, OF=0$
- If there is a negative overflow ($a > b$) → $SF=1, OF=1$
- If there is a positive overflow ($a < b$) → $SF=0, OF=1$

a < b in assembly

- Translate the line: return (a<b);
- Suppose a is in %edx, b is in %eax:

```
    cmpl    %eax,%edx                # compare a to b
```

<code>setl D</code>	<code>setnge</code>	<code>D ← SF ^ OF</code>	<code>Less (Signed <)</code>
------------------------	---------------------	--------------------------	---------------------------------

```
    movzbl   %al,%eax                # set %eax to 0 or 1
```

Unconditional Jump

■ Direct jump

```
jmp L1
```

■ Indirect jump

```
□ jmp *%rax
```

```
□ jmp *(%rax)
```

```
L1:
```

```
movq %rdx, %rcx
```

```
movq %rdx, %rcx
```

```
jmp L1
```

Unconditional Jump

■ Direct jump

```
jmp L1
```

■ Indirect jump

```
□ jmp *%rax
```

```
□ jmp *(%rax)
```

```
%rax = L2
```

```
L1:
```

```
    movq %rdx, %rcx
```

```
L2:
```

```
    movq %rdx, %rcx
```

```
    jmp *%rax
```

Unconditional Jump

■ Direct jump

```
jmp L1
```

■ Indirect jump

```
□ jmp *%rax
```

```
□ jmp *(%rax)
```

```
%rax = 1000
```

```
L1:
```

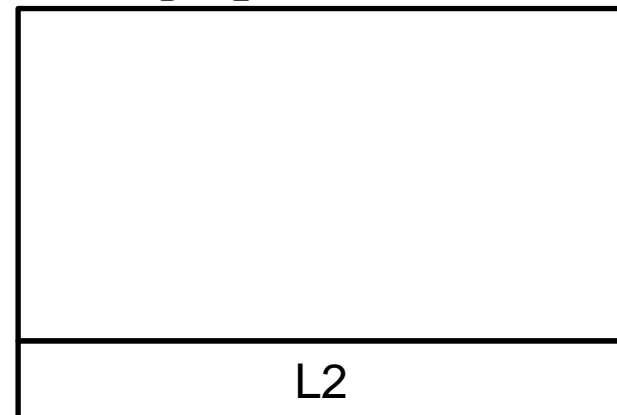
```
movq %rdx, %rcx
```

```
L2:
```

```
movq %rdx, %rcx
```

```
jmp *(%rax)
```

1000



Jump Instructions

Instruction	Synonym	Jump Condition	Description
<code>jmp Label</code>		1	Direct Jump
<code>jmp *Operand</code>		1	Indirect Jump
<code>jz Label</code>	<code>jz</code>	ZF	Equal / Zero
<code>jne Label</code>	<code>jnz</code>	$\sim ZF$	Not Equal / Not Zero
<code>js Label</code>		SF	Negative
<code>jns Label</code>		$\sim SF$	Nonnegative
<code>jg Label</code>	<code>jnle</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed >)
<code>jge Label</code>	<code>jnl</code>	$\sim (SF \wedge OF)$	Greater or Equal (Signed >=)
<code>jl Label</code>	<code>jnge</code>	$SF \wedge OF$	Less (Signed <)
<code>jle Label</code>	<code>jng</code>	$(SF \wedge OF) \mid ZF$	Less or Equal (Signed <=)
<code>ja Label</code>	<code>jnbe</code>	$\sim CF \ \& \ \sim ZF$	Above (Unsigned >)
<code>jae Label</code>	<code>jnb</code>	$\sim CF$	Above or Equal (Unsigned >=)
<code>jb Label</code>	<code>jnae</code>	CF	Below (Unsigned <)
<code>jbe Label</code>	<code>jna</code>	$CF \ \& \ \sim ZF$	Below or Equal (Unsigned <=)

Figure 3.10: **The jump Instructions.** These instructions jump to a labeled destination when the jump condition holds. Some instructions have “synonyms,” alternate names for the same machine instruction.



Conditional Jump

- Can't use indirect jump
- Use it to implement
 - if conditions
 - loops
 - switch statements

Goto in C

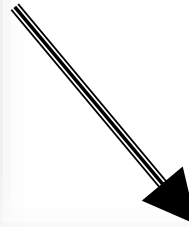
```
if (test-expr)
    then-statement
else
    else-statement
```



```
t = test-expr;
if (t)
    goto true;
else-statement
goto done;
true:
    then-statement
done:
```

If Condition in Assembly

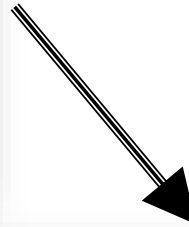
```
1 int absdiff(int x, int y)
2 {
3     if (x < y)
4         return y - x;
5     else
6         return x - y;
7 }
```



```
1  # x in %edi, y in %esi
2  cmpl    %esi,%edi    # Compare x:y
3  jl     .L3           # if <, go to less:
4  subl    %esi,%edi    # Compute x-y as return value
5  movl    %edi,%eax    # Set as return value
6  jmp     .L5           # Goto done:
7  .L3:                # less:
8  subl    %edi,%esi    # Compute y-x as return value
9  movl    %esi,%eax    # Set as return value
10 .L5:                # done: Begin completion code
```

If Condition in Assembly

```
1 int absdiff(int x, int y)
2 {
3     if (x < y)
4         return y - x;
5     else
6         return x - y;
7 }
```




```
1  # x in %edi, y in %esi
2  cmpl    %esi,%edi    # Compare x:y
3  

|     |       |      |         |                 |
|-----|-------|------|---------|-----------------|
| jnl | Label | jnge | SF ^ OF | Less (Signed <) |
|-----|-------|------|---------|-----------------|


4  subl    %esi,%edi    # Compute x-y as return value
5  movl    %edi,%eax    # Set as return value
6  jmp     .L5          # Goto done:
7  .L3:                # less:
8  subl    %edi,%esi    # Compute y-x as return value
9  movl    %esi,%eax    # Set as return value
10 .L5:                # done: Begin completion code
```

Do-While Loops

<pre>do <i>body-statement</i> while (<i>test-expr</i>) ;</pre>		<pre>loop: <i>body-statement</i> t = <i>test-expr</i>; if (t) goto loop;</pre>
--	--	--

Do-While Loops in Assembly

```
1 long fib_dw(long n)
2 {
3     long i = 0;
4     long val = 0;
5     long nval = 1;
6
7     do {
8         long t = val + nval;
9         val = nval;
10        nval = t;
11        i++;
12    } while(i < n);
13
14    return val;
15 }
```



```
1 .L6:
2     leaq    (%rdx,%rbx),%rax
3     movq    %rdx,%rbx
4     movq    %rax,%rdx
5     incq    %rcx
6     cmpq    %rdi,%rcx
7     jl      .L6
8     movq    %rbx,%rax
```

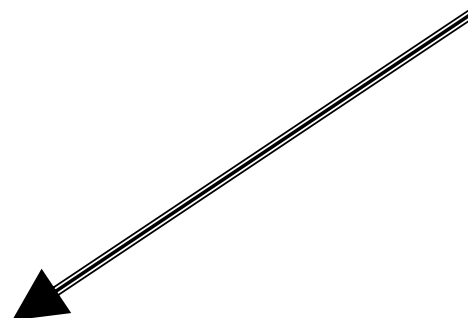
Register Usage		
Register	Variable	Initially
%rcx	i	0
%rdi	n	n
%rbx	val	0
%rdx	nval	1
%rax	t	— ₂₃

While Loops

```
while (test-expr)  
  body-statement
```



```
if (!test-expr)  
  goto done;  
do  
  body-statement  
  while (test-expr);  
done:
```



```
  t = test-expr;  
  if (!t)  
    goto done;  
loop:  
  body-statement  
  t = test-expr;  
  if (t)  
    goto loop;  
done:
```


Exercise

```
1 int loop_while(int a, int b)
2 {
3     int i = 0;
4     int result = a;
5     while (i < 256) {
6         result += a;
7         a -= b;
8         i += b;
9     }
10    return result;
11 }
```



```
1  # a in %edi, b in %esi
2  xorl    %ecx,%ecx
3  movl    %edi,%edx
4  .p2align 4,,7
5  .L5:
6  addl    %edi,%edx
7  subl    %esi,%edi
8  addl    %esi,%ecx
9  cmpl    $255,%ecx
10 jle .L5
```

Register	Variable	Initially
%eax		
%ebx		
%ecx		
%edx		

Exercise

```
1 int loop_while(int a, int b)
2 {
3     int i = 0;
4     int result = a;
5     while (i < 256) {
6         result += a;
7         a -= b;
8         i += b;
9     }
10    return result;
11 }
```



```
1  # a in %edi, b in %esi
2  xorl    %ecx,%ecx
3  movl    %edi,%edx
4  .p2align 4,,7
5  .L5:
6  addl    %edi,%edx
7  subl    %esi,%edi
8  addl    %esi,%ecx
9  cmpl    $255,%ecx
10 jle .L5
```

Register	Variable	Initially
%eax	a	a
%ebx	b	b
%ecx	i	0
%edx	result	a

Exercise's Solution

```
1  # a in %edi, b in %esi
2  xorl    %ecx,%ecx    # i = 0
3  movl    %edi,%edx    # result = 0
4  .p2align 4,,7
5  .L5:                # loop:
6  addl    %edi,%edx    # result += a
7  subl    %esi,%edi    # a -= b
8  addl    %esi,%ecx    # i += b
9  cmpl    $255,%ecx    # Compare i:255
10 jle .L5             # If <= goto loop
11 movl    %edx,%eax    # Set result as return value
```

Note the optimization done by the compiler!

For Loops

```
for (init-expr; test-expr; update-expr)  
    body-statement
```



```
init-expr;  
while (test-expr) {  
    body-statement  
    update-expr;  
}
```



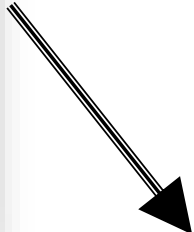
```
    init-expr;  
    t = test-expr;  
    if (!t)  
        goto done;  
loop:  
    body-statement  
    update-expr;  
    t = test-expr;  
    if (t)  
        goto loop;  
done:
```



```
init-expr;  
if (!test-expr)  
    goto done;  
do {  
    body-statement  
    update-expr;  
} while (test-expr);  
done:
```

Exercise

```
1  # x in %rdi, y in %rsi, n in %rdx
2  xorq    %rax,%rax
3  decq    %rdx
4  js      .L4
5  imulq   %rdi,%rsi
6  .L6:
7  addq    %rsi,%rax
8  subq    %rdi,%rdx
9  jns     .L6
10 .L4:
```



```
1 long loop(long x, long y, long n)
2 {
3     long result = 0;
4     long i;
5     for (i = ____ ; i ____ ; i = ____) {
6         result += ____ ;
7     }
8     return result;
9 }
```

Exercise's Solution

```
1 long loop(long x, long y, long n)
2 {
3     long result = 0;
4     long i;
5     for (i = n - 1 ; i >= 0 ; i = i - x) {
6         result += y * x ;
7     }
8     return result;
9 }
```

Note the optimization done by the compiler!

Switch Statements in C

```
long switch_eq(long x)
{
    long result = x;
    switch(x) {
```

```
    case 100:
        result *= 13;
        break;
```

```
    case 102:
        result += 10;
        /* Fall through */
```

```
    case 103:
        result += 11;
        break;
```

```
        case 104:
        case 106:
            result *= result;
            break;
```

```
        default:
            result = 0;
```

```
    }
    return result;
}
```



A diagram consisting of a vertical line on the left and a horizontal line at the top, connected by a downward-pointing arrow. This arrow points from the closing brace of the switch statement to the start of the case 104: case 106: block, illustrating a jump in the code execution flow.

Switch Statements in Assembly

Building the jump table:

```
.section .rodata
.align 8 # Align address to multiple of 8
.L10:
.quad .L4 # Case 100: loc_A
.quad .L9 # Case 101: loc_def
.quad .L5 # Case 102: loc_B
.quad .L6 # Case 103: loc_C
.quad .L8 # Case 104: loc_D
.quad .L9 # Case 105: loc_def
.quad .L8 # Case 106: loc_D
```


Switch Statements in Assembly

Set up the jump table access

leaq -100(%rdi),%rsi # Compute xi = x-100

cmpq \$6,%rsi # Compare xi:6

ja .L9 # if >, goto **default-case**

jmp *.L10(%rsi,8) # Goto jt[xi]

Case 100

.L4: # loc_A:

leaq (%rdi,%rdi,2),%rax # Compute 3*result

leaq (%rdi,%rax,4),%rdi # Compute x+4*3*result

jmp .L3 # Goto **done**

Case 102

.L5: # loc_B:

addq \$10,%rdi # result += 10, Fall through:

Case 103

.L6: # loc_C:

addq \$11,%rdi # result += 11

jmp .L3 # Goto **done**

Switch Statements in Assembly

Cases 104,106

```
.L8:          # loc_D:
    imulq %rdi,%rdi    # result *= result
    jmp .L3           # Goto done
```

Default case

```
.L9:          # loc_D:
    xorq %rdi,%rdi     # result = 0
```

Return result

```
.L3:          # done:
    movq %rdi,%rax     # Set result as return value
```

Exercise

```
1 long switch2(long x)
2 {
3     long result = 0;
4     switch(x) {
5         /* Body of switch statement omitted */
6     }
7     return result;
8 }
```



```
1 # Setting up the jump table access
2 leaq $2(%rdi),%rsi # x in %rdi
3 cmpq $6,%rsi
4 ja .L10
5 jmp *.L11(,%rsi,8)
```

```
1 # Jump table for switch2
2 .L11:
3     .quad .L4
4     .quad .L10
5     .quad .L5
6     .quad .L6
7     .quad .L8
8     .quad .L8
9     .quad .L9
```

Q & A

■ Q:

- What were the values of the case labels in the switch statement body?
- What cases had multiple labels in the C code?

■ A:

- The case labels had values -2,0,1,2,3,4
- The case with the labels 2 and 3