

Information Theory: Problem Set

General instructions:

- Please hand in your homework via Gradescope (entry code: KDPE8G) before 11:59 PM.
- Numbered exercises are taken from the book “Information Theory: From Coding to Learning” by Y. Polyanskiy and Y. Wu, available online at <https://people.lids.mit.edu/yp/homepage/data/itbook-export.pdf>.
- Unless otherwise specified, all logarithms (including those in entropy, mutual information, and KL divergence) are in base e .

Homework 1 (Due on Oct 1, 2025)

Required problems:

R1. I.13

R2. III.19 (Note: in Part (b), $\sqrt{2E_0}$ should be $\sqrt{2E_0/\log e}$.)

R3. (a) Show that $I(X; Y) \geq I(X; Y|U)$ for a Markov chain $U - X - Y$. Conclude that $I(X; Y)$ is concave in P_X for fixed $P_{Y|X}$.

(b) Show that $I(X; Y) \leq I(X; Y|U)$ if X and U are independent. Conclude that $I(X; Y)$ is convex in $P_{Y|X}$ for fixed P_X .

R4. Prove Tao’s inequality: for random variables X, Y, Z with $X \in [-1, 1]$ almost surely,

$$\mathbb{E}|\mathbb{E}[X|Y] - \mathbb{E}[X|Y, Z]| \leq \sqrt{2I(X; Z|Y)}.$$

Optional problems (solve three of them):

O1. I.49 (Note: the claimed limit $1/\sqrt{1-\tau}$ is incorrect and should be replaced by

$$\frac{e^{-\tau/2-\tau^2/4}}{\sqrt{1-\tau}} - 1. \quad)$$

O2. I.51

O3. I.53

O4. I.59 (Note: there is an obvious typo in Part (d): $\varepsilon^{\lambda-1}$ should be $\varepsilon^{1-\lambda}$)

O5. I.63

O6. III.28

O7. *Shearer for sums.* Let X, Y, Z be independent random integers. Prove that

$$2H(X + Y + Z) \leq H(X + Y) + H(X + Z) + H(Y + Z).$$

O8. *Pinning lemma.* Let (X_1, \dots, X_n) be $\{\pm 1\}^n$ -valued random vector. For $2 \leq k \leq n$, let S be a uniformly random subset of $[n]$ of size k , and $i, j \in S$ be two uniformly random draws from S without replacement. Define the quantity

$$f_k = \mathbb{E}[I(X_i; X_j | X_{S \setminus \{i, j\}})].$$

(a) Prove that $\sum_{k=2}^n f_k \leq \log 2$.

(b) Deduce that for $m \geq 0$, there exists a subset $T \subseteq [n]$ with $|T| \leq m$ such that

$$\mathbb{E}[\text{Cov}(X_i, X_j | X_T)^2] \leq \frac{2 \log 2}{m + 1}.$$

Here the expectation is taken over the randomness in both the uniformly random pair $(i, j) \in \binom{[n]}{2}$ and X_T .

O9. *Coin weighing.* There is an unknown subset $X \subseteq [n]$. You must choose in advance k subsets $S_1, \dots, S_k \subseteq [n]$, and receive the cardinalities $|X \cap S_i|$ for all $i \in [k]$. We wish to determine the smallest number k needed to recover the unknown subset X .

(a) Prove that if $k \leq 1.99n / \log_2 n$ and n is sufficiently large, any strategy cannot guarantee the recovery of every subset $X \subseteq [n]$.

(b) Propose a successful strategy if $k \geq 3.17n / \log_2 n$ and n is sufficiently large (here the constant is chosen such that $3.17 > 2 \log_2 3$).

(c) (*challenging and not graded*) Prove (b) with 3.17 replaced by 2.01.

O10. *Information bound on variance.* Let X_1, \dots, X_n be i.i.d., and $(\phi_t)_{t \in \mathcal{T}}$ be a collection of functions $\phi_t : \mathcal{X} \rightarrow [0, 1]$. For $t \in \mathcal{T}$, let $\sigma^2(\phi_t) = \text{Var}(\phi_t(X_1))$ be the true variance, and

$$s_n^2(\phi_t) = \frac{1}{n} \sum_{i=1}^n \phi_t(X_i)^2 - \left(\frac{1}{n} \sum_{i=1}^n \phi_t(X_i) \right)^2$$

be the sample variance. Show that for any $C > 0$ and random index T , it holds that

$$\mathbb{E} \left[\frac{s_n^2(\phi_T)}{\max\{C, \sigma^2(\phi_T)\}} \right] \leq \frac{I(T; X^n)}{nC} + 2.$$

(Hint: use $\mathbb{E}[e^X] \leq \mathbb{E}[1 + 2X] \leq e^{2\mathbb{E}[X]}$ for $X \in [0, 1]$.)

Homework 2 (Due on Nov 1, 2025)

Required problems:

R1. VI.8

R2. VI.14, Part (a) - (c)

R3. Suppose $X_1, \dots, X_n \sim \text{Bern}(p)$ with unknown $p \in [0, 1]$. Using the two-point method, argue that if there is an estimator T such that

$$\sup_{p \in [0,1]} \mathbb{P}_p(|T(X) - p| > \varepsilon) \leq \delta$$

with $\varepsilon, \delta \in (0, 1/8)$, then

$$n \geq c \cdot \frac{\log(1/\delta)}{\varepsilon^2}$$

for a universal constant $c > 0$. (*Hint: $1 - \text{TV} \geq \frac{1}{2} \exp(-\text{KL})$.*)

R4. Let X_1, \dots, X_n be i.i.d. drawn from a discrete distribution $P = (p_1, \dots, p_k)$, the learner aims to estimate the entropy $H(P) = -\sum_{i=1}^k p_i \log p_i$. With a slight abuse of notation, we also use the same letter P to denote the free parameter (p_1, \dots, p_{k-1}) , which belongs to the parameter set $\mathcal{P}_k = \{(p_1, \dots, p_{k-1}) \in \mathbb{R}_+^{k-1} : \sum_{i=1}^{k-1} p_i \leq 1\}$ with a non-empty interior in \mathbb{R}^{k-1} .

(a) For a fixed P in the interior of \mathcal{P}_k , find the expression of the Fisher information $I(P)$ and the inverse Fisher information $I(P)^{-1}$ in the above model with $n = 1$. (*Hint: for $I(P)^{-1}$, use Woodbury matrix identity.*)

(b) Use the local asymptotic minimax theorem to show that for any P_0 in the interior of \mathcal{P}_k and any sequence of estimators \hat{H}_n based on n samples, it holds that

$$\lim_{C \rightarrow \infty} \liminf_{n \rightarrow \infty} n \cdot \sup_{P \in \mathcal{P}_k : \|P - P_0\|_2 \leq C/\sqrt{n}} \mathbb{E}_P[(\hat{H}_n - H(P))^2] \geq \text{Var}_{X \sim P_0}(\log P_0(X)).$$

(c) Find a suitable P_0 in (b) to conclude that

$$\liminf_{n \rightarrow \infty} n \cdot \inf_{\hat{H}_n} \sup_{P \in \mathcal{P}_k} \mathbb{E}_P[(\hat{H}_n - H(P))^2] \geq c \cdot \log^2 k,$$

where $c > 0$ is a universal constant.

Optional problems (solve three of them):

O1. I.65 (In Part (c), n should be d)

O2. I.66

O3. VI.14, Part (d)

O4. VI.16

O5. VI.17

O6. *Monotonic CLT.* Let X_1, X_2, \dots be i.i.d. random variables with $\mathbb{E}[X_1] = 0$, $\text{Var}(X_1) = 1$, and $h(X_1) > -\infty$. Let $S_n = \sum_{i=1}^n X_i$ and $T_n = \frac{1}{\sqrt{n}} S_n$.

- (a) Let $S_{\sim i} = S_n - X_i$, and ρ_i be the score function of $S_{\sim i}$, where we recall that the score function for a random variable X with density f is $\rho(x) = (\log f)' = \frac{f'(x)}{f(x)}$. Show that the score function ρ of S_n is $\rho(S_n) = \mathbb{E}[\rho_i(S_{\sim i})|S_n]$.
- (b) Prove the following lemma: for independent Z_1, \dots, Z_n and functions f_1, \dots, f_n such that f_i depends only on $Z_{\sim i}$ and $\mathbb{E}[f_i(Z_{\sim i})] = 0$, it holds that

$$\mathbb{E} \left[\left(\sum_{i=1}^n f_i(Z_{\sim i}) \right)^2 \right] \leq (n-1) \sum_{i=1}^n \mathbb{E}[f_i(Z_{\sim i})^2].$$

(Hint: ANOVA decomposition.)

- (c) Show that if $J(X_1) < \infty$, the Fisher information satisfies $J(T_n) \leq J(T_{n-1})$.
- (d) Show that the differential entropy satisfies $h(T_n) \geq h(T_{n-1})$.
- O7. *Bernoulli EPI: Mrs. Gerber's Lemma.* Let $h_2(p) = -p \log p - (1-p) \log(1-p)$ be the binary entropy function, and $h_2^{-1} : [0, \log 2] \rightarrow [0, \frac{1}{2}]$ be its inverse.
- (a) Show that for any fixed $p \in [0, 1]$, the function $v \mapsto h_2(h_2^{-1}(v) * p)$ is convex, where $p * q = p(1-q) + (1-p)q$ denotes the convolution.
- (b) Use (a) to show that for any (X, U) with $X \in \{0, 1\}$ and $Y = X \oplus \text{Bern}(p)$,

$$H(Y|U) \geq h_2(h_2^{-1}(H(X|U)) * p).$$

- (c) Use (b) to show that for any $X^n \in \{\pm 1\}^n$ and $Y^n = X^n \oplus \text{Bern}(p)^{\otimes n}$,

$$\frac{H(Y^n)}{n} \geq h_2 \left(h_2^{-1} \left(\frac{H(X^n)}{n} \right) * p \right).$$

O8. *Tree-based lower bound.* This problem proves another lower bound for the test error of testing multiple hypotheses. Let $T = ([m], E)$ be an undirected graph with vertex set $[m]$ and edge set E , and be a tree in the sense that T is both connected and acyclic.

- (a) Show that for any real numbers x_1, \dots, x_m , it holds that

$$\sum_{i=1}^m x_i - \max_{i \in [m]} x_i \geq \sum_{(i,j) \in E} \min\{x_i, x_j\}.$$

(b) Use the result in Part (a), show that for probability distributions P_1, \dots, P_m ,

$$\min_{\Psi} \frac{1}{m} \sum_{i=1}^m P_i(\Psi \neq i) \geq \frac{1}{m} \sum_{(i,j) \in E} (1 - \text{TV}(P_i, P_j)),$$

where the minimum is over all possible tests $\Psi : \mathcal{X} \rightarrow [m]$.

(c) Evaluate the terms on both sides of (b) under $P_i = \mathcal{N}(i\Delta, 1)$ and a line tree with edge set $E = \{(1, 2), (2, 3), \dots, (m-1, m)\}$, and show that they are equal.

O9. *VC class with small oracle risk.* We have a function class \mathcal{F} with VC dimension d , and n training data $(x_1, y_1), \dots, (x_n, y_n)$ drawn from an unknown joint distribution P_{XY} , with $\mathcal{Y} = \{0, 1\}$. Define the following class $\mathcal{P}(\mathcal{F}, \varepsilon)$ of joint distributions where the best classifier has an error at most ε :

$$\mathcal{P}(\mathcal{F}, \varepsilon) = \left\{ P_{XY} : \inf_{f^* \in \mathcal{F}} P_{XY}(Y \neq f^*(X)) \leq \varepsilon \right\}.$$

So $\varepsilon = 0$ corresponds to the well-specified case, and $\varepsilon = 1$ corresponds to the misspecified case. Define the minimax excess risk $R^*(\mathcal{F}, \varepsilon)$ over $\mathcal{P}(\mathcal{F}, \varepsilon)$ as

$$R^*(\mathcal{F}, \varepsilon) = \inf_{\hat{f}} \sup_{P_{XY} \in \mathcal{P}(\mathcal{F}, \varepsilon)} \mathbb{E} \left[P_{XY}(Y \neq \hat{f}(X)) - \inf_{f^* \in \mathcal{F}} P_{XY}(Y \neq f^*(X)) \right].$$

Show that for all $\varepsilon \in [0, 1]$,

$$R^*(\mathcal{F}, \varepsilon) = \Omega \left(\min \left\{ \sqrt{\frac{d}{n}} \cdot \varepsilon + \frac{d}{n}, 1 \right\} \right).$$

O10. *Bias-variance analysis for orthogonal polynomials.* The concept of orthogonal polynomials is useful not only in proving lower bounds, but also in constructing and analyzing unbiased estimators. Let $(P_\theta)_{\theta \in [\theta_0 - \varepsilon, \theta_0 + \varepsilon]}$ be a one-dimensional family of probability distributions with the following local expansion:

$$\frac{P_{\theta_0+u}(x)}{P_{\theta_0}(x)} = \sum_{m=0}^{\infty} p_m(x; \theta_0) \frac{u^m}{m!}, \quad \forall |u| \leq \varepsilon, x \in \mathcal{X}.$$

In addition, assume that the quantity $\sum_{x \in \mathcal{X}} P_{\theta_0+u}(x) P_{\theta_0+v}(x) / P_{\theta_0}(x)$ depends only on θ_0 and uv , for all $u, v \in [-\varepsilon, \varepsilon]$. In class we showed that $\{p_m(x; \theta_0)\}_{m \geq 0}$ are orthogonal in $L^2(P_{\theta_0})$, i.e.,

$$\mathbb{E}_{X \sim P_{\theta_0}} [p_m(X; \theta_0) p_n(X; \theta_0)] = A_m(\theta_0) \cdot \mathbb{1}(m = n)$$

for some constants $\{A_m(\theta_0)\}_{m \geq 0}$.

(a) Show that for $X \sim P_{\theta_0+u}$ with $u \in [-\varepsilon, \varepsilon]$,

$$\mathbb{E}_{X \sim P_{\theta_0+u}} [p_m(X; \theta_0)] = c_m u^m$$

for some constant c_m independent of u . Find the expression of c_m using $A_m(\theta_0)$.

- (b) Suppose the same local expansion and orthogonality condition also hold under $\theta_0 + u$, and

$$p_m(x; \theta_0) = \sum_{\ell=0}^m b(m, \ell, \theta_0, u) \cdot p_\ell(x; \theta_0 + u), \quad \forall |u| \leq \varepsilon, x \in \mathcal{X},$$

Show that

$$\mathbb{E}_{X \sim P_{\theta_0+u}}[p_m(X; \theta_0)^2] = \sum_{\ell=0}^m b(m, \ell, \theta_0, u)^2 \cdot A_\ell(\theta_0 + u).$$

- (c) Show that in the Poisson model $\mathcal{X} = \mathbb{N}$, $P_\theta = \text{Poi}(\theta)$,

$$b(m, \ell, \theta_0, u) = \binom{m}{\ell} \frac{(\theta_0 + u)^\ell u^{m-\ell}}{\theta_0^m}.$$

Homework 3 (Due on Dec 1, 2025)

Required problems:

R1. II.20

R2. VI.25

- R3. Let $A \in \mathbb{R}^{m \times n}$ be a random matrix with i.i.d. $\mathcal{N}(0, 1)$ entries. The operator norm of A is defined as $\|A\|_{\text{op}} = \max_{v \in S^{n-1}} \|Av\|_2$, where S^{n-1} denotes the unit sphere in \mathbb{R}^n .

- (a) Let $\mathcal{U} = \{u_1, \dots, u_M\}$ and $\mathcal{V} = \{v_1, \dots, v_N\}$ be an ε -net (under the ℓ_2 norm) for S^{m-1} and S^{n-1} respectively. Show that for $\varepsilon < 1/2$,

$$\|A\|_{\text{op}} \leq \frac{1}{1 - 2\varepsilon} \max_{u \in \mathcal{U}, v \in \mathcal{V}} u^\top Av.$$

- (b) Deduce from (a) that $\mathbb{E}\|A\|_{\text{op}} \lesssim \sqrt{m} + \sqrt{n}$.

- (c) Show a matching lower bound $\mathbb{E}\|A\|_{\text{op}} \gtrsim \sqrt{m} + \sqrt{n}$ using Sudakov minoration.

- R4. Recall that $M(A, d, \varepsilon)$ denotes the maximum number m of points x_1, \dots, x_m such that $d(x_i, x_j) \geq \varepsilon$ for every $i \neq j \in [m]$.

- (a) Let A be the set of all non-decreasing functions $f : [0, 1] \rightarrow [0, 1]$. Show that for $\varepsilon \in [0, 1]$, there exist universal constants $c_1, c_2 > 0$ such that

$$\log M(A, L_2([0, 1]), c_1 \varepsilon) \geq \frac{c_2}{\varepsilon}.$$

- (b) Now let A be the set of all convex functions $f : [0, 1] \rightarrow [0, 1]$. Show that for $\varepsilon \in [0, 1]$, there exist universal constants $c_1, c_2 > 0$ such that

$$\log M(A, L_2([0, 1]), c_1 \varepsilon) \geq \frac{c_2}{\sqrt{\varepsilon}}.$$

Hint: you may try to break into several small intervals, find two possible function constructions in each interval, and concatenate them. Use the Gilbert–Varshamov bound in class.

Optional problems (solve three of them):

O1. VI.19

O2. VI.24

O3. VI.26 (*A typo in the problem: the inequality should be*

$$\frac{1}{n - |T| - 1} \sum_{i \neq j \in T^c} D(\nu_{X_i, X_j} \| \mu_{X_i, X_j}^{(\sigma_T)}) \geq \left(2 - \frac{c}{n - |T| - 1}\right) \sum_{i \in T^c} D(\nu_{X_i} \| \mu_{X_i}^{(\sigma_T)}).$$

Additional hint: for any function $h : 2^{[n]} \rightarrow \mathbb{R}$ and $S \sim \text{Bern}(\tau)^{\otimes n}$, show that

$$\begin{aligned} \frac{d}{d\tau} \mathbb{E}[h(S)] &= \sum_{i=1}^n \mathbb{E}[h(S \cup \{i\}) - h(S)], \\ \frac{d^2}{d\tau^2} \mathbb{E}[h(S)] &= \sum_{i \neq j} \mathbb{E}[h(S \cup \{i, j\}) - h(S \cup \{i\}) - h(S \cup \{j\}) + h(S)]. \end{aligned}$$

O4. *Redundancy bound with general Beta mixing.* Let $P_\theta = \text{Bern}(\theta)$, and $\theta \sim \text{Beta}(\alpha, \beta)$ follow a Beta distribution. For $x^n \in \{0, 1\}^n$, let $Q(x^n) = \mathbb{E}_\theta[P_\theta^{\otimes n}(x^n)]$, and n_0, n_1 be the number of 0's and 1's in x^n (with $n_0 + n_1 = n$).

(a) Let $B(\alpha, \beta)$ be the Beta function. Show that

$$\max_{\theta \in [0, 1]} \frac{P_\theta^{\otimes n}(x^n)}{Q(x^n)} = \frac{\frac{n_0^{n_0} n_1^{n_1}}{n^n}}{\frac{B(\alpha + n_0, \beta + n_1)}{B(\alpha, \beta)}}.$$

(b) By Stirling's approximation $1 \leq \Gamma(z)/(\sqrt{2\pi} z^{z-1/2} e^{-z}) \leq e^{\frac{1}{12z}}$ for $z > 0$, show that

$$\max_{\theta \in [0, 1]} D_{\text{KL}}(P_\theta^{\otimes n} \| Q) \leq \max_{\theta \in [0, 1]} \max_{x^n} \log \frac{P_\theta^{\otimes n}(x^n)}{Q(x^n)} \leq \max \left\{ \frac{1}{2}, \alpha, \beta \right\} \log n + O_{\alpha, \beta}(1).$$

In other words, there is a range of parameters for the Beta prior that can achieve the optimal regret up to first order.

O5. *Last-iterate convergence for Hellinger.* For probability distributions P_1, \dots, P_m on the same space \mathcal{X} , let $Q_{X^n} = \frac{1}{m} \sum_{i=1}^m P_i^{\otimes n}$ be the Yang–Barron type mixture with a uniform prior. Show that for $X_1, \dots, X_{n-1} \sim P_1$ and a universal constant C ,

$$\mathbb{E}[H^2(P_1, Q_{X_n|X^{n-1}})] \leq \frac{C \log m}{n}.$$

(*Hint: Express $Q_{X_n|X^{n-1}} = \sum_i w_i P_i$ and upper bound $\mathbb{E}[w_i]$ based on $H^2(P_1, P_i)$.)*

O6. *Jeffreys prior.*

(a) I.57, Part (a)

(b) Let π be a prior on $\Theta \subseteq \mathbb{R}^d$ with density $\pi(\theta)$, and $\theta_0 \in \text{int}(\Theta)$. Let $(P_\theta)_{\theta \in \Theta}$ be a family of distributions, with Fisher information matrix $I(\theta)$ at $\theta = \theta_0$. Show that for $Q = \mathbb{E}_{\theta \sim \pi}[P_\theta^{\otimes n}]$, by Laplace's method and the local behavior of KL divergence by Fisher information, one has the approximate upper bound

$$D_{\text{KL}}(P_{\theta_0}^{\otimes n} \| Q) \leq \frac{d}{2} \log \frac{n}{2\pi} - \log \frac{1}{\pi(\theta_0)} + \frac{1}{2} \log \det I(\theta_0) + O(1).$$

Therefore, choosing $\pi(\theta) \propto (\log \det I(\theta))^{-1/2}$ (known as *the Jeffreys prior*) achieves a redundancy upper bound $\frac{d}{2} \log n + O(1)$ assuming regularity conditions.

O7. *Redundancy of uniform family.* Let $\mathcal{P} = \{\text{Unif}(0, \theta) : \theta \in [\frac{1}{2}, 1]\}$. Show that $\text{Red}(\mathcal{P}^{\otimes n}) \sim \log n$ by proving redundancy upper and lower bounds. Explain the difference from the usual relation $\text{Red}(\mathcal{P}^{\otimes n}) \sim \frac{1}{2} \log n$.

O8. *Branching number.* For a countable rooted tree T where each vertex has a finite degree, a *flow* is a function $f : V(T) \rightarrow \mathbb{R}_+$ such that $f_u = \sum_{v \in \text{children}(u)} f_v$ for all vertices u . The *branching number* $\text{br}(T)$ is defined as the supremum of $\lambda \in \mathbb{R}$ such that there exists a flow f with $f_u > 0$ for some u , and $f_u \leq \lambda^{-d(u)}$ for all vertices u , where $d(u)$ denotes the depth of u .

(a) Show that for broadcasting on tree T , if each edge represents a channel $P_{Y|X}$ with $\eta_{\text{KL}}(P_{Y|X}) \leq \eta$, then the model has non-reconstruction when $\text{br}(T)\eta < 1$.

(b) For $p \in [0, 1]$, let T_p be the connected component containing the root in a random graph where each edge of T is removed independently with probability $1 - p$. Let $p_c = p_c(T) \in [0, 1]$ be the critical (percolation) probability:

$$p_c = \sup\{p \in [0, 1] : \mathbb{P}(T_p \text{ has infinitely many vertices}) = 0\}.$$

Show that $\text{br}(T) \leq p_c^{-1}$. (*Hint: for $p > \text{br}(T)^{-1}$, construct a flow $\{f_u\}$ and define $M_n = \sum_{u \in V(T_p) : \text{depth}(u)=n} f_u p^{-n}$. Show that $\{M_n\}$ is a martingale, and that $\sup_n \mathbb{E}[M_n^2] < \infty$. Therefore $M_n \rightarrow M_\infty$ in L^1 .)*

(c) Show that $\text{br}(T) \geq p_c^{-1}$. (*Hint: show that if $p < \text{br}(T)^{-1}$, find a sequence of cuts $\{C_n\}$ of T such that*

$$\lim_{n \rightarrow \infty} \mathbb{E}[\text{number of edges in } T_p \text{ crossed by the cut } C_n] = 0;$$

the max-flow min-cut theorem might be useful.)

(d) Conclude that if the offspring distribution of a branching process has mean $m > 1$, then given the event that this process does not become extinct, the corresponding Galton–Watson tree T has branching number m almost surely. (*Hint: recall that the extinction probability in a branching process with $m > 1$ is the unique solution in $[0, 1)$ to the equation $x = \sum_{i=0}^{\infty} p_i x^i$. How about the probability that T has branching number at least λ , for different choices of λ ?)*

O9. F_I curve. For a joint distribution P_{XY} and $t > 0$, define

$$F_I(P_{XY}, t) = \sup\{I(U; Y) : I(U; X) \leq t, U - X - Y \text{ forms a Markov chain}\}.$$

- (a) Show that $t \mapsto F_I(P_{XY}, t)$ is concave, and $\frac{d}{dt}F_I(P_{XY}, t)|_{t=0} = \eta_{\text{KL}}(P_X, P_{Y|X})$.
- (b) Using EPI and Bernoulli EPI (Mrs. Gerber's lemma in HW2 O7), find the expression of $F_I(P_{XY}, t)$ in the following two scenarios:
 - i. (X, Y) is zero-mean and jointly Gaussian with correlation $\rho \in [-1, 1]$;
 - ii. (X, Y) is zero-mean and jointly Bernoulli with correlation $\rho \in [-1, 1]$.
- (c) Conclude that in both scenarios, the maximal correlation between X and Y is $|\rho|$.

O10. *SDPI for Fisher information.* Let $(P_\theta)_{\theta \in \Theta \subseteq \mathbb{R}^n}$ be a family of distributions, with score function $s_\theta(\cdot)$ and Fisher information matrix $I(\theta)$.

- (a) Show that for $\theta - X - Y$, the score function for Y is $s_\theta^Y(y) = \mathbb{E}[s_\theta(X)|Y = y]$.
- (b) If P_θ is a discrete pmf $(\frac{1}{2n} + \theta_1, \frac{1}{2n} - \theta_1, \dots, \frac{1}{2n} + \theta_n, \frac{1}{2n} - \theta_n)$, show that

$$\sup\{\text{trace}(I^Y(0)) : \theta - X - Y, |\mathcal{Y}| \leq \ell\} \asymp \min\{n(\ell - 1), n^2\},$$

where $I^Y(\theta)$ denotes the Fisher information matrix for Y , and $\ell \in \mathbb{N}$.

- (c) If $P_\theta = \mathcal{N}(\theta, I_n)$, show that

$$\sup\{\text{trace}(I^Y(0)) : \theta - X - Y, |\mathcal{Y}| \leq \ell\} \asymp \min\{\log \ell, n\},$$