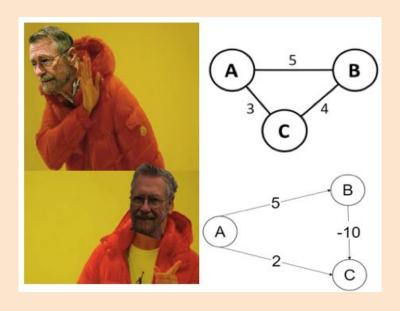
# Bellman-Ford Algorithm

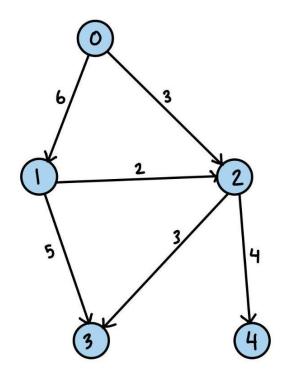


Presented by Group 4: Zadie, Yan, Peng

## Algorithm Introduction

- Shortest path algorithm
- Performs shortest path traversal on a source node from a weighted digraph
- Edges can have negative weights and the algorithm can detect negative weight cycles
- Real life application: Negative weights can exist in an transportation graph or a discount or rebate is offered for a specific flight or route

Review



Node	MinDist	lastStep
0	0	[0,0]
1	6	[0, 1]
2	3	[0,2]
3	6	[2,3]
4	7	(2,4)

Dijkstra's Algorithm

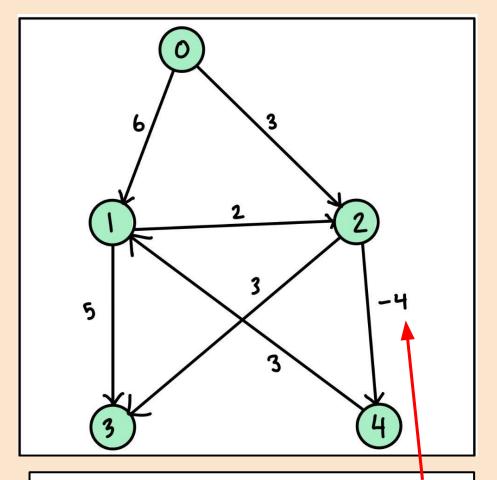
### BMF:

- 1) Initalize all weights to  $\infty$
- 2) Source = 0
- 3) Relax all edges Inner loop runs Etimes
- 4) Outer 100p runs V-1 times
- 5) Negative Weight check

Algorithm should detect Shortest path in V-1 100ps

Otherwise:

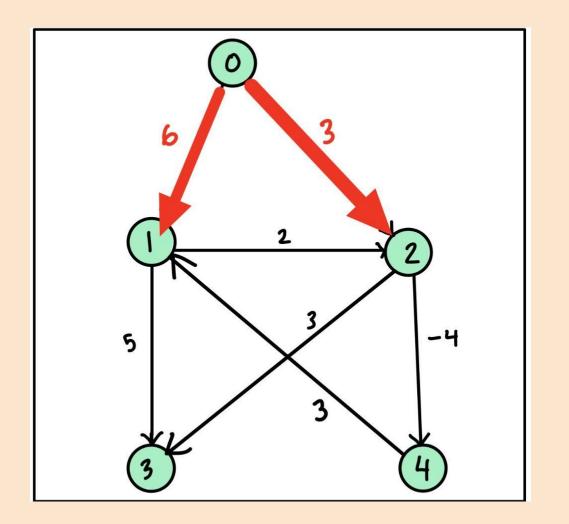
A negative cycle exists



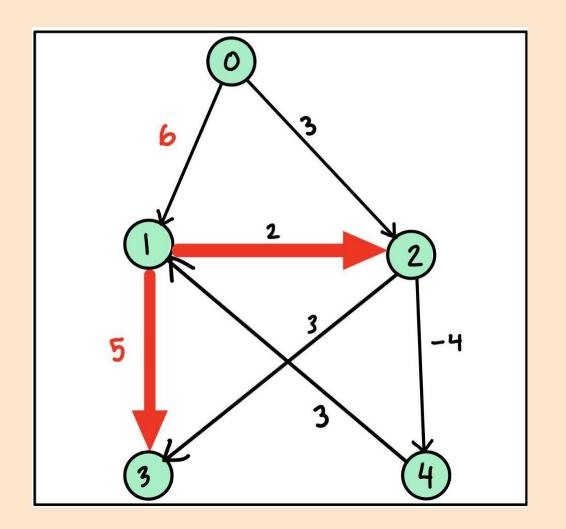
**New: Handles negative edge weights** 

### Initialization

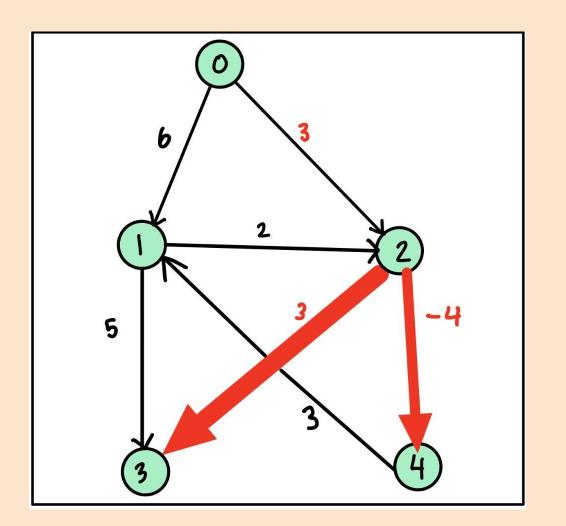
Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	∞
[0,2] [1,2]	2	∞
[1,3] [2,3]	3	∞
[2,4]	4	∞



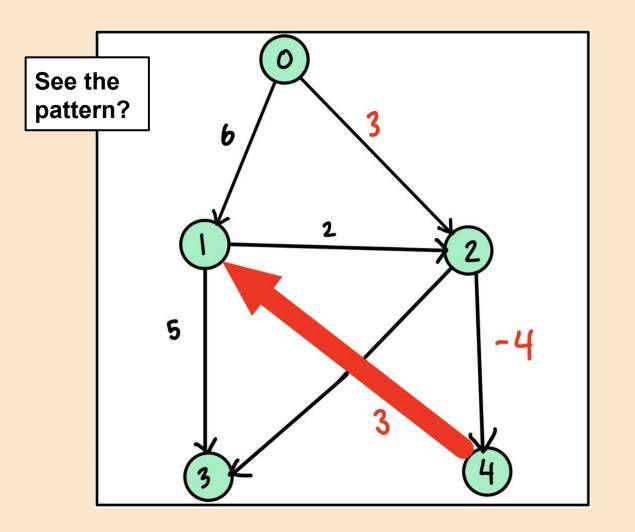
Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	6
[0,2] [1,2]	2	3
[1,3] [2,3]	3	∞
[2,4]	4	∞



Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	6
[0,2] [1,2]	2	3
[1,3] [2,3]	3	11
[2,4]	4	∞



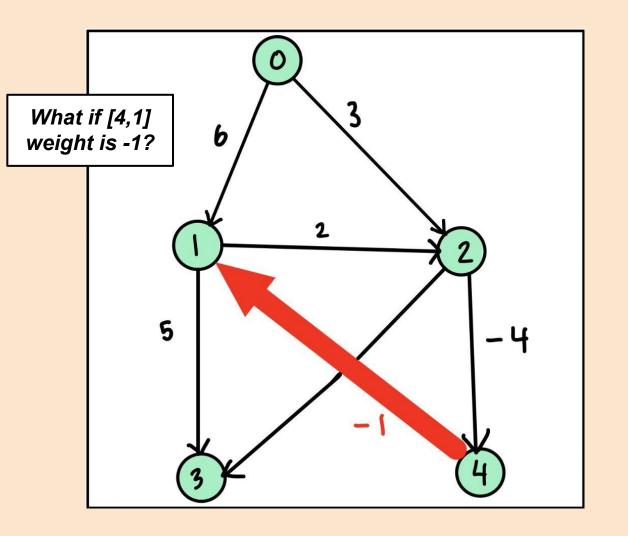
Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	6
[0,2] [1,2]	2	3
[1,3] [2,3]	3	6
[2,4]	4	-1



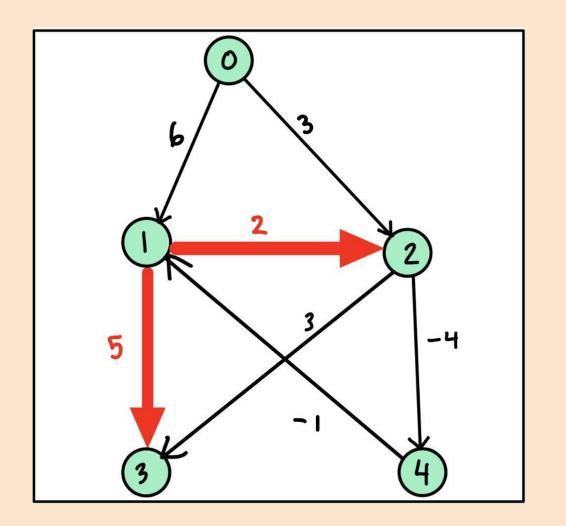
Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	2
[0,2] [1,2]	2	3
[1,3] [2,3]	3	6
[2,4]	4	-1

### Now let's look at a <u>bad</u> implementation...



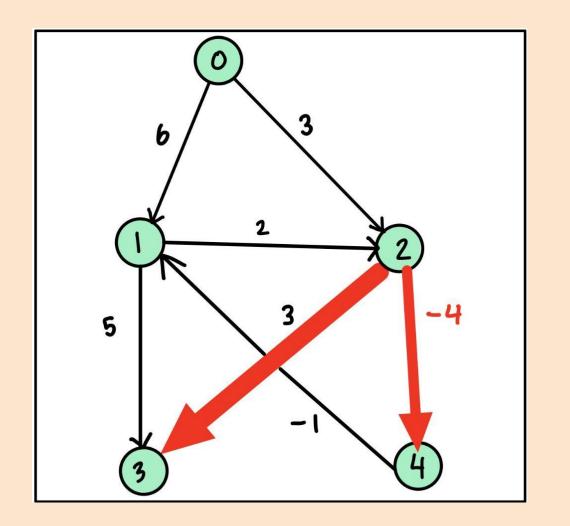


Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	-2
[0,2] [1,2]	2	3
[1,3] [2,3]	3	6
[2,4]	4	-1



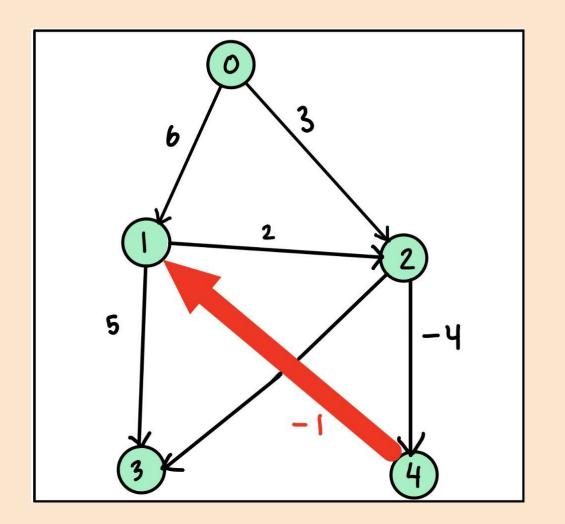
### 2nd Iteration

Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	-2
[0,2] [1,2]	2	0
[1,3] [2,3]	3	3
[2,4]	4	-1



### 2nd Iteration

Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	-2
[0,2] [1,2]	2	0
[1,3] [2,3]	3	3
[2,4]	4	-4



### 2nd Iteration

Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	-5
[0,2] [1,2]	2	0
[1,3] [2,3]	3	3
[2,4]	4	-4

# to infinity negative weight cycle

### **3rd Iteration**

Edges	Node	Min Dist
	0	0
[0,1] [4,1]	1	-5
[0,2] [1,2]	2	-3
[1,3] [2,3]	3	0
[2,4]	4	-4

### Bellman-Ford Pseudocode:

```
function BellmanFord(vertices, edges, source):
// Step 1: Initialize distances from source to all other vertices as infinity
dist = \{\}
for each vertex in vertices:
   dist[vertex] = infinity
dist[source] = 0
// Step 2: Relax edges repeatedly
for i from 1 to |vertices|-1:
   for each edge (u, v, w) in edges:
      if dist[u] + w < dist[v]:
        dist[v] = dist[u] + w
// Step 3: Check for negative-weight cycles
for each edge (u, v, w) in edges:
   if dist[u] + w < dist[v]:
      throw "Graph contains a negative-weight cycle"
return dist
```

# Time and Space Complexity

Worst case time complexity: O (V³)

Average case time complexity: O(|V| \* |E|)

Best case time complexity: O(E)

Space Complexity: O(V + E) for all cases

# Thanks for learning with us!



### References

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- 2) R. Bellman, "On a routing problem," Quarterly of Applied Mathematics, vol. 16, no. 1, pp. 87–90, 1958.
- 3) L. R. Ford Jr., "Network Flow Theory," Paper P-923, Santa Monica, CA, USA: RAND Corporation, 1956.
- 4) E. F. Moore, "The shortest path through a maze," in Proc. Internat. Sympos. Switching Theory 1957, Part II, Cambridge, MA, USA: Harvard Univ. Press, 1959, pp. 285-292.
- 5) J. Y. Yen, "An algorithm for finding shortest routes from all source nodes to a given destination in general networks," Quarterly of Applied Mathematics, vol. 27, no. 4, pp. 526–530, 1970.
- 6) M. J. Bannister and D. Eppstein, "Randomized speedup of the bellman–Ford algorithm," 2012 Proceedings of the Ninth Workshop on Analytic Algorithmics and Combinatorics (ANALCO), 2012.
- 7) "Bellman–Ford algorithm," Wikipedia, 27-Feb-2023. [Online]. Available: https://en.wikipedia.org/wiki/Bellman%E2%80%93Ford\_algorithm. [Accessed: 23-Mar-2023].