Intervals Intersection (medium)

We'll cover the following ^

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Problem Statement

Given two lists of intervals, find the **intersection of these two lists**. Each list consists of **disjoint intervals sorted on their start time**.

Example 1:

```
Input: arr1=[[1, 3], [5, 6], [7, 9]], arr2=[[2, 3], [5, 7]]
Output: [2, 3], [5, 6], [7, 7]
Explanation: The output list contains the common intervals between the two lists.
```

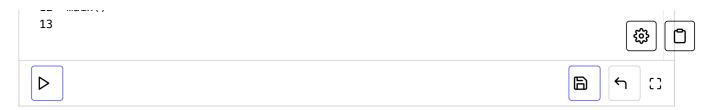
Example 2:

```
Input: arr1=[[1, 3], [5, 7], [9, 12]], arr2=[[5, 10]]
Output: [5, 7], [9, 10]
Explanation: The output list contains the common intervals between the two lists.
```

Try it yourself

Try solving this question here:

```
👙 Java
           Python3
                                    ③ C++
                        JS JS
 1 def merge(intervals_a, intervals_b):
      result = []
 2
 3
      # TODO: Write your code here
 4
      return result
 7 def main():
      print("Intervals Intersection: " + str(merge([[1, 3], [5, 6], [7, 9]], [[2, 3], [5, 7]])
      print("Intervals Intersection: " + str(merge([[1, 3], [5, 7], [9, 12]], [[5, 10]])))
educative
12 main()
```

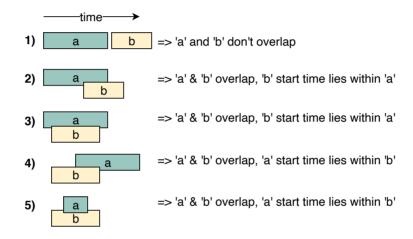


Solution

This problem follows the Merge Intervals

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/565201724243 9680/) pattern. As we have discussed under Insert Interval

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/571831435762 0736/), there are five overlapping possibilities between two intervals 'a' and 'b'. A close observation will tell us that whenever the two intervals overlap, one of the interval's start time lies within the other interval. This rule can help us identify if any two intervals overlap or not.



Now, if we have found that the two intervals overlap, how can we find the overlapped part?

Again from the above diagram, the overlapping interval will be equal to:

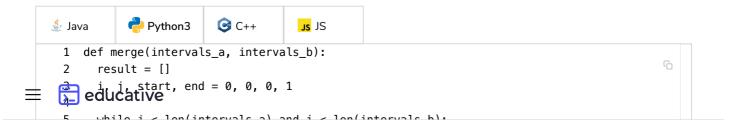
```
start = max(a.start, b.start)
end = min(a.end, b.end)
```

That is, the highest start time and the lowest end time will be the overlapping interval.

So our algorithm will be to iterate through both the lists together to see if any two intervals overlap. If two intervals overlap, we will insert the overlapped part into a result list and move on to the next interval which is finishing early.

Code

Here is what our algorithm will look like:



```
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        # check if intervals overlap and intervals_a[i]'s start time lies within the other
 6
 7
        a_overlaps_b = intervals_a[i][start] >= intervals_b[j][start] and \
 8
                         intervals_a[i][start] <= intervals_b[j][end]</pre>
 9
10
        # check if intervals overlap and intervals_a[j]'s start time lies within the other int
11
        b overlaps a = intervals b[j][start] >= intervals a[i][start] and \
                         intervals_b[j][start] <= intervals_a[i][end]</pre>
12
13
14
        # store the the intersection part
15
        if (a_overlaps_b or b_overlaps_a):
           result.append([max(intervals_a[i][start], intervals_b[j][start]), min(
16
             intervals_a[i][end], intervals_b[j][end])])
17
18
        # move next from the interval which is finishing first
19
20
        if intervals_a[i][end] < intervals_b[j][end]:</pre>
21
           i += 1
22
        else:
23
           j += 1
24
25
       return result
26
27
    def main():
28
\triangleright
                                                                                     []
```

Time complexity

As we are iterating through both the lists once, the time complexity of the above algorithm is O(N+M), where 'N' and 'M' are the total number of intervals in the input arrays respectively.

Space complexity

Ignoring the space needed for the result list, the algorithm runs in constant space O(1).

