



Solution Review: Trace the Complete Path of a Journey

This review provides a detailed analysis of the solution to the Trace the Complete Path of a Journey Challenge.

We'll cover the following

- Solution: A Hash Table to Deduce The Starting Point
 - Time Complexity

Solution: A Hash Table to Deduce The Starting Point

```
def trace_path(my_dict):
 2
        result = []
 3
        # Create a reverse dict of the given dict i.e if the given dict has (N,C)
        # then reverse dict will have (C,N) as key-value pair
        # Traverse original dict and see if it's key exists in reverse dict
 5
 6
        # If it doesn't exist then we found our starting point.
 7
        # After the starting point is found, simply trace the complete path
 8
        # from the original dict.
 9
        reverse_dict = dict()
10
        # To fill reverse dict, iterate through the given dict
11
        keys = my_dict.keys()
12
        for key in keys:
13
            reverse_dict[my_dict.get(key)] = key
14
        # Find the starting point of itinerary
15
        from_loc = None
16
        keys_rev = reverse_dict.keys()
17
        for key in keys:
18
            if key not in reverse_dict:
19
                from_loc = key
20
                break
21
                # Trace complete path
22
        to = my dict.get(from loc)
23
        while to is not None:
24
            result.append([from_loc, to])
25
            from_loc = to
26
            to = my_dict.get(to)
27
        return result
28
\triangleright
                                                                                              []
```

The first thing we need to do is find the starting point of the journey. A reverse_dict is created to switch the sources and destinations in the original map.

The key which does not appear in reverse_dict has never been a destination in map. Hence, it is the starting city.

From here, we simply traverse from city to city based on the previous destination.

Time Complexity





Although a hash table is created and traversed, both take the same amount of time. The complexity for this algorithm is O(n) where **n** is the number of source-destination pairs.

