

No-repeat Substring (hard)

We'll cover the following ^

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Problem Statement

Given a string, find the **length of the longest substring** which has **no repeating characters**.

Example 1:

```
Input: String="aabccbb"
Output: 3
Explanation: The longest substring without any repeating characters is "abc".
```

Example 2:

```
Input: String="abbbb"
Output: 2
Explanation: The longest substring without any repeating characters is "ab".
```

Example 3:

```
Input: String="abccde"
Output: 3
Explanation: Longest substrings without any repeating characters are "abc" & "cde"
.
```

Try it yourself

Try solving this question here:





Solution

This problem follows the **Sliding Window** pattern and we can use a similar dynamic sliding window strategy as discussed in Longest Substring with K Distinct Characters (https://www.educative.io/collection/page/5668639101419520/5671464854355968/569821771281 2032/). We can use a **HashMap** to remember the last index of each character we have processed. Whenever we get a repeating character we will shrink our sliding window to ensure that we always have distinct characters in the sliding window.

Code

Here is what our algorithm will look like:

```
Js JS
👙 Java
           🯓 Python3
                         G C++
 1 def non_repeat_substring(str1):
 2
      window start = 0
 3
      max length = 0
 4
      char_index_map = {}
 5
      # try to extend the range [windowStart, windowEnd]
 6
 7
      for window_end in range(len(str1)):
 8
        right char = str1[window end]
 9
        # if the map already contains the 'right_char', shrink the window from the beginning s
10
        # we have only one occurrence of 'right_char'
11
        if right_char in char_index_map:
          # this is tricky; in the current window, we will not have any 'right_char' after its
12
13
          # and if 'window_start' is already ahead of the last index of 'right_char', we'll ke
          window_start = max(window_start, char_index_map[right_char] + 1)
14
15
        # insert the 'right_char' into the map
        char_index_map[right_char] = window_end
16
17
        # remember the maximum length so far
18
        max_length = max(max_length, window_end - window_start + 1)
19
      return max length
20
21
22 def main():
      print("Length of the longest substring: " + str(non_repeat_substring("aabccbb")))
23
      print("Length of the longest substring: " + str(non_repeat_substring("abbbb")))
24
      print("Length of the longest substring: " + str(non_repeat_substring("abccde")))
25
26
27
28
   main()
\triangleright
                                                                                             :3
```

Time Complexity

The time complexity of the above algorithm will be O(N) where 'N' is the number of characters in the input string.







The space complexity of the algorithm will be O(K) where K is the number of distinct characters in the input string. This also means K <= N, because in the worst case, the whole string might not have any repeating character so the entire string will be added to the **HashMap**. Having said that, since we can expect a fixed set of characters in the input string (e.g., 26 for English letters), we can say that the algorithm runs in fixed space O(1); in this case, we can use a fixed-size array instead of the **HashMap**.

