

Red-Black Tree Insertion

This lesson will cover the insertion operation in Red-Black trees, discussing all the four insertion cases.

We'll cover the following



- Insertion in Red-Black Tree
- Rebalancing the Tree
 - Case 1: Left-Left
 - Case 2: Left-Right
 - Case 3: Right-Right
 - Case 4: Right-Left

Insertion in Red-Black Tree

Here is a high-level description of the algorithm involved in inserting a value in a Red-Black Tree:

1. Insert the given node using the standard BST Insertion technique that we studied earlier and color it **Red**.
2. If the given node is the root, then change its color to **Black**
3. If the given node is not the root, then we will have to perform some operations to make the tree follow the Red-Black property.

Rebalancing the Tree

There are two ways to balance an unbalanced tree:

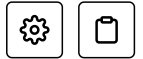
1. Recoloring Nodes
2. Rotating Nodes (left or right)

But before the details are explained, let's define the structure of the Red-Black Tree and some nodes relative to the given node, which is the node that we inserted in the Tree.

- Node C – the newly inserted node
- Node P – the parent of the newly inserted node
- Node G – the grandparent of the newly inserted node
- Node U – the sibling of the parent of the newly inserted node, i.e., the sibling of Node P / child of Node G

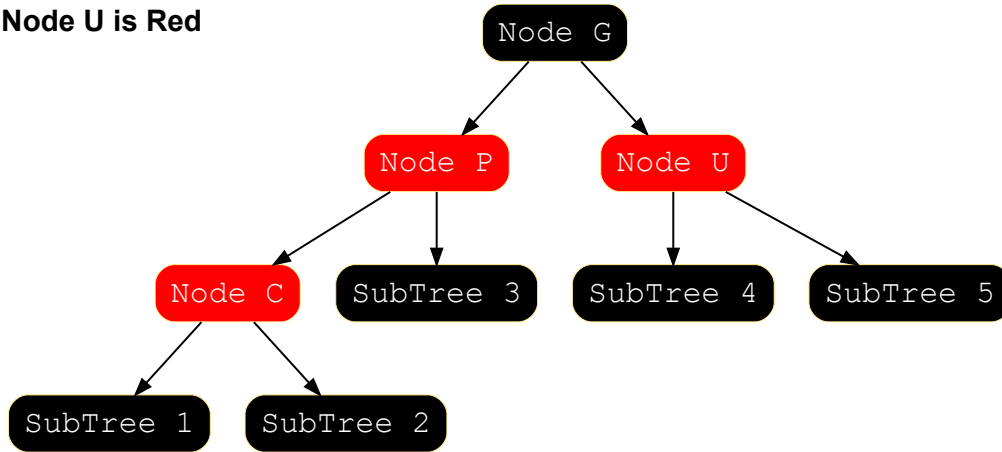
If the newly inserted node is not a root and the parent of the newly inserted node is not **Black**, first, we will check Node U which is the and based on Node U's color, we balance the tree. If

Node U is **Red**, do the following:



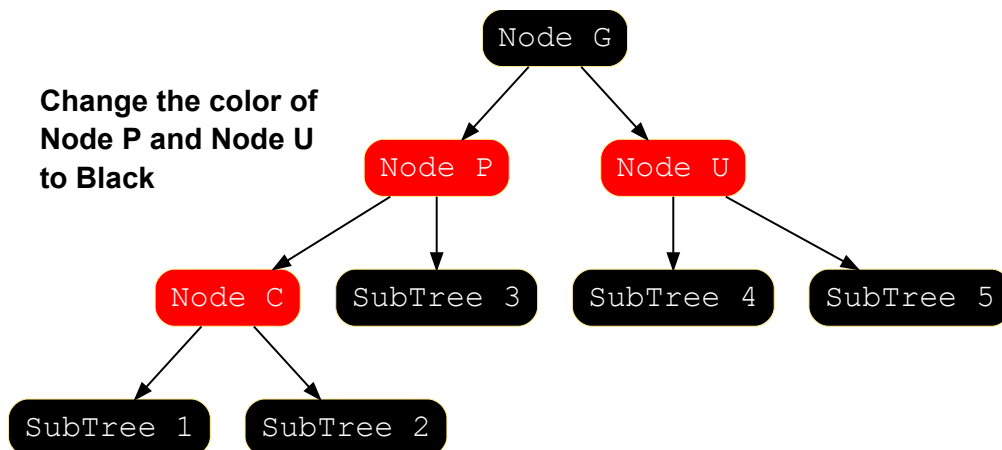
1. Change color of Node P and U to **Black**
2. Change color of Node G to **Red**
3. Make Node G the new *currentNode* and repeat the same process from step two

Node U is Red



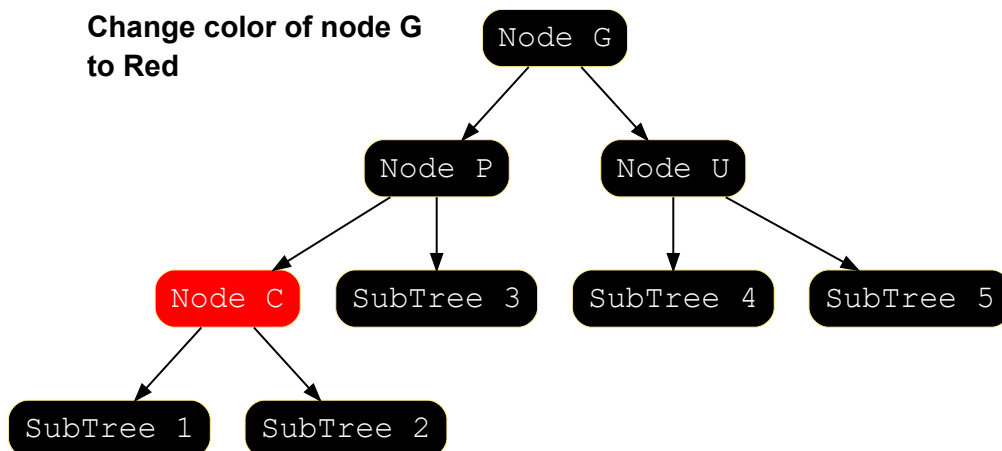
1 of 6

**Change the color of
Node P and Node U
to Black**

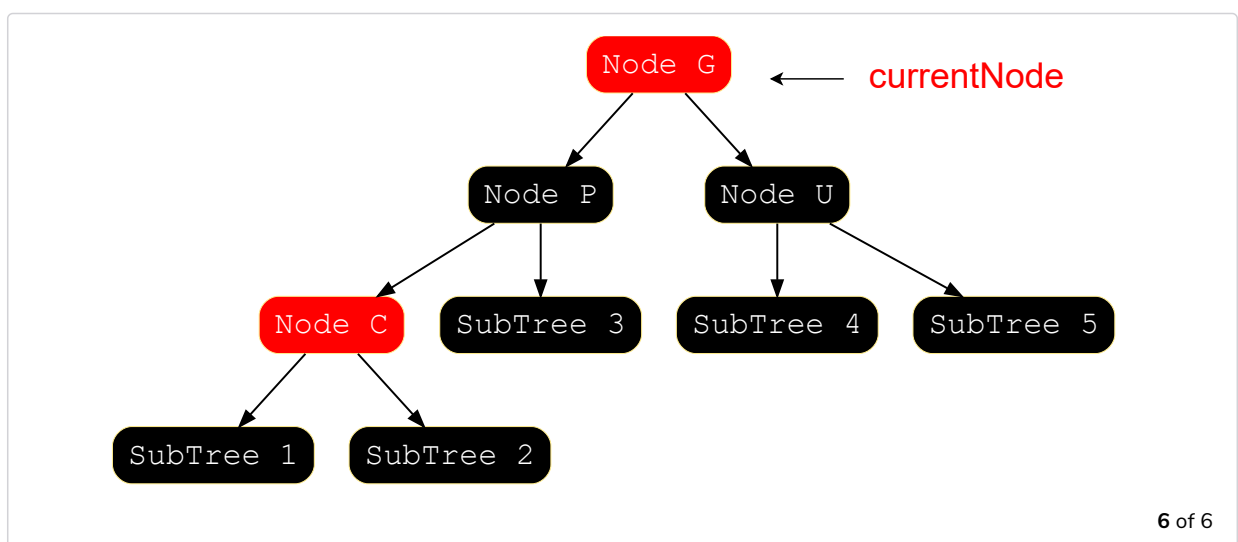
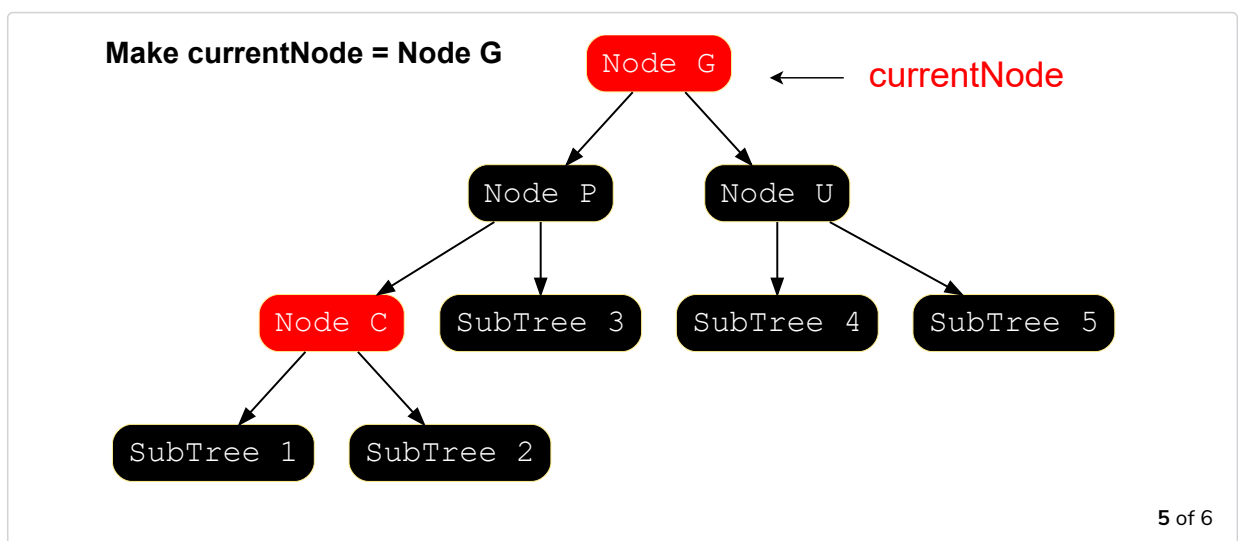
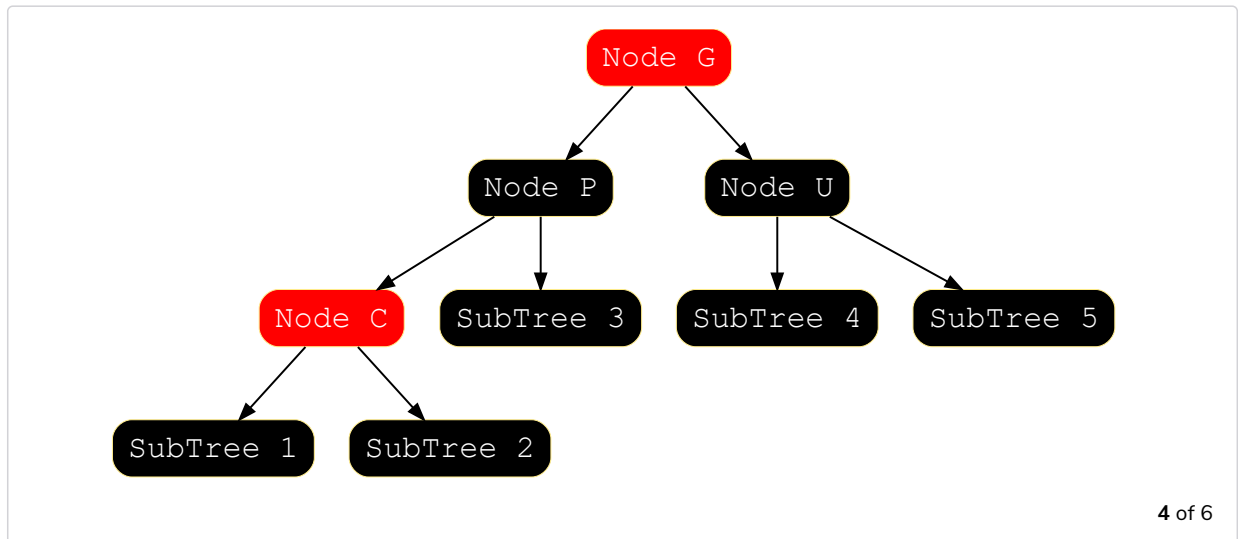


2 of 6

**Change color of node G
to Red**



3 of 6



— []

If Node U is **Black**, then we come across four different scenarios based on the arrangements of Node P and G, just like we did in AVL trees. We will cover each of these scenarios and try to help you understand through illustrations.

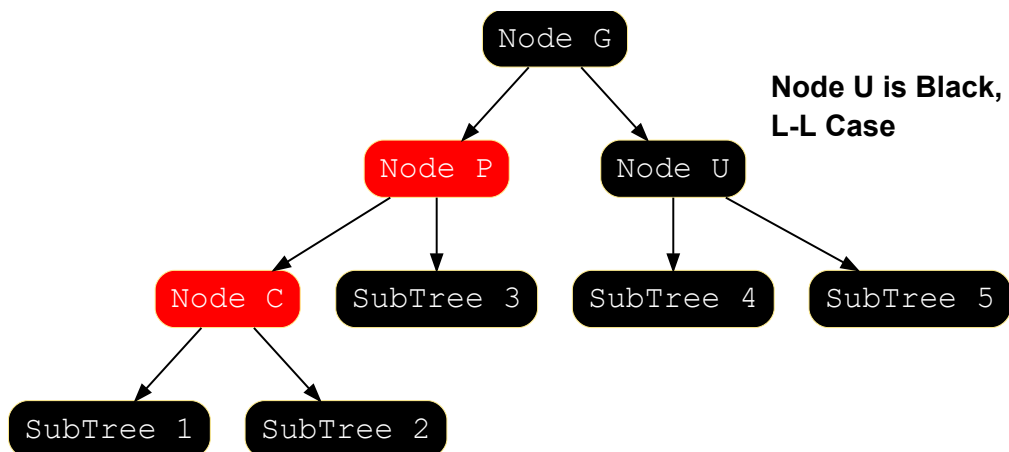
- Left-Left: Node P is the *leftChild* of Node G and *currentNode* is the *leftChild* of Node P

- Left-Right: Node P is the *leftChild* of Node G and *currentNode* is the *rightChild* of Node P
- Right-Right: Node P is the *rightChild* of Node G and *currentNode* is the *rightChild* of Node P
- Right-Left: Node P is the *rightChild* of Node G and *currentNode* is the *leftChild* of Node P

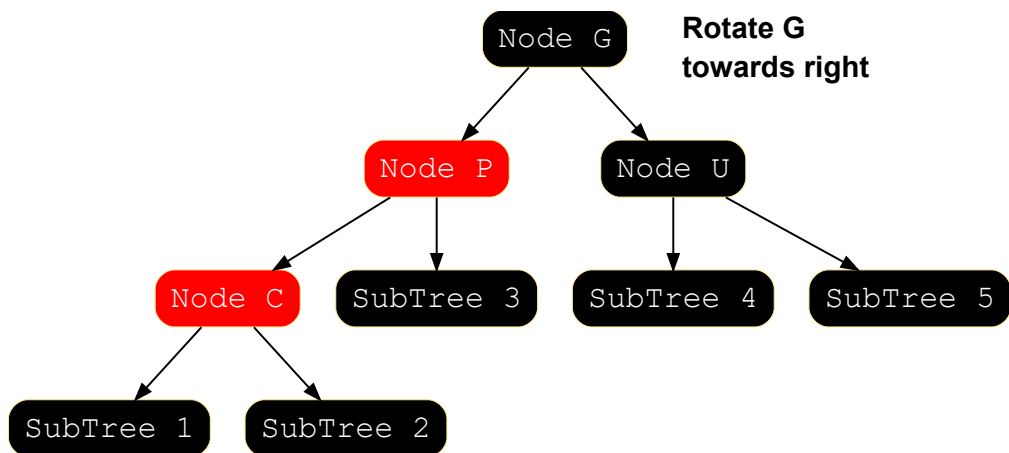
Case 1: Left-Left

In case when the Node P is the *leftChild* of Node G and *currentNode* is the *leftChild* of Node P, we perform the following steps. Look at the illustration below for a better understanding.

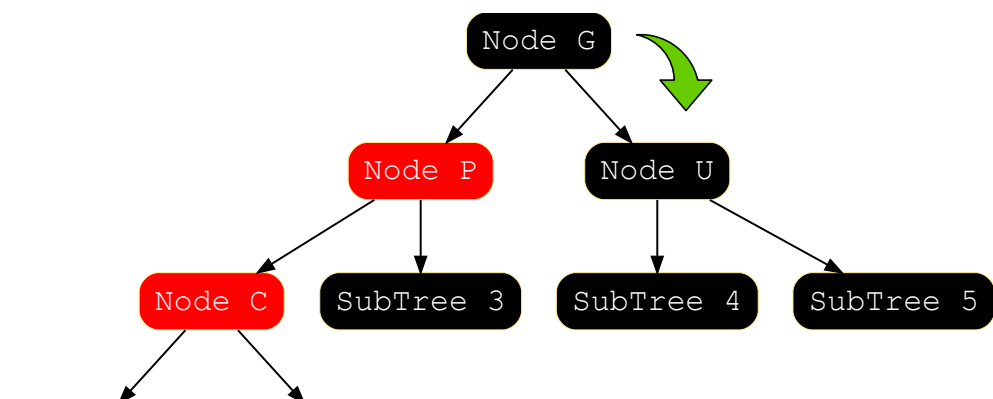
1. Right Rotate Node G
2. Swap the colors of Nodes G and P



1 of 5



2 of 5

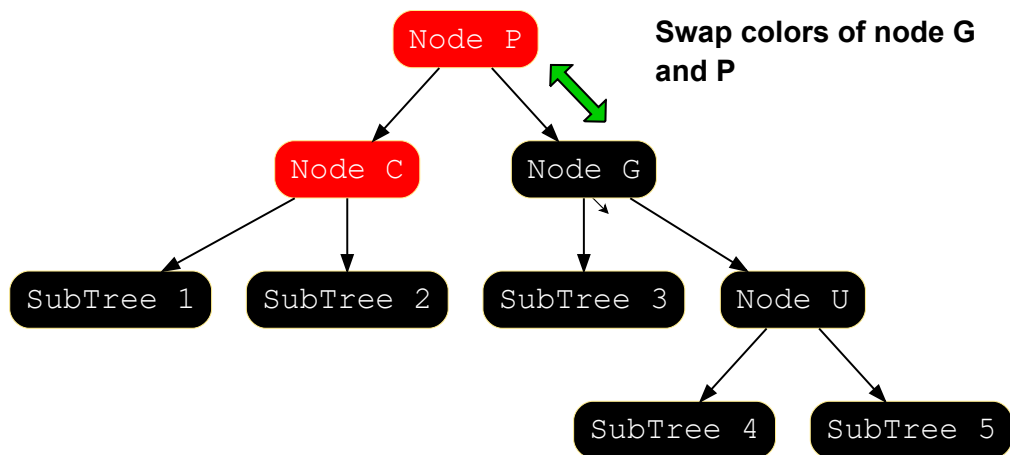


SubTree 1

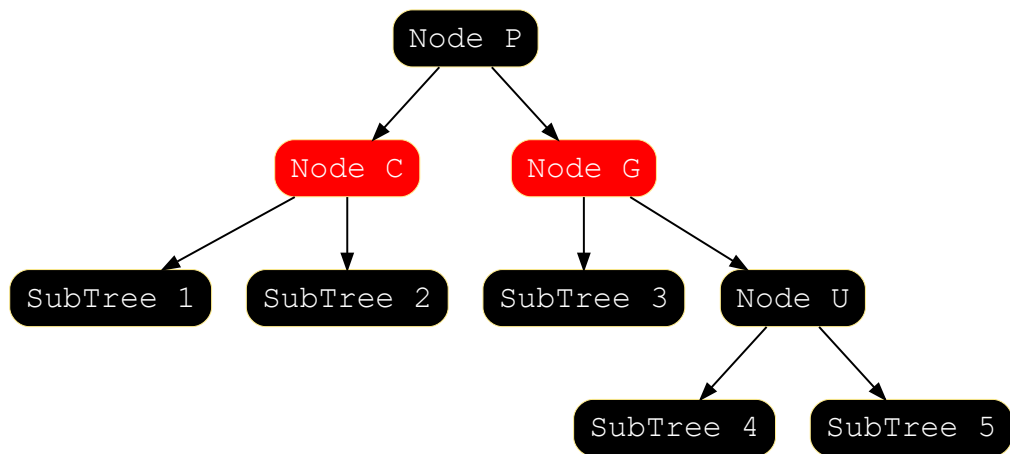
SubTree 2



3 of 5



4 of 5



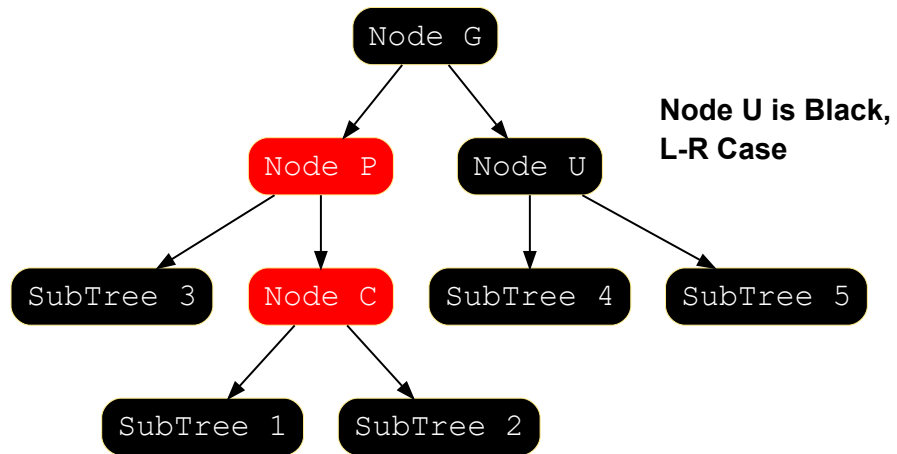
5 of 5

— {}

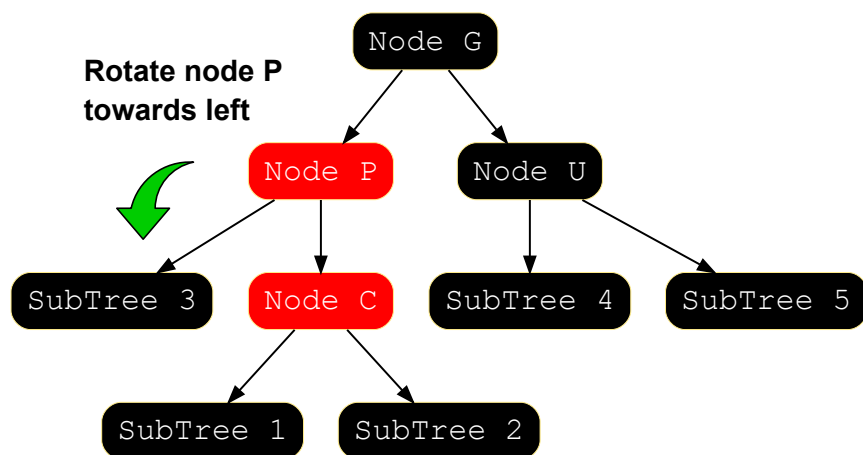
Case 2: Left-Right

In the case when the Node P is the *leftChild* of Node G and the *currentNode* is the *rightChild* of Node P, we perform the following steps. Look at the illustration below for better understanding:

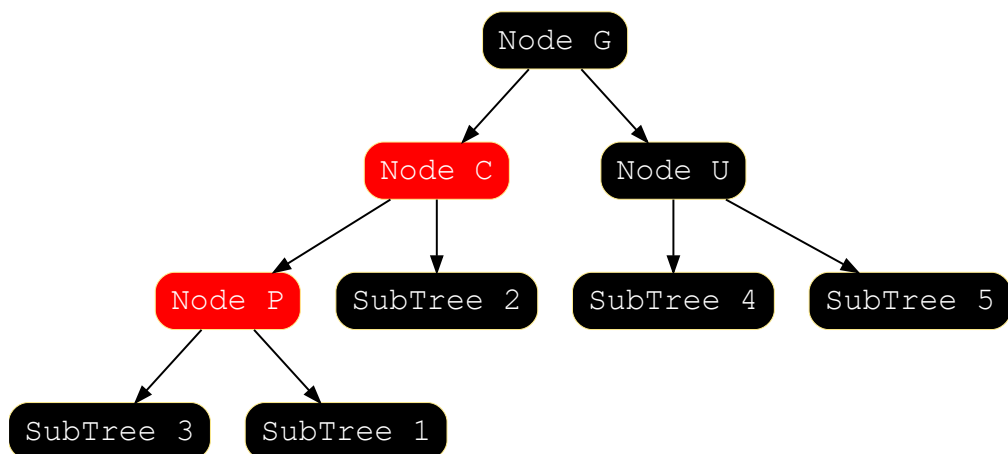
1. Left Rotate Node P
2. After that, repeat the steps that we covered in the Left-Left case



1 of 8



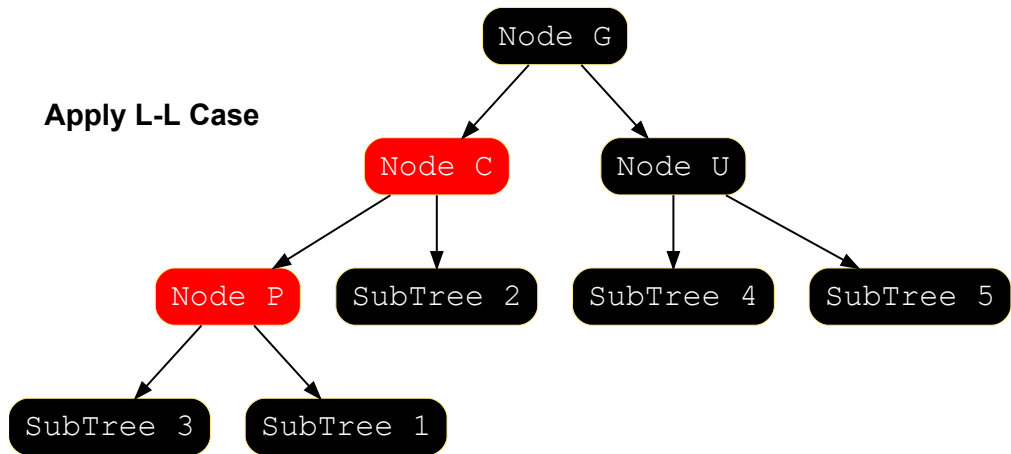
2 of 8



3 of 8

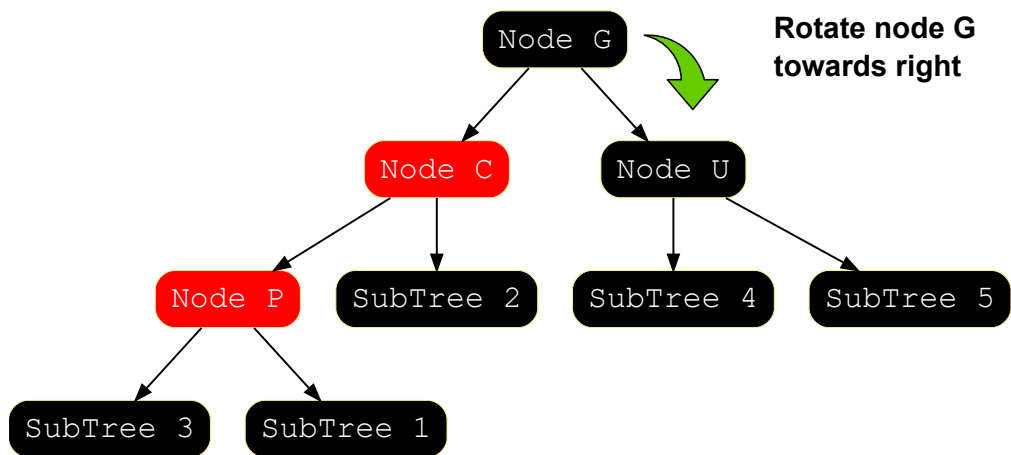


Apply L-L Case

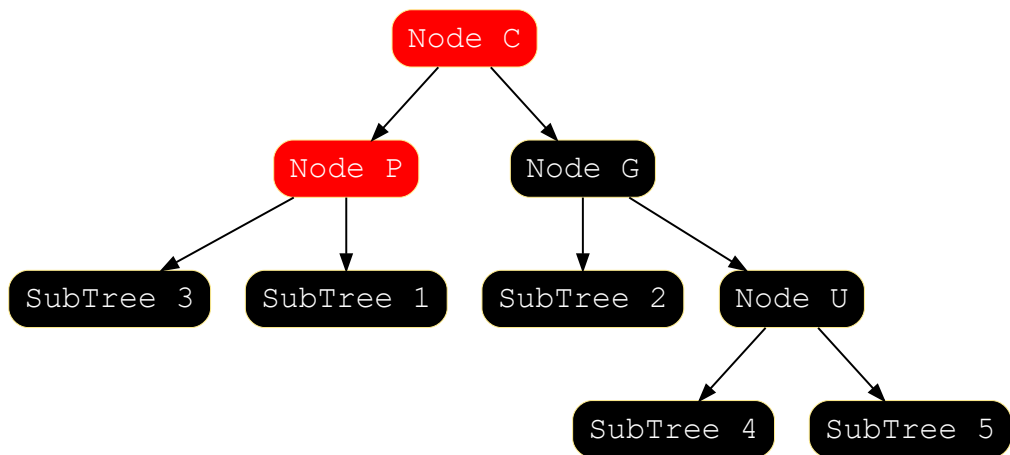


4 of 8

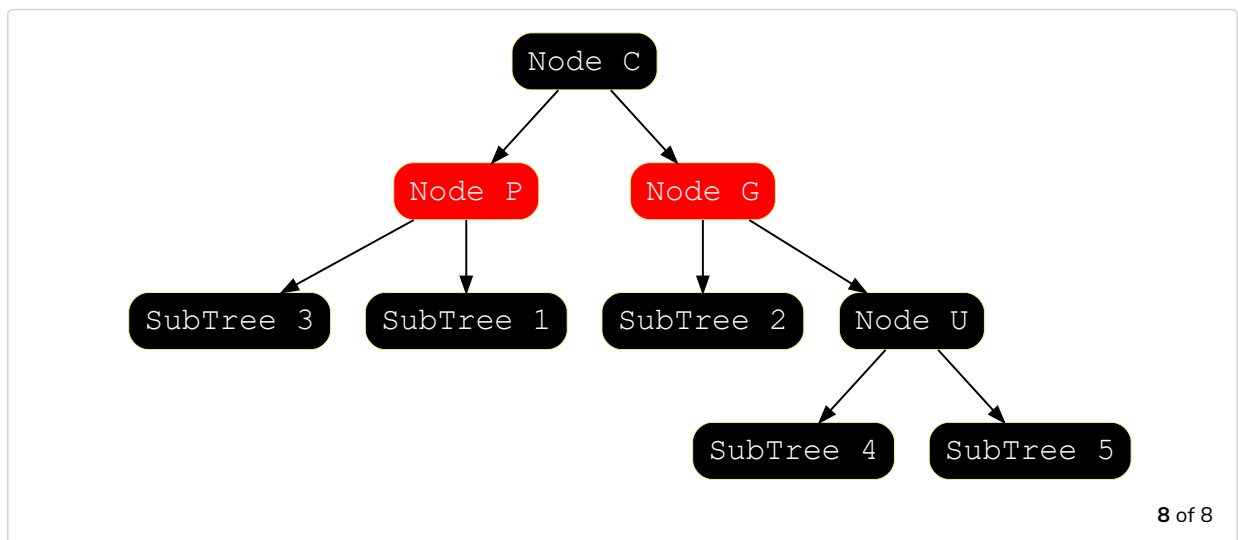
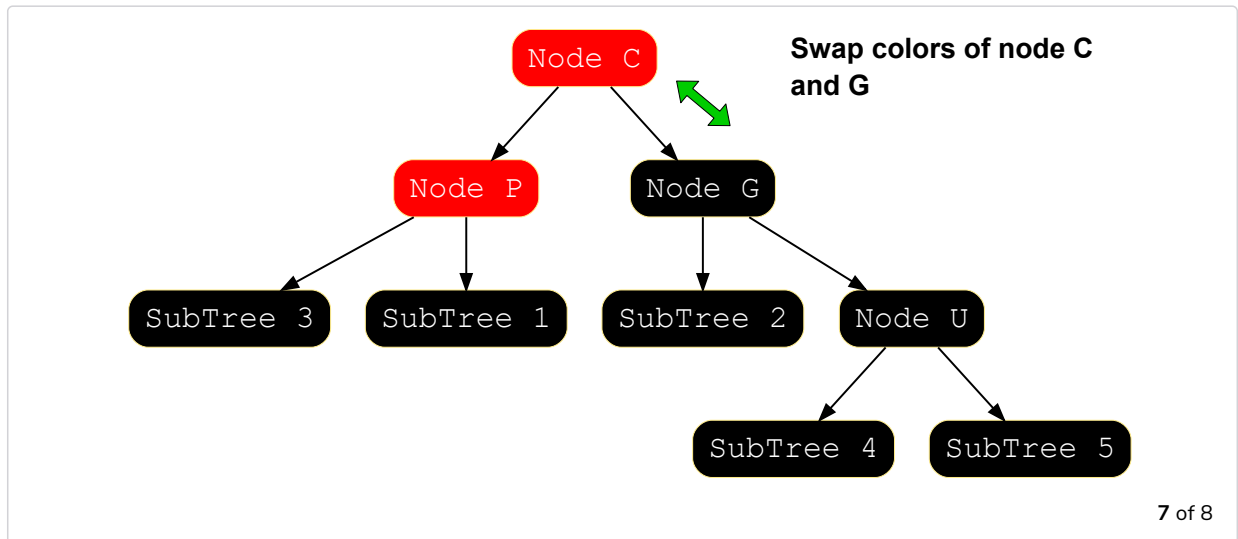
Rotate node G
towards right



5 of 8



6 of 8

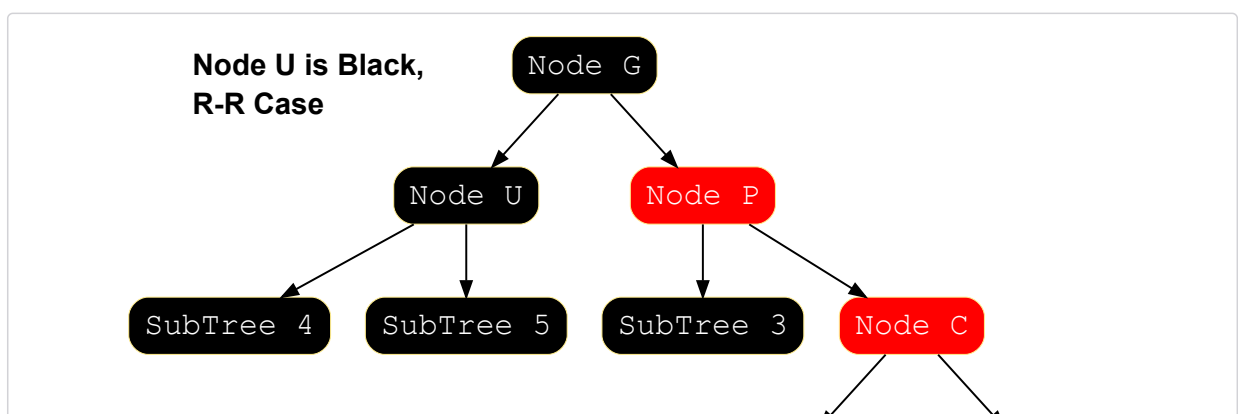


— []

Case 3: Right-Right

In the case when Node P is the *rightChild* of Node G and the *currentNode* is the *rightChild* of Node P, we perform the following steps. Look at the illustration below for a better understanding.

1. Left Rotate Node G
2. Swap colors of Node G and P



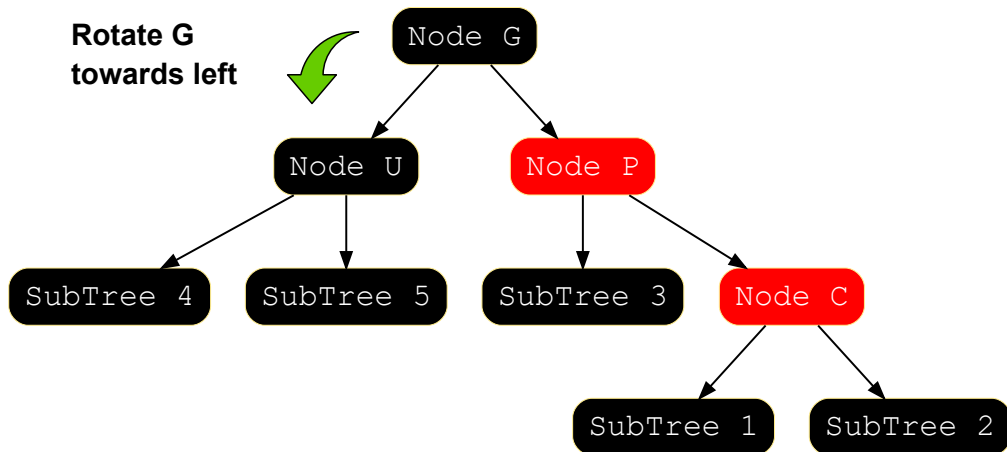
SubTree 1

SubTree 2

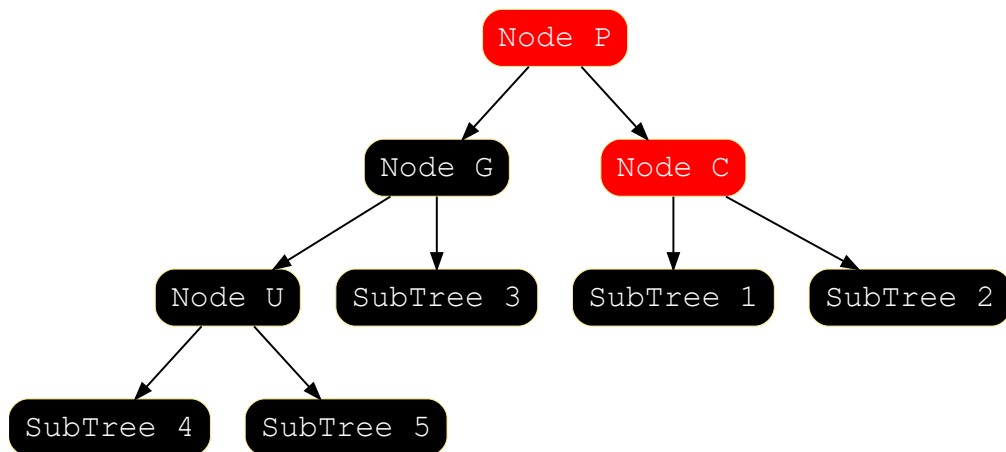


1 of 5

Rotate G
towards left

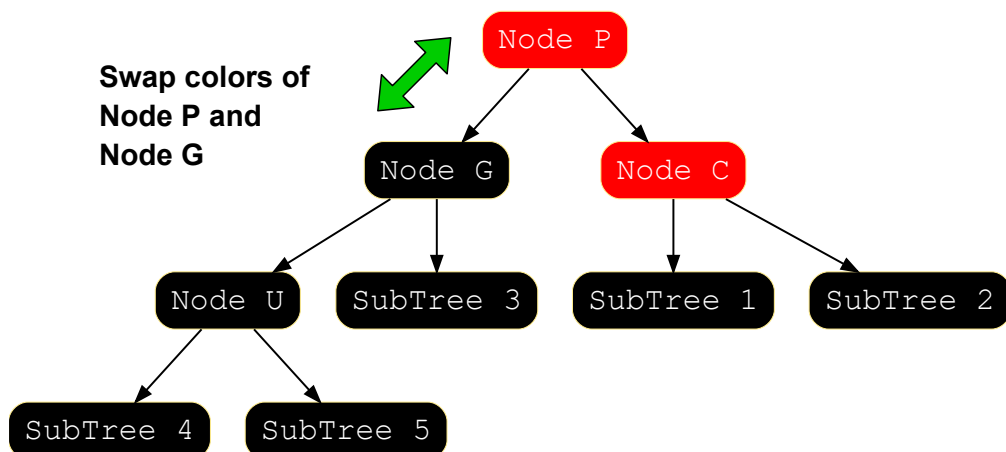


2 of 5

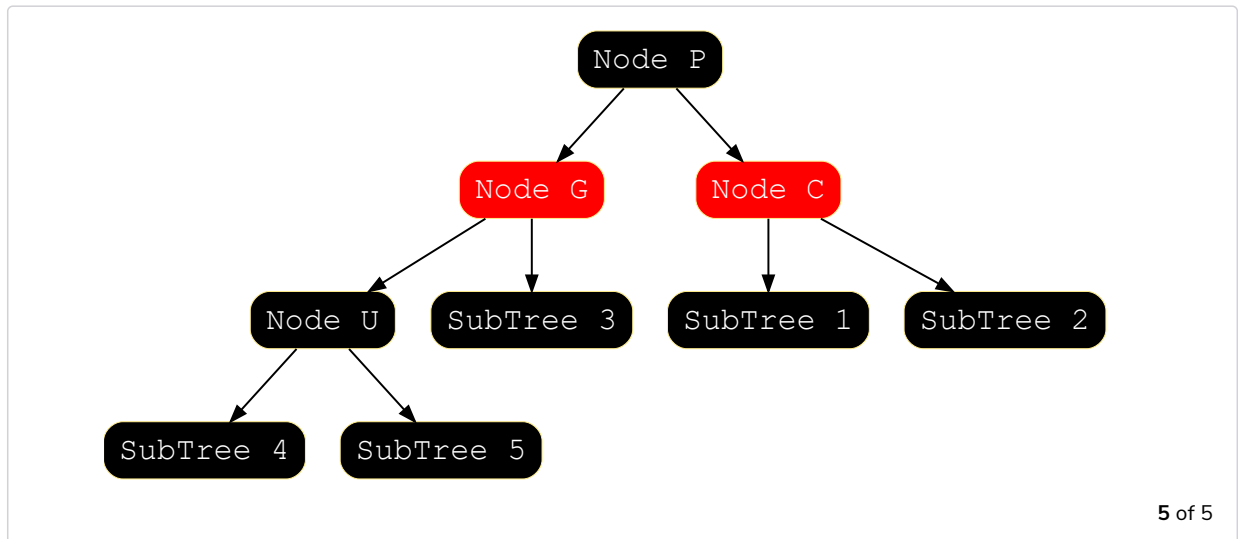


3 of 5

Swap colors of
Node P and
Node G



4 of 5

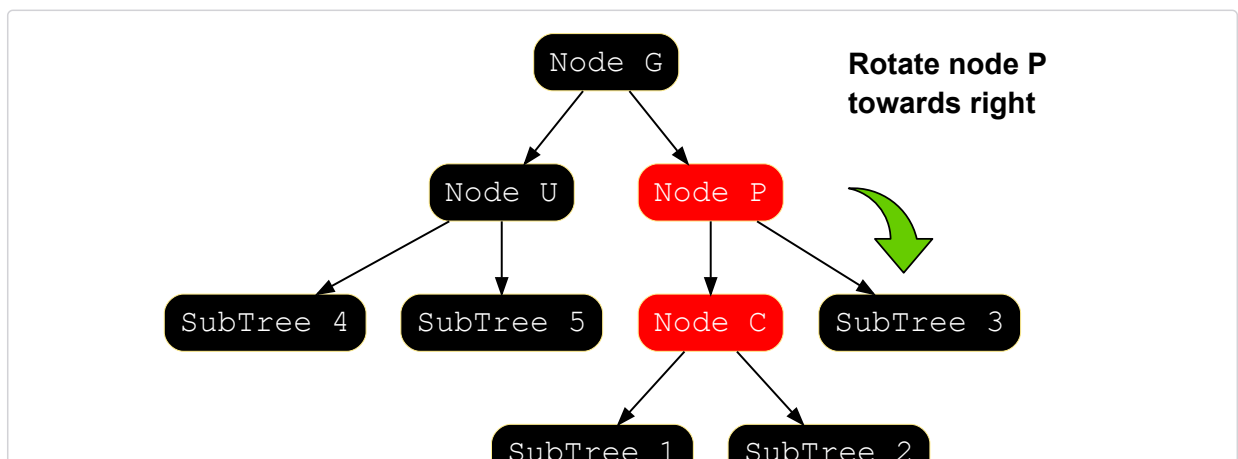
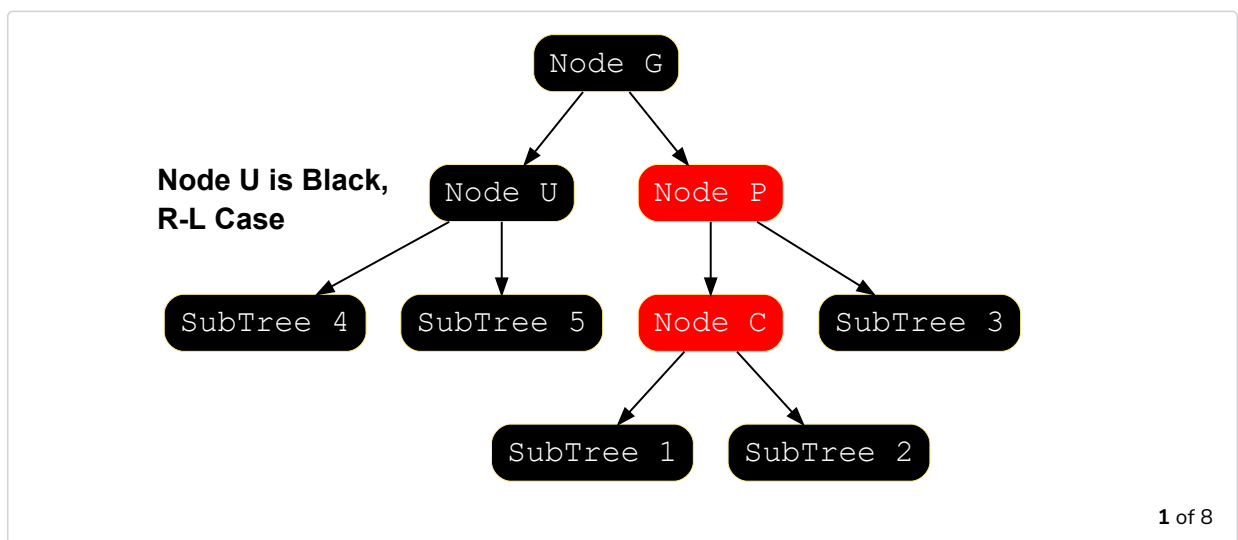


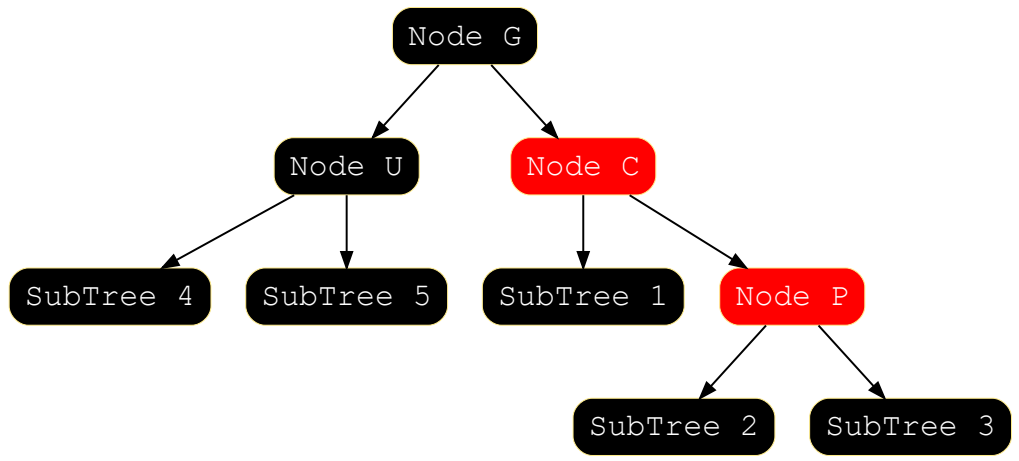
— []

Case 4: Right-Left

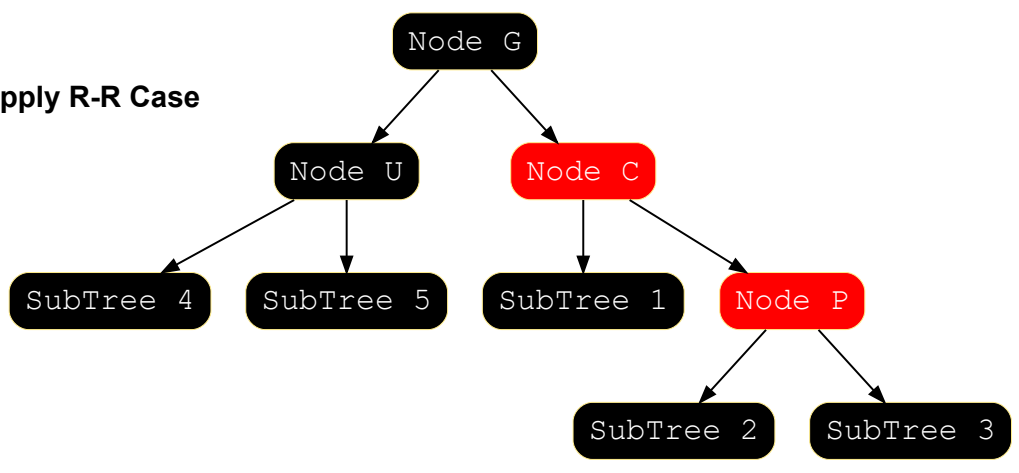
In the case when Node P is the *rightChild* of Node G and the *currentNode* is the *leftChild* of Node P, we perform the following steps. Look at the illustration below for a better understanding.

1. Right Rotate Node P
2. After that, repeat the steps that we covered in Right-Right case

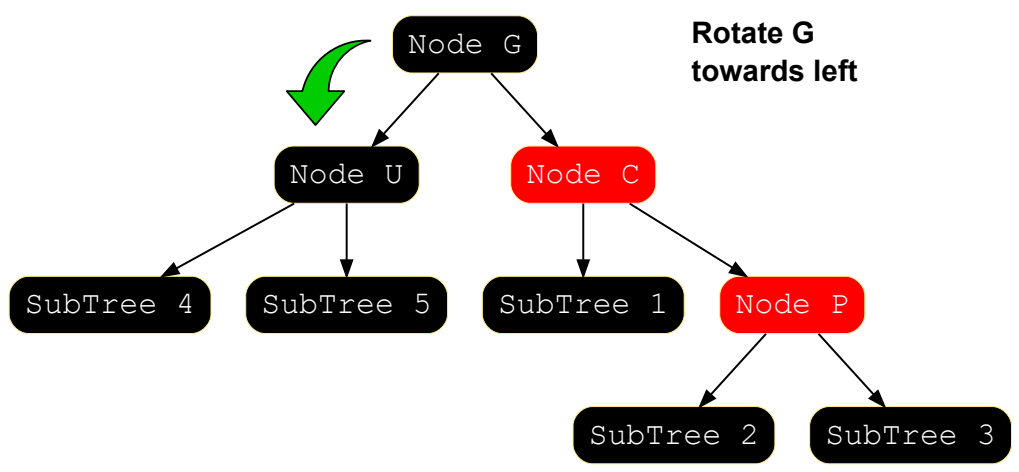


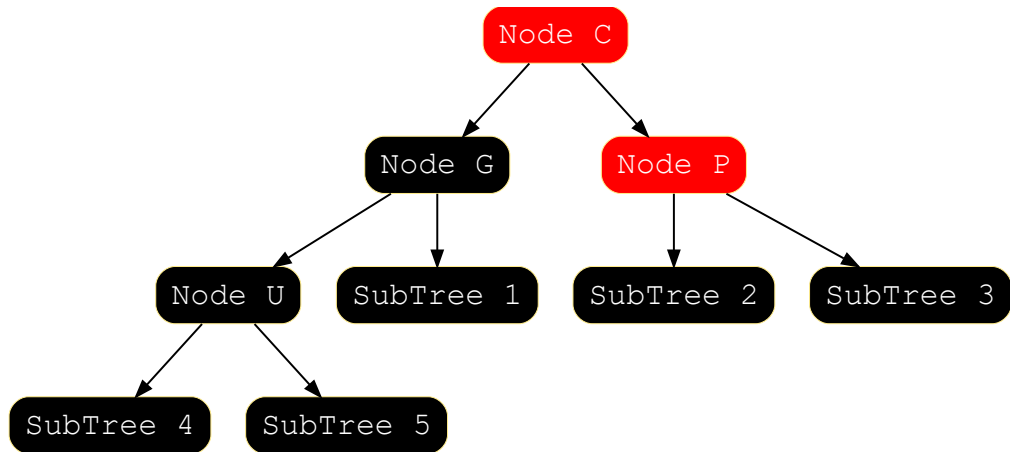


Apply R-R Case

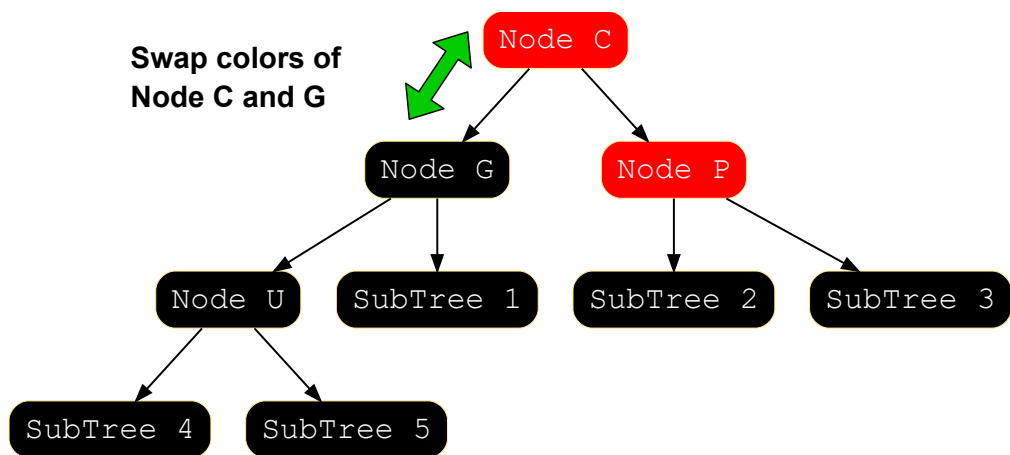


Rotate G towards left

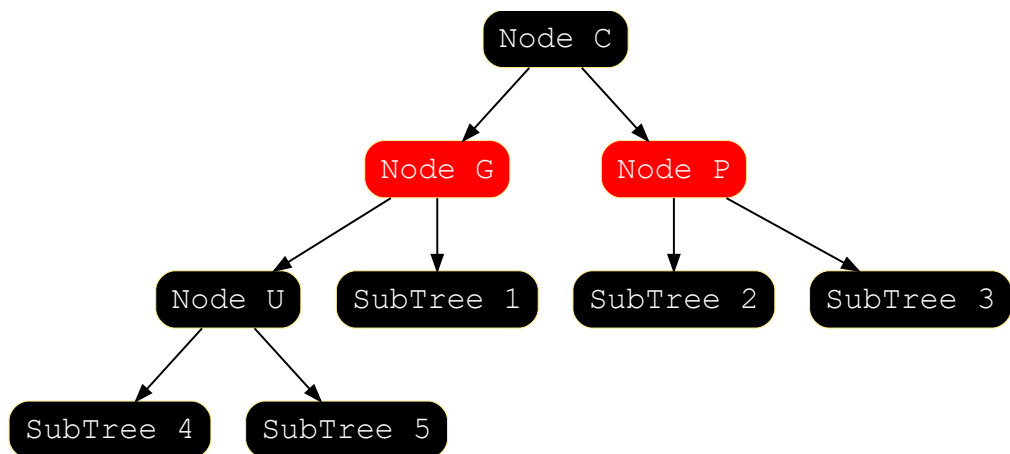




6 of 8



7 of 8



8 of 8

In the next lesson, we'll study red-black tree deletion!



Mark as Completed



Report an Issue



Ask a Question

(https://discuss.educative.io/tag/red-black-tree-insertion__introduction-to-trees__data-structures-for-coding-interviews-in-python)