

## All Tasks Scheduling Orders (hard)

# We'll cover the following

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  - Time and Space Complexity

#### Problem Statement #

There are 'N' tasks, labeled from '0' to 'N-1'. Each task can have some prerequisite tasks which need to be completed before it can be scheduled. Given the number of tasks and a list of prerequisite pairs, write a method to print all possible ordering of tasks meeting all prerequisites.

#### Example 1:

```
Input: Tasks=3, Prerequisites=[0, 1], [1, 2]
Output: [0, 1, 2]
Explanation: There is only possible ordering of the tasks.
```

#### Example 2:

```
Input: Tasks=4, Prerequisites=[3, 2], [3, 0], [2, 0], [2, 1]
Output:
1) [3, 2, 0, 1]
2) [3, 2, 1, 0]
Explanation: There are two possible orderings of the tasks meeting all prerequisit es.
```

#### Example 3:



```
Input: Tasks=6, Prerequisites=[2, 5], [0, 5], [0, 4], [1, 4], [3, 2], [1, 3]
Output:
1) [0, 1, 4, 3, 2, 5]
2) [0, 1, 3, 4, 2, 5]
3) [0, 1, 3, 2, 4, 5]
4) [0, 1, 3, 2, 5, 4]
5) [1, 0, 3, 4, 2, 5]
6) [1, 0, 3, 2, 4, 5]
7) [1, 0, 3, 2, 5, 4]
8) [1, 0, 4, 3, 2, 5]
9) [1, 3, 0, 2, 4, 5]
10) [1, 3, 0, 2, 5, 4]
11) [1, 3, 0, 4, 2, 5]
12) [1, 3, 2, 0, 5, 4]
13) [1, 3, 2, 0, 4, 5]
```

### Try it yourself #

Try solving this question here:

```
Python3
👙 Java
                                       ⊘ C++
                          JS JS
    def print_orders(tasks, prerequisites):
 2
       # TODO: Write your code here
 3
       print()
 4
 5
 6
 7
    def main():
 8
       print("Task Orders: ")
 9
       print_orders(3, [[0, 1], [1, 2]])
10
       print("Task Orders: ")
11
       print_orders(4, [[3, 2], [3, 0], [2, 0], [2, 1]])
12
13
       print("Task Orders: ")
14
15
       print_orders(6, [[2, 5], [0, 5], [0, 4], [1, 4], [3, 2], [1, 3]])
16
17
18 main()
19
                                                                                            \leftarrow
                                                                                                 []
\triangleright
```

#### Solution #

This problem is similar to Tasks Scheduling Order

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/5066018374287360/), the only difference is that we need to find all the topological orderings of the tasks.

At any stage, if we have more than one source available and since we can choose any source, therefore, in this case, we will have multiple orderings of the tasks. We can use a recursive



#### Code #

Here is what our algorithm will look like:

```
👙 Java
            🤁 Python3
                         ⊘ C++
                                      JS JS
    from collections import deque
 2
 3
 4
    def print_orders(tasks, prerequisites):
 5
      sortedOrder = []
      if tasks <= 0:
 6
 7
         return False
 8
 9
      # a. Initialize the graph
      inDegree = {i: 0 for i in range(tasks)} # count of incoming edges
10
      graph = {i: [] for i in range(tasks)} # adjacency list graph
11
12
13
      # b. Build the graph
14
      for prerequisite in prerequisites:
15
        parent, child = prerequisite[0], prerequisite[1]
        graph[parent].append(child) # put the child into it's parent's list
16
17
         inDegree[child] += 1 # increment child's inDegree
18
19
      # c. Find all sources i.e., all vertices with 0 in-degrees
      sources = deque()
20
21
      for key in inDegree:
22
        if inDegree[key] == 0:
23
          sources.append(key)
24
25
      print_all_topological_sorts(graph, inDegree, sources, sortedOrder)
26
27
28 def print all topological sorts(graph, inDegree, sources, sortedOrder):
\triangleright
                                                                                          \leftarrow
                                                                                               []
```

#### Time and Space Complexity #

If we don't have any prerequisites, all combinations of the tasks can represent a topological ordering. As we know, that there can be N! combinations for 'N' numbers, therefore the time and space complexity of our algorithm will be O(V!\*E) where 'V' is the total number of tasks and 'E' is the total prerequisites. We need the 'E' part because in each recursive call, at max, we remove (and add back) all the edges.

