



Count of Palindromic Substrings

We'll cover the following ^

- Problem Statement
- Solution

Problem Statement

Given a string, find the total number of palindromic substrings in it. Please note we need to find the total number of substrings and not subsequences.

Example 1:

```
Input: "abdbca"
Output: 7
Explanation: Here are the palindromic substrings, "a", "b", "d", "b", "c", "a", "b
db".
```

Example 2:

```
Input: = "cddpd"
Output: 7
Explanation: Here are the palindromic substrings, "c", "d", "d", "p", "d", "dd",
"dpd".
```

Example 3:

```
Input: = "pqr"
Output: 3
Explanation: Here are the palindromic substrings,"p", "q", "r".
```

Solution

This problem follows the Longest Palindromic Subsequence (https://www.educative.io/collection/page/5668639101419520/5633779737559040/574811928317 1328/) pattern, and can be easily converted to Longest Palindromic Substring (https://www.educative.io/collection/page/5668639101419520/5633779737559040/566160146196 0704/). The only difference is that instead of calculating the longest palindromic substring, we will instead count all the palindromic substrings.

Let's jump directly to the bottom-up dynamic programming solution:



```
∣ 🌄 JS
                        Python3
 Java
    def count_PS(st):
 1
 2
      n = len(st)
 3
      # dp[i][j] will be 'true' if the string from index 'i' to index 'j' is a palindrome
      dp = [[False for _ in range(n)] for _ in range(n)]
      count = 0
 5
 6
 7
      # every string with one character is a palindrome
 8
      for i in range(n):
        dp[i][i] = True
 9
10
        count += 1
11
      for startIndex in range(n - 1, -1, -1):
12
        for endIndex in range(startIndex + 1, n):
13
14
          if st[startIndex] == st[endIndex]:
            # if it's a two character string or if the remaining string is a palindrome too
15
16
             if endIndex - startIndex == 1 or dp[startIndex + 1][endIndex - 1]:
              dp[startIndex][endIndex] = True
17
18
               count += 1
19
20
       return count
21
22
23
    def main():
      print(count_PS("abdbca"))
24
25
      print(count_PS("cddpd"))
26
      print(count_PS("pqr"))
27
      print(count_PS("qqq"))
28
29
30
    main()
\triangleright
                                                                                               []
```

The time and space complexity of the above algorithm is $O(n^2)$, where 'n' is the length of the input string.

