

Solution Review: Problem Challenge 4

We'll cover the following

- Words Concatenation (hard)
- Solution
- Code
 - Time Complexity
 - Space Complexity

Words Concatenation (hard)

Given a string and a list of words, find all the starting indices of substrings in the given string that are a **concatenation of all the given words** exactly once **without any overlapping** of words. It is given that all words are of the same length.

Example 1:

```
Input: String="catfoxcat", Words=["cat", "fox"]
Output: [0, 3]
Explanation: The two substring containing both the words are "catfox" & "foxcat".
```

Example 2:

```
Input: String="catcatfoxfox", Words=["cat", "fox"]
Output: [3]
Explanation: The only substring containing both the words is "catfox".
```

Solution

This problem follows the **Sliding Window** pattern and has a lot of similarities with Maximum Sum Subarray of Size K

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/517704302723 0720/). We will keep track of all the words in a **HashMap** and try to match them in the given string. Here are the set of steps for our algorithm:

- 1. Keep the frequency of every word in a **HashMap**.
- 2. Starting from every index in the string, try to match all the words.
- 3. In each iteration, keep track of all the words that we have already seen in another **HashMap**.
- 4. If a word is not found or has a higher frequency than required, we can move on to the next character in the string.
- 5. Store the index if we have found all the words.
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Here is what our algorithm will look like:



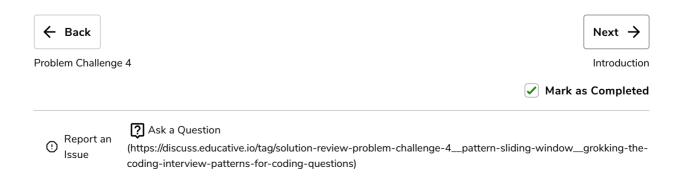
```
👙 Java
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                                     JS JS
    def find_word_concatenation(str1, words):
      if len(words) == 0 or len(words[0]) == 0:
 3
        return []
 4
 5
      word_frequency = {}
 6
      for word in words:
 7
 8
        if word not in word frequency:
 9
          word frequency[word] = 0
10
        word_frequency[word] += 1
11
12
      result_indices = []
13
      words_count = len(words)
14
      word_length = len(words[0])
15
16
      for i in range((len(str1) - words_count * word_length)+1):
17
        words_seen = {}
18
        for j in range(0, words count):
19
          next_word_index = i + j * word_length
20
          # Get the next word from the string
21
          word = str1[next_word_index: next_word_index + word_length]
          if word not in word frequency: # Break if we don't need this word
22
23
            break
24
          # Add the word to the 'words_seen' map
25
26
          if word not in words_seen:
27
            words_seen[word] = 0
28
          words seen[word] += 1
                                                                                            []
D
```

Time Complexity

The time complexity of the above algorithm will be O(N*M*Len) where 'N' is the number of characters in the given string, 'M' is the total number of words, and 'Len' is the length of a word.

Space Complexity

The space complexity of the algorithm is O(M) since at most, we will be storing all the words in the two **HashMaps**. In the worst case, we also need O(N) space for the resulting list. So, the overall space complexity of the algorithm will be O(M+N).









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