Conflicting Appointments (medium)

We'll cover the following ^

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Problem Statement

Given an array of intervals representing 'N' appointments, find out if a person can **attend all the appointments**.

Example 1:

```
Appointments: [[1,4], [2,5], [7,9]]
Output: false
Explanation: Since [1,4] and [2,5] overlap, a person cannot attend both of these a ppointments.
```

Example 2:

```
Appointments: [[6,7], [2,4], [8,12]]
Output: true
Explanation: None of the appointments overlap, therefore a person can attend all o f them.
```

Example 3:

```
Appointments: [[4,5], [2,3], [3,6]]
Output: false
Explanation: Since [4,5] and [3,6] overlap, a person cannot attend both of these a ppointments.
```

Try it yourself

Try solving this question here:









```
def can_attend_all_appointments(intervals):
 2
      # TODO: Write your code here
 3
      return False
 4
 5
 6 def main():
 7
      print("Can attend all appointments: " + str(can_attend_all_appointments([[1, 4], [2, 5],
      print("Can attend all appointments: " + str(can_attend_all_appointments([[6, 7], [2, 4],
 8
      print("Can attend all appointments: " + str(can_attend_all_appointments([[4, 5], [2, 3],
 9
10
11
12
    main()
13
\triangleright
                                                                                    \leftarrow
                                                                                                 []
```

Solution

The problem follows the Merge Intervals

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/565201724243 9680/) pattern. We can sort all the intervals by start time and then check if any two intervals overlap. A person will not be able to attend all appointments if any two appointments overlap.

Code

Time complexity #

Here is what our algorithm will look like:

```
👙 Java
           Pvthon3
                         G C++
                                     Js JS
 1 ∨ def can_attend_all_appointments(intervals):
       intervals.sort(key=lambda x: x[0])
 3
      start, end = 0, 1
 4 \vee for i in range(1, len(intervals)):
        if intervals[i][start] < intervals[i-1][end]:</pre>
           # please note the comparison above, it is "<" and not "<="</pre>
 6
           # while merging we needed "<=" comparison, as we will be merging the two
 7
           \# intervals having condition "intervals[i][start] == intervals[i - 1][end]" but
 9
           # such intervals don't represent conflicting appointments as one starts right
10
           # after the other
11
           return False
       return True
12
13
14
15 ∨def main():
      print("Can attend all appointments: " + str(can_attend_all_appointments([[1, 4], [2, 5],
16
       print("Can attend all appointments: " + str(can_attend_all_appointments([[6, 7], [2, 4],
17
       print("Can attend all appointments: " + str(can_attend_all_appointments([[4, 5], [2, 3],
18
19
20
21
    main()
22
                                                                                              []
   educative
```





The time complexity of the above algorithm is O(N*logN), where 'N' is the total number of appointments. Though we are iterating the intervals only once, our algorithm will take O(N*logN) since we need to sort them in the beginning.

Space complexity

The space complexity of the above algorithm will be O(N), which we need for sorting. For Java, Arrays.sort() uses Timsort (https://en.wikipedia.org/wiki/Timsort), which needs O(N) space.

Similar Problems

Problem 1: Given a list of appointments, find all the conflicting appointments.

coding-interview-patterns-for-coding-questions)

Example:



