## Subsets With Duplicates (easy)

# We'll cover the following ^

- Problem Statement
- Try it yourself
- Solution
- Code
  - Time complexity
  - Space complexity

#### Problem Statement #

Given a set of numbers that might contain duplicates, find all of its distinct subsets.

#### Example 1:

```
Input: [1, 3, 3]
Output: [], [1], [3], [1,3], [1,3,3]
```

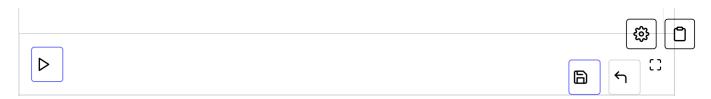
#### Example 2:

```
Input: [1, 5, 3, 3]
Output: [], [1], [5], [3], [1,5], [1,3], [5,3], [1,5,3], [3,3], [1,3,3], [1,5,3,3]
```

### Try it yourself #

Try solving this question here:

```
Python3
👙 Java
                        Js JS
                                    G C++
 1 def find_subsets(nums):
      subsets = []
      # TODO: Write your code here
 3
      return subsets
 5
 6
 7
    def main():
 8
      print("Here is the list of subsets: " + str(find_subsets([1, 3, 3])))
 9
      print("Here is the list of subsets: " + str(find_subsets([1, 5, 3, 3])))
10
11
    educative
```



#### Solution #

This problem follows the Subsets

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/567024937861 1200) pattern and we can follow a similar **Breadth First Search (BFS)** approach. The only additional thing we need to do is handle duplicates. Since the given set can have duplicate numbers, if we follow the same approach discussed in Subsets

(https://www.educative.io/collection/page/5668639101419520/5671464854355968/567024937861 1200), we will end up with duplicate subsets, which is not acceptable. To handle this, we will do two extra things:

- 1. Sort all numbers of the given set. This will ensure that all duplicate numbers are next to each other.
- 2. Follow the same BFS approach but whenever we are about to process a duplicate (i.e., when the current and the previous numbers are same), instead of adding the current number (which is a duplicate) to all the existing subsets, only add it to the subsets which were created in the previous step.

Let's take Example-2 mentioned above to go through each step of our algorithm:

```
Given set: [1, 5, 3, 3]
Sorted set: [1, 3, 3, 5]
```

- 1. Start with an empty set: [[]]
- 2. Add the first number (1) to all the existing subsets to create new subsets: [[], [1]];
- 3. Add the second number (3) to all the existing subsets: [[], [1], [3], [1,3]].
- 4. The next number (3) is a duplicate. If we add it to all existing subsets we will get:

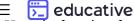
```
[[], [1], [3], [1,3], [3], [1,3], [3,3], [1,3,3]]
```

```
We got two duplicate subsets: [3], [1,3]
Whereas we only needed the new subsets: [3,3], [1,3,3]
```

To handle this instead of adding (3) to all the existing subsets, we only add it to the new subsets which were created in the previous (3rd) step:

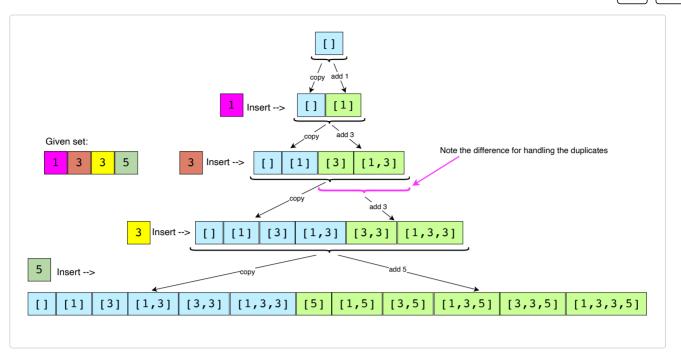
```
[[], [1], [3], [1,3], [1,3,3]]
```

5. Finally, add the forth number (5) to all the existing subsets: [[], [1], [3], [1,3], [1,3,3], [5], [1,5], [3,5], [1,3,5], [1,3,3,5]]



Here is the visual representation of the above steps:





#### Code #

Here is what our algorithm will look like:

```
Python3
                         ⊘ C++
👙 Java
                                     Js JS
 1 def find_subsets(nums):
 2
      # sort the numbers to handle duplicates
 3
      list.sort(nums)
 4
      subsets = []
 5
      subsets.append([])
      startIndex, endIndex = 0, 0
 6
 7
      for i in range(len(nums)):
 8
        startIndex = 0
 9
        # if current and the previous elements are same, create new subsets only from the subs
10
        # added in the previous step
        if i > 0 and nums[i] == nums[i - 1]:
11
           startIndex = endIndex + 1
12
13
        endIndex = len(subsets) - 1
        for j in range(startIndex, endIndex+1):
14
15
          # create a new subset from the existing subset and add the current element to it
          set = list(subsets[j])
16
17
           set.append(nums[i])
18
          subsets.append(set)
19
       return subsets
20
21
    def main():
22
23
      print("Here is the list of subsets: " + str(find_subsets([1, 3, 3])))
24
      print("Here is the list of subsets: " + str(find_subsets([1, 5, 3, 3])))
25
26
27
28
    main()
                                                                                  \leftarrow
                                                                                              []
     educative
```





Since, in each step, the number of subsets could double (if not duplicate) as we add each element to all the existing subsets, the time complexity of the above algorithm is  $O(2^N)$ , where 'N' is the total number of elements in the input set. This also means that, in the end, we will have a total of  $O(2^N)$  subsets at the most.

#### Space complexity #

All the additional space used by our algorithm is for the output list. Since at most we will have a total of  $O(2^N)$  subsets, the space complexity of our algorithm is also  $O(2^N)$ .



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