Binary Tree Level Order Traversal (easy)

We'll cover the following ⌃

- Problem Statement
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- Solution
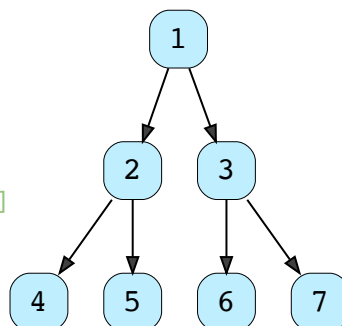- Code
  - Time complexity
  - Space complexity

Problem Statement

Given a binary tree, populate an array to represent its level-by-level traversal. You should populate the values of all **nodes of each level from left to right** in separate sub-arrays.

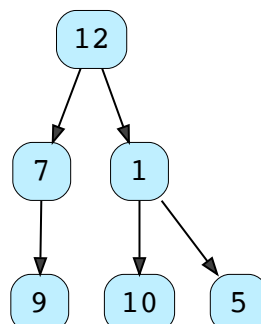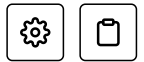**Example 1:**

Level Order Traversal: [[1], [2,3], [4,5,6,7]]

Example 2:

Level Order Traversal: [[12], [7,1], [9,10,5]]

Try it yourself

Try solving this question here:

Java	Python3	JS	C++

```python
1   from collections import deque
2
3
4   class TreeNode:
5       def __init__(self, val):
6           self.val = val
7           self.left, self.right = None, None
8
9
10  def traverse(root):
11      result = []
12      # TODO: Write your code here
13      return result
14
15
16  def main():
17      root = TreeNode(12)
18      root.left = TreeNode(7)
19      root.right = TreeNode(1)
20      root.left.left = TreeNode(9)
21      root.right.left = TreeNode(10)
22      root.right.right = TreeNode(5)
23      print("Level order traversal: " + str(traverse(root)))
24
25
26  main()
27
```

Solution

Since we need to traverse all nodes of each level before moving onto the next level, we can use the **Breadth First Search (BFS)** technique to solve this problem.

We can use a Queue to efficiently traverse in BFS fashion. Here are the steps of our algorithm:
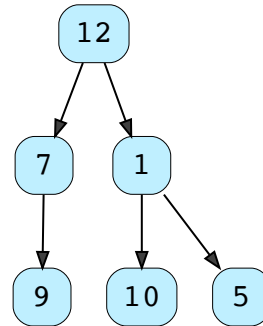
1. Start by pushing the `root` node to the queue.

2. Keep iterating until the queue is empty.

3. In each iteration, first count the elements in the queue (let's call it `levelSize`). We will have these many nodes in the current level.

4. Next, remove `levelSize` nodes from the queue and push their `value` in an array to represent the current level.

5. After removing each node from the queue, insert both of its children into the queue.

6. If the queue is not empty, repeat from step 3 for the next level.

Let's take the example-2 mentioned above to visually represent our algorithm:
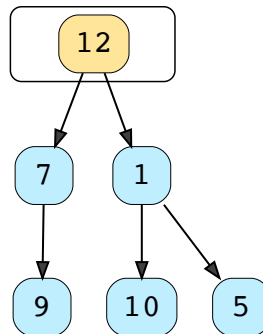
Level Size: 1

Queue: 12

Start by pushing the root to the queue.

Level Size: 1

Queue: 12

Count the elements of the queue (levelSize = 1), they all will be in the first level. Since the levelSize is "1" there will be one element in the first level.
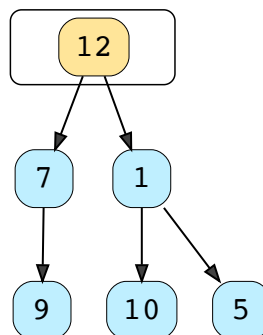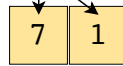
Level Size: 1

Queue: 12 7 1

← insert

Level 1: 12

Move "one" element to the the output array representing the first level and push its children to the queue.
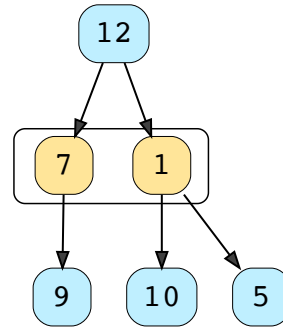
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Level Size: 2

Queue: 7 | 1

Level 1: 12

Count the elements of the queue (levelSize = 2), they all will be in the second level. Since the levelSize is "2" there will be two elements in the second level.

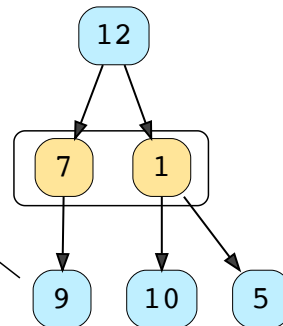**Level Size: 2**

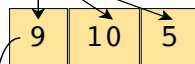Queue: 7 | 1 | 9 | 10 | 5

Level 1: 12

Level 2: 7 | 1

insert

Move "two" elements to the the output array representing the second level and push their children to the queue in the same order.
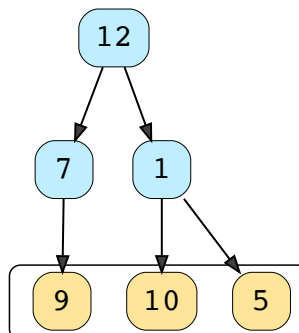

Level Size: 3

Queue: 9 | 10 | 5

Level 1: 12

Level 2: 7 | 1

Level 3:

Count the elements of the queue (levelSize = 3), they all will be in the third level. Since the levelSize is "3" there will be three elements in the third level.

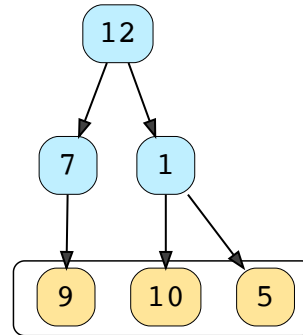educative

Level Size: 3

Queue:

Level 1: `12`

Level 2: `7` `1`

Level 3: `9` `10` `5`

Move "three" elements to the the output array representing third level.

Code

Here is what our algorithm will look like:

| Java	Python3	C++	JS

```python
from collections import deque


class TreeNode:
  def __init__(self, val):
    self.val = val
    self.left, self.right = None, None


def traverse(root):
  result = []
  if root is None:
    return result

  queue = deque()
  queue.append(root)
  while queue:
    levelSize = len(queue)
    currentLevel = []
    for _ in range(levelSize):
      currentNode = queue.popleft()
      # add the node to the current level
      currentLevel.append(currentNode.val)
      # insert the children of current node in the queue
      if currentNode.left:
        queue.append(currentNode.left)
      if currentNode.right:
        queue.append(currentNode.right)
```

Time complexity #

The time complexity of the above algorithm is $O(N)$, where 'N' is the total number of nodes in the tree. This is due to the fact that we traverse each node once.

Space complexity #

The space complexity of the above algorithm will be $O(N)$ as we need to return a list containing the level order traversal. We will also need $O(N)$ space for the queue. Since we can have a maximum of $N/2$ nodes at any level (this could happen only at the lowest level), therefore we will need $O(N)$ space to store them in the queue.

✓ Mark as Completed

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