## MTH 4320 Homework 3

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### Problem 1

Solution. Let the graph on the left be G and the graph of the drawing be H. Notice that the set of vertices  $\{a,c,f,g,h\}$  in G is an independent set, where no two vertices in the set are adjacent. We can label the vertices in H as  $\{a,c,f,h\}$  on the left side and  $\{g,e,d,b\}$  on the right side from top to bottom respectively. Then we observe that vertex a is adjacent to the vertices  $\{b,d,e\}$  denoted by  $a \sim \{b,d,e\}$ . Similarly, we observe that  $c \sim \{b,d,g\}, f \sim \{b,e,g\}$ , and  $h \sim \{d,e,g\}$ . Hence we have that any two vertices with the same labels are adjacent in G if and only if they are adjacent in G. Therefore, it is shown that G and G are identical or  $G \cong H$ , G is isomorphic to G.

### Problem 2

Solution. The algorithm is

- 1. Label the vertices in graph G as  $V = \{v_1, \ldots, v_n\}$ . This takes O(|V|) time.
- 2. For every vertex  $v_i$  in V starting from  $v_1$ 
  - Assign a direction to every edge incidental to  $v_i$  that has not been assigned a direction.
  - If the other endpoint has a greater value of *i* then we assign the edge to that direction else we assign it to the opposite direction.
- 3. Labeling all of the edges takes O(|E|) time.

The algorithm make sure that the directions of the edges of every vertex go to the adjacent vertices with a label of higher value. Therefore, there must not be any cycles in the directed graph. The running time of the algorithm is O(|V|) + O(|E|) = O(|V| + |E|).

# Problem 3

Solution. For a binary sequence of length k = 3, G can be a cube with 8 vertices being the corners having edges connecting the other 3 adjacent corners. In general, for any positive

integer n we have that G(n) is a connected graph with only one component. We can prove this proposition using proof by induction. Let the statement be P(k) for all positive integer k where k is the length of the binary sequence.

Base case: For k = 1 we have two vertices 0 and 1 connected by an edge so P(1) is true. Induction steps: Assume that P(k) is true so G(k) is connected. When we add one more bit to the binary sequence of G(k), we can choose any position and insert the same bit to maintain the graph G(k). Then we can change the new bit of each new binary sequence to create some new vertices connected to the existing vertices of G(k) to form the graph G(k+1). Hence G(k+1) is constructed from G(k) and G(k) is connected so G(k+1) is connected and G(k+1) is true.

Therefore, it is proved by mathematical induction that for any positive integer n the graph G(n) is connected so it has only one component.