

Review

* Markov decision process (MDP)

$$\text{MDP} = \{ \mathcal{S}, \mathcal{A}, P(\mathcal{S}' | \mathcal{S}, a), R(\mathcal{S}) \}$$

states, actions, transitions, rewards

* Policy = deterministic mapping $\pi : \mathcal{S} \rightarrow \mathcal{A}$

* Value functions

$$V^\pi(\mathcal{S}) = E^\pi \left[\sum_{t=0}^{\infty} \gamma^t R(\mathcal{S}_t) \mid \mathcal{S}_0 = \mathcal{S} \right] \quad (\text{state})$$

$$Q^\pi(\mathcal{S}, a) = E^\pi \left[\sum_{t=0}^{\infty} \gamma^t R(\mathcal{S}_t) \mid \mathcal{S}_0 = \mathcal{S}, a_0 = a \right] \quad (\text{action})$$

* Policy evaluation

Solve linear equations:

$$\sum_{\mathcal{S}'} \left[\mathbb{I}(\mathcal{S}, \mathcal{S}') - \gamma P(\mathcal{S}' | \mathcal{S}, \pi(\mathcal{S})) \right] V^\pi(\mathcal{S}') = R(\mathcal{S})$$

* Policy improvement

greedy policy $\pi'(\mathcal{S}) = \underset{a}{\operatorname{argmax}} Q^\pi(\mathcal{S}, a)$

Theorem: $V^{\pi'}(\mathcal{S}) \geq V^\pi(\mathcal{S})$ for all states \mathcal{S}

* Policy iteration

$$\pi_0 \xrightarrow[\text{improve}]{\text{evaluate}} \frac{V^{\pi_0}(\mathcal{S})}{Q^{\pi_0}(\mathcal{S}, a)} \xrightarrow[\text{improve}]{\text{evaluate}} \pi_1 \xrightarrow[\text{improve}]{\text{evaluate}} \frac{V^{\pi_1}(\mathcal{S})}{Q^{\pi_1}(\mathcal{S}, a)} \xrightarrow[\text{improve}]{\text{evaluate}} \pi_2 \rightarrow \dots$$

* Is policy iteration guaranteed to converge? yes

* Does it always converge to an optimal policy π^* ? yes.

* Theorem: Suppose $\pi'(\mathcal{S}) = \pi(\mathcal{S})$ for all states \mathcal{S}

(or even more generally, that $V^{\pi'}(\mathcal{S}) = V^\pi(\mathcal{S})$)

Then: $V^\pi(\mathcal{S}) = V^*(\mathcal{S})$.

Note: optimal value function $V^*(\mathcal{S})$ is unique, even if there are many optimal policies.

* proof strategy:

1) Derive "Bellman optimality equation"

Satisfied by $V^{\pi^*}(s)$ when $V^{\pi^*}(s) = V^{\pi}(s)$.

2) show that $V^{\pi}(s) \geq V^{\tilde{\pi}}(s)$ for all policies $\tilde{\pi}$ and states s in MDP.

Hence: $V^{\pi}(s) = V^*(s)$

Step 1.

From Bellman equation for $\pi'(s)$

$$V^{\pi'}(s) = R(s) + \gamma \sum_{s'} p(s'|s, \pi'(s)) V^{\pi'}(s')$$

By assumption, $V^{\pi'}(s) = V^{\pi}(s)$ at convergence

$$\text{Hence: } V^{\pi}(s) = R(s) + \gamma \sum_{s'} p(s'|s, \pi'(s)) V^{\pi}(s')$$

By assumption, $\pi'(s)$ is greedy w.r.t. $V^{\pi}(s)$.

$$\text{Hence, } \boxed{V^{\pi}(s) = R(s) + \gamma \max_a \sum_{s'} p(s'|s, a) V^{\pi}(s')}$$

"Bellman optimality equation"

(set of n non-linear equations for $s=1, 2, \dots, n$)

non-linear b/c max operation is not linear.

* different than linear Bellman equation

Step 2

Iterate right hand side:

$$V^{\pi}(s) = R(s) + \gamma \max_a \sum_{s'} p(s'|s, a) \left[R(s') + \gamma \max_{a'} \sum_{s''} p(s''|s', a') V^{\pi}(s'') \right]$$

$\nwarrow V^{\pi}(s')$

Iterate again and again

Now show that this iterated expression (taken out an infinite # terms) implies optimality.

Let $\tilde{\pi}(s)$ be any other policy with Bellman equation:

$$\begin{aligned}
 V^{\tilde{\pi}}(s) &= R(s) + \gamma \sum_{s'} P(s'|s, \tilde{\pi}(s)) V^{\tilde{\pi}}(s') \\
 &\leq R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V^{\tilde{\pi}}(s') \quad \left. \begin{array}{l} \text{"be greedy"} \\ \text{"use Bellman equation"} \end{array} \right\} \\
 &= R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) [R(s') + \gamma \sum_{s''} P(s''|s', \tilde{\pi}(s')) V^{\tilde{\pi}}(s'')] \\
 &\leq R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) [R(s') + \gamma \max_{a'} \sum_{s''} P(s''|s', \tilde{\pi}(s')) V^{\tilde{\pi}}(s'')] \quad \left. \begin{array}{l} \text{greedy} \end{array} \right\}
 \end{aligned}$$

Consider upper bound on $V^{\tilde{\pi}}(s)$ from iterating above t times (being greedy, then applying Bellman equation)

Compare this to equality after t iterations for $V^{\pi}(s)$.

As $t \rightarrow \infty$, RHS of upper bound on $V^{\tilde{\pi}}(s)$ converges to RHS of equality for $V^{\pi}(s)$

Thus as $t \rightarrow \infty$:

$$V^{\tilde{\pi}}(s) \leq \lim_{t \rightarrow \infty} [] = \lim_{t \rightarrow \infty} [] = V^{\pi}(s)$$

Thus for all policies $\tilde{\pi}$ and states s , we have

$$V^{\pi}(s) \geq V^{\tilde{\pi}}(s)$$

$$V^{\pi}(s) = \max_{\tilde{\pi}} V^{\tilde{\pi}}(s) \quad \text{or} \quad V^{\pi}(s) = V^*(s).$$

To compute π^* :

$$\begin{aligned}
 \pi^*(s) &= \operatorname{argmax}_a Q^*(s, a) \\
 &= \operatorname{argmax}_a \sum_{s'} P(s'|s, a) V^*(s')
 \end{aligned}$$

pros/cons of policy evaluation:

(+) Converges quickly (in few steps)

(-) Each step requires policy evaluations $O(n^3)$

Value iteration

* How to compute $V^*(s)$ directly?

$$\begin{aligned} V^*(s) &= \max_a Q^*(s, a) \\ &= \max_a \left[R(s) + \gamma \sum_{s'} p(s'|s, a) V^*(s') \right] \end{aligned}$$

$$V^*(s) = R(s) + \gamma \max_a \sum_{s'} p(s'|s, a) V^*(s')$$

* n nonlinear equations for n unknowns $V^*(s)$ for $s=1, \dots, n$
How to solve?

* Algorithm: Value iteration.

(1) Initialize $V_0(s) = 0$ for all s

(2) Iterate

$$V_{k+1}(s) = R(s) + \gamma \max_a \left[\sum_{s'} p(s'|s, a) V_k(s') \right]$$

for all $s = 1, 2, \dots, n$

current estimate of $V^*(s')$
at k -th iteration



Note: this algorithm works directly on value functions,
no policies.

But incremental policies can be computed from:

$$\begin{aligned} \pi_{k+1}(s) &= \text{greedy}[V_k(s)] \\ &= \arg\max_a \left[\sum_{s'} p(s'|s, a) V_k(s') \right] \end{aligned}$$

(3) Suppose this converges: $\lim_{k \rightarrow \infty} V_k(s) = V^*(s)$

$$\text{then compute } \pi^*(s) = \arg\max_a \left[\sum_{s'} p(s'|s, a) V^*(s') \right]$$

Does algorithm converge?

Clearly, $V^*(s)$ is fixed point of iteration. But are there other fixed points? No. Does it always reach $V^*(s)$? yes.

* Lemma:

for any functions $f(a)$ and $g(a)$:

$$\left| \max_a f(a) - \max_a g(a) \right| \leq \max_a |f(a) - g(a)|$$

Proof of lemma:

$$\text{for all } a: f(a) - \max_{a'} g(a') \leq f(a) - g(a)$$

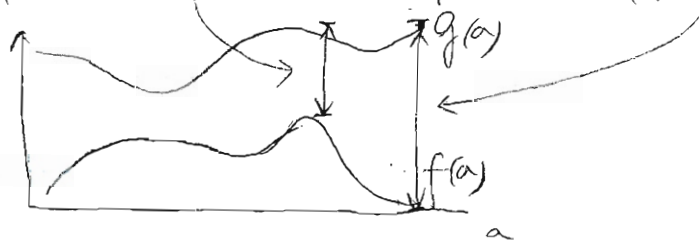
$$\begin{aligned} \text{max over } a: \max_a f(a) - \max_a g(a) &\leq \max_a [f(a) - g(a)] \\ &\leq \max_a |f(a) - g(a)| \end{aligned}$$

By symmetry, exchanging $f \leftrightarrow g$ everywhere:

$$\max_a g(a) - \max_a f(a) \leq \max_a |g(a) - f(a)|$$

Combining last two inequalities:

$$\left| \max_a f(a) - \max_a g(a) \right| \leq \max_a |g(a) - f(a)|$$



Thm: Value iteration Converges.

$$\lim_{k \rightarrow \infty} [V_k(s)] \rightarrow V^*(s) \text{ for all states } s$$

Proof: let $\Delta_k = \max_s |V_k(s) - V^*(s)|$ error at k -th iteration

$$\begin{aligned} \Delta_{k+1} &= \max_s |V_{k+1}(s) - V^*(s)| \\ &= \max_s \left| \left[\cancel{R(s)} + \gamma \max_a \sum_{s'} p(s'|s, a) V_k(s') \right] - \left[\cancel{R(s)} + \gamma \max_a \sum_{s'} p(s'|s, a) V^*(s') \right] \right| \end{aligned}$$

plug in def for value iteration and Bellman optimality equation for $V^*(s)$.

(cont.)

$$\Delta_{k+1} = \gamma \max_s \left| \max_a \underbrace{\sum_{s'} P(s'|s, a) V_k(s')}_{f(a)} - \max_a \underbrace{\sum_{s'} P(s'|s, a) V^*(s')}_{g(a)} \right|$$

apply lemma: $\left| \max_a f(a) - \max_a g(a) \right| \leq \max_a |f(a) - g(a)|$

$$\begin{aligned} \Delta_{k+1} &\leq \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) [V_k(s') - V^*(s')] \right| \\ &\leq \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) \left(\max_{s''} |V_k(s'') - V^*(s'')| \right) \right| \\ &= \gamma \max_s \max_a \left| \left\{ \sum_{s'} P(s'|s, a) \right\} \Delta_k \right| \quad \text{Worst case bound on difference.} \\ &= \gamma \Delta_k \max_s \max_a (1) \\ &= \gamma \Delta_k \end{aligned}$$

Hence: $\Delta_{k+1} \leq \gamma \Delta_k$

By iteration: $\Delta_k \leq \gamma^k \Delta_0 \rightarrow 0$ as $k \rightarrow \infty$ for $\gamma < 1$

Assume rewards are bounded:

$$\begin{aligned} \Delta_0 &= \max_s |V_0(s) - V^*(s)| = \max_s |V^*(s)| \\ &\leq \left[\max_s |R(s)| \right] (1 + \gamma + \gamma^2 + \gamma^3 + \dots) \\ &= \max_s |R(s)| \frac{1}{1-\gamma} \end{aligned}$$

Thm: $\Delta_k \leq \left(\frac{\gamma^k}{1-\gamma} \right) \max_s |R(s)| \rightarrow 0$ as $k \rightarrow \infty$

convergence rate depends on γ .

Suggests that more iterations are required as $\gamma \rightarrow 1$.