

# *Algorithm Design*

## *(Algorithms as High-Order-Primitives)*

### *Video 6.3e*

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*Algorithm is Cool. Learn Algorithms.*

# Quick review

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**Problem-1: Algorithm to Compute the sum of  
(1 + 2 + 3 + ... + 99 + 100)**

Give no-brainer calculated algorithm  
(BAD-Sum-to-Hundred)

Evolved to first algorithm with **a loop**  
**(Sum-1-to-100)**  
(given in flowchart and in pseudo-code)

# Algorithm Sum-1-to-100

**Problem-1: Algorithm to Compute the sum of  
(1 + 2 + 3 + ... + 99 + 100)**

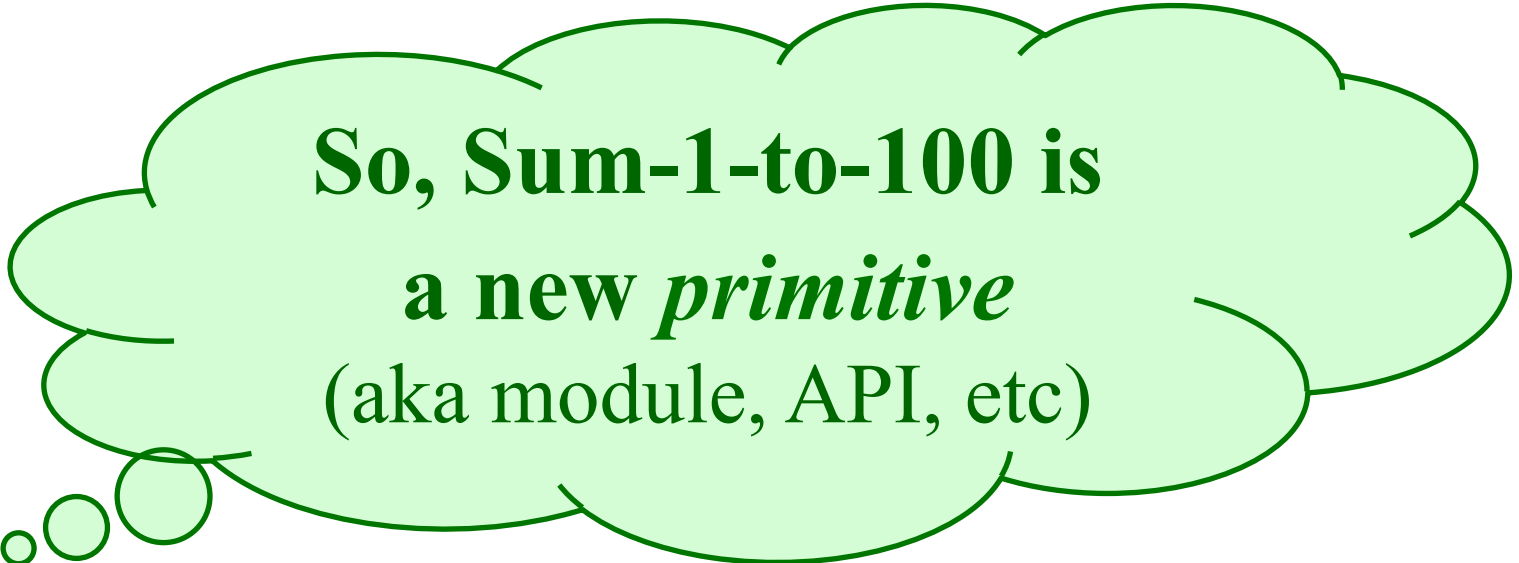
## ALGORITHM Sum-1-to-100;

1. Let  $\text{Sum} \leftarrow 0$  ;
2. Let  $k \leftarrow 1$  ;
3. While ( $k \leq 100$ ) repeat Steps 4-6
4.    $\text{Sum} \leftarrow \text{Sum} + k$
5.    $k \leftarrow k + 1$
6. end-of-while-block;
7. Print out the value of Sum
8. End

Algorithm Sum-1-to-100  
(in pseudo-code)

# An algorithm is a new *primitive*

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So, Sum-1-to-100 is  
a new *primitive*  
(aka module, API, etc)

Anyone can use Sum-1-to-100 to solve Prob-1.

If implemented in code, aka **module**,  
**component**, or **API**, *sharable* with others.

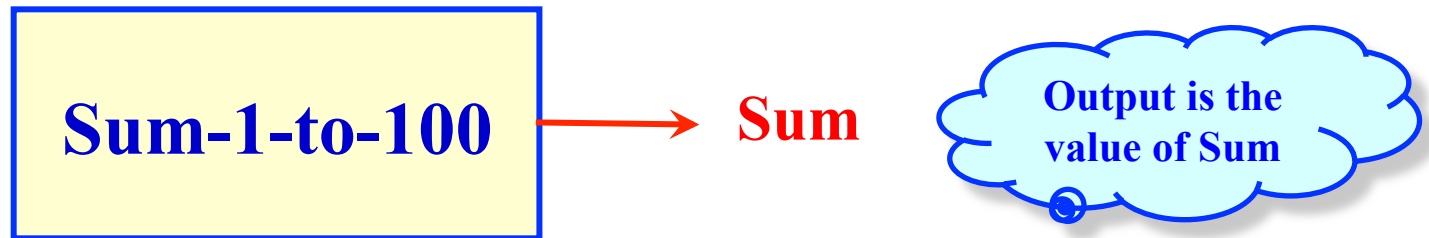
# Define new primitive carefully

**So**, we have an algorithm for computing sum of integers from 1 to 100.

We call it **Sum-1-to-100**

**Definition:** A new primitive called **Sum-1-to-100**

The (high-level) primitive **Sum-1-to-100** computes the sum of integers 1 to 100 and returns the total via variable **sum**.



# Analysis of Sum-1-to-100 primitive

**Definition:** A new primitive called **Sum-1-to-100**  
The (high-level) primitive **Sum-1-to-100** computes the sum of integers 1 to 100 and returns the total via variable sum.

**Primitive Sum-1-to-100 has no input... WHY?**

If I execute/run Sum-1-to-100 many time,  
it always produce the same answer. (5050)

Does not matter what “input” you give to it.

**Occam's Razor:**

Fewest parameters rules.

In this case, 0 parameters!

**Abstraction**

Module: CT,

# How to make this more useful?

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## Reformulate the Problem Evolve to Problem-2

**Problem-2:** Want to compute the sum of integers from 1 to  $n$ ?  
Namely, calculate  $(1 + 2 + 3 + \dots + n)$

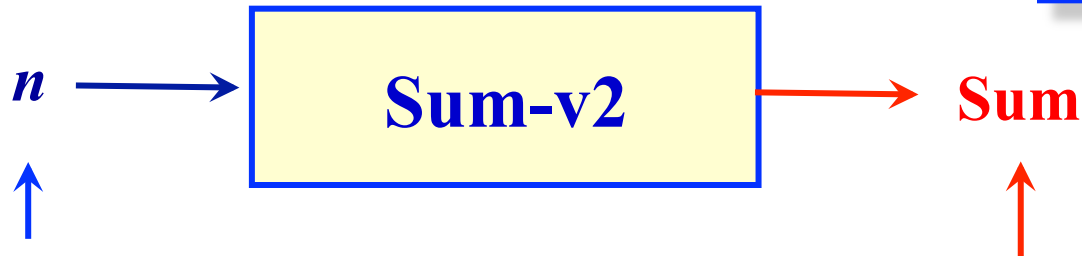
(But we want to specify  $n$  later!)

**Why is this useful?**

# Call it primitive **Sum-v2( $n$ )**

- Sum-v2( $n$ ) is a *high-level primitive* with an input parameter  $n$

**Abstraction**



Input to **Sum-v2**: variable  $n$

Output is variable **Sum**

## **Definition: Sum-v2 ( $n$ )**

The high-level primitive Sum-v2 takes in as input a integer  $n$ , and it computes and returns the sum of  $(1 + 2 + \dots + n)$ .



# Sum-v2( $n$ ) more useful?

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**We use it to compute many different sums**

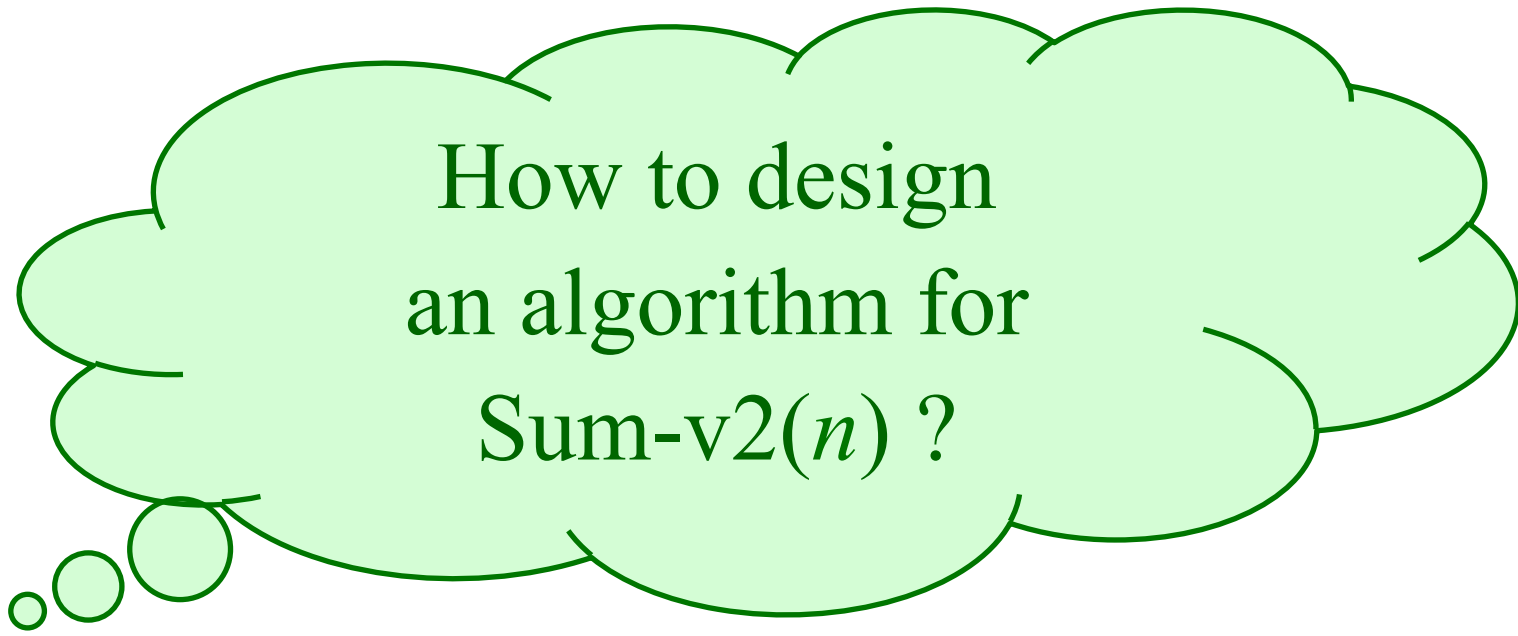
- just by sending different value of  $n$

**Examples:**

- Sum-v2(100) gives  $(1+2+3+\dots+100)$
- Sum-v2(24) gives  $(1+2+3+\dots+24)$
- Can also call Sum-v2(1024)

**SO, Sum-v2( $n$ )** is more general than Sum-1-to-100.

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How to design  
an algorithm for  
 $\text{Sum-v2}(n)$  ?

# Appeal to Polya Step-2

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**PQ: Have we seen a similar problem?**

**PQ: Can we reuse the result** (black box reuse)  
**or the method** (white box reuse).

**Answer: Can reuse Sum-1-to-100 ?**

**White box reuse:** Reuse the algorithm (method) for Sum-1-to-100 with minor changes

**Black box reuse:** Reuse (result) without any changes to the primitive Sum-1-to-100

# White box reuse (small change)

Algorithm Sum-1-to-100

## ALGORITHM Sum-1-to-100;

1. Let  $\text{Sum} \leftarrow 0$  ;
2. Let  $k \leftarrow 1$  ;
3. While ( $k \leq 100$ ) repeat Steps 4-6
4.    $\text{Sum} \leftarrow \text{Sum} + k$
5.    $k \leftarrow k + 1$
6. end-of-while-block;
7. Print out the value of Sum
8. End

Algorithm Sum-v2( $n$ )

## ALGORITHM Sum-v2( $n$ );

1. Let  $\text{Sum} \leftarrow 0$  ;
2. Let  $k \leftarrow 1$  ;
3. While ( $k \leq n$  ) repeat Steps 4-6
4.    $\text{Sum} \leftarrow \text{Sum} + k$
5.    $k \leftarrow k + 1$
6. end-of-while-block;
7. Print out the value of Sum
8. End

# Black box reuse (no change)


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**Q: Can we reuse Sum-1-to-100 as black box to help solve Problem-2 ?**

**Answer: NO.**

**Every time we use (run) Sum-1-to-100, it always gives 5050.**

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**Problem  
Reformulation  
again.**

# Reformulate again...

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## Reformulate the Problem Evolve to Problem-3

**Problem-3:** We want Sum-Range( $p$ ,  $q$ ) that computes the sum of integers from  $p$  to  $q$ .  
Namely,  $(p + (p+1) + \dots + q)$

Eg: Sum-Range(25,100) =  $(25 + 26 + \dots + 100)$

**PQ: Can we reuse the result** (black box reuse)  
**or the method** (white box reuse).

# White box reuse

**White box reuse:** Reuse the algorithm (method) for Sum-v2(n) with minor changes to get algorithm for Sum-Range(p,q)

## ALGORITHM Sum-v2(n);

1. Let  $\text{Sum} \leftarrow 0$  ;
2. Let  $k \leftarrow 1$  ;
3. While ( $k \leq n$ ) repeat Steps 4-6
4.    $\text{Sum} \leftarrow \text{Sum} + k$
5.    $k \leftarrow k + 1$
6. end-of-while-block;
7. Print out the value of Sum
8. End

**Your DIY HW.**  
(you remember  
what is HW?)



# Black box reuse of Sum-v2( $n$ )

**Black box reuse: Reuse without any changes to the primitive Sum-v2( $n$ )**

**Example: Sum-Range(25,100)**

- Sum-v2(100) gives (**1+2+...+24**+25+...100)
- Sum-v2(24) gives (**1+2+...+24**)

So, Sum-v2(100) – Sum-v2(24) gives  
Sum-Range(25,100)

**COOL!**  
**Did not change  
any code!**

$$\text{Sum-Range}(p,q) = \text{Sum-v2}(q) - \text{Sum-v2}(p-1)$$

# Summary of 6.3d, 6.3e

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- ❑ Design first algorithm with a loop;
  - ❖ Magic power of *loop (iterations)*
- ❑ Algorithms for 3 related problems
- ❑ Learned how to re-use algorithms
  - ❖ White box reuse (modify method slightly)
  - ❖ Black box reuse (cannot change method)
  - ❖ Different thinking skills are needed

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***(End of video 6.3e)***

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