Sorting

```
Click the sequence of integers that represent the content of the array A=[42, 8, 26, 35, 23, 32, 12, 39] with n integers after 5 passes of this version of Insertion Sort.

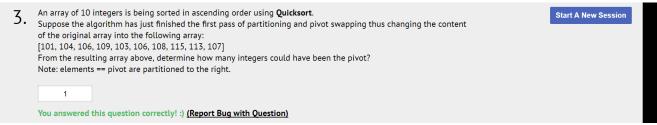
for (i = 1; i <= 5; i++) { // 5 passes}
e = A[i]; j = i;
while (j > 0) {
if (A[j-1] > e)
A[j] = A[j-1];
else
break;
j--;
}

A[j] = e;
}

Your answer is: 8, 23, 26, 32, 35, 42, 12, 39

You answered this question correctly!:) (Report Bug with Question)
```

First n+1 numbers will be sorted. Insertion sort is **STABLE**, so the rest of the elements is unaffected.



Find number such that every number to the left is lesser than (or equal to) itself and every number to the right is greater than (or equal to) itself

```
How many swap(s) is/are required to sort an array of n=6 integers: [16, 28, 9, 13, 26, 15] using this version of Bubble Start A New Session

For (j = 0; j < n-1; j++)
for (i = 0; i < n-j-1; i++)
if (A[i] > A[i+1]);

swap(A[i], A[i+1]);

8

You answered this question correctly!:) (Report Bug with Question)
```

Brute force count swaps

```
Click the sequence of integers that represent the content of the array A=[6, 38, 18, 22, 20, 19, 29, 34] with n integers after 4 passes of this version of Selection Sort.

for (i = 0; i < 4; i++) { // 4 passes

cur_min = i;

for (j = i+1; j < n; j++)

if (A[j] < A[cur_min])

cur_min = j;

swap (A[i], A[cur_min]);
}

Your answer is: 6,18,19,20,38,22,29,34

You answered this question wrongly.

The correct answer is: 6,18,19,20,22,38,29,34
```

Selection sort is **UNSTABLE**. Keep track of where elements are swapped to.

```
Click the sequence of integers that represent the content of the array A=[38, 4, 6, 5, 29, 34, 32, 26] with n integers after 7 passes of this version of Bubble Sort.

for (j = 0; j < 7; j++) // 7 passes

for (i = 0; i < n-j-1; i++)

if (A[i] > A[i+1])

swap (A[i], A[i+1]);

Your answer is: 4,5,6,26,29,32,34,38

You answered this question correctly!:) (Report Bug with Question)
```

Biggest 7 elements are bubbled right. So, this sequence of 8 integers will be fully sorted lol

Become the algorithm

```
How many pass(es) is/are required to sort an array of n=5 integers: [35, 26, 37, 9, 7] using this version of Bubble

Sort?

j = 0;
do {
    swapped = false; j++;
    for (i = 0; i < n-j; i++) // each completion of this for-loop is counted as one pass
    if (A[i] > A[i+1]) { // each execution of this if-statement is counted as one comparison
        swapped = true;
    }
} while (swapped);

4

You answered this question wrongly.

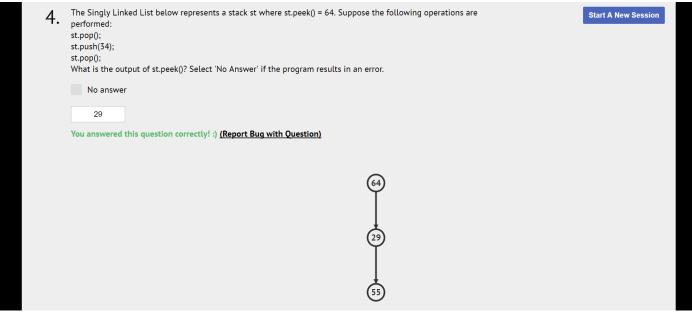
The correct answer is: 5
```

Become the algorithm. This is the "smart" version which breaks if no new swaps are made in the pass-through.

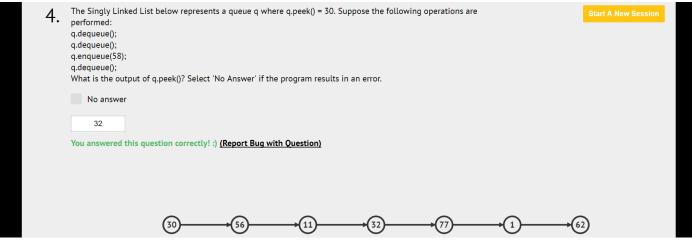
Linked list



Just be careful. Tail pointer is usually not updated. Final_element.next will lead to NULL (still valid), but attempting to NULL.next or dereference NULL will lead to compilation error -> NO ANSWER



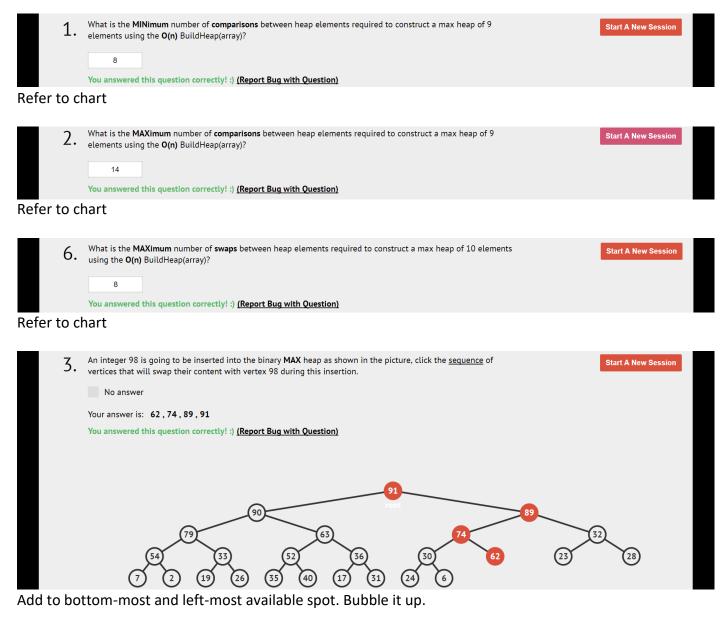
Add and remove stuff from the TOP (64 is the element at the top in this example)

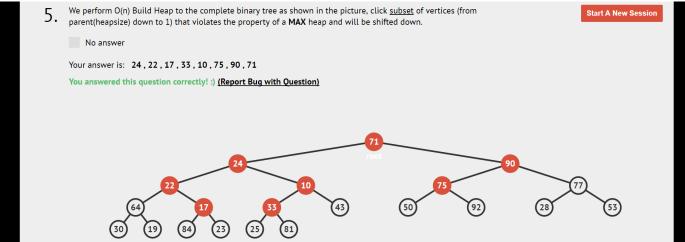


Enqueue adds at the RIGHT, removes from the LEFT

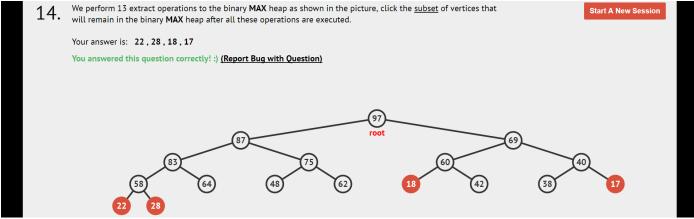
Binary Heap

Valid Max Heap: Make sure parent bigger than child

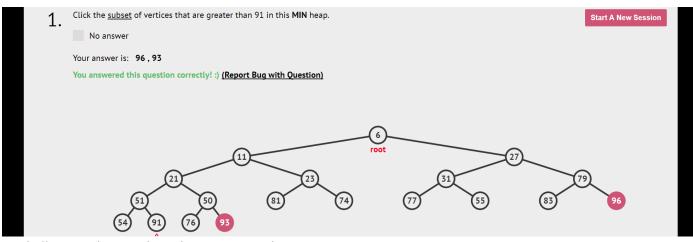




Vertices that are smaller than any of their descendants will be shifted down.



Smallest 17 – 13 = 4 elements are remaining



Find all vertex bigger than the query number

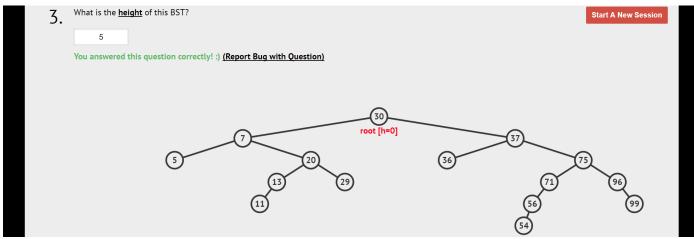
Binary Search Tree

BST: Left child < current < Right child.

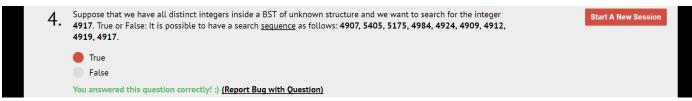
Successor: Most direct way to the next element bigger than the node of interest Predecessor: Most direct way to the next element smaller than the node of interest

Preorder: Print first, then check children

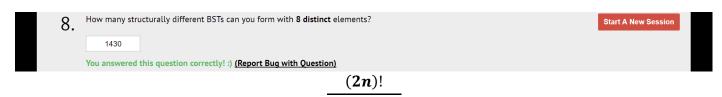
Inorder: Check left, **then print**, then check right **Post**order: Check children first, **then print**.

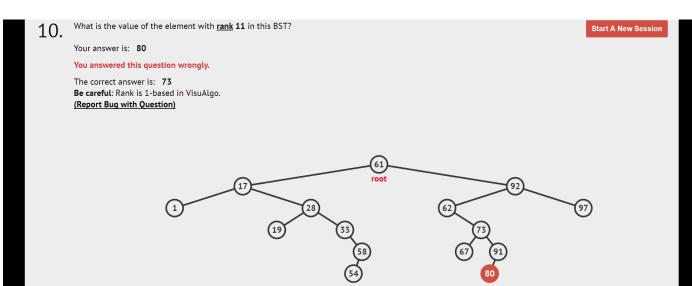


Largest number of edges from root to leaf



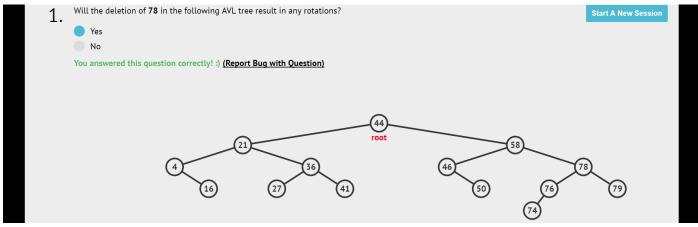
If target is bigger, next number should be bigger; if target smaller, next number should be smaller.



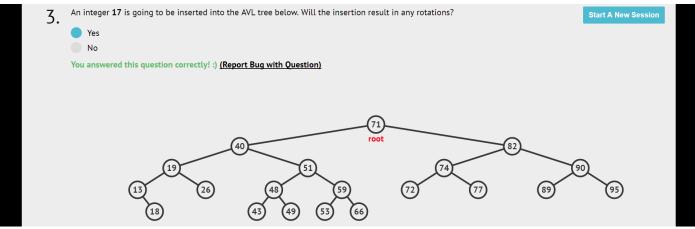


indexed-1. Leftmost is 1. Flatten the tree and count until rank.

AVL Tree



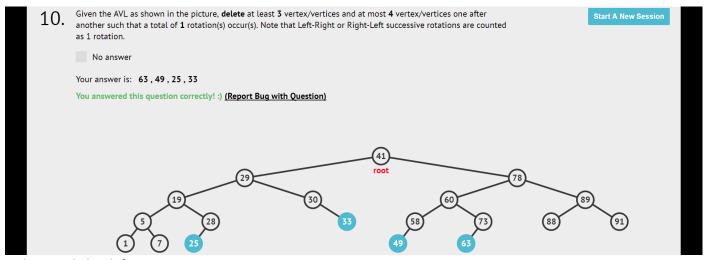
Search upwards starting from the replaced node all the way to root and see if anyone complains



Insert into correct position and check for rotations. Max 1 "fix" required for insertion.



Height Vertices



Delete and check for rotation.

Graph DS

Directed acyclic, just make sure small number -> bigger number.

Tree: V vertices, V-1 edges, acyclic.

Complete graph: An edge between any pair of vertices.

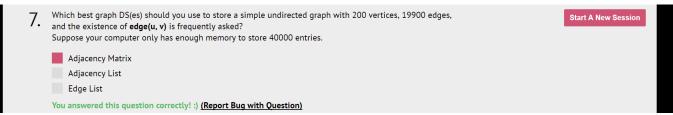
Connected graph: From any vertex, can get to all other vertex

2.	Which best graph DS(es) should you use to store a simple undirected graph with 10000 vertices, 10000 edges, and the neighbours are frequently enumerated? Suppose your computer only has enough memory to store 40000 entries.	New Session
	Adjacency Matrix	
	Adjacency List	
	Edge List	
	You answered this question correctly! :) (Report Bug with Question)	
 _		

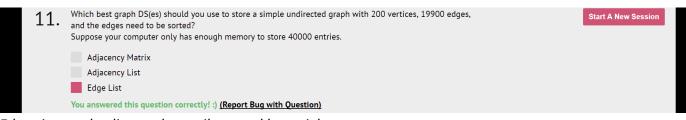
Need to enumerate neighbours, **BUT NOT ENOUGH MEMORY (we need V^2 entries for AM)**

3.	Which best graph DS(es) should you use to store a simple undirected graph with 200 vertices, 19900 edges, and the neighbours are frequently enumerated? Suppose your computer only has enough memory to store 40000 entries.	Start A New Session
	Adjacency Matrix Adjacency List Edge List	
	You answered this question correctly! :) (Report Bug with Question)	

Adjacency Matrix enumerate not as fast but still accepted, so it depends on memory.

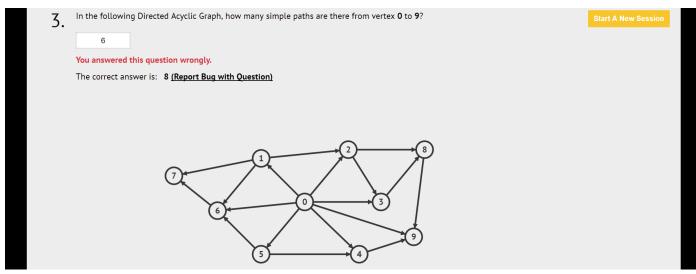


Ask about existence of edge -> AM (Good for Floyd-Warshall's 4 line wonder :D)

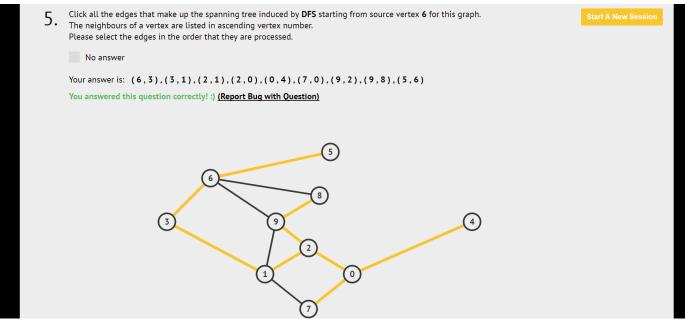


Edges in an edge list can be easily sorted by weight

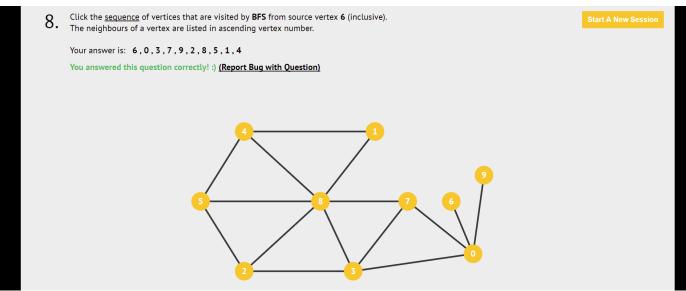
Graph traversal



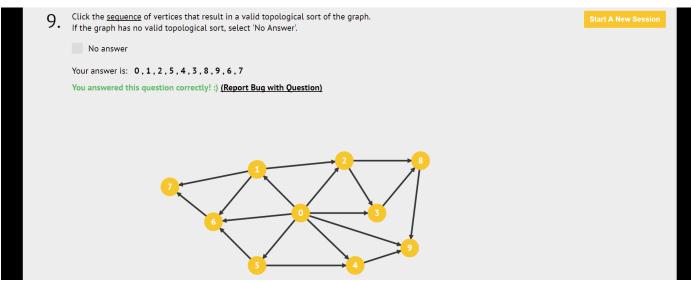
Write down all the paths



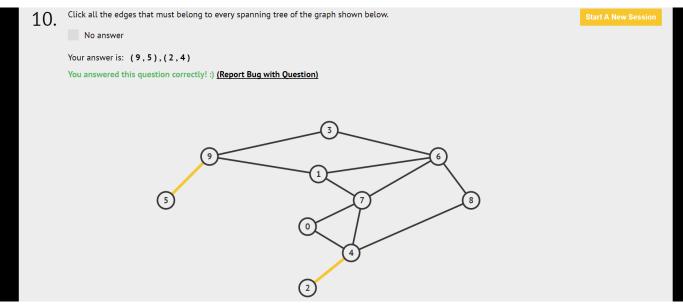
DFS: Go deep, recurse back and go deep for neighbour. Take node of order of neighbours.



BFS: Visit all neighbours, then all neighbours of neighbours, and so on. Keep track of the order of exploration of neighbours



Topological sort



Vertex with only 1 edge

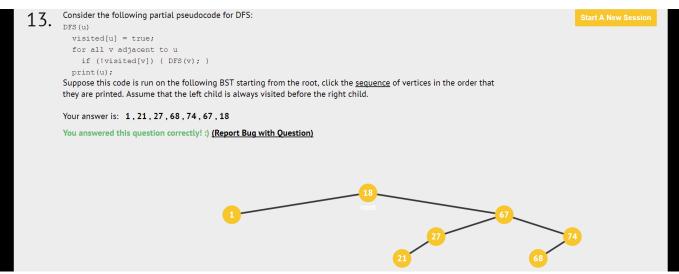
 ${\bf 11.} \ \ {}^{\hbox{\scriptsize How many different spanning trees are there in a complete graph with 8 vertices?}}$

Start A New Session

262144

You answered this question correctly! :) (Report Bug with Question)

 n^{n-2}

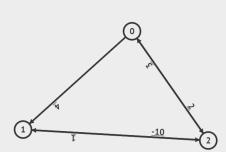


Print comes after exploring all neighbours -> Post-order traversal



1. Draw a simple connected weighted directed graph with 3 vertices and 5 directed edges such that running Modified Dijkstra's algorithm from source vertex 0 makes the algorithm run indefinitely. The weights of the edges have to be distinct.

You answered this question correctly! :) (Report Bug with Question)



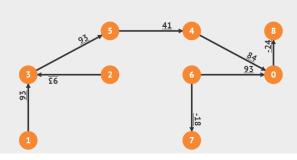
Create a negative cycle. MAKE SURE THE WEIGHTS ARE DISTINCT

7. Click the <u>sequence</u> of vertices such that when all the outgoing edges of these vertices are relaxed in this order using One-Pass Bellman-Ford's algorithm, the SSSP problem can be solved in O(V+E) time.

Start A New Session

Your answer is: 6,7,1,2,3,5,4,0,8

You answered this question correctly! :) (Report Bug with Question)

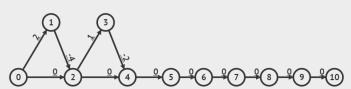


Topological sort

Draw a simple connected weighted directed graph with 11 vertices and at most 12 edges such that running Modified Dijkstra's algorithm from source vertex 0 successfuly relaxes ≥31 edges to get the correct shortest paths. We count 1 successful relaxation if relax(u, v, w_u_v) decreases D[v]. Your graph cannot contain a negative weight cycle.

Start A New Session

You answered this question correctly! :) (Report Bug with Question)



Dijkstra Killer. THERE ARE (E-V+1) TRIANGLES.

Draw a simple connected weighted directed graph with 7 vertices such that running
Optimized Bellman-Ford's algorithm from source vertex 0 uses at least 3 rounds to get the correct shortest paths.
Your graph cannot contain a negative weight cycle.
Note that all edges are processed according to the Adjacency List
i.e. all outgoing edges from vertex 0 is processed then edges from vertex 1 and so on.

Start A New Session

You answered this question correctly! :) (Report Bug with Question)



Bellman-Ford Killer. 0 -> (N-1) -> (N-2) -> ... -> 2 -> 1

SSSP- Path Weight Criteria

	Bellman-Ford	Original Dijkstra	Modified Dijkstra
Terminate	Always	Always	Does not terminate
			when there's negative
			weight cycle
Wrong when	Negative eight cycle	Negative edges and	Negative weight cycle- it
there's		Negative weight cycle	doesn't even terminate
			in the first place 🙁

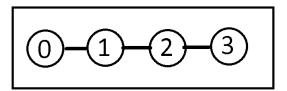
Note:

If the graph contains negative edges but does not contain any negative cycles, Bellman Ford and modified Dijkstra will give the correct answer all the time.

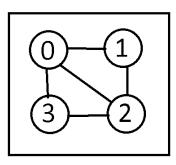
Original Dijkstra *could* still give the correct answer (run Bellman Ford and original Dijkstra and compare the answers)

Bonus: Graph creation-exploration with 4 nodes

Draw a connected undirected unweighted <u>tree</u> with 4 vertices such that running either DFS/BFS from the source 0 will result in the same sequence of visited vertices.



Draw a connected undirected unweighted graph with 4 vertices such that running either DFS/BFS from the source 0 will result in the same sequence of visited vertices. This graph cannot be a tree.



Tables

Max Heap max swaps

Elements	6	7	8	9	10	11	12	13	14
Swaps	4	4	7	7	8	8	10	10	11

Max Heap max comparisons

Elements	6	7	8	9	10	11	12	13	14
Compare	7	8	11	14	15	16	18	20	21

Max Heap Min Comparisons: N-1

Elements	6	7	8	9	10	11	12	13	14
Compare	5	6	7	8	9	10	11	12	13

Binary Search Tree: how many permutations

Elements	1	2	3	4	5	6	7	8	9
Catalan #	1	2	5	14	42	132	429	1430	4862

AVL Tree: Minimum # of vertices (Height(n) = height(n-1) + height(n-2) + 1)

						•	_	` '	•		,	•	•	•		
Height	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Vertices	1	2	4	7	12	20	33	54	88	143	232	376	609	986	1596	2583

Graph Structures

Neighbours Frequently Enumerated AND **HAS ENOUGH MEMORY: AL + AM** (We need V^2 memory for AM)

Neighbours Frequently Enumerated AND NOT ENOUGH MEMORY: AL ONLY (memory < V^2)

Existence of edge(u,v): AM ONLY

Edges need to be sorted: EDGE LIST ONLY

Number of spanning trees in complete graph n^(n-2)

Vertices	3	4	5	6	7	8	9	10	11
Ans	3	16	625	1296	16807	262144	4782969	100000000	2357947691

Path Weight Criteria

	Bellman-Ford	Original Dijkstra	Modified Dijkstra
Terminate	Always	Always	Does not terminate when there's negative weight cycle
Wrong when there's	Negative weight cycle	Negative weight (might still be correct if the negative edge doesn't mess up any of the nodes)	Negative weight cycle- it doesn't even terminate in the first place

Modulo Tables

%11

0	1	2	3	4	5	6	7	8	9	10
0	1	2	3	4	5	6	7	8	9	10
11	12	13	14	15	16	17	18	19	20	21
22	23	24	25	26	27	28	29	30	31	32
33	34	35	36	37	38	39	40	41	42	43
44	45	46	47	48	49	50	51	52	53	54
55	56	57	58	59	60	61	62	63	64	65
66	67	68	69	70	71	72	73	74	75	76
77	78	79	80	81	82	83	84	85	86	87
88	89	90	91	92	93	94	95	96	97	98
99	100	101	102	103	104	105	106	107	108	109

%13

0	1	2	3	4	5	6	7	8	9	10	11	12
0	1	2	3	4	5	6	7	8	9	10	11	12
13	14	15	16	17	18	19	20	21	22	23	24	25
26	27	28	29	30	31	32	33	34	35	36	37	38
39	40	41	42	43	44	45	46	47	48	49	50	51
52	53	54	55	56	57	58	59	60	61	62	63	64
65	66	67	68	69	70	71	72	73	74	75	76	77
78	79	80	81	82	83	84	85	86	87	88	89	90
91	92	93	94	95	96	97	98	99	100	101	102	103
104	105	106	107	108	109	110	111	112	113	114	115	116

%12

0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	5	6	7	8	9	10	11
12	13	14	15	16	17	18	19	20	21	22	23
24	25	26	27	28	29	30	31	32	33	34	35
36	37	38	39	40	41	42	43	44	45	46	47
48	49	50	51	52	53	54	55	56	57	58	59
60	61	62	63	64	65	66	67	68	69	70	71
72	73	74	75	76	77	78	79	80	81	82	83
84	85	86	87	88	89	90	91	92	93	94	95
96	97	98	99	100	101	102	103	104	105	106	107
108	109	110	111	112	113	114	115	116	117	118	119

%14

	0	1	2	3	4	5	6	7	8	9	10	11	12	13
	0	1	2	3	4	5	6	7	8	9	10	11	12	13
1	14	15	16	17	18	19	20	21	22	23	24	25	26	27
2	28	29	30	31	32	33	34	35	36	37	38	39	40	41
4	12	43	44	45	46	47	48	49	50	51	52	53	54	55
5	56	57	58	59	60	61	62	63	64	65	66	67	68	69
7	70	71	72	73	74	75	76	77	78	79	80	81	82	83
8	34	85	86	87	88	89	90	91	92	93	94	95	96	97
ç	98	99	100	101	102	103	104	105	106	107	108	109	110	111