CS2102 Tutorial 1: Relational Model Wk 3, Sem 2, 2020/21

- 1. **Solution:** In this question (and also Question 2), we are mainly inferring properties of the relation schema from an instance of the schema. Hence the superkeys/candidate keys identified here are generally only possibilities unless we have further information to derive more definite answers.
 - (a) Possible superkeys of R are $\{A,C\}$, $\{A,D\}$, $\{A,B,C\}$, $\{A,B,D\}$, $\{A,C,D\}$, and $\{A,B,C,D\}$.
 - (b) If $\{A,C\}$ is indeed a superkey of R, then $\{A,C\}$ must also be a candidate key of R since neither $\{A\}$ nor $\{C\}$ is a superkey of R (based on r). Moreover, any superkey which is a superset of $\{A,C\}$ can't be a candidate key of R. The two remaining possible superkeys, $\{A,D\}$ and $\{A,B,D\}$, are possible candidate keys of R. Note that at most one of $\{A,D\}$ and $\{A,B,D\}$ could be a candidate key.
- 2. Solution: The possible foreign keys are W, Y, and Z.

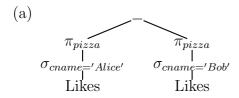
3. Solution:

- (a) Equivalent.
- (b) Not equivalent. Q_2 is an invalid relational algebra expression as the selection condition refers to a non-existent attribute.
- (c) Equivalent.
- (d) Equivalent.
- (e) Equivalent.

Note that binary operators in relational algebra are left associative. Hence, in the absence of parenthesis, the expression $R \times S \times T$ is evaluated as $(R \times S) \times T$. As the cross product operator is associative, $(R \times S) \times T$ is equivalent to $R \times (S \times T)$.

- (f) Equivalent.
- (g) Not equivalent. Consider the following database instance d: $r = \{(10, 10)\}$ and $s = \{(10, 20)\}$. The result of Q_1 on d is $\{(10)\}$ while the result of Q_2 on d is \emptyset .

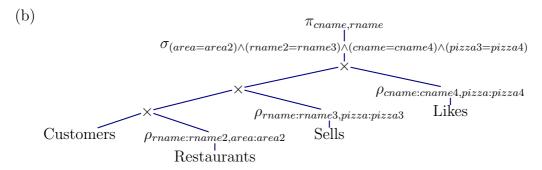
4. Solution:



The following answer is incorrect:

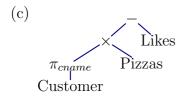
 $\pi_{pizza}(\sigma_{(pizza=pizza2)and(cname='Alice')and(cname2<>'Bob')}(Likes \times \rho_{cname:cname2.pizza:pizza2}(Likes)))$

The above answer will compute exactly the set of pizzas that Alices likes (independent of whether Bob likes them) since for each pizza p that Alice likes, we have $('Alice', p, 'Alice', p) \in Likes \times Likes$.

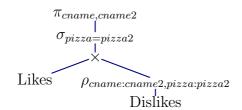


 $\pi_{cname,rname}(Customers \bowtie Restaurants \bowtie Sells \bowtie Likes)$

 $\pi_{cname,rname}(Customers \bowtie Restaurants) \cap \pi_{cname,rname}(Sells \bowtie Likes)$

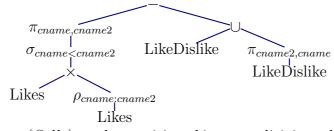


(d) Let Dislike(cname, pizza) denote the output computed by part (c).

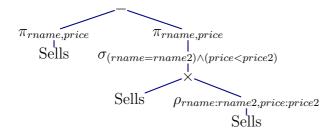


(e) Two customers C and C2 likes exactly the same set of pizzas if there does not exist any pizza that (1) C likes but C2 does not like or (2) C does not like but C2 likes.

Let LikeDislike(cname, cname2) denote the output computed by part (d).



(f) The set $\pi_{rname,price}(Sells)$ can be partitioned into two disjoint subsets $\pi_{rname,price}(Sells)$ = $S_1 \cup S_2$ where S_1 is the set of restaurant name and price pairs (r, pr) such that p is the price of some pizza sold by r that is not the most expensive pizze sold by r, and S_2 is the set of restaurant name and price pairs (r, pr) such that p is the price of the most expensive pizza sold by r. Thus, the required answer is given by $S_2 = \pi_{rname,price}(Sells) - S_1$. Note that $(r, pr) \in S_1$ if there exists some tuple $(r_2, p_2, pr_2) \in Sells$ where $r_2 = r$ and $pr_2 > pr$.



(g)
$$\pi_{cname.nizza}(Customers \to (Restaurants \bowtie Sells))$$

The following answer is incorrect:

$$\pi_{cname,pizza}(Customers \rightarrow Restaurants \bowtie Sells)$$

The above answer is equivalent to

$$\pi_{cname,pizza}((Customers \rightarrow Restaurants) \bowtie Sells)$$

However, outer joins are generally not associative with inner joins. In particular,

$$(Customers \rightarrow Restaurants) \bowtie Sells \not\equiv Customers \rightarrow (Restaurants \bowtie Sells)$$

To see this, consider a customer tuple $c \in Customers$ who is not co-located with any restaurant. Therefore, c will appear exactly once in the result of $Customers \to Restaurants$ with a null value for attributes rname. Due to the null value for rname, this tuple with c will be a dangling tuple w.r.t. the natural join with Sells. Thus, c will not be preserved in the result of $(Customers \to Restaurants) \bowtie Sells$. In contrast, by performing the outer join after the natural join, c will be preserved in the result of $Customers \to (Restaurants \bowtie Sells)$.

You should pay attention to the join ordering when using outer joins!

The following is another incorrect answer:

$$\pi_{cname,pizza}(Customers \leftrightarrow (Restaurants \leftrightarrow Sells))$$

By using full outer joins for both joins, this answer will also preserve dangling tuples from Sells and Restaurants relations; thus, the query result could have tuples with a null value for cname, but such tuples are not required by the question.

To see this, consider a tuple $(r, p, pr) \in Sells$ where the restaurant r is not colocated with any customer. Thus, (r, p, pr) will join with exactly one tuple from Restaurants (say (r, a)), but the resultant tuple (r, a, p, pr) from the first join computation will not join with any tuple from Customers. This will produce (null, p) in the query result, which is incorrect.

Moreover, if there is some tuple $r \in Restaurants$ where r is not co-located with any customer and r does not sell any pizza, then this will produce (null, null) in the query result, which is also incorrect.

The type of joins matters: you should use appropriate outer joins to preserve the required dangling tuples only when necessary!

5. Solution:

- R_1 is the set of pizzas that Maggie likes.
- R_2 is the set of restaurant-pizza pairs (r, p) where r is a restaurant that sells some pizza and p is some pizza that Maggie likes.
- \bullet R_3 is the set of restaurants that don't sell some pizza that Maggie likes.
- R_4 is the set of restaurants that sell all the pizzas that Maggie likes.
- R_5 is the set of pizzas that Ralph likes.
- R_6 is the set of restaurants that sell some pizza that Ralph likes.
- R_7 is the set of restaurants that sell all the pizzas that Maggie likes and don't sell any pizza that Ralph likes.