CS2100 Computer Organisation Tutorial #4: Instruction Set Architecture Answers

1. Given below are the contents of some registers and memory locations, where mem(A) refers to the data word stored in the memory location at address A:

mem (100) =	1000
mem (160) =	1600
mem (200) =	2000
mem (240) =	2400
mem (300) =	3000
mem (10000) =	100
mem (15000) =	150
mem (20000) =	200
mem (25000) =	250
mem (30000) =	300
	mem (100) = mem (160) = mem (200) = mem (240) = mem (300) = mem (10000) = mem (15000) = mem (25000) = mem (30000) =

For each of the MIPS instructions below, 'op' is an unknown opcode, which is not of our concern here. For each of the data operands in the instruction, if the format is legal, give its **memory address** (if applicable) and the content inside. Note that \$s0, ..., \$s4, \$t1, ..., \$t4 are symbolic register names, \$zero is register 0 with content zero inside, and only the three addressing modes (register, immediate, and displacement) mentioned in class are considered as legal.

- (a) op \$t1, \$t2
- (b) op \$s2, 100(\$zero)
- (c) op \$t4, 40(\$s2)
- (d) op \$s3, 200(\$zero)
- (e) op \$t3, \$zero(\$t1)
- (f) op \$s1, 140(\$s1)

Notes to students:

- 1. Do <u>not</u> use pseudo-instructions (move, bge, etc.) unless the question explicitly permits so.
- 2. Do not use non-standard MIPS instruction formats. QTSpim is a simulator; it accepts some non-standard format such as lw, \$t1, \$t2(0) or beq \$t3, 4, label. Use only standard MIPS instruction syntax in your answers.

The answers for (a) have been done for you.

	Operand	Target Memory Address	Content
(a)	\$t1	Not applicable	15000
	\$t2	Not applicable	20000
(b)	\$s2	Not applicable	200
	100(\$zero)	100	1000
(c)	\$t4	Not applicable	30000
	40(\$s2)	240	2400
(d)	\$s3	Not applicable	240
	200(\$zero)	200	2000
(e)	\$t3	Not applicable	25000
	\$zero(\$t1)	Illegal	Illegal
(f)	\$s1	Not applicable	160
	140(\$s1)	300	3000

For question 2 below, we are exploring four different types of ISAs. MIPS belongs to the *Register-Register* style.

2. You are tasked to design a dedicated 16-bit processor to perform simple addition and would like to evaluate the design using the four different instruction set architecture (ISA) styles. For each of these architectures, the corresponding data movement and arithmetic operations are shown below. Note that the operands for the instructions are annotated with either @ for address, or \$ for register.

ISA	Instructions	Explanation
Stack	push @src	Load value in @src onto top of stack.
	pop @dest	Transfer value at top of stack to @dest.
	add	Remove top two values in stack, add
		them, and load the sum onto top of stack.
Accumulator	load @src	Load value in @src into accumulator.
	add @src	Add value in @src and value in
	_	accumulator, and put sum back into
		accumulator.
	store @dest	Store the value in accumulator into
		@dest.
Memory-Memory	add @dest, @src1, @src2	Add values in @src1 and @src2, and put
, ,	_ , _ , _	the sum into @dest.
Register-Register	load \$reg, @src	Load value in @src into \$reg.
	add \$dest, \$src1, \$src2	Add values in $\$src1$ and $\$src2$, and put sum into $\$dest$.
	store \$reg, @dest	Store value in \$reg into @dest.

Consider the following three C statements with integer variables a0, a1 and a2:

a0 = a1 + a2; a1 = a0 + a2; a2 = a0 + a1;

For each of the four architectural styles above, write the assembly code corresponding to the above C code. Assume that the values of a0, a1, a2 are pre-assigned to memory in addresses @a0, @a1, @a2 respectively. The registers are denoted by \$r0 to \$r4. First part of each code is given.

Stack	Accumulator	Memory-Memory	Register-Register
push @a1	load @a1	add @a0, @a1, @a2	load \$r1, @a1
push @a2	add @a2	add @a1, @a0, @a2	load \$r2, @a2
add	store @a0	add @a2, @a0, @a1	add \$r0, \$r1, \$r2
pop @a0	add @a2		store \$r0, @a0
push @a0	store @a1		add \$r1, \$r0, \$r2
push @a2	add @a0		store \$r1, @a1
add	store @a2		add \$r2, \$r0, \$r1
pop @a1			store \$r2, @a2
push @a0			
push @a1			
add			
pop @a2			

3. This is a follow up on question 2. We studied the assembly code generated for four different storage architectures, namely stack, accumulator, memory-memory and register-register. For your reference, the code fragment corresponds to the following high level statements:

```
a0 = a1 + a2;
a1 = a0 + a2;
a2 = a0 + a1;
```

(a) Let us study the instruction encoding for this question. Assume that 3 bits will be used for the opcode, and minimal space will be used to represent 128 bytes of addressable memory. Moreover, for the Register-Register ISA, there are only 5 general-purpose registers available. Assume also a fixed-length instruction format, and that the memory is byte-addressable.

For each of the four architectures, what is the number of bits required for the longest instruction? Hence, what is the size, in number of bytes, of the instructions?

Partial answers are given below. Discuss.

	Number of bits for longest instruction	Number of bytes
Stack	10	2
Accumulator	10	2
Memory-Memory	24	4
Register-Register	13	2

- 3 bits needed to represent 5 registers;
- For a memory space of 128 bytes, 7 bits needed for a byte-addressable machine.
- A common question is why 4 bytes instead of 3 bytes for the 24-bit instruction?
 It could have been 3, but 4 bytes would permit a simpler design to avoid
 alignment problem (at the expense of more space). An instruction is usually
 fetched in a word, and a word usually contains 2, 4, or 8 bytes. It is uncommon
 to have a word that contains 3 bytes.

Stack	Accumulator	Memory-Memory	Register-Register
push @src 10	load @src 10	add @dest, @src1, @src2 24	load \$reg, @src 13
pop @dest 10	add @src 10		add \$dest1,\$src1,\$src212
add 3	store @dest 10		store \$reg, @dest 13

The longest instruction for each of the four ISAs from left to right are 10, 10, 24 and 13 bits respectively.

To compute the instruction code size for a 16-bit processor, we consider the minimum number of 16-bit words necessary to encompass the bits. The sizes are therefore 2, 2, 4 and 2 bytes respectively. This information will be needed for the next question.

(b) What is the size (in number of bytes) of the code in question 2 for each of the four architectures? Which architecture is most efficient in terms of code size for this code?

Code size = number of instructions \times size of each instruction

Stack: $12 \times 2 = 24$ bytes

Accumulator: $7 \times 2 = 14$ bytes Memory-Memory: $3 \times 4 = 12$ bytes Register-Register: $8 \times 2 = 16$ bytes

Hence Memory-Memory is the most efficient in terms of code size for this code.

4. [Past-year's exam question]

An ISA has 16-bit instructions and 5-bit addresses. There are two classes of instructions: class *A* instructions have one address, while class *B* instructions have two addresses. Both classes exist and the encoding space for opcode is completely utilized.

- (a) What is the minimum total number of instructions?
- (b) What is the maximum total number of instructions?

Answers

Class *A*: ppppppppppp xxxxx Class *B*: qqqqqq xxxxx yyyyy

'ppppppppppp' and 'qqqqqq' are opcodes, 'xxxxx' and 'yyyyy' are addresses.

To obtain the minimum total number of instructions, we give class B ($2^6 - 1$) opcodes, leaving one 6-bit opcode as the prefix for class A 11-bit opcodes. Hence class A has 2^5 opcodes. Total = ($2^6 - 1$) + (2^5) = 63 + 32 = 95.

To obtain the maximum total number of instructions, we give class B only 1 opcode, leaving $(2^6 - 1)$ prefixes for class A opcodes. Hence, class A has $(2^6 - 1) \times (2^5) = 63 \times 32 = 2016$ opcodes. Hence a total of 2017 instructions.