# CS2100: Computer Organisation Tutorial #2: C and MIPS (3 February – 7 February 2020)

Discussion questions D1 and D2 are meant for your own exploration and discussion on LumiNUS Forum. They will not be discussed in class.

### D1. Exploration: C to MIPS

Go to this website <a href="http://reliant.colab.duke.edu/c2mips/">http://reliant.colab.duke.edu/c2mips/</a> and copy the C code below into the box. Check that the 'Optimization' option is "-O0 – No optimization (default)" and click the 'Submit' button. Scroll down and study the Assembly output.

```
int main(void) {
   int a, b, c;

a = 3;
b = 5;
c = a + b;
return 0;
}
```

The following is extracted from the Assembly output.

```
addiu $sp, $sp, -32
      $fp,28($sp)
      $fp,$sp
move
li
      $2,3
                           # 0x3
      $2,8($fp)
SW
      $31,$31,$0
movz
      $2,5
                           # 0x5
li
      $2,12($fp)
SW
      $3,8($fp)
      $2,12($fp)
lw
nop
      $2,$3,$2
addu
      $2,16($fp)
move $2,$0
move $sp,$fp
      $fp,28($sp)
addiu $sp,$sp,32
j
       $31
nop
```

We covered **sw**, **lw** and **j** in lecture, but not the rest. Find out what they are. (Note: **li** will be used in the labs later and **nop** will be mentioned in the topic on Pipelining. **move**, like **li**, is a pseudo-instruction. **\$sp**, **\$fp**, **movz**, **addiu**, and **addu** are not in the syllabus.)

D2. For each of the following instructions, indicate if it is valid or not. If not, explain why and suggest a correction. Note that the | in (d) is the bitwise OR operation.

a. add \$t1, \$t2, \$t3 #\$t3 = \$t1 + \$t2

b. addi \$t1, \$0, 0x25 #\$t1 = 0x25

c. **subi** \$t2, \$t1, 3 #\$t2 = \$t1 - 3

d. ori \$t3, \$t4, 0xAC120000 #\$t3 = \$t4 | 0xAC120000

e. **sl1** \$t5, \$t2, 0x21 # shift left \$t2 33 bits and store in \$t5

# **Tutorial questions**

## 1. Bitwise operations

Find out about the following bitwise operations in C and explain and illustrate each of them with an example.

| (bitwise OR)
& (bitwise AND)
^ (bitwise XOR)
~ (one's complement)
<< (left shift)</li>
>> (right shift)

You may use the following code template for your illustration. Variables of the data type **char** take up 8 bits of memory.

```
#include <stdio.h>
typedef unsigned char byte t;
void printByte(byte t);
int main(void) {
   byte t a, b;
   a = 5;
   b = 22;
   printf("a = "); printByte(a); printf("\n");
printf("b = "); printByte(b); printf("\n");
   printf("a|b = "); printByte(a|b); printf("\n");
   return 0;
}
void printByte(byte t x) {
   printf("%c%c%c%c%c%c%c%c",
           (x \& 0x80 ? '1': '0'),
           (x \& 0x40 ? '1': '0'),
           (x \& 0x20 ? '1': '0'),
           (x \& 0x10 ? '1': '0'),
           (x \& 0x08 ? '1': '0'),
           (x \& 0x04 ? '1': '0'),
           (x \& 0x02 ? '1': '0'),
           (x \& 0x01 ? '1': '0'));
}
```

#### 2. Swapping

Your friend is challenging you with a magic trick. He has a deck of 17 cards. All cards are either  $\clubsuit$  or  $\spadesuit$  suit with numbers from 2 to 10. Since there are only 17 cards, it means one of the cards is missing. He says that he can find the missing card by only looping through the deck exactly once! To prove his point, he produces the following code.

```
int findCard(int deck[]) {
   int res, i;
   int count[9] = {0, 0, 0, 0, 0, 0, 0, 0, 0};

for (i=0; i<17; i++) {
     count[deck[i]-2]++;
   }

for (i=2; i<=10; i++) {
     if (count[i-2] == 1) {
        res = i;
     }
   }
}

return res;
}</pre>
```

As a CS2100 student who has learned about bitwise operator, you say that you can produce a better code! In particular, your friend is using a separate array to count the number of occurrences and you believe you can do without it. To do that, you will need to use one of the bitwise operators from Question 1.

Improve the function findCard(int deck[]) such that it can find the missing card without using a separate array to store the number of occurrences.

**Hint:** You will need a bitwise operator  $\otimes$  such that  $x \otimes x = 0$  and  $0 \otimes x = x$ .

3. For this question, show the minimal sequence of MIPS instructions needed to translate the given C statement. Use only MIPS instructions covered in MIPS Part 1: Introduction. Assume that variables **a** and **b** are mapped to registers \$s0 and \$s1 respectively.

```
    (a) b = a % 16; // b is the remainder of a divides by 16
    (b) b = (a / 8) * 8; // b is the nearest multiple of 8 that is smaller than a. // Note that integer division is used.
```

4. The mysterious MIPS code below assumes that \$s0 is a 16-bit binary sequence, i.e. the upper 16-bits of \$s0 and \$s1 are assumed to be zero at the start of the code.

```
add $t0, $s0, $zero
add $t1, $s1, $zero
add $t3, $zero, $zero

lp: beq $t1, $zero, e
andi $t2, $t1, 1
beq $t2, $zero, s
add $t3, $t3, $t0

s: sll $t0, $t0, 1
srl $t1, $t1, 1
j lp
e:
```

- a) Given the initial value of \$s0 = 3 and \$s1 = 5. What is the value of \$t3 in hexadecimal at the end of execution?
- b) Given the initial value of \$s0 = 5 and \$s1 = 13. How many times do statement "beq \$t2, \$zero, s" results in a branch to the label s?
- c) Explain the purpose of the code in one sentence.

[AY2012/13 Semesters 2 Term Test #2]
 Take a look at the following MIPS program and the memory contents. For simplicity, all values (memory addresses and contents) are in decimal.

Address	Content
100	120
104	132
108	128
112	108
116	124
120	116
124	104
128	100
132	136
136	112

Suppose we execute the code for **3 iterations**, answer the following:

- (a) Give the value of register \$s1 at the end of the third iteration.
- (b) If we want to terminate the loop at the fourth check of **beq**, some memory contents need to be modified. Give the address(es) and the modified content(s) to achieve this. You should find the minimum changes required.
- (c) What do you think the code does? Give a high-level code fragment that works in a similar way.