



The x86 PC

assembly language, design, and interfacing

fifth
edition

Prentice Hall

Chapter TWO

Assembly Language Programming

The x86 PC

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OBJECTIVES

this chapter enables the student to:

- Explain the difference between Assembly language instructions and pseudo-instructions.
- Identify the segments of an Assembly program.
- Assemble, link, and run a simple Assembly program.
- Code control transfer instructions such as **conditional jump** and **unconditional jump** and **call**.
- Data directives for binary, hex, decimal, or ASCII.
- Write an Assembly language program using either the simplified/full segment definition.
- Explore the use of the (MASM and emu8086) assemblers.

2.0: ASSEMBLY LANGUAGE

- An Assembly language program is a series of statements, or lines.
 - Either Assembly language instructions, or statements called *directives*.
 - Directives (pseudo-instructions) give directions to the assembler about how it should translate the Assembly language instructions into machine code.
- Assembly language instructions consist of four fields:
`[label:] mnemonic [operands][;comment]`
 - Brackets indicate that the field is optional.
 - Do not type in the brackets.

2.1: DIRECTIVES AND A SAMPLE PROGRAM

assembly language instructions

[label:] **mnemonic** [operands] [;comment]

- Examples of directives are **DB**, **END**, and **ENDP**.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX,@DATA ;load the data segment address
MOV DS,AX ;assign value to DS
MOV AL,DATA1 ;get the first operand
MOV BL,DATA2 ;get the second operand
ADD AL,BL ;add the operands
MOV SUM,AL ;store the result in location SUM
MOV AH,4CH ;set up to return to OS
INT 21H ;
MAIN ENDP
END MAIN ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

model definition

- After the first two comments is the MODEL directive.
 - This directive selects the size of the memory model.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
        .MODEL SMALL
        .STACK 64
        .DATA
DATA1    DB    52H
DATA2    DB    29H
SUM      DB    ?
        .CODE
MAIN     PROC   FAR           ;this is the program entry point
        MOV    AX,@DATA      ;load the data segment address
        MOV    DS,AX         ;assign value to DS
        MOV    AL,DATA1      ;get the first operand
        MOV    BL,DATA2      ;get the second operand
        ADD    AL,BL         ;add the operands
        MOV    SUM,AL        ;store the result in location SUM
        MOV    AH,4CH        ;set up to return to OS
        INT    21H           ;
MAIN     ENDP
        END    MAIN         ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

model definition

- Among the options for the memory model are **SMALL**, **MEDIUM**, **COMPACT**, and **LARGE**.

<code>.MODEL SMALL</code>	<code>;this directive defines the model as small</code>
<code>.MODEL MEDIUM</code>	<code>;the data must fit into 64K bytes</code> <code>;but the code can exceed 64K bytes of memory</code>
<code>.MODEL COMPACT</code>	<code>;the data can exceed 64K bytes</code> <code>;but the code cannot exceed 64K bytes</code>
<code>.MODEL LARGE</code>	<code>;both data and code can exceed 64K</code> <code>;but no single set of data should exceed 64K</code>
<code>.MODEL HUGE</code>	<code>;both code and data can exceed 64K</code> <code>;data items (such as arrays) can exceed 64K</code>
<code>.MODEL TINY</code>	<code>;used with COM files in which data and code</code> <code>;must fit into 64K bytes</code>

2.1: DIRECTIVES AND A SAMPLE PROGRAM

segment definition

- Every line of an Assembly language program must correspond to one an x86 CPU segment register.
 - CS (code segment); DS (data segment).
 - SS (stack segment); ES (extra segment).
- The simplified segment definition format uses three simple directives: ".CODE" ".DATA" ".STACK"
 - Which correspond to the CS, DS, and SS registers.

```
.STACK      ;marks the beginning of the stack segment
.DATA       ;marks the beginning of the data segment
.CODE       ;marks the beginning of the code segment
```

- The stack segment defines storage for the stack.
- The data segment defines the data the program will use.
- The code segment contains Assembly language instructions.

2.1: DIRECTIVES AND A SAMPLE PROGRAM

stack segment

- This directive reserves 64 bytes of memory for the stack:

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1    DB    52H
DATA2    DB    29H
SUM       DB    ?
.CODE
MAIN      PROC   FAR           ;this is the program entry point
          MOV    AX,@DATA      ;load the data segment address
          MOV    DS,AX         ;assign value to DS
          MOV    AL,DATA1      ;get the first operand
          MOV    BL,DATA2      ;get the second operand
          ADD    AL,BL         ;add the operands
          MOV    SUM,AL        ;store the result in location SUM
          MOV    AH,4CH         ;set up to return to OS
          INT    21H           ;
MAIN      ENDP
          END     MAIN         ;this is the program exit point
```


2.1: DIRECTIVES AND A SAMPLE PROGRAM

data segment

- The data segment defines three data items:
 - DATA1, DATA2, and SUM.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX,@DATA ;load the data segment address
MOV DS,AX ;assign value to DS
MOV AL,DATA1 ;get the first operand
MOV BL,DATA2 ;get the second operand
ADD AL,BL ;add the operands
MOV SUM,AL ;store the result in location SUM
MOV AH,4CH ;set up to return to OS
INT 21H ;
MAIN ENDP
END MAIN ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- The first line of the segment after the `.CODE` directive is the `PROC` directive.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
        .MODEL SMALL
        .STACK 64
        .DATA
DATA1    DB    52H
DATA2    DB    29H
SUM      DB    ?
        .CODE
MAIN     PROC FAR                ;this is the program entry point
        MOV    AX,@DATA          ;load the data segment address
        MOV    DS,AX             ;assign value to DS
        MOV    AL,DATA1          ;get the first operand
        MOV    BL,DATA2          ;get the second operand
        ADD    AL,BL             ;add the operands
        MOV    SUM,AL            ;store the result in location SUM
        MOV    AH,4CH            ;set up to return to OS
        INT    21H              ;
MAIN     ENDP
        END    MAIN             ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- *A procedure* is a group of instructions designed to accomplish a specific function.
 - A code segment is organized into several small procedures to make the program more structured.
- Every procedure must have a name defined by the PROC directive.
 - Followed by the assembly language instructions, and closed by the ENDP directive.
 - The PROC and ENDP statements must have the same label.
 - The PROC directive may have the option FAR or NEAR.
 - The OS requires the entry point to the user program to be a FAR procedure.

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- The program loads **AL** & **BL** with **DATA1** & **DATA2**, ADDs them together, and stores the result in **SUM**.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION

.MODEL SMALL
.STACK 64
.DATA
DATA1    DB    52H
DATA2    DB    29H
SUM       DB    ?
.CODE
MAIN     PROC   FAR           ;this is the program entry point
        MOV    AX,@DATA      ;load the data segment address
        MOV    DS,AX         ;assign value to DS
        MOV    AL,DATA1      ;get the first operand
        MOV    BL,DATA2      ;get the second operand
        ADD    AL,BL         ;add the operands
        MOV    SUM,AL        ;store the result in location SUM
        MOV    AH,4CH         ;set up to return to OS
        INT    21H           ;
MAIN     ENDP
        END    MAIN         ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- The first instructions, "**MOV AX, @DATA**" & "**MOV DS, AX**" initialize DS register

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX, @DATA ;load the data segment address
MOV DS, AX ;assign value to DS
MOV AL, DATA1 ;get the first operand
MOV BL, DATA2 ;get the second operand
ADD AL, BL ;add the operands
MOV SUM, AL ;store the result in location SUM
MOV AH, 4CH ;set up to return to OS
INT 21H ;
MAIN ENDP
END MAIN ;this is the program exit point
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- The last instructions, "**MOV AH, 4CH**" & "**INT 21H**" return control to the operating system.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX,@DATA ;load the data segment address
MOV DS,AX ;assign value to DS
MOV AL,DATA1 ;get the first operand
MOV BL,DATA2 ;get the second operand
ADD AL,BL ;add the operands
MOV SUM,AL ;store the result in location SUM
MOV AH,4CH ;set up to return to OS
INT 21H ;
MAIN ENDP
END MAIN ;this is the program exit point
```


2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- The last two lines end the procedure & program.
 - The label for **ENDP(MAIN)** matches the label for **PROC**.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX,@DATA ;load the data segment address
MOV DS,AX ;assign value to DS
MOV AL,DATA1 ;get the first operand
MOV BL,DATA2 ;get the second operand
ADD AL,BL ;add the operands
MOV SUM,AL ;store the result in location SUM
MOV AH,4CH ;set up to return to OS
INT 21H ;
MAIN ENDP ;this is the program exit point
END MAIN
```

2.1: DIRECTIVES AND A SAMPLE PROGRAM

code segment definition

- It is handy to keep a sample shell & fill it in with the instructions and data for your program.

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
; USING SIMPLIFIED SEGMENT DEFINITION
.MODEL SMALL
.STACK 64
.DATA
;
;place data definitions here
;
.CODE
MAIN PROC FAR           ;this is the program entry point
MOV AX,@DATA           ;load the data segment address
MOV DS,AX              ;assign value to DS
;
;place code here
;
MOV AH,4CH             ;set up to
INT 21H                ;return to OS
MAIN ENDP
END MAIN               ;this is the program exit point
```

2.2: ASSEMBLE, LINK, AND RUN A PROGRAM

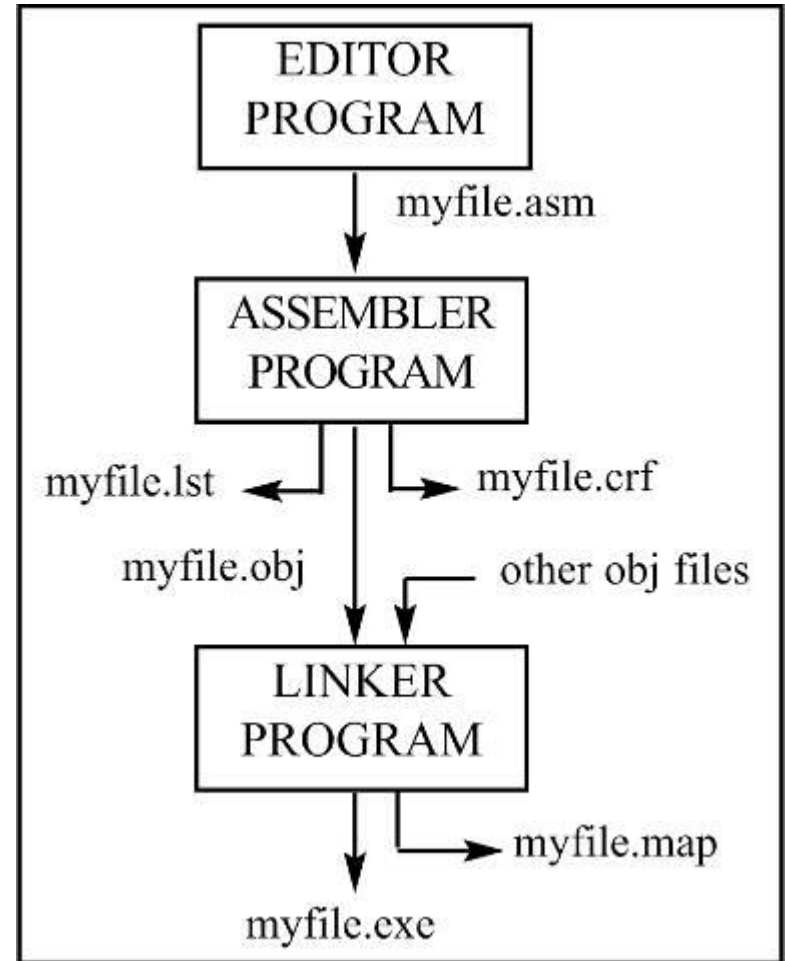
- MASM & LINK are the assembler & linker programs.
 - Many editors or word processors can be used to create and/or edit the program, and produce an ASCII file.
 - The steps to create an executable Assembly language program are as follows:

Step	Input	Program	Output
1. Edit the program	keyboard	editor	myfile.asm
2. Assemble the program	myfile.asm	MASM or TASM	myfile.obj
3. Link the program	myfile.obj	LINK or TLINK	myfile.exe

2.2: ASSEMBLE, LINK, AND RUN A PROGRAM

Before feeding the ".obj" file into LINK, all syntax errors must be corrected.

Fixing these errors will not guarantee the program will work as intended, as the program may contain conceptual errors.



2.4: CONTROL TRANSFER INSTRUCTIONS

FAR and NEAR

- In the sequence of instructions, it is often necessary to transfer program control to a different location.
 - If control is transferred to a memory location within the current code segment, it is NEAR.
 - Sometimes called intrasegment. (within segment)
 - the IP is updated and CS remains the same
 - If control is transferred outside the current code segment, it is a FAR jump.
 - Or intersegment. (between segments)
 - both CS and IP have to be updated to the new values

2.4: CONTROL TRANSFER INSTRUCTIONS

conditional jumps

- Conditional jumps have mnemonics such as JNZ (jump not zero) and JC (jump if carry).
 - In the conditional jump, control is transferred to a new location if a certain condition is met.
 - The flag register indicates the current condition.
- For example, with "JNZ label", the processor looks at the zero flag to see if it is raised.
 - If not, the CPU starts to fetch and execute instructions from the address of the label.
 - If $ZF = 1$, it will not jump but will execute the next instruction below the JNZ.

2.4: CONTROL TRANSFER INSTRUCTIONS

conditional jumps

**Table 2-1: 8086
Conditional
Jump Instructions**

Note: “Above” and “below” refer to the relationship of two unsigned values; “greater” and “less” refer to the relationship of two signed values.

Mnemonic	Condition Tested	“Jump IF ...”
JA/JNBE	$(CF = 0) \text{ and } (ZF = 0)$	above/not below nor zero
JAE/JNB	$CF = 0$	above or equal/not below
JB/JNAE	$CF = 1$	below/not above nor equal
JBE/JNA	$(CF \text{ or } ZF) = 1$	below or equal/not above
JC	$CF = 1$	carry
JE/JZ	$ZF = 1$	equal/zero
JG/JNLE	$((SF \text{ xor } OF) \text{ or } ZF) = 0$	greater/not less nor equal
JGE/JNL	$(SF \text{ xor } OF) = 0$	greater or equal/not less
JL/JNGE	$(SF \text{ xor } OF) = 1$	less/not greater nor equal
JLE/JNG	$((SF \text{ xor } OF) \text{ or } ZF) = 1$	less or equal/not greater
JNC	$CF = 0$	not carry
JNE/JNZ	$ZF = 0$	not equal/not zero
JNO	$OF = 0$	not overflow
JNP/JPO	$PF = 0$	not parity/parity odd
JNS	$SF = 0$	not sign
JO	$OF = 1$	overflow
JP/JPE	$PF = 1$	parity/parity equal
JS	$SF = 1$	sign

2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- All conditional jumps are short jumps.
 - The address of the target must be within -128 to +127 bytes of the IP.
- The conditional jump is a two-byte instruction.
 - One byte is the opcode of the J condition.
 - The second byte is a value between 00 and FF.
 - An offset range of 00 to FF gives 256 possible addresses.
- In a jump backward, the second byte is the 2's complement of the displacement value

2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- To calculate the target address, the second byte is added to the IP of the instruction after the jump.

```
1067:0000 B86610      MOV    AX,1066
1067:0003 8ED8      MOV    DS,AX
1067:0005 B90500      MOV    CX,0005
1067:0008 BB0000      MOV    BX,0000
1067:000D 0207      ADD    AL,[ BX]
1067:000F 43          INC    BX
1067:0010 49          DEC    CX
1067:0011 75FA      JNZ    000D
1067:0013 A20500      MOV    [ 0005] ,AL
1067:0016 B44C      MOV    AH,4C
1067:0018 CD21      INT    21
```

- "JNZ AGAIN" was assembled as "JNZ 000D", and 000D is the address of the instruction with the label AGAIN.
 - "JNZ 000D" has the opcode 75 and the target address FA.

2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- This is followed by "MOV SUM, AL", which is located beginning at offset address 0013.

```
1067:0000 B86610      MOV    AX,1066
1067:0003 8ED8      MOV    DS,AX
1067:0005 B90500      MOV    CX,0005
1067:0008 BB0000      MOV    BX,0000
1067:000D 0207      ADD    AL,[ BX]
1067:000F 43          INC    BX
1067:0010 49          DEC    CX
1067:0011 75FA      JNZ    000D
1067:0013 A20500      MOV    [ 0005] ,AL
1067:0016 B44C      MOV    AH,4C
1067:0018 CD21      INT    21
```

- The IP value of MOV, 0013, is added to FA to calculate the address of label AGAIN, and the carry is dropped.
 - FA is the 2's complement of -6.

2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- Calculate a forward jump target address by adding the IP of the following instruction to the operand.
 - The displacement value is positive, as shown.

0005	8A 47 02	AGAIN:	MOV	AL,[BX] +2
0008	3C 61		CMP	AL,61H
000A	72 06		JB	NEXT
000C	3C 7A		CMP	AL,7AH
000E	77 02		JA	NEXT
0010	24 DF		AND	AL,ODFH
0012	88 04	NEXT:	MOV	[SI],AL

- "**JB NEXT**" has the opcode **72**, the target address **06** and is located at **IP = 000A** and **000B**.
 - The jump is 6 bytes from the next instruction, is IP = 000C.
 - Adding gives us 000CH + 0006H = 0012H, which is the exact address of the **NEXT** label.

2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- For conditional jumps, the address of the target address can never be more than -128 to +127 bytes away from the IP associated with the instruction following the jump.
 - Any attempt is made to violate this rule will generate a "relative jump out of range" message.

2.4: CONTROL TRANSFER INSTRUCTIONS

unconditional jumps

- An unconditional jump transfers control to the target location label unconditionally, in the following forms:
 - **SHORT JUMP** - in the format "JMP SHORT label".
 - A jump within -128 to +127 bytes of memory relative to the address of the current IP, opcode EB.
 - **NEAR JUMP** - the default, has the format "JMP label".
 - A jump within the current code segment, opcode E9.
 - The target address can be any of the addressing modes of direct, register, register indirect, or memory indirect:
 - **Direct JUMP** - exactly like the short jump.
 - Except that the target address can be anywhere in the segment in the range +32767 to -32768 of the current IP.

2.4: CONTROL TRANSFER INSTRUCTIONS

unconditional jumps

- An unconditional jump transfers control to the target location label unconditionally, in the following forms:
 - **Register indirect JUMP** - target address is in a register.
 - In "JMP BX", IP takes the value BX.
 - **Memory indirect JMP** - target address is the contents of two memory locations, pointed at by the register.
 - "JMP [DI]" will replace the IP with the contents of memory locations pointed at by DI and DI+1.
 - **FAR JUMP** - in the format "JMP FAR PTR label".
register.
 - A jump out of the current code segment
 - IP and CS are both replaced with new values.

2.4: CONTROL TRANSFER INSTRUCTIONS

CALL statements

- The CALL instruction is used to call a procedure, to perform tasks that need to be performed frequently.
 - The target address could be in the current segment, in which case it will be a NEAR call or outside the current CS segment, which is a FAR call.
- The microprocessor saves the address of the instruction following the call on the stack.
 - To know where to return, after executing the subroutine.
 - In the NEAR call only the IP is saved on the stack.
 - In a FAR call both CS and IP are saved.

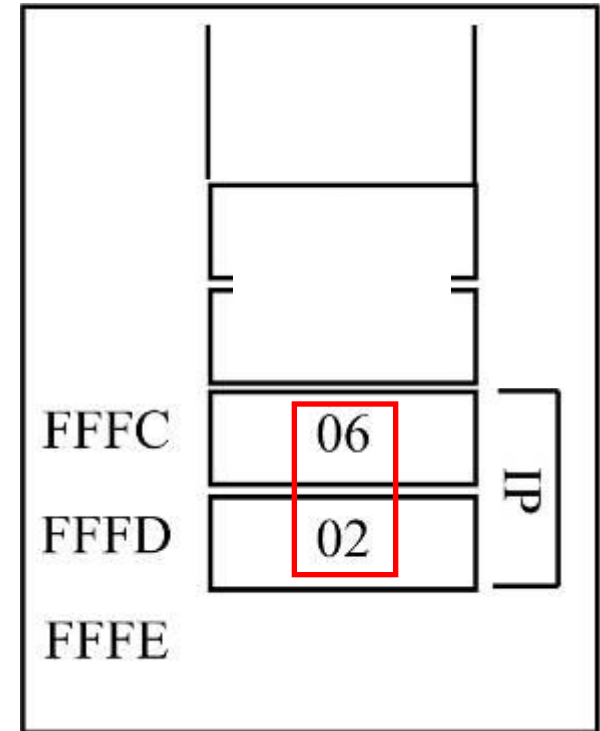
2.4: CONTROL TRANSFER INSTRUCTIONS

CALL statements

- For control to be transferred back to the caller, the last subroutine instruction must be RET (return).
 - For NEAR calls, the IP is restored.
 - For FAR calls, CS & IP are restored.
- Assume SP = FFFEh:

```
12B0:0200  BB1295  MOV BX, 9512
12B0:0203  E8FA00  CALL 0300
12B0:0206  B82F14  MOV AX, 142F
```

- Since this is a NEAR call, only IP is saved on the stack.
 - The IP address **0206**, which belongs to the "**MOV AX, 142F**" instruction, is saved on the stack.

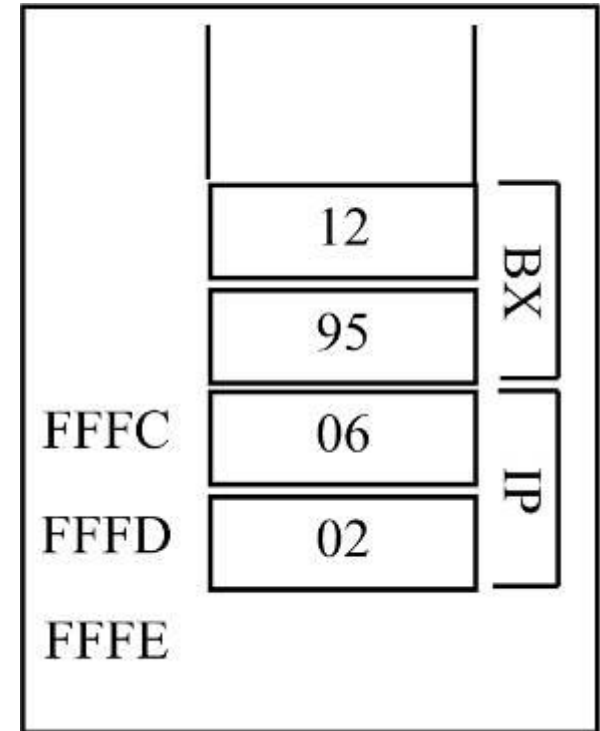


2.4: CONTROL TRANSFER INSTRUCTIONS

short jumps

- The last instruction of the called subroutine must be a RET instruction that directs the CPU to POP the top 2 bytes of the stack into the IP and resume executing at offset address 0206.
 - The number of PUSH and POP instructions (which alter the SP) **must match**.
 - For every **PUSH** there must be a **POP**.

```
12B0:0300  53  PUSH BX
12B0:0301  ...  ....
...
12B0:0309  5B  POP  BX
12B0:030A  C3  RET
```



2.5: DATA TYPES AND DATA DEFINITION

- DUP will duplicate a given number of characters.
- DW is used to allocate memory 2 bytes (one word)
- EQU associates a constant value with a data label.
 - Defines a constant without occupying a memory location.
- The DD directive is used to allocate memory locations that are 4 bytes (two words) in size.
 - Data is converted to hex & placed in memory locations
 - Low byte to low address and high byte to high address.
- DQ is used to allocate memory 8 bytes (four words) in size, to represent any variable up to 64 bits wide

2.5: DATA TYPES AND DATA DEFINITION

- DQ is used to allocate memory 8 bytes (four words) in size, to represent any variable up to 64 bits wide
- DT is used for memory allocation of packed BCD numbers.
 - This directive allocates 10 bytes.
 - A maximum of 18 digits can be entered.
 - The "H" after the data is not needed.

2.6: FULL SEGMENT DEFINITION

→ segment definition

<code>;FULL SEGMENT DEFINITION</code>	<code>;SIMPLIFIED FORMAT</code>
<code>;--- stack segment ---</code>	<code>.MODEL SMALL</code>
<code>name1 SEGMENT</code>	<code>.STACK 64</code>
<code>DB 64 DUP (?)</code>	<code>;</code>
<code>name1 ENDS</code>	<code>;</code>
<code>;--- data segment ---</code>	<code>;</code>
<code>name2 SEGMENT</code>	<code>. DATA</code>
<code>;place data definitions here</code>	<code>;place data definitions here</code>
<code>name2 ENDS</code>	<code>;</code>
<code>;--- code segment ---</code>	<code>;</code>
<code>name3 SEGMENT</code>	<code>.CODE</code>
<code>MAIN PROC FAR</code>	<code>MAIN PROC FAR</code>
<code>ASSUME ...</code>	<code>MOV AX, @DATA</code>
<code>MOV AX, name2</code>	<code>MOV DS, AX</code>
<code>MOV DS, AX</code>	<code>...</code>
<code>...</code>	<code>...</code>
<code>MAIN ENDP</code>	<code>MAIN ENDP</code>
<code>name3 ENDS</code>	<code>END MAIN</code>
<code>END MAIN</code>	

Figure 2-8

ENDS ; TWO



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