

Q2

a) Assume CFLs closed under intersections then

$$A = \{a^m b^n c^n \mid m, n \geq 0\} \rightarrow A \cap B = \{a^n b^n c^n \mid n \geq 0\}$$
$$B = \{a^n b^n c^m \mid m, n \geq 0\}$$

We know that $A \cap B$ is not CFL therefore we reach to a contradiction. Our assumption was wrong. \square

b) Assume CFLs closed under complementation, Then \bar{A} and \bar{B} are CFLs then we can write

$$\overline{(\bar{A} \cup \bar{B})}$$

CFL

should be CFL since we assumed CFLs are closed under complement

$$\overline{(\bar{A} \cup \bar{B})} = A \cap B \text{ which is not CFL, contradiction.}$$

Q3)

a) Pick a word $w \in B$ s.t. $|w| > p$ where p is the pumping length $w = 0^p 1^{2p} 0^p \in B$ then

we can split w as $uvxyz$ where $|xy| > 0$

and $|vxy| \leq p$ then we have 5 cases

* vxy contains 0s from left then
 uv^2xy^2z contains more 0s than z

* vxy contains 0s and 1s from left then
 uv^2xy^2z contains $00 \dots 00 \text{---} 11 \dots 11 0 \dots 0$
 0s followed 1s on the left and not a palindrome

- * vxy contains 1s only then uv^2xy^2z contains more 1s than 0s
- * vxy contains 1s and 0s then similar to case M *

- * vxy contains 0s from right then uv^2xy^2z contains more 0s than 1s

b) Similarly let p be p.l. and pick $w = 1^p 3^p 2^p 4^p$ then we can split w as $uvxy^2z$ where $|vy| > 0$ $|vxy| \leq p$

- * vxy contains only one of the digits (e.g. 1) then uv^2xy^2z contains more 1s than 2s.
- * vxy contains 2 of the digits (e.g. 2 and 6) then uv^2xy^2z contains less 2s and 6s than 1s and 3s

c) Let p be p.l. s be $1^{2p} 1^p 0^p 1^{2p} \in A$ $|w| > p$ then we can split s as $uvxy^2z$ where $|vy| > 0$ $|vxy| \leq p$ $s = \underbrace{1^{2p}}_w \underbrace{1^p}_t \underbrace{0^p}_{w^R} 1^{2p}$ we have

- * vxy contains only 1s from w then uv^2xy^2z $|w| > 2p$ but only has 1s w^R should also have cardinality $2p$ but can only do it by borrowing 0s from t meaning $(w^R)^R \neq w$
- * vxy contains 1s from w and t then uv^2xy^2z will have $|w| > 2p$ and similarly $|w^R| > 2p$ can only be achieved by borrowing 0s from t .
- * vxy containing 1s, 1s and 0s or 0s only from t then uv^2xy^2z will make it impossible to obtain $(w^R)^R = w$ by breaking the balance of 1s and 0s.

end Os.
 * vxy contains Os from t end Is from w^R

** If amount of Os in vxy is more than Is then w^R will borrow Os from t end
 (with t)

** If amount of Os in vxy is equal to amount of Is in vxy

then in wxy^Rz we will have
 $|w^R| = |t| \neq |x|$

** If amount of Is in vxy is larger than amount of Os in vxy then in wxy^Rz we will have w^R with Os but we won't have any Os.

* vxy contains only Is from w^R then in wxy^Rz we will have w^R with Os but we won't have any Os.

