

# Operating Systems and Algorithms

## Assignment-2

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Q.1. Explanation: Given an array of integers. If we see the problem closely here we need to find two parts of the subarray such that the difference of the two parts is minimum.

Algorithm: Every element has two choices either it will be part of  $s_1$  or  $s_2$ . So accordingly I will call the recursive call and try to minimise the sum and  $\text{total\_sum} - \text{sum}$ .

```
This is the recursive call
int find(vector<int>&A, vector<vector<int>>&dp, int ind, int sum, int tsum)

and two options
int take=find(A, dp, ind+1, sum + A[ind], tsum);
int n_take=find(A, dp, ind+1, sum, tsum);
```

Time complexity:  $O(N * \text{total\_sum})$ . Basically I run two nested for loops that is defining the time complexity.

Space complexity:  $O(N * \text{total\_sum})$ . This is the total space for the size of dp matrix.

Q.2. This is a classical knapsack problem. Implementation is just take the item or not and according try all the combination to maximize the values.

Algorithm: I implemented using dynamic programming where  $dp[i][j]$  basically denoting the maximum value obtained up to the  $i$ th index of the array with  $j$  weight.

```

this is the main working of the code and according to it the dp states are updated.
if(wieght[i-1]<=j){
    dp[i][j]=max(value[i-1]+dp[i-1][j-wieght[i-1]],dp[i-1][j]);
}
else{
    dp[i][j]=dp[i-1][j];
}

```

Time complexcity:  $O(N \times \text{wight})$  . Basically I run two nested for loop that is defining the time complexcity.

Space complexcity:  $O(N \times \text{total\_sum})$  . this is the total space for the size of dp matrix.

Q.4.

Explanation: Here we need to find minimum wastage of area . wastage is defined as any rectangular piece which is not listed.

Algorithm: I used a recursion with memorization algoritihm to solve the question. for a given w and h i am initilizing the dp array with maximum wastage that is  $w \times h$  . Then for all the valued of w and h i am calling the recursive function and update the dp state whenever a less wastage is obtained.

dp state:  $dp[w][h]$  what is the minimum wasteage with rectangle of width w and hieght h.

```

main algorithm:
for (int i = 1; i <= h / 2; i++) {
    temp = find(w, i);
    if (temp + find(w, h - i) < dp[w][h]) {
        dp[w][h]= temp + find(w, h - i);
    }
}
for (int i = 1; i <= w / 2; i++) {
    temp = find(i, h);
    if (temp + find(w - i, h) < dp[w][h]) {
        dp[w][h] = temp + find(w - i, h);
    }
}

```

```
}  
}
```

Time complexity:  $O(w \cdot h)$  . maximum recursive call will be  $w \cdot h$  times with using the memorization for subproblems.

space complexity is  $O(w \cdot h)$  storing the value in dp state.