

Let dist. of X be $\{p_1, p_2, \dots, p_n\}$

$$\Rightarrow H(X) = - \sum_{i=1}^n p_i \log(p_i)$$

$$X = \left\{ \begin{array}{l} a \rightarrow \frac{1}{2} \\ b \rightarrow \frac{1}{4} \\ c \rightarrow \frac{1}{8} \\ d \rightarrow \frac{1}{8} \end{array} \right\} \Rightarrow H(X) = \frac{1}{2} \log 2 + \frac{1}{4} \log 4 + \frac{1}{8} \times 2 \log 8$$

$$= \frac{1}{2} + \frac{1}{2} + \frac{3}{4} = 1 + \frac{3}{4} = \frac{7}{4} \text{ bits}$$

P.T. $H(X|Y) \leq H(X)$

$$P(X|Y) = \frac{P(X, Y)}{P(Y)} \text{ and } P(X) \quad P(X, Y) \leq P(X).$$



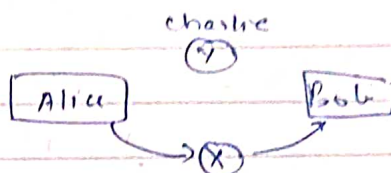
Conditional Entropy

X : R.V. that A will visit B

Y : R.V. that A's car got punched.

A and B are now dependent

Let $q_{k\ell} = P(B_\ell | A_k)$, and $\pi_{k\ell} = P(A_k, B_\ell)$



$\Rightarrow \pi_{k\ell} = p_k \cdot q_{k\ell}$

$$-H(AB) = \sum_k \sum_\ell \pi_{k\ell} \log(\pi_{k\ell})$$

$$= \sum_k \sum_\ell p_k q_{k\ell} \{ \log(p_k) + \log(q_{k\ell}) \}$$

$$= \underbrace{\sum_k p_k \log(p_k)}_{-H(A)} \cdot \underbrace{\sum_\ell q_{k\ell}}_1 + \sum_k p_k \sum_\ell q_{k\ell} \log(q_{k\ell}).$$

$$-H(AB) = -H(A) + \sum_k p_k \sum_\ell q_{k\ell} \log(q_{k\ell}).$$

\Rightarrow

$$H(AB) = H(A) + \sum_k p_k \left(- \sum_\ell q_{k\ell} \log(q_{k\ell}) \right).$$

$$H(AB) = H(A) + \sum_k p_k (H_k(B))$$

$$H(AB) = H(A) + H(B|A). \quad \text{--- (1)}$$

P.T. $H(B|A) \leq H(B).$

$$H(BA) = H(B) + H(A|B).$$

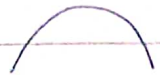
$$- \sum_k p_k \log p_k + \sum_k \sum_\ell p_k q_{k\ell}$$

$$H(B|A) = H(B) + H(A|B) - H(A).$$

$$= H(B) - (H(A) - H(A|B))$$

$$H(B) - H(B|A)$$

$$= -\sum_i q_i \log q_i + \sum_k \sum_i p_k q_{ki} \log(q_{ki})$$



concave
for convex funct. (cont.),

$$\sum_k \lambda_k f(x_k) \geq f(\sum_k \lambda_k x_k), \text{ if } \lambda_k \geq 0 \text{ and } \sum \lambda_k = 1$$

$$\begin{cases} f(x) = x \log x \\ \lambda_k = p_k, x_k = q_{ki} \end{cases}$$

$$\sum_k p_k q_{ki} \log(q_{ki}) \geq \sum_k p_k q_{ki} \log\left(\sum_k p_k q_{ki}\right)$$

$$\Rightarrow \sum_k p_k q_{ki} \log(q_{ki}) \geq q_i \log(q_i)$$

$$\Rightarrow -\sum_k p_k H_k(B) \geq -H(B)$$

$$\Rightarrow -H(B|A) \geq -H(B)$$

$$\boxed{H(B) \geq H(B|A)}$$

X \ Y	1	2	3	4	
1	$\frac{1}{8}$	$\frac{1}{16}$	$\frac{1}{32}$	$\frac{1}{32}$	$\left(\frac{1}{4}\right)$
2	$\frac{1}{16}$	$\frac{1}{8}$	$\frac{1}{32}$	$\frac{1}{32}$	$\left(\frac{1}{4}\right)$
3	$\frac{1}{16}$	$\frac{1}{16}$	$\frac{1}{16}$	$\frac{1}{16}$	$\left(\frac{1}{4}\right)$
4	$\frac{1}{4}$	0	0	0	$\left(\frac{1}{4}\right)$
	$\left(\frac{1}{2}\right)$	$\left(\frac{1}{4}\right)$	$\left(\frac{1}{8}\right)$	$\left(\frac{1}{8}\right)$	

$$\frac{3}{16} + \frac{1}{16}$$

find $H(X)$, $H(X|Y)$, $H(Y|X)$
 $H(XY)$

$$H(XY) = H(X) + H(Y|X)$$

$$H(X) = -\sum_x p_x \log(p_x)$$

$$= \frac{1}{2} \cdot 1 + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 = \frac{1}{2} + \frac{1}{2} + \frac{3}{8} = \frac{7}{8}$$

$$H(XY) = -\sum_x \sum_y p_{xy} \log(p_{xy})$$

$$= H(Y|X) = \sum_x p_x \left(-\sum_y p_{y|x} \log(p_{y|x}) \right)$$

$$H_1(Y) = \frac{1}{8} \times 3 + \frac{1}{16} \times 2 + \frac{1}{16} \times 2 = \frac{1}{8} + \frac{1}{8} = \frac{1}{4}$$

$$H_2(Y) =$$

$$H(XY) = H(Y) + H(X|Y)$$

$$= H(XY) - H(Y) = \frac{27}{8} - \frac{2 \times 8}{8} = \frac{17}{8}$$

$$H(Y|X) = \sum_x \sum_y p(x,y) \log \left(\frac{p(x)}{p(x,y)} \right)$$

$$= H(Y)_1 = \frac{1}{8} \log \left(\frac{4}{2} \right) + \frac{1}{16} \log \left(\frac{5}{2} \right) + \frac{1}{16} \log \left(\frac{16}{2} \right) + \frac{1}{4} \log \left(\frac{2}{2} \right)$$

$$= \frac{1}{8} + \frac{3}{16} + \frac{3}{16} + \frac{1}{4} = \frac{1}{2} + \frac{3}{8} = \left(\frac{7}{8} \right)$$

$$H(Y)_2 = \frac{1}{16} \log \left(\frac{16}{4} \right) \times 2 + \frac{1}{8} \log \left(\frac{8}{4} \right) = \frac{2}{4} + \frac{1}{8} = \left(\frac{5}{8} \right)$$

$$H(Y)_3 = \frac{1 \times 2 \times \log \left(\frac{32}{8} \right)}{32 \times 16 \times 8} + \frac{1}{16} \log \left(\frac{16}{8} \right) = \frac{1}{8} + \frac{1}{16} = \left(\frac{3}{16} \right)$$

$$H(Y)_4 = \frac{1 \times 2 \times \log \left(\frac{32}{8} \right)}{32 \times 16 \times 8} + \frac{1}{16} \log \left(\frac{16}{8} \right) = \frac{1}{16} + \frac{1}{8} = \left(\frac{3}{16} \right)$$

$$P(Y|X) = \frac{7}{8} \times \frac{1}{2} + \frac{3}{8} \times \frac{1}{4} + \frac{1}{8} \times \frac{3}{16} + \frac{1}{8} \times \frac{3}{16}$$

$$= \frac{7}{16} + \frac{3}{32} + \frac{2 \times 3}{4 \times 8 \times 16} = \frac{7}{16} + \frac{3}{32} + \frac{3}{64}$$

$$= \frac{28 + 6 + 3}{64} = \frac{37}{64}$$

$$\checkmark H(X|Y) = \sum_y p_y H(X|Y=y)$$

$$H(X|Y=1) = H\left(\frac{1}{2} \frac{1}{4} \frac{1}{8} \frac{1}{8}\right) = H(X|Y=2) = \frac{1}{4} \frac{1}{2} \frac{1}{8} \frac{1}{8}$$

$$H(X|Y=3) = H\left(\frac{1}{4} \frac{1}{4} \frac{1}{4} \frac{1}{4}\right) \quad H(X|Y=4) = (1, 0, 0, 0)$$

$$\frac{1}{2} + \frac{1}{4} + \frac{3}{8} = \left(\frac{7}{4} \right) \quad \frac{1}{4} + \frac{1}{2} + \frac{1 \times 2 \times 3}{8 \times 4} = \left(\frac{7}{4} \right)$$

(2)

(6)

$$H(X|Y) = \frac{1}{4} \cdot \frac{7}{4} + \frac{1}{4} \cdot \frac{7}{4} + \frac{1}{4} \cdot \frac{2 \cdot 4}{4} = \frac{7}{16} + \frac{7}{16} + \frac{2}{4} = \frac{22}{16} = \left(\frac{11}{8} \right)$$

$$H(XY) = H(Y) + H(X|Y) = H(X) + H(Y|X)$$

$$= 2 + \frac{11}{8} = \frac{27}{8} = H(Y|X)$$

$$H(Y|X) = \frac{16 + 11}{8} - \frac{7}{4} = \frac{27 - 14}{8} = \left(\frac{13}{8} \right)$$

$$H(Y|X) = \sum_x p_x H(Y|X=x)$$

$$H_1(Y) = H\left(\frac{1}{4}, \frac{1}{8}, \frac{1}{8}, \frac{1}{2}\right) = \frac{2}{4} + \frac{3}{8} + \frac{1}{2} = \left(\frac{7}{4}\right)$$

$$H_2(Y) = H\left(\frac{1}{4}, \frac{1}{2}, \frac{1}{4}, 0\right) = \frac{1}{2} + 2 \times \frac{2}{4} = \left(\frac{3}{2}\right)$$

$$H_3(Y) = H\left(\frac{1}{4}, \frac{1}{4}, \frac{1}{2}, 0\right) = \left(\frac{3}{2}\right)$$

$$H_4(Y) = H\left(\frac{1}{4}, \frac{1}{4}, \frac{1}{2}, 0\right) = \left(\frac{3}{2}\right)$$

$$H(Y|X) = \frac{1}{2} \cdot \frac{7}{4} + \frac{1}{4} \cdot \frac{3}{2} + \frac{1}{8} \times 2 \times \frac{3}{2}$$

$$= \frac{7}{8} + \frac{3}{8} + \frac{3}{8} = \left(\frac{13}{8}\right)$$

$$\frac{16}{8} - \frac{13}{8} = \frac{3}{8}$$

Remember $H(X, Y) = H(X) + H(Y|X) = H(Y) + H(X|Y)$

$$H(X, Y)$$

* The joint entropy $H(X, Y)$ corresponding to a pair of discrete random variables (X, Y) with prob. dist. $p(x, y)$ is given by

$$H(X, Y) = - \sum_x \sum_y p(x, y) \cdot \log(p(x, y)) = - \mathbb{E}_{x, y} [\log(p(x, y))]$$

and $H(X, Y) = H(X) + H(Y)$ when X, Y are mut. ind.

* Conditional Entropy: If there are 2 random variables $X, Y \sim p(x, y)$ then $H(Y|X)$ is given by

$$H(Y|X) = \sum_x p(x) \cdot H(Y|X=x) = - \sum_x p(x) \sum_y p(y|x) \log p(y|x)$$

$$= - \sum_x \sum_y p(x) \cdot p(y|x) \log(p(y|x))$$

$$\rightarrow = - \sum_x \sum_y p(x, y) \cdot \log(p(y|x))$$

$$\rightarrow = - \mathbb{E}_{x, y} \log(p(y|x))$$

P.T. $H(X, Y|Z) = H(X|Z) + H(Y|X, Z)$

$$H(X, Y|Z) = \sum_z P(z) \cdot H(X, Y|Z=z)$$

$$= - \sum_x \sum_y \sum_z P(z) \cdot P(x, y|z) \cdot \log(P(x, y|z))$$

$$= - \sum_x \sum_y \sum_z P(z) \cdot P(x, y, z) / P(z) \cdot \{ \log P(x, y, z) + \log P(z) \}$$

$$= - \sum_x \sum_y \sum_z P(z) P(x, y, z) P(z) \log P(x, y, z)$$

$$- \sum_x \sum_y \sum_z P(z) P(x, y, z) P(z) \log(P(z))$$

$$= - \sum_z P(z) \cdot P(x, y)$$

$$\text{RHS} = - \sum_x \sum_z P(x, z) \log(P(x|z))$$

$$= - \sum_x \sum_y \sum_z P(x, y, z) \cdot \log\left(\frac{P(x, y, z)}{P(x, z)}\right)$$

Soln. $H(X, Y|Z) = H(X, Y, Z) - H(Z)$

$$H(X|Z) = H(X, Z) - H(Z)$$

$$H(Y|X, Z) = H(X, Y, Z) - H(X, Z)$$

$$= H(X, Y, Z) - \{H(X, Z)\}$$

$$\Rightarrow H(X, Y|Z) = H(X, Z) + H(Y|X, Z)$$

(K.L. Divergence)

► Relative Entropy: It is the measure of distance between 2 probability distributions. It is also the measure of inefficiency of assuming one dist. while other dist. is ~~given~~ true.

$D(p||q) \rightarrow$ 'q' is assumed while 'p' is the true distribution

$$D(p||q) = \sum_{x \in X} p(x) \cdot \log \frac{p(x)}{q(x)} = \mathbb{E}_{p(x)} \left\{ \log \left\{ \frac{p(x)}{q(x)} \right\} \right\}$$

1) When $p(x) = 0$ $D(p||q) = 0$

2) When $q(x) = 0$ $D(p||q) \rightarrow \infty$

↳ can $D(p||q) = \infty$ even if $q(x) \neq 0$?

$$p: \{0,1\} = \{1-r, r\}$$

$$q: \{0,1\} = \{1-s, s\}$$

$$\left. \begin{aligned} D(p||q) &= (1-r) \cdot \log\left(\frac{1-r}{1-s}\right) + r \cdot \log\left(\frac{r}{s}\right) \\ D(q||p) &= (1-s) \cdot \log\left(\frac{1-s}{1-r}\right) + s \cdot \log\left(\frac{s}{r}\right) \end{aligned} \right\} D(p||q) \neq D(q||p)$$

* Relative Entropy is a non-negative quantity

► Information Inequality ($D(p||q) \geq 0$)

We have 2 probability distributions $p(x)$ & $q(x)$ (pmf)

Then $D(p||q) \geq 0$

with equality iff $p(x) = q(x) \forall x \in X$

Let $A = \{x : p(x) > 0\}$ be the support set of $p(x)$

Then

$$-D(p||q) = -\sum_{x \in A} p(x) \log\left(\frac{p(x)}{q(x)}\right) = \sum_{x \in A} p(x) \log\left(\frac{q(x)}{p(x)}\right)$$

for concave function $\log(x)$, $f'(x) = 1/x \leftarrow$ dec

$$E[f(x)] \leq f(E[x]) \text{ --- (I)}$$

using (I).

$$-D = \sum_{x \in A} p(x) \log\left(\frac{q(x)}{p(x)}\right) \leq \log\left(\sum_{x \in A} p(x) \cdot \frac{q(x)}{p(x)}\right) \leq \log\left(\sum_{x \in X} q(x)\right) \leq 0$$

$$\Rightarrow \boxed{D \geq 0}$$

P.T. $\ln x \leq x - 1$

$$e^x \geq 1+x \Rightarrow \text{let } x = \ln x$$

$$\Rightarrow e^{\ln x} \geq 1 + \ln x \Rightarrow \boxed{\ln(x) \leq x - 1} \text{ --- (II)}$$

We can prove $D \geq 0$ using (II) also

► Mutual Information ($I(X:Y)$)

Consider X and Y to be 2 r.v. with a joint prob. mass function $p(x, y)$. Let $p(x)$ & $p(y)$ be marginal prob. dist. (pmf).
The mutual information $I(X:Y)$ is given as the relative entropy between the joint prob. dist. $p(x, y)$ & $p(x)p(y)$.

$$I(X:Y) = \sum_x \sum_y p(x, y) \cdot \log \left\{ \frac{p(x, y)}{p(x)p(y)} \right\}$$

$$I(X:Y) = \mathbb{E}_{p(x, y)} \left[\log \left\{ \frac{p(x, y)}{p(x)p(y)} \right\} \right] = D(p(x, y) \| p(x)p(y))$$

$$= \sum_x \sum_y p(x, y) \log \left(\frac{p(x, y)}{p(x)} \right).$$

$$= - \sum_x \sum_y p(x, y) \cdot \log p(x) - \left(- \sum_x \sum_y p(x, y) \cdot \log (p(x|y)) \right)$$

$$= H(X) - H(X|Y)$$

$$\Rightarrow (*) \quad \boxed{I(X:Y) = H(X) - H(X|Y)}$$

Similarly $I(X:Y) = H(Y) - H(Y|X)$.

$$\Rightarrow H(X, Y) = H(X) + H(Y|X)$$

$$I(X:Y) = H(Y) - (H(X, Y) - H(X))$$

$$(*) \quad \boxed{I(X:Y) = H(X) + H(Y) - H(X, Y)} \quad \left. \vphantom{\boxed{I(X:Y) = H(X) + H(Y) - H(X, Y)}} \right\} \text{sth. like Inclusion exclusion principle.}$$

$$(*) \quad \boxed{I(X:X) = H(X) - H(X|X) = H(X) - 0 = H(X)}$$

\Rightarrow we can call entropy ($H(X)$) as self information $I(X:X)$

Theorem: Chain Rule of Entropy:

We have 'n' r.v. $X_1, X_2, X_3, \dots, X_n$ drawn acc. to prob. dist. $p(x_1, x_2, \dots, x_n)$. Then,

$$H(x_1, x_2, \dots, x_n) = \sum_{i=1}^n H(X_i | X_1, \dots, X_{i-1})$$

Proof

$$\hookrightarrow H(X_1, X_2, \dots, X_n) = \sum_{i=1}^n H(X_i | X_1, \dots, X_{i-1})$$

$$H(X_1, X_2) = H(X_2 | X_1) + H(X_1) \rightarrow \checkmark \text{ Base case}$$

$$\text{Let } H(X_1, \dots, X_k) = \sum_{i=1}^k H(X_i | X_1, \dots, X_{i-1})$$

$$\begin{aligned} H(X_1, \dots, X_k, X_{k+1}) &= H(X_{k+1} | X_1, X_2, \dots, X_k) + H(X_1, X_2, \dots, X_k) \\ &= H(X_{k+1} | X_1, \dots, X_k) + \sum_{i=1}^k H(X_i | X_1, \dots, X_{i-1}) \\ &= \sum_{i=1}^{k+1} H(X_i | X_1, \dots, X_{i-1}) \rightarrow \text{Ind. case} \end{aligned}$$

Hence Proved

► Conditional Mutual Information $I(X; Y | Z)$

Let X, Y, Z be 3 random variables, then conditional mutual info. b/w X & Y given Z is given as:-

$$\begin{aligned} I(X; Y | Z) &= H(X | Z) - H(X | Y, Z) \\ &= \mathbb{E} \left[\log \frac{P(X, Y | Z)}{P(X | Z) \cdot P(Y | Z)} \right] \end{aligned}$$

► Chain Rule for Mutual Info.

$$I(X_1, X_2, \dots, X_n; Y) = \sum_{i=1}^n I(X_i; Y | X_1, \dots, X_{i-1})$$

$$\begin{aligned} I(X_1, X_2, \dots, X_n; Y) &= H(X_1, X_2, \dots, X_n) - H(X_1, X_2, \dots, X_n | Y) \\ &= \sum_{i=1}^n H(X_i | X_1, \dots, X_{i-1}) - \sum_{i=1}^n H(X_i | X_1, \dots, X_n, Y) \\ &= \sum_{i=1}^n H(X_i; Y | X_1, \dots, X_{i-1}) \end{aligned}$$

► Conditional Relative Entropy : between 2 condⁿ prob dist $P(y|x)$ & $q(y|x)$ is given by

$$D(P(y|x) \parallel q(y|x)) = \sum_x p(x) \left\{ \sum_y p(y|x) \log \left[\frac{p(y|x)}{q(y|x)} \right] \right\}$$

$$= \sum_{x,y} p(x,y) \log \left[\frac{p(y|x)}{q(y|x)} \right] = E_x \left[\log \frac{p(y|x)}{q(y|x)} \right]$$

Show: $D(P(x,y) \parallel q(x,y))$

$$= D(P(x) \parallel q(x)) + D(P(y|x) \parallel q(y|x))$$

$$\text{LHS} = \sum_{x,y} p(x,y) \log \frac{p(x,y)}{q(x,y)} \quad \sum_y p(y) \log \frac{p(y)}{q(y)}$$

$$= \sum_{x,y} p(x,y) \log \left(\frac{p(y|x) p(x)}{q(y|x) q(x)} \right)$$

$$= \sum_{x,y} p(x,y) \log \left(\frac{p(x)}{q(x)} \right) + \sum_{x,y} p(x,y) \log \left(\frac{p(y|x)}{q(y|x)} \right)$$

$$= D(P(x) \parallel q(x)) + D(P(y|x) \parallel q(y|x))$$

Ex

$y \backslash x$	1	2
1	0	$3/4$
2	$1/8$	$1/8$

$\left(\frac{3}{4} \right)$
 $\left(\frac{1}{8} \right)$ $\left(\frac{2}{8} \right)$

$$H(x), H(x|y),$$

$$H(x) = \frac{1}{8} \times 3 + \frac{7}{8} \log \left(\frac{8}{7} \right) = 0.544 \text{ bits}$$

$$H(y) = \frac{3}{4} \log \left(\frac{4}{3} \right) + \frac{1}{4} \times 2$$

$$H(x,y) = H(x|y) + H(y)$$

$$H(x,y) = \frac{3}{4} \log \frac{4}{3} + 2 \times \frac{1}{8} \times 3 = \frac{3}{4} \left(1 + \log \frac{4}{3} \right)$$

$$H(x|y=1) = 0$$

$$H(x|y=2) = 1$$

$$H(x|y) = \frac{3}{4} + \frac{3}{4} \log \frac{4}{3} - \frac{3}{4} \log \frac{4}{3} - \frac{1}{2} \times 2$$

$$= \left(\frac{1}{4} \right)$$

Note that $H(x|y=1) > H(x)$

but $H(x|y)$ will always be $\leq H(x)$.

① for any 2 R.V. X, Y $I(X; Y) \geq 0$
(equality holds when X & Y are independent).

② ~~$D(p(y|x) || q(y|x))$~~ $D(p(y|x) || q(y|x)) \geq 0$
(equality when $p(y|x) = q(y|x)$ with $p(x) > 0$).

③. $I(X; Y | Z) \geq 0$ with equality iff X & Y are cond. independent given Z .

Create examples for equalities in ①, ② & ③.

Q. I am tossing a coin twice. Let X be the number of heads that occur. Then I toss it twice again, Y be no. of heads, now, this time

i) find $P(X < 2, Y > 1)$.

ii) find $I(X; Y)$.

$X \backslash Y$	0	1	2
0	$1/4$	$1/2$	$1/4$
1	$1/2$	$1/2$	$1/4$
2	$1/4$	$1/4$	$1/4$

i) $P(X < 2) \times P(Y > 1)$
 $= (1 - \frac{1}{4}) \cdot \frac{1}{4} = \frac{3}{16}$

$X: \quad 0 \quad 1 \quad 2$
 $\quad \frac{1}{4} \quad \frac{1}{2} \quad \frac{1}{4}$

ii) $I(X; Y) = D(p(x, y) || p(x) \cdot p(y))$
 $= 0$

Show that: $I(X; Y, Z) = I(X; Z) + I(X; Y | Z)$ * ①
 $I(X; Y, Z) = I(X; Y) + I(X; Z | Y)$ ②

①. $I(X; Y, Z) = \sum p(x, y, z) \cdot \log \frac{p(x, y, z)}{p(x) \cdot p(y, z)}$

$I(X; Z) = \sum p(x, z) \cdot \log \frac{p(x, z)}{p(x) \cdot p(z)}$

$I(X; Y | Z) = \sum p(x, y, z) \cdot \log \frac{p(x, y | z)}{p(x | z) \cdot p(y | z)}$

$$\textcircled{i} \quad I(X:Z) = H(X) - H(X|Z)$$

$$I(X:Y|Z) = H(X|Z) - H(X|Y,Z)$$

addly

$$I(X:Z) + I(X:Y|Z) = H(X) - H(X|Y,Z)$$

$$= I(X:Y,Z) \rightarrow \text{Thus proved.}$$

Similarly prove (ii)

$$\textcircled{ii} \quad I(X:Y) = H(X) - H(X|Y)$$

$$I(X:Z|Y) = H(X|Y) - H(X|Y,Z)$$

$$\Sigma = H(X) - H(X|Y,Z) = I(X:Y,Z) \quad \text{Hence Proved}$$

Let there are 3 r.v. X, Y, Z . They will form a Markov chain in the order $X \rightarrow Y \rightarrow Z$ if the conditional distribution of Z depends only upon Y and ~~is~~ conditionally independent on X .

$$p(x, y, z) = p(x) \cdot p(y|x) \cdot p(z|y).$$

$$\Rightarrow p(x, z|y) = \frac{p(x, y, z)}{p(y)} = \frac{p(x) \cdot p(y|x) \cdot p(z|y)}{p(y)}$$

$$= \frac{p(x, y)}{p(y)} \cdot p(z|y) = p(x|y) \cdot p(z|y).$$

Data Processing Inequality: If we have a Markov chain $X \rightarrow Y \rightarrow Z$ then $I(X:Y) \geq I(X:Z)$.

Proof.

$$\left. \begin{aligned} I(X:Y) &= I(X:Y,Z) - I(X:Z|Y) \quad \text{--- (i)} \\ I(X:Z) &= I(X:Y,Z) - I(X:Y|Z) \quad \text{--- (ii)} \end{aligned} \right\}$$

Note $Z|Y$ and X are ind. (since Markov chain).

$$\Rightarrow I(X:Z|Y) = 0 \text{ \& } I(X:Y|Z) \text{ exists } \geq 0$$

$$\Rightarrow I(X:Y) \geq I(X:Z)$$

HW

(1) Let $X = x_1, x_2$ and $Y = y_1, y_2$ are random bits such that $x_1 \oplus x_2 = y_1 \oplus y_2 = z$.

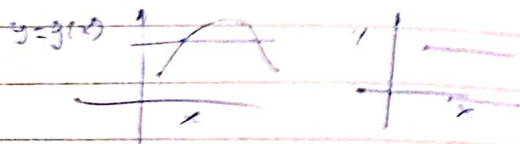
$z = 0$	$z = 1$
$\begin{pmatrix} x_1 & x_2 \\ (0,0) \\ (0,1) \\ (1,1) \\ (1,0) \end{pmatrix}$	$\begin{pmatrix} x_1 & x_2 \\ (0,0) \\ (0,1) \\ (1,0) \\ (1,1) \end{pmatrix}$
\textcircled{X}	\textcircled{Y}

Find $I(X:Y|Z)$

$I(X:Z|Y)$

$I(Y:Z|X)$

$$H(Y) \leq H(X)$$



HW

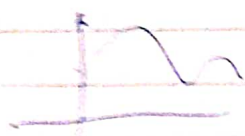
Let X & Y be 2 r.v. $Y = g(X)$ which of $H(X)$ or $H(Y)$ is greater

$$H(X) = \sum_x p(x) \log \frac{1}{p(x)}$$

$$\text{Let } g(x_1) = g(x_2) = \dots = g(x_n) = y$$

$$H(Y) = \sum_y \left(\sum_{x: y=g(x)} p(x) \log \frac{1}{\sum_{x: y=g(x)} p(x)} \right)$$

$$\sum_y \sum_{x: y=g(x)} p(x) \log \left(\frac{1}{\sum_{x: y=g(x)} p(x)} \right)$$



$$H(Y) = \sum_y \left(\sum_{x: y=g(x)} p(x) \log \frac{1}{\sum_{x: y=g(x)} p(x)} \right) \leq H(X) ?$$

Prove

► Jensen's Inequality

A function $f(x)$ will be called a convex function over an interval (a, b) if for $x_1, x_2 \in (a, b)$ and $0 \leq \lambda < 1$

$$f(\lambda x_1 + (1-\lambda)x_2) \leq \lambda f(x_1) + (1-\lambda)f(x_2)$$

Then $-f$ is concave.

Thm: If the function's 2nd derivative is ≥ 0 everywhere, then we call function to be convex.

Proof: $f(x) = f(x_0) + f'(x_0)(x-x_0) + \frac{f''(x^*)}{2}(x-x_0)^2 \rightarrow$ Taylor exp. around x_0
 $x_0 < x^* < x$

$$\Rightarrow f(x) - f(x_0) - f'(x_0)(x-x_0)$$

$$f(x) \geq f(x_0) + f'(x_0)(x-x_0)$$

1st case
take
 $x = x_1$

$$f(x_1) \geq f(x_0) + f'(x_0)[x_1 - \lambda x_1 - (1-\lambda)x_2]$$

$$x_0 = \lambda x_1 + (1-\lambda)x_2$$

$$x_1 - \lambda x_1 - x_2 + \lambda x_1 = (x_1 - x_2) - \lambda(x_1 - x_2) = (1-\lambda)(x_1 - x_2)$$

$$* f(x_1) \geq f(x_0) + f'(x_0)[(1-\lambda)(x_1 - x_2)] \quad \text{--- (1)}$$

2nd case:

$$x = x_2$$

$$x_0 = \lambda x_1 + (1-\lambda)x_2$$

$$\rightarrow f(x_2) \geq f(x_0) + f'(x_0)[\lambda(x_2 - x_1)] \quad \text{--- (2)}$$

multiplying (1) by λ & (2) by $(1-\lambda)$ & then adding

$$\lambda f(x_1) + (1-\lambda)f(x_2) \geq \lambda f(x_0) + \lambda(1-\lambda)f'(x_0)(x_1 - x_2) + (1-\lambda)f(x_0) + \lambda(1-\lambda)f'(x_0)(x_2 - x_1)$$

$$\Rightarrow \lambda f(x_1) + (1-\lambda)f(x_2) \geq f(x_0)$$

$$= f(\lambda x_1 + (1-\lambda)x_2)$$

$$\Rightarrow \boxed{\lambda f(x_1) + (1-\lambda)f(x_2) \geq f(\lambda x_1 + (1-\lambda)x_2)} \rightarrow f \text{ is convex.}$$

7)

We have a convex function f and x is a random variable. Then

$$\boxed{E[f(x)] \geq f(E[x])} \rightarrow \text{Jensen's Inequality.}$$

Proof: for 2 mass point distribution, it directly follows from being convex.
 $[p_1 f(x_1) + p_2 f(x_2) \geq f(p_1 x_1 + p_2 x_2)]$

Let this be true for $k-1$ mass point. & let $p_i' = \frac{p_i}{1-p_k}$

$$\Rightarrow \sum_{i=1}^k p_i f(x_i) = p_k f(x_k) + (1-p_k) \left\{ \sum_{i=1}^{k-1} p_i' f(x_i) \right\}$$

$$\geq p_k f(x_k) + (1-p_k) f\left(\sum_{i=1}^{k-1} p_i' x_i\right)$$

$$\geq f\left(p_k x_k + (1-p_k) \sum_{i=1}^{k-1} p_i' x_i\right)$$

$$\geq f\left(\sum_{i=1}^k p_i x_i\right)$$

Now since

$$\sum_{i=1}^k p_i f(x_i) \geq f\left(\sum_{i=1}^k p_i x_i\right)$$

$$\Rightarrow \boxed{E[f(x)] \geq f(E[x])} \checkmark$$

Note

$$\sum p_i' = \frac{p_1 + p_2 + \dots + p_{k-1}}{1-p_k} = \frac{1-p_k}{1-p_k} = 1$$

$$\frac{p_1 + p_2 + \dots + p_{k-1}}{1-p_k} = \frac{1-p_k}{1-p_k} = 1$$

Q

Show that $H(x) \leq \log |X|$

where

$|X|$ denotes no. of elements x in the range of X

with equality iff X has uniform distribution.

Let X have uniform dist:

$$H(x) = -\left(\frac{1}{n} \log\left(\frac{1}{n}\right) + \frac{1}{n} \log\left(\frac{1}{n}\right) + \dots + \frac{1}{n} \log\left(\frac{1}{n}\right)\right)$$

$$\text{At } \left. \begin{matrix} p_1 \rightarrow p_1 + \epsilon \\ p_2 \rightarrow p_2 - \epsilon \end{matrix} \right\} \Rightarrow -H'(x) = -\left(\left(\frac{1}{n} + \epsilon\right) \log\left(\frac{1}{n} + \epsilon\right) + \left(\frac{1}{n} - \epsilon\right) \log\left(\frac{1}{n} - \epsilon\right) + \dots\right)$$

$$\leq \frac{p_1 + \epsilon}{n} \log \frac{n}{p_1 + \epsilon} + \frac{p_2 - \epsilon}{n} \log \frac{n}{p_2 - \epsilon} + \dots$$

$$\leq \log n - H(x)$$

Sol Let $u(x)$ be uniform dist $u(x) = \frac{1}{n} \quad \forall x \in X$ ($n = |X|$)
 $p(x)$ be any prob dist of X .

$$\begin{aligned} \rightarrow D(p \| u) &= \sum_{x \in X} p(x) \cdot \log \frac{p(x)}{u(x)} = \sum p(x) \cdot \log(n \cdot p(x)) \\ &= \sum p(x) \log n + \sum p(x) \log(p(x)) \end{aligned}$$

$$D(p \| u) = \log n - H(x) \geq 0$$

$$\Rightarrow \boxed{H(x) \leq \log n}$$