

Name : Ajay ray

course : Entropy & Information

Roll No : 2021102032

M	T	W	T	F	S	S
Page No :					YOUVA	
Date :						

Q1.

Q1.

for unbiased Dice

- probability of each outcome = $1/6$

- favourable outcome = 6

= no. of trials req. of 6 = n

we have to find.

$\Rightarrow H(n) = ?$

Now

$P(n=1)$: prob. of getting 6 in 1st chance = $1/6$

for

$n=2$, $P(n)$ = prob. of getting 6 in 2nd chance = $\frac{5}{6} \times \frac{1}{6}$

for $n=3$

$P(n)$ = " " " " in 3rd = $\frac{5}{6} \times \frac{5}{6} \times \frac{1}{6}$

for $n=4$

$P(n)$ = " " " " 4th = $\frac{5}{6} \times \frac{5}{6} \times \frac{5}{6} \times \frac{1}{6}$

$P(n=n)$ = " " " " n^{th} = $\left(\frac{5}{6}\right)^{n-1} \times \frac{1}{6}$

we know that

$$\therefore H(n) = - \sum_{n} P(n) \log_2 P(n)$$

$$= - \left[\frac{1}{6} \log_2 \frac{1}{6} + \frac{5}{6} \times \frac{1}{6} \log_2 \frac{5}{6} \times \frac{1}{6} + \dots + \left(\frac{5}{6}\right)^{n-1} \frac{1}{6} \log_2 \left(\frac{5}{6}\right)^{n-1} \frac{1}{6} \right]$$

$$= - \frac{1}{6} \left[\frac{5}{6} \log_2 \frac{5}{6} + 2 \left(\frac{5}{6}\right)^2 \log_2 \frac{5}{6} + \dots + n \left(\frac{5}{6}\right)^{n-1} \log_2 \frac{5}{6} \right]$$

now we can say that

from here we have two geometric series

$H(n) = a + b$ where

$$a = - \frac{1}{6} \log_2 \frac{1}{6} \left[1 + \frac{5}{6} + \left(\frac{5}{6}\right)^2 + \dots + \left(\frac{5}{6}\right)^{n-1} \right]$$

$$b = - \frac{1}{6} \log_2 \frac{5}{6} \left[\frac{5}{6} + 2 \left(\frac{5}{6}\right)^2 + 3 \left(\frac{5}{6}\right)^3 + \dots + (n-1) \left(\frac{5}{6}\right)^{n-1} \right]$$

from here we can conclude that

a = geometric series

b = arithmetic geometric series

Now for a,

first term = 1

and common ratio = $5/6$

$$a = -\frac{1}{6} \log \frac{1}{6} \times 6 = -\frac{1}{6} \log \left(\frac{1}{6}\right)^{\times 6}$$

$$a = -\log_2 \frac{1}{6} = \log_2 6 = \boxed{\log_2 6}$$

Now for b

common ratio, $r = 5/6$

$$\text{constant } k = -\frac{1}{6} \log \frac{5}{6}$$

$$\text{so } b = k [1r^1 + 2r^2 + 3r^3 + \dots + n-1 r^{n-1}] \quad (1)$$

after ~~now~~ dividing b by r

$$\frac{b}{r} = k [1 + 2r + 3r^2 + 4r^3 + \dots + n-1 r^{n-2}] \quad (2)$$

the sum of RHS of eqn (2) is derivative of geometric series summed from $k=1$ to $n-1$.
and sum of geometric series $\frac{r^{n-1} - 1}{r-1}$

$$\Rightarrow \frac{b}{r} = k \frac{d}{dr} \left(\left[\frac{r^{(n-1)-1} - 1}{r-1} \right] - 1 \right)$$

$$= k \frac{d}{dr} \left[\frac{r^n - 1}{r-1} \right]$$

$$= k [(n-1)r^n - nr^{n-1} + 1] / (r-1)^2$$

$$b = k[(n-1)r^{n-1} - nr^n + r] / (r-1)^2$$

$\therefore r = \frac{5}{6} < 1$ so infinite sum will converge to $\frac{r}{(r-1)^2}$

$$\therefore b = -\frac{1}{6} \log \frac{5}{6} \times \frac{5/6}{1/36}$$

$$b = -\log 5/6 \times 5$$

$$b = 5 \log_2 6/5$$

$$\therefore H(n) = \log_2 6 + 5 \log_2 6/5$$

$$= 6 \log_2 6 - 5 \log_2 5$$

$$= \log_2 \frac{6^6}{5^5} = \log_2 \left(\frac{46656}{3125} \right)$$

$$= \log_2 (14.89952)$$

$$= 3.91185 \text{ bits}$$

$$\therefore \boxed{H(n) = 3.91185 \text{ bits}}$$

for probability $P(n) = \frac{1}{m} \& n$ is no of trials for getting n

we can generalise $H(n) = m \log_2 m - (m-1) \log_2 m$

(b) now die is biased
given that

	1	2	3	4	5	6
distribution	$\frac{1}{4}$	$\frac{1}{4}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{1}{8}$

n_i = no. of trials needed to get i th value

we have to find $H(n_i)$

now, as prob. distribution of (1,2) and (3,4,5,6)
have same $\frac{1}{4}$, $\frac{1}{8}$ respectively

we can say that

$$H(n_1) = H(n_2)$$

$$H(n_3) = H(n_4) = H(n_5) = H(n_6)$$

now

let's calculate for x_1

$p(n_1=1)$ = prob. of getting 1 in 1st chance = $\frac{1}{4}$

$$p(n_1=2) = \dots \dots \dots \text{2nd} = \frac{1}{4} \times \frac{3}{4}$$

$$p(n_1=n) = \dots \dots \dots \text{nth} = \frac{1}{4} \left(\frac{3}{4}\right)^{n-1}$$

$$\therefore H(n_1) = - \sum_{n \in X_1} p(n) \log_2 p(n) \quad \text{where } n \in [1, \infty)$$

$$= - \sum_n \frac{1}{4} \left(\frac{3}{4}\right)^{n-1} \log_2 \left[\frac{1}{4} \left(\frac{3}{4}\right)^{n-1} \right]$$

now, for of form our generalised eqⁿ for prob

$p(n) = \frac{1}{m}$ if n is no. of trials
required to get x , then

$$H(n) = m \log_2 m - (m-1) \log_2 (m-1) \quad \text{eq ①}$$

(derived in last part)

$$\begin{aligned}
 3. \quad H(X_1) &= H(X_2) = 4 \log_2 4 - 3 \log_2 3 \\
 &= 8 - 3(1.58) \\
 &= 3.2153 \text{ bits}
 \end{aligned}$$

$$\therefore \boxed{H(X_1) = H(X_2) = 3.2153 \text{ bits}}$$

Now let's calculate entropy of X_3

$$P(X_3=1) = \text{prob. of getting 3 in 1st chance} = \frac{1}{8}$$

$$\text{for } n_3=2 \quad \text{prob. of getting 3 in 2nd chance} = \frac{1}{8} \times \frac{7}{8}$$

$$\text{for } n_3=n \quad P(n_3=n) = \frac{1}{8} \left(\frac{7}{8}\right)^{n-1}$$

$$\therefore H(X_3) = - \sum_{x \in X_3} P(x) \log_2 P(x); \quad x \in [1, \infty)$$

$$= - \sum_{x \in X_3} \frac{1}{8} \left(\frac{7}{8}\right)^{n-1} \log_2 \left[\frac{1}{8} \left(\frac{7}{8}\right)^{n-1} \right]$$

now putting $m=8$ in eqⁿ (1) (Generalised for \log_2)

$$= 8 \log_2 8 - 7 \log_2 7 = 24 - 7(2.8)$$

$$\boxed{= H(X_4) = H(X_5) = H(X_6) = 4.9 \text{ bits}}$$

Q.2 we know that

the axioms of uniqueness theorem are

cond. 1: for a given n and prob. dist p . ($\sum_{i=1}^n P_i = 1$). ~~the~~
for the function $H(P_1, P_2, \dots, P_n)$ which gives at
it highest value at $P_k = \frac{1}{n} \forall k \in \{1, 2, \dots, n\}$

cond. 2: let $A \in B$ be two r.v
then

$$H(A, B) = H(A) + \sum_a P(A=a) \cdot H(B/A=a)$$

cond. 3: $H(P_1, P_2, \dots, P_n) = H(P_1, P_2, \dots, P_n, 0)$.

now let S_1 and S_2 satisfy the condition for $H(n)$
but S_2 does not. then function will have
its maximum ~~distance~~ at uniform distance.

now

$$\text{let } H(n) = f(x)$$

we can find maxima of $f(x)$ by partially
differentiating w.r.t P_1, P_2, \dots, P_n and
will set derivative as 0

$$\Rightarrow \frac{\partial f}{\partial P_1} = \frac{\partial f}{\partial P_2} = \frac{\partial f}{\partial P_3} = \frac{\partial f}{\partial P_4} = \dots = \frac{\partial f}{\partial P_n} = 0$$

when $P_1 = P_2 = \dots = P_n = \frac{1}{n}$ let $\frac{\partial f}{\partial P_i}$ be a linear eqⁿ

then

$$\frac{\partial f}{\partial P_i} = a \cdot P_i + b = 0 \Rightarrow P_i = -\frac{b}{a}$$

let $b = +1, a = -n$

$$\Rightarrow \frac{\partial f}{\partial P_i} = -n P_i + 1 = 0 \quad (\text{by substituting the value})$$

using eqⁿ we can choose $a = -2, b = 2/n$
which satisfy $P_i = -b/a = 1/n$

$$\Rightarrow \boxed{\frac{\partial f}{\partial p_i} = \frac{2}{n} - 2 p_i}$$

we have to note that this condition must hold for all i ($\forall i$)

and value of f can easily be calculated by using

$$\boxed{f(n) = \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right)} \quad \text{--- ①}$$

we can verify eq ① for condition 1

$$\text{since } \frac{\partial f}{\partial p_i} = \frac{\partial}{\partial p_i} \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right)$$

$$= \frac{2}{n} - 2 p_i = 0 \Rightarrow \frac{2}{n} = 2 p_i \text{ or } \boxed{p_i = \frac{1}{n}}$$

also here we can check condition

$$\text{or } f(p_1, p_2, \dots, p_n) = \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right)$$

$$f(p_1, p_2, \dots, p_n, 0) = \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right) + \left(\frac{2(0)}{n} - (0)^2 \right) =$$

$$f(p_1, p_2, \dots, p_n)$$

$$\Rightarrow \text{since } f(p_1, p_2, \dots, p_n) = \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right) \text{ satisfy}$$

both the conditions axioms condition 1 and condition 3 of uniqueness theorem but does not satisfy the axiom 2 (cond 2)

thus

$$H(p_1, p_2, \dots, p_n) = \sum_{k=1}^n \left(\frac{2 p_k}{n} - p_k^2 \right) \text{ can be one}$$

such measure of entropy.