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Section :- E

Branch :- CSE

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Sheet :- Tuit-4

Ans-1  $T(n) = 3T(n/2) + n^2$

$$a = 3 \quad b = 2$$

$$k = \log_2 3 = 1.58$$

$$n^2 > n^{1.58}$$

$$\therefore \underline{O(n^2)}$$

Ans-2  $T(n) = 4T(n/2) + n^2$

$$k = \log_2 4 = 2$$

$$n^2 = n^2$$

$$\therefore \underline{O(n^2 \log n)}$$

Ans-3  $T(n) = T(n/2) + 2^n$

here  $F(n)$  is not a polynomial

$\therefore$  we can not apply Master theorem.

Ans-4  $T(n) = 2^n T(n/2) + n^n$

yes

Here, this recurrence relation can't be solved using Master's Method.

Ans-5  $T(n) = 16T(n/4) + n$

$$k = \log_4 16 = 2$$

$$n^2 > n$$

$$\therefore \underline{\underline{O(n^2)}}$$

Ans-6  $T(n) = 2T(n/2) + n \log n$

$$\therefore \text{Case } k = 1 \quad (\text{by normal Master theorem})$$

Now by using extended Master's theorem.

$$T(n) = aT(n/b) + O(n^k \log^p n)$$

$$a = b^k \cdot 2 = 2$$

$$T(n) = O(n^{\log_2 2} \log^{1+1} n)$$

$$O(n \log^2 n)$$

Ans-7  $T(n) = 2T(n/2) + n \log^{-1} n$

using extended Master's theorem.

$$T(n) = aT(n/b) + O(n^k \log^{p_0} n)$$

$$a = 2, \quad b = 2, \quad k = 1, \quad p_0 = -1$$

$$n = -1$$

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$$\therefore T(n) = O(n^{\log_2 2} \log \log n)$$

$$= O(n \log \log n)$$

Ans-8  $T(n) = 2T(n/4) + n^{0.51}$

$$a = 2, b = 4$$

$$\text{Case } k = \log_4 2 = 0.5$$

$$n^{0.5} < n^{0.51}$$

$$\therefore O(n^{0.51})$$

Ans-9  $T(n) = 0.5T(n/2) + n^{-1}$

As  $a < 1$ ,  $\therefore$  Master's theorem's can not apply here.

Ans-10  $T(n) = 16T(n/4) + n!$

$$k = \log_4 16 = 2$$

$$n^2 < n!$$

$$\therefore \underline{O(n!)}$$

Ans-11  $T(n) = 4T(n/2) + \log n$

here,  $a = 4, b = 2, k = 0, n = 1$

$\therefore$  using Extended Master's theorem.

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$$T(n) = a T(n/b) + O(n^k \log^h n)$$

$$a > b^k$$

$$4 > 2^0$$

$$\therefore O(n \log_2^4)$$

$$\therefore O(n^2)$$

Ans-12  $T(n) = \sqrt{n} T(n/2) + \log n$

here  $a$  is not constant, so Master theorem's can not apply here.

Ans-13  $T(n) = 3T(n/2) + n$

$$k = \log_2 3 = 1.58$$

$$n^{1.58} > n$$

$$= O(n^{1.58})$$

Ans-14  $T(n) = 3T(n/3) + \sqrt{n}$

$$k = \log_3 3 = 1$$

$$n^1 > \sqrt{n}$$

$$= O(n)$$

Ans-15  $T(n) = 4T(n/2) + n$

$$k = \log_2 4 = 2$$

$$n^2 > n$$

$$\therefore \underline{O(n^2)}$$

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Ans-16  $T(n) = 3T(n/4) + n \log n$

using extended master's theorem

$$T(n) = a T(n/b) + \Theta(n^k \log^h n)$$

here  $a = 3, b = 4, k = 2, h = 2$

$$a < b^k$$

$$3 < 4^2$$

$$h > 0$$

$$\therefore T(n) = \Theta(n^k \log^h n)$$

$$= \underline{\underline{\Theta(n \log n)}}$$

Ans-17  $T(n) = 3T(n/3)$

$$c = \log_3 3 = 1$$

$$m^1 = n$$

$$\therefore \underline{\underline{\Theta(n \log n)}}$$

Ans-18  $T(n) = 6T(n/3) + n^2 \log n$

using extended master's theorem

$$T(n) = a T(n/b) + \Theta(n^k \log^h n)$$

$a = 6, b = 3, k = 2, h = 1$

$$a < b^k$$

$$6 < 3^2$$

$$h > 0$$

$$\therefore T(n) = \Theta(n^k \log^h n)$$

$$= \underline{\underline{\Theta(n^2 \log n)}}$$

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Ans-19  $T(n) = .4T(n/2) + n \log^{-1} n$   
using Master theorem's extended.

$$T(n) = aT(n/b) + O(n^k \log^h n)$$

$$a = .4, b = .2, k = 1, h = 1$$

$$a > b^k$$

$$.4 > 2^1$$

$$\therefore T(n) = O(n^{\log_2 .4})$$

$$= O(n^2)$$

Ans-20  $T(n) = .64T(n/8) - n^2 \log n$

As  $f(n)$  is negative.

$\therefore$  we can not apply Master's theorem.

Ans-21  $T(n) = .7T(n/3) + n^2$

using Master theorem

$$k = \log_3 7 = 1.77$$

$$n^{1.77} < n^2$$

$$\therefore \underline{O(n^2)}$$

Ans-22

here, we can not apply Master's theorem.

yes