



# CHANDIGARH UNIVERSITY

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## UNIVERSITY INSTITUTE OF ENGINEERING

### Bachelor of Engineering (Computer Science & Engineering)

### Operating System (20CST/ITT-313)

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Introduction to Operating System  
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# System Protection and Security



# Cryptography as a Security Tool

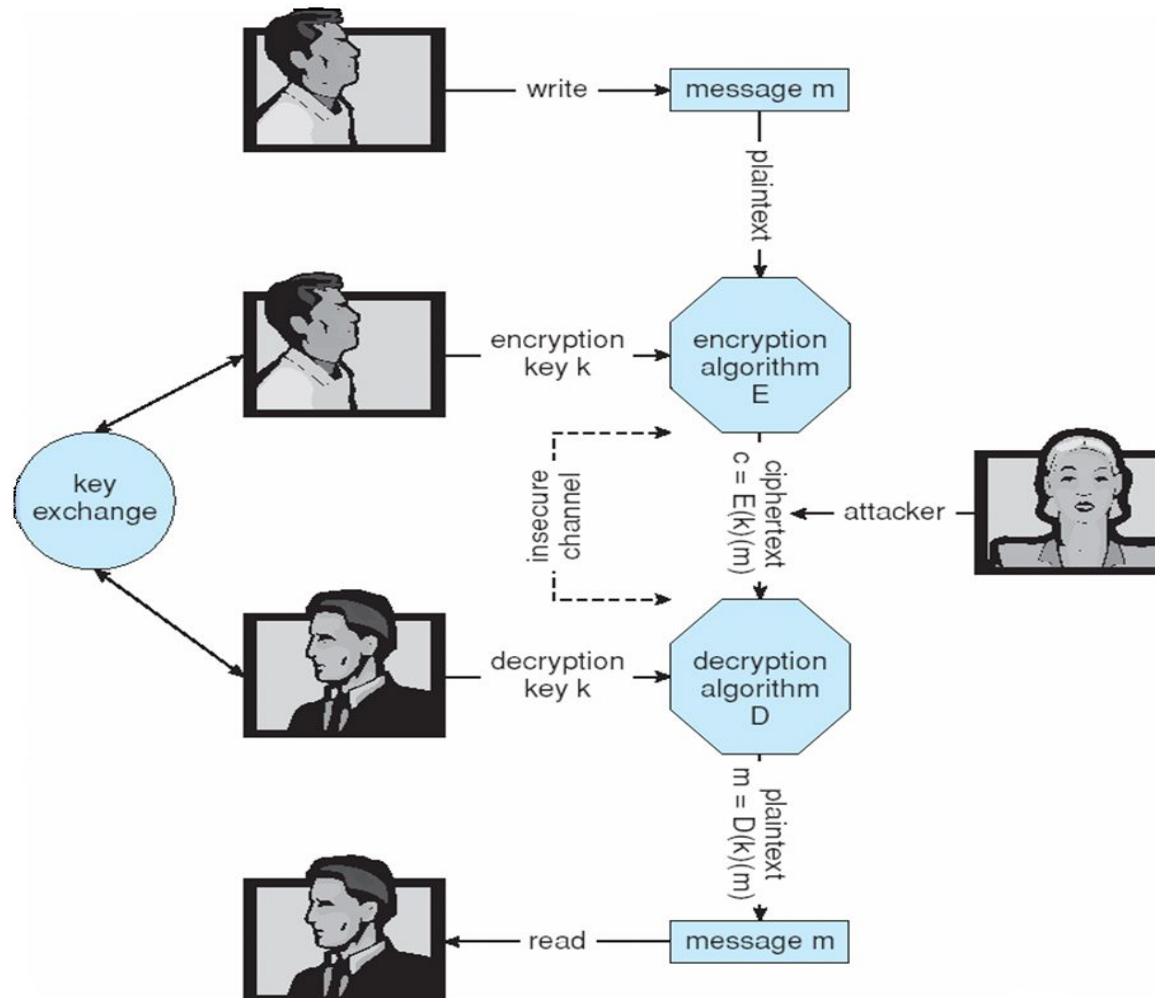
- Broadest security tool available
  - Internal to a given computer, source and destination of messages can be known and protected
    - OS creates, manages, protects process IDs, communication ports
  - Source and destination of messages on network cannot be trusted without cryptography
    - Local network – IP address?
      - Consider unauthorized host added
    - WAN / Internet – how to establish authenticity
      - Not via IP address



# Cryptography

- Means to constrain potential senders (*sources*) and / or receivers (*destinations*) of *messages*
  - Based on secrets (**keys**)
  - Enables
    - Confirmation of source
    - Receipt only by certain destination
    - Trust relationship between sender and receiver

# Secure Communication over Insecure Medium





# Encryption

- Encryption algorithm consists of
  - Set  $K$  of keys
  - Set  $M$  of Messages
  - Set  $C$  of ciphertexts (encrypted messages)
  - A function  $E : K \rightarrow (M \rightarrow C)$ . That is, for each  $k \in K$ ,  $E(k)$  is a function for generating ciphertexts from messages
    - Both  $E$  and  $E(k)$  for any  $k$  should be efficiently computable functions
  - A function  $D : K \rightarrow (C \rightarrow M)$ . That is, for each  $k \in K$ ,  $D(k)$  is a function for generating messages from ciphertexts
    - Both  $D$  and  $D(k)$  for any  $k$  should be efficiently computable functions
- An encryption algorithm must provide this essential property: Given a ciphertext  $c \in C$ , a computer can compute  $m$  such that  $E(k)(m) = c$  only if it possesses  $D(k)$ 
  - Thus, a computer holding  $D(k)$  can decrypt ciphertexts to the plaintexts used to produce them, but a computer not holding  $D(k)$  cannot decrypt ciphertexts
  - Since ciphertexts are generally exposed (for example, sent on the network), it is important that it be infeasible to derive  $D(k)$  from the ciphertexts



# Symmetric Encryption

- Same key used to encrypt and decrypt
  - $E(k)$  can be derived from  $D(k)$ , and vice versa
- DES is most commonly used symmetric block-encryption algorithm (created by US Govt)
  - Encrypts a block of data at a time
- Triple-DES considered more secure
- Advanced Encryption Standard (**AES**),
- RC4 is most common symmetric stream cipher, but known to have vulnerabilities
  - Encrypts/decrypts a stream of bytes (i.e., wireless transmission)
  - Key is a input to psuedo-random-bit generator
    - Generates an infinite **keystream**



# Asymmetric Encryption

- Public-key encryption based on each user having two keys:
  - public key – public key used to encrypt data
  - private key – key known only to individual user used to decrypt data
- Must be an encryption scheme that can be made public without making it easy to figure out the decryption scheme
  - Most common is RSA block cipher
  - No efficient algorithm is known for finding the prime factors of a number which is product of two large prime numbers.





# Asymmetric Encryption (Cont.)

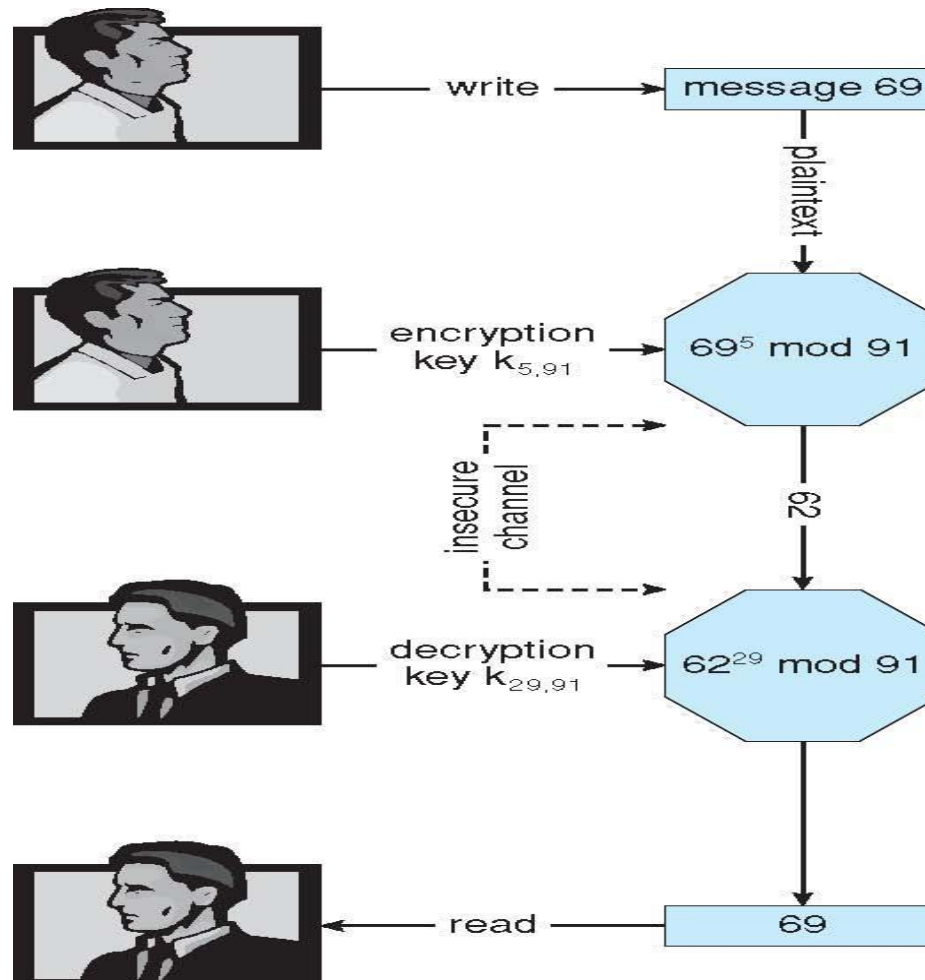
- Formally, it is computationally infeasible to derive  $D(k_d, N)$  from  $E(k_e, N)$ , and so  $E(k_e, N)$  need not be kept secret and can be widely disseminated
  - $E(k_e, N)$  (or just  $k_e$ ) is the **public key**
  - $D(k_d, N)$  (or just  $k_d$ ) is the **private key**
  - $N$  is the product of two large, randomly chosen prime numbers  $p$  and  $q$  (for example,  $p$  and  $q$  are 512 bits each)
  - Encryption algorithm is  $E(k_e, N)(m) = m^{k_e} \bmod N$ , where  $k_e$  satisfies  $k_e k_d \bmod (p-1)(q-1) = 1$
  - The decryption algorithm is then  $D(k_d, N)(c) = c^{k_d} \bmod N$



# Asymmetric Encryption Example

- For example. make  $p = 7$  and  $q = 13$
- We then calculate  $N = 7 * 13 = 91$  and  $(p-1)(q-1) = 72$
- We next select  $k_e$  relatively prime to 72 and  $< 72$ , yielding 5
- Finally, we calculate  $k_d$  such that  $k_e k_d \bmod 72 = 1$ , yielding 29
- We now have our keys
  - Public key,  $k_e, N = 5, 91$
  - Private key,  $k_d, N = 29, 91$
- Encrypting the message 69 with the public key results in the cyphertext 62
- Cyphertext can be decoded with the private key
  - Public key can be distributed in cleartext to anyone who wants to communicate with holder of public key

# Encryption and Decryption using RSA Asymmetric Cryptography





## Cryptography (Cont.)

- Note symmetric cryptography based on transformations, asymmetric based on mathematical functions
  - Asymmetric much more compute intensive
  - Typically not used for bulk data encryption

## Video Links

<https://www.edureka.co/blog/what-is-cryptography/>

<https://book.cyberyozh.com/comprehensive-encryption-of-operating-system-or-hard-disk-drive/>

# References

- <https://www.unf.edu/public/cop4610/ree/Notes/PPT/PPT8E/CH15-OS8e.pdf>
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- <https://www.coursehero.com/file/19323929/Operating-System-Threats-and-Vulnerabilities/>
- [https://www.cs.uic.edu/~jbell/CourseNotes/OperatingSystems/15\\_Security.html](https://www.cs.uic.edu/~jbell/CourseNotes/OperatingSystems/15_Security.html)
- <https://devqa.io/security-threats-attack-vectors/>
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