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Assignment-8

1. Explain log based recovery technique with example.

→ log-based recovery is a technique where all changes made to the database are recorded in a log file. If a crash occurs the system uses the log to redo and undo transactions key components:

- log: records transaction ID, operation old value, new value
- write:- Ahead logging (WAL): log must be written before actual data is changed.

Example:- Assume transaction T_1 updates A from 100 to 150.

Log

$\langle T_1, \text{start} \rangle$

$\langle T_1, A, 100, 150 \rangle$

$\langle T_1, \text{commit} \rangle$

If the system crashes after T_1 commit we redo the change. If it crashed before T_1 commit, we undo it using old value.

2. Explain shadow paging with suitable example and diagram.

→ shadow paging is a recovery technique where two copies of the database are maintained

3. Write a note on serializability
→ serializability ensures that the current execution of transaction result in a database state that would be same as if the transactions were executed one after another (serially)

Types:-

Conflict serializability:-
Based on reordering non conflicting operations
view serializability:-
Based on final state even if operation are reordered differently.
Based on final state even if operations are reordered prevents issues like lost updates and inconsistent analysis

4. Explain the 2 phase locking protocol with respect to DBMS.

→ 2PL ensures serializability by dividing a transaction into two phases:

1. Growing phase:-

- Transaction can acquire locks
- No locks can be released.

2. Shrinking phase:-

- Transaction can release locks
- Cannot acquire any new locks

ex:-

T_1 :- Read and write data item A

T_2 :- Read and write data item B.

If both follow 2PL, deadlocks and inconsistency can be avoided.

5. Describe deadlock avoidance using suitable example.
→ Deadlock occurs when two or more transactions wait forever for each other to release locks.

Deadlock avoidance techniques -

- wait-die scheme
older transactions wait, younger ones are rolled back
- wound-wait.
older transactions wound (force roll back of) younger ones

Example (wait die)

- T_1 (timestamp = 10) wants a lock held by T_2 (timestamp = 20)
- T_1 is older → T_1 waits.
- T_2 wants a lock held by T_1
 T_2 is younger → T_2 did rollback