#### CS 4530: Fundamentals of Software Engineering

Module 10.4: Case Studies

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# Learning Goals for this Lesson

- By the end of this lesson you should be able to:
  - Briefly describe several typical examples of distributed systems
  - Briefly describe how each of them deals with scalability, fault tolerance, etc.

# Case Study 1: the Network File System NFS

- NFS is a distributed file system: multiple clients can read/write the same files
- Created in 1984, still widely used
- In a UNIX (POSIX-compliant) operating system, files are stored in a tree from "/"
- Access a remote file by name like /username@remotehost/path/to/remote/file
- Or you could "mount" a remote filesystem to access it as if it were local.

# NFS is a Monolithic Shared Filesystem

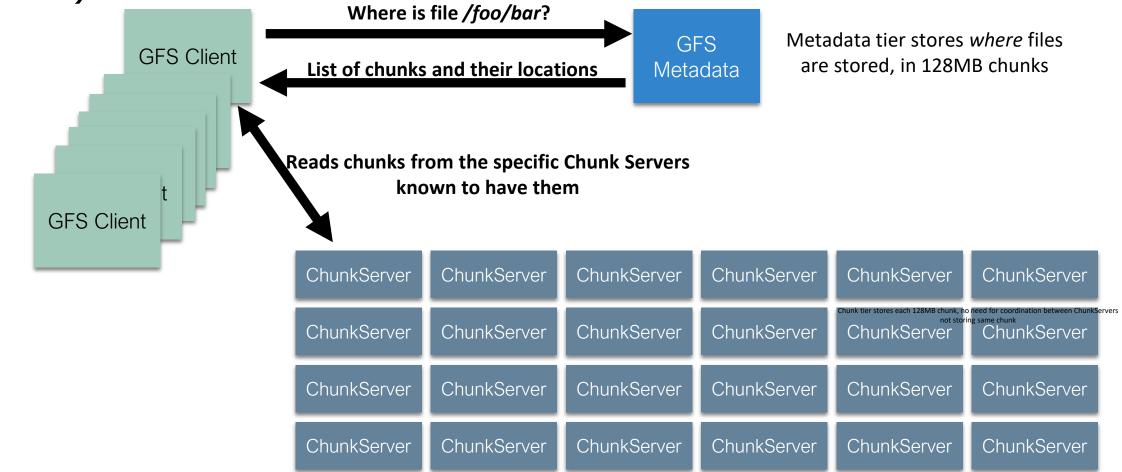
- All files are stored on a single server
- To list files in a directory, clients make request to server
- To read or write files, clients make request to server
- Clients might "lock" files to prevent concurrent updates
- Assuming that scale, throughput, fault tolerance requirements are relatively low, this is an acceptable architecture
- This architecture is the easiest to build fast and correctly

# Case Study 2: GFS (Google File System, ~2010)

- Stated requirements:
- "High sustained bandwidth is more important than low latency. Most of our target applications place a premium on processing data in bulk at a high rate, while few have stringent response time requirements for an individual read or write."

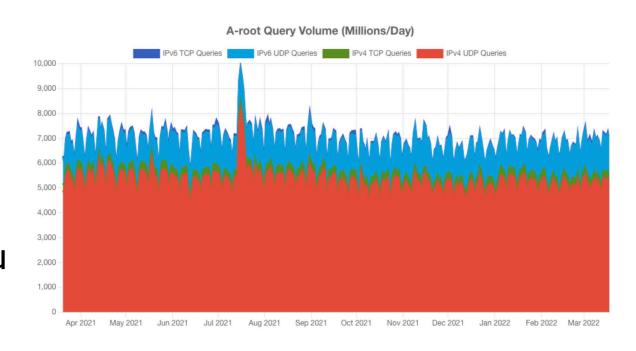
#### GFS is a tiered filesystem with two tiers: Metadata and File Chunks

• Example: GFS (Google File System, c 2010)



# Case Study 3: Domain Name System (DNS)

- Nodes (hosts) on a network are identified by IP addresses
- E.g.: 142.251.41.4
- We humans prefer something easier to remember: calendar.google.com, facebook.com, www.khoury.northeastern.edu
- We need to keep a directory of domain names and their addresses
- We also need to make sure everybody gets directed to the correct host



https://a.root-servers.org/metrics

# Requirements for the DNS system

- Need to handle millions of DNS queries per second
- Not immediately obvious how to scale: how do we maintain replication, some measure of consistency?



# DNS distributed system requirements

- We need a scalable solution
  - New hosts keep being added
  - Number of users increases
  - Need to maintain speed/responsiveness
- We need our service to be available and fault tolerant
  - It is a crucial basic service
  - A problematic node shouldn't "crash the internet"
  - Reads are more important that writes: far more queries to resolve records than to update them
- Global in scope
  - Domain names mean the same thing everywhere

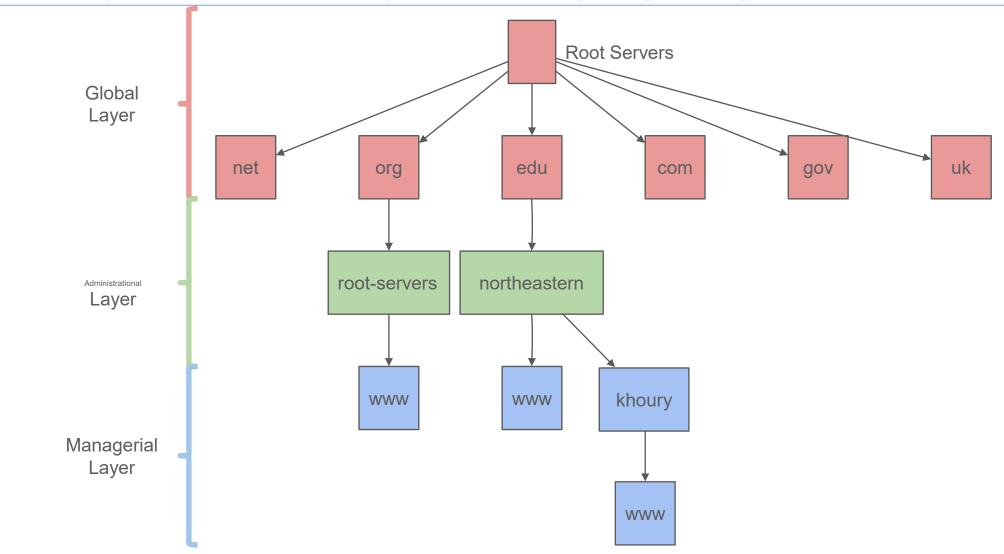
#### Strawman solution: Use a local file

- Keep local copy of mapping from all hosts to all IPs (e.g., /etc/hosts)
- Space would be reasonable: a few dozen Gbytes.
- BUT hosts change IPs regularly, so need to download file frequently
- Lot of constant internet bandwidth use
- Not scalable!

# A tiered architecture yields a scalable solution

- Idea: break apart responsibility for each part of a domain name (zone) to a different group of servers
- Each zone is a continuous section of the name space, eg \*.northeastern.edu
- Each zone has an associated set of name servers.

# DNS partitions responsibility by "layers".



# This is an example of a tiered architecture

- Each server need only needs to know about its immediate descendants in its zone.
- It only processes requests about a single zone.
- Both data and processing are limited to requests about this zone— any other requests are delegated to this server's parent server.

# DNS uses Tree search algorithm

\*.northeastern.edu

\*.khoury.northeastern.edu.

etc.

\*.registrar.northeastern.edu.

\*.law.northeastern.edu.

# Updating name servers

\*.edu
\*.northeastern.edu

\*.khoury.northeastern.edu.

\*.registrar.northeastern.edu.

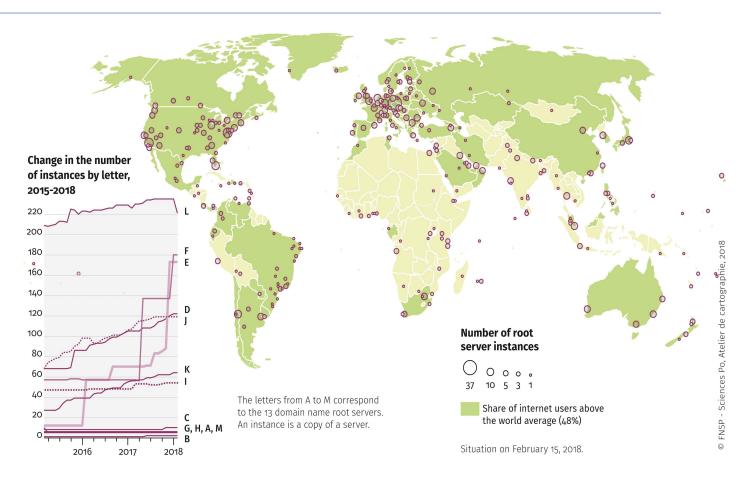
course.khoury.northeastern.edu.

# But some zones are too big and too busy to be handled by a single server

- Eg, .edu, .com, .gov, etc.
- So these servers are replicated.

# There is replication even within the root servers

- 13 root servers
  - [a-m].root-servers.org
  - E.g., d.root-servers.org
- But each root server has multiple copies of the database, which need to be kept in sync.
- Somewhere around 1500 replicas in total.



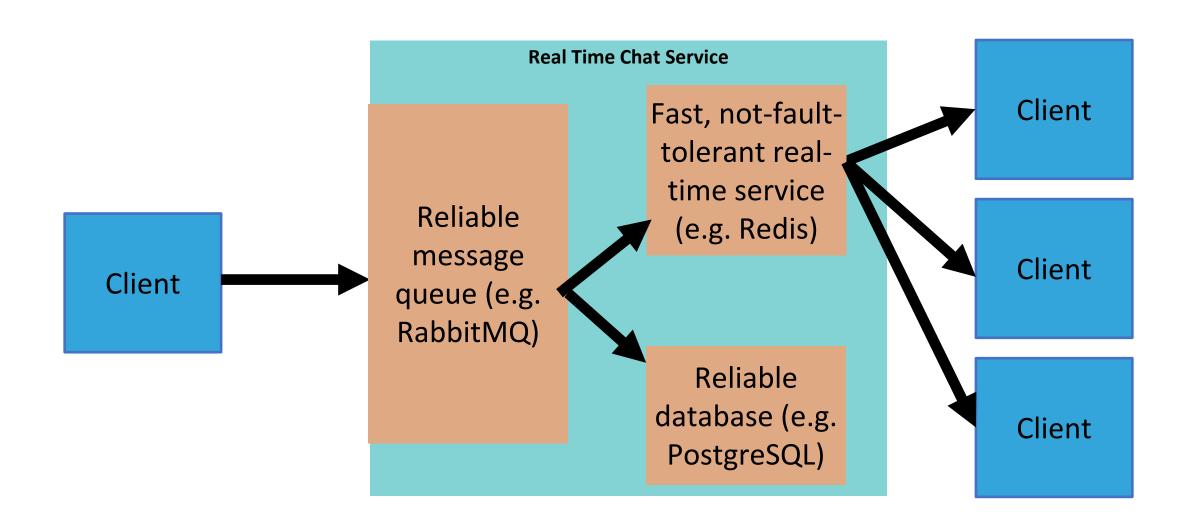
# Case Study 4: Reliable Real-Time Chat

- Requirements:
  - Must support real-time text chat for 2,000 users exchanging messages.
  - Must have best-effort delivery in real-time
  - Must guarantee that all messages acknowledged are preserved in the central database"

# Possible solution: use separate processing units for each requirement.

- Allocate separate processing units for these requirements:
  - "Real time" component optimizes for speed and availability (sacrificing faulttolerance)
  - "Persistence" component optimizes for fault-tolerance, sacrificing speed and availability
  - Event queue service receives events, dispatches to both processing units and is fault tolerant

#### Block diagram for a real-time chat service



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