

Algorithm

Root Array

Array Value

Array Index

Here is a sample quick find implementation of the Disjoint Set.

0 0 0

2

0

```
Run Playground
                Python3
 C++
        Java
 1 ⋅ class UnionFind {
2 public:
        UnionFind(int sz) : root(sz) {
    for (int i = 0; i < sz; i++) {
4 =
                root[i] = i;
9 +
        int find(int x) {
10
            return root[x];
11
12
         void unionSet(int x, int y) {
13 *
14
            int rootX = find(x);
15
             int rootY = find(y);
             if (rootX != rootY) {
16 *
17 *
                 for (int i = 0; i < root.size(); i++) {</pre>
18 *
                     if (root[i] == rootY) {
19
                         root[i] = rootX;
20
21
22
23
24
        bool connected(int x, int y) \{
25 *
26
            return find(x) == find(y);
```

0 0

Root Vertex

\$ 22

Time Complexity

	Union-find Constructor	Find	Union	Connected
Time Complexity	O(N)	O(1)	O(N)	O(1)

Note: N is the number of vertices in the graph.

- ullet When initializing a union-find constructor, we need to create an array of size N with the values equal to the corresponding array indices; this requires linear time.
- Each call to $\frac{1}{1}$ will require O(1) time since we are just accessing an element of the array at the given index.
- Each call to union will require O(N) time because we need to traverse through the entire array and update the root vertices for all the vertices of the set that is going to be merged into another set.
- ullet The connected operation takes O(1) time since it involves the two find calls and the equality check operation.

Space Complexity

We need ${\cal O}(N)$ space to store the array of size N.