### Recap: Markov Chain

- Let  $X_n$ ,  $n = 0, 1, \cdots$  be a sequence of discrete random variables taking values in S.
- ▶ We say it is a Markov chain if

$$Pr[X_{n+1} = x_{n+1} | X_n = x_n, X_{n-1} = x_{n-1} \cdots X_0 = x_0] = Pr[X_{n+1} = x_{n+1} | X_n = x_n]$$

► We can write it as

$$f_{X_{n+1}|X_n,\cdots X_0}(x_{n+1}|x_n,\cdots,x_0) = f_{X_{n+1}|X_n}(x_{n+1}|x_n), \ \forall x_i$$

► For a Markov chain, given the current state, the future evolution is independent of the history of how you reached the current state

# Recap: Transition Probabilities

▶ Transition probability function is  $P: S \times S \rightarrow [0, 1]$ 

$$P(x,y) = Pr[X_{n+1} = y | X_n = x]$$

The chain is said to be homogeneous when this is not a function of time.

► For a homogeneous chain

$$Pr[X_{n+1} = y | X_n = x] = Pr[X_1 = y | X_0 = x], \ \forall n$$

- ightharpoonup P satisfies
  - $P(x,y) \ge 0, \ \forall x,y \in S$
  - $\triangleright \sum_{y \in S} P(x, y) = 1, \forall x \in S$
- ▶ If S is finite then P can be represented as a matrix

# Recap: Initial State Probabilities

▶ Initial state probabilities  $\pi_0: S \to [0, 1]$ 

$$\pi_0(x) = Pr[X_0 = x]$$

It satisfies

- $\blacktriangleright \pi_0(x) > 0, \ \forall x \in S$
- $\sum_{x \in S} \pi_0(x) = 1$
- ▶ The P and  $\pi_0$  together determine all joint distributions:

$$Pr[X_0 = x_0, X_1 = x_1, \dots, X_m = x_m] = \pi_0(x_0)P(x_0, x_1) \cdots P(x_{m-1}, x_m)$$

Using this we can find joint distribution of any finite number of  $X_i$ 's

### Recap

► The Markov property implies

$$Pr[X_{m+n} = y | X_m = x, X_{m-1} \cdots] = Pr[X_{m+n} = y | X_m = x]$$
  
=  $Pr[X_n = y | X_0 = x]$ 

► Or, in general,

$$f_{X_{m+n}|X_m,\cdots,X_0}(y|x,\cdots) = f_{X_{m+n}|X_m}(y|x)$$

► Further, we can show

$$\Pr[X_{m+n} = y | X_m = x, X_{m-k} \in A_k, k = 1, \dots, m] = Pr[X_{m+n} = y | X_m = x]$$

$$Pr[X_{m+n+r} \in B_r, \ r = 0, \cdots, s \mid X_m = x, \ X_{m-k} \in A_k, \ k = 1, \cdots, m]$$
  
=  $Pr[X_{m+n+r} \in B_r, \ r = 0, \cdots, s \mid X_m = x]$ 

► The transition probabilities we defined earlier are also called one step transition probabilities

$$P(x,y) = Pr[X_{n+1} = y | X_n = x] = Pr[X_1 = y | X_0 = x]$$

- We can define transition probabilities for multiple steps, that is,  $Pr[X_n = y | X_0 = x]$
- ► Now we get

$$Pr[X_{m+n} = y | X_0 = x] = \sum_{z} Pr[X_{m+n} = y, X_m = z | X_0 = x]$$

$$= \sum_{z} Pr[X_{m+n} = y | X_m = z, X_0 = x] Pr[X_m = z | X_0 = x]$$
$$= \sum_{z} Pr[X_{m+n} = y | X_m = z] Pr[X_m = z | X_0 = x]$$

$$= \sum_{n=0}^{\infty} \Pr[X_n = y | X_0 = z] \Pr[X_m = z | X_0 = x]$$

# Chapman-Kolmogorov Equations

- ▶ Define:  $P^n(x,y) = Pr[X_n = y | x_0 = x]$
- ▶ These are called *n*-step transition probabilities.
- From what we showed,

$$P^{m+n}(x,y) = Pr[X_{m+n} = y | X_0 = x]$$

$$= \sum_{z} Pr[X_n = y | X_0 = z] Pr[X_m = z | X_0 = x]$$

$$= \sum_{z} P^m(x,z) P^n(z,y)$$

- ► These are known as Chapman-Kolmogorov equations
- ► This relationship is intuitively clear

► Specifically, using Chapman-Kolmogorov equations,

$$P^{2}(x,y) = \sum P(x,z)P(z,y)$$

- For a finite chain, P is a matrix
- ► Thus  $P^2(x,y)$  is the  $(x,y)^{th}$  element of the matrix,  $P \times P$
- ightharpoonup That is why we use  $P^n$  for n-step transition probabilities

- ▶ Define:  $\pi_n(x) = Pr[X_n = x]$ .
- ► Then we get

$$\pi_n(y) = \sum_x Pr[X_n = y | X_0 = x] Pr[X_0 = x]$$

$$= \sum_x \pi_0(x) P^n(x, y)$$

► In particular

$$\pi_{n+1}(y) = \sum_{x} Pr[X_{n+1} = y | X_n = x] Pr[X_n = x]$$

$$= \sum_{x} \pi_n(x) P(x, y)$$

### Hitting times

- ▶ Let y be a state.
- ▶ We define hitting time for y as the random variable

$$T_y = \min\{n > 0 : X_n = y\}$$

- ▶  $T_y$  is the first time that the chain is in state y (after t = 0 when the chain is initiated).
- lt is easy to see that  $Pr[T_y = 1 | X_0 = x] = P(x, y)$ .
- Notation:  $P_z(A) = Pr[A|X_0 = z]$
- ► We write the above as  $P_x(T_y = 1) = P(x, y)$ Note that

$$P_x[T_y = n] = Pr[T_y = n | X_0 = x] \neq P^n(x, y)$$

$$T_{y} = \min\{n > 0 : X_{n} = y\}$$

► We can now get

 $z\neq y$ 

$$P_x(T_y = 2) = \sum_{z \neq y} P(x, z) P(z, y) = \sum_{z \neq y} P(x, z) P_z(T_y = 1)$$

$$P_x(T_y = m) = Pr[T_y = m | X_0 = x]$$

$$= \sum_{z \neq y} P(x, z) P(z, y) = \sum_{z \neq y} P(x, z) P_z(T_y = 1)$$

$$= \sum_{z \neq y} P(x, z) P(z, y) = \sum_{z \neq y} P(x, z) P_z(T_y = 1)$$

$$= \sum_{z \neq y} P(x, z) Pr[T_y = m | X_1 = z]$$
$$= \sum_{z \neq y} P(x, z) P_z(T_y = m - 1)$$

 $z \neq y$ 

Similarly we can get:

$$P^{n}(x,y) = \sum_{y=0}^{n} P_{x}(T_{y} = m)P^{n-m}(y,y)$$

We can derive this as

$$P^n(x,y) = Pr[X_n = y | X_0 = x]$$

$$= \sum_{n=0}^{n} \Pr[T_{y} = m, X_{n} = y | X_{0} = x]$$

$$= \sum_{m=1} Pr[T_y = m, X_n = y | X_0 = x]$$

$$\frac{1}{m-1} - \sum_{i=1}^{n} Pr[X_i - y|T_i - y_i] Pr[T_i -$$

$$= \sum_{m=1}^{n} \Pr[X_n = y | T_y = m, X_0 = x] \Pr[T_y = m | X_0 = x]$$

 $= \sum Pr[X_n = y | X_m = y] \ Pr[T_y = m | X_0 = x]$ 

$$= \sum P^{n-m}(y,y) P_x(T_y = m)$$

### transient and recurrent states

- ▶ Define  $\rho_{xy} = P_x(T_y < \infty)$ .
- $\blacktriangleright$  It is the probability that starting in x you will visit y
- Note that

$$\rho_{xy} = \lim_{n \to \infty} P_x(T_y < n) = \lim_{n \to \infty} \sum_{m=1}^{n-1} P_x(T_y = m) = \sum_{m=1}^{\infty} P_x(T_y = m)$$

**Definition**: A state y is called **transient** if  $\rho_{yy} < 1$ ; it is called **recurrent** if  $\rho_{uy} = 1$ .

- Intuitively, all transient states would be visited only finitely many times while recurrent states are visited infinitely often.
- ► For any state y define

$$I_y(X_n) = \begin{cases} 1 & \text{if } X_n = y \\ 0 & \text{otherwise} \end{cases}$$

▶ Now, the total number of visits to y is given by

$$N_y = \sum_{n=1}^{\infty} I_y(X_n)$$

 $\blacktriangleright$  We can get distribution of  $N_u$  as

$$P_{x}(N_{y} \ge 1) = P_{x}(T_{y} < \infty) = \rho_{xy}$$

$$P_{x}(N_{y} \ge 2) = \sum_{m} P_{x}(T_{y} = m)P_{y}(T_{y} < \infty)$$

$$= \rho_{yy} \sum_{m} P_{x}(T_{y} = m) = \rho_{yy} \rho_{xy}$$

$$P_{x}(N_{y} \ge m) = \rho_{yy}^{m-1} \rho_{xy}$$

$$P_{x}(N_{y} = m) = P_{x}(N_{y} \ge m) - P_{x}(N_{y} \ge m + 1)$$

$$= \rho_{yy}^{m-1} \rho_{xy} - \rho_{yy}^{m} \rho_{xy} = \rho_{xy} \rho_{yy}^{m-1} (1 - \rho_{yy})$$

$$P_{x}(N_{y} = 0) = 1 - P_{x}(N_{y} \ge 1) = 1 - \rho_{xy}$$

- Notation:  $E_x[Z] = E[Z|X_0 = x]$
- Define

$$G(x,y) \triangleq E_x[N_y]$$

$$= E_x \left[ \sum_{n=1}^{\infty} I_y(X_n) \right]$$

$$= \sum_{n=1}^{\infty} E_x [I_y(X_n)]$$

$$= \sum_{n=1}^{\infty} P^n(x,y)$$

▶ G(x,y) is the expected number of visits to y for a chain that is started in x.

#### Theorem:

(i). Let y be transient. Then

$$P_x(N_y < \infty) = 1, \ \forall x \ \text{ and } \ G(x,y) = \frac{\rho_{xy}}{1 - \rho_{yy}} < \infty, \ \forall x$$

(ii) Let y be recurrent. Then

$$P_{u}[N_{u}=\infty]=1$$
, and  $G(y,y)=E_{u}[N_{u}]=\infty$ 

$$P_x[N_y = \infty] = \rho_{xy}$$
, and  $G(x,y) = \begin{cases} 0 & \text{if } \rho_{xy} = 0 \\ \infty & \text{if } \rho_{xy} > 0 \end{cases}$ 

#### **Proof of (i)**: y is transient; $\rho_{yy} < 1$

$$\begin{split} G(x,y) &=& E_x[N(y)] = \sum_m m P_x[N(y) = m] \\ &=& \sum_m m \; \rho_{xy} \; \rho_{yy}^{m-1} (1 - \rho_{yy}) \\ &=& \rho_{xy} \; \sum_{m=1}^\infty m \; \rho_{yy}^{m-1} (1 - \rho_{yy}) \\ &=& \rho_{xy} \; \frac{1}{1 - \rho_{yy}} < \infty, \; \text{ because } \; \rho_{yy} < 1 \\ &\Rightarrow \; P_x[N(y) < \infty] = 1 \end{split}$$

### Proof of (ii):

 $y \text{ recurrent } \Rightarrow \rho_{yy} = 1. \text{ Hence}$ 

$$P_y[N(y) \ge m] = \rho_{yy}^m = 1, \forall m$$

$$\Rightarrow P_y[N(y) = \infty] = \lim_{m \to \infty} P_y[N(y) \ge m] = 1$$

$$\Rightarrow G(y, y) = E_y[N(y)] = \infty$$

$$P_x[N(y) \ge m] = \rho_{xy} \ \rho_{yy}^{m-1} = \rho_{xy}, \ \forall m$$

Hence 
$$P_x[N(y) = \infty] = \rho_{xy}$$

$$\rho_{xy} = 0 \implies P_x[N(y) > m] = 0, \forall m > 0 \implies G(x, y) = 0$$

$$\rho_{xy} > 0 \implies P_x[N(y) = \infty] > 0 \implies G(x, y) = \infty$$

- ► Transient states are visited only finitely many times while recurrent states are visited infinitely often
- ▶ If *S* is finite, it should have at least one recurrent state
- ightharpoonup If y is transient, then, for all x

$$G(x,y) = \sum_{n=1}^{\infty} P^n(x,y) < \infty \quad \Rightarrow \quad \lim_{n \to \infty} P^n(x,y) = 0$$

- ▶ However,  $\sum_{y} P^{n}(x, y) = 1, \ \forall n, \ \forall x$
- ▶ If all  $y \in S$  are transient, then we get a contradiction

$$1 = \lim_{n \to \infty} \sum_{y \in S} P^n(x, y) = \sum_{y \in S} \lim_{n \to \infty} P^n(x, y) = 0$$

- ▶ A finite chain has to have at least one recurrent state
- ► An infinite chain can have only transient states

- ▶ We say, x leads to y if  $\rho_{xy} > 0$ 
  - **Theorem**: If x is recurrent and x leads to y then y is recurrent and  $\rho_{xy} = \rho_{yx} = 1$ .

#### Proof:

- ► Take  $x \neq y$ , wlog. Since  $\rho_{xy} > 0$ ,  $\exists n$  s.t.  $P^n(x,y) > 0$
- ▶ Take least such n. Then we have states  $y_1, \dots, y_{n-1}$ , none of which is x (or y) such that

$$P(x, y_1) P(y_1, y_2) \cdots P(y_{n-1}, y) > 0$$

Now suppose,  $\rho_{ux} < 1$ . Then

$$P(x, y_1) P(y_1, y_2) \cdots P(y_{n-1}, y) (1 - \rho_{yx}) > 0$$

is the probability of starting in x but not returning to x.

- $\blacktriangleright$  But this cannot be because x is recurrent and hence  $\rho_{xx}=1$
- Hence, if x is recurrent and x leads to y then  $\rho_{ux} = 1$

Now,  $\exists n_0, n_1 \text{ s.t. } P^{n_0}(x, y) > 0, P^{n_1}(y, x) > 0.$ 

$$P^{n_1+n+n_0}(y,y) = P_y[X_{n_1+n+n_0} = y]$$

$$\geq P_y[X_{n_1} = x, X_{n_1+n} = x, X_{n_1+n+n_0} = y]$$

$$= P^{n_1}(y,x)P^n(x,x)P^{n_0}(x,y), \forall n$$

▶ We know  $G(x,x) = \sum_{m=1}^{\infty} P^m(x,x) = \infty$ 

$$\sum_{m=1}^{\infty} P^m(y,y) \geq \sum_{m=n_0+n_1+1}^{\infty} P^m(y,y) = \sum_{n=1}^{\infty} P^{n_1+n+n_0}(y,y)$$

$$\geq \sum_{n=1}^{\infty} P^{n_1}(y,x) P^n(x,x) P^{n_0}(x,y)$$

$$= \infty, \text{ because } x \text{ is recurrent}$$

$$\Rightarrow y \text{ is recurrent}$$

- Mhat we showed so far is: if x leads to y and x is recurrent, then  $\rho_{ux}=1$  and y is recurrent.
- Now, y is recurrent and y leads to x and hence  $\rho_{xy} = 1$ .
- ► This completes proof of the theorem

### equivalence relation

- $\blacktriangleright$  let R be a relation on set A. Note  $R \subset A \times A$
- ightharpoonup R is called an equivalence relation if it is
  - 1. reflexive, i.e.,  $(x, x) \in R$ ,  $\forall x \in A$
  - 2. symmetric, i.e.,  $(x,y) \in R \implies (y,x) \in R$
  - 3. transitive, i.e.,  $(x,y),(y,z)\in R \Rightarrow (x,z)\in R$

### example

- ▶ Let  $A = \{\frac{m}{n} \mid m, n \text{ are integers}\}$
- ightharpoonup Define relation R by

$$\left(\frac{m}{n}, \frac{p}{q}\right) \in R \text{ if } mq = np$$

- ► This is the usual equality of fractions
- Easy to check it is an equivalence relation.

# Equivalence classes

- ▶ Let R be an equivalence relation on A.
- ightharpoonup Then, A can be partitioned as

$$A = C_1 + C_2 + \cdots$$

Where  $C_i$  satisfy

- $\triangleright x, y \in C_i \Rightarrow (x, y) \in R, \forall i$
- $ightharpoonup x \in C_i, \ y \in C_j, \ i \neq j \Rightarrow (x,y) \notin R$
- In our example, each equivalence class corresponds to a rational number.
- ightharpoonup Here,  $C_i$  contains all fractions that are equal to that rational number

The state space of any Markov chain can be partitioned into the transient and recurrent states:  $S = S_T + S_R$ :

$$S_T = \{ y \in S : \rho_{yy} < 1 \}$$
  $S_R = \{ y \in S : \rho_{yy} = 1 \}$ 

- ▶ On  $S_R$ , consider the relation: 'x leads to y' (i.e., x is related to y if  $\rho_{xy} > 0$ )
- ► This is an equivalence relation
  - $ightharpoonup 
    ho_{xx} > 0, \ \forall x \in S_R$
  - $\rho_{xy} > 0 \implies \rho_{yx} > 0, \ \forall x, y \in S_R$
  - $ightharpoonup 
    ho_{xy} > 0, \ 
    ho_{yz} > 0 \ \Rightarrow \ 
    ho_{xz} > 0$
- ▶ Hence we get a partition:  $S_R = C_1 + C_2 + \cdots$  where  $C_i$  are equivalence classes.

- $\triangleright$  On  $S_R$ , "x leads to y" is an equivalence relation.
- ▶ This gives rise to the partition  $S_R = C_1 + C_2 + \cdots$
- ightharpoonup Since  $C_i$  are equivalence classes, they satisfy:
  - $\blacktriangleright x, y \in C_i \Rightarrow x \text{ leads to } y$
  - $ightharpoonup x \in C_i, y \in C_j, i \neq j \Rightarrow \rho_{xy} = 0$
- lacktriangle All states in any  $C_i$  lead to each other or communicate with each other
- ▶ If  $i \neq j$  and  $x \in C_i$  and  $y \in C_j$ , then,  $\rho_{xy} = \rho_{yx} = 0$ . x and y do not communicate with each other.

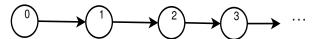
- ▶ A set of states,  $C \subset S$  is said to be irreducible if x leads to y for all  $x, y \in C$
- ▶ An irreducible set is also called a communicating class
- A set of states,  $C \subset S$ , is said to be closed if  $x \in C$ ,  $y \notin C$  implies  $\rho_{xy} = 0$ .
- Once the chain visits a state in a closed set, it cannot leave that set.
- ▶ We get a partition of recurrent states

$$S_R = C_1 + C_2 + \cdots$$

where each  $C_i$  is a closed and irreducible set of states.

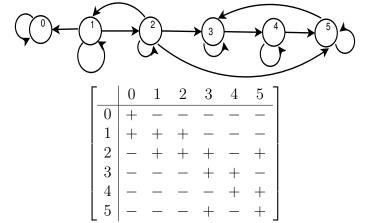
▶ If S is irreducible then the chain is said to be irreducible. (Note that S is trivially closed)

- ▶ In an irreducible set of states, if one state is recurrent, then, all states are recurrent.
- ► We saw that a finite chain has to have at least one recurrent state.
- ▶ Thus, a finite irreducible chain is recurrent.
- ► For example, in the umbrellas problem, the chain is irreducible and hence all states are recurrent.
- An infinite irreducible chain may be wholly transient
- ▶ Here is a trivial example of non-irreducible transient chain:



- ► The state space of any Markov chain can be partitioned into transient and recurrent states.
- We need not calculate  $\rho_{xx}$  to do this partition.
- ▶ By looking at the structure of transition probability matrix we can get this partition

### Example



- ▶ State 0 is called an absorbing state. {0} is a closed irreducible set.
- ▶ 1,2 are transient states.
- We get:  $S_T = \{1, 2\}$  and  $S_R = \{0\} + \{3, 4, 5\}$