

int a = 5; // Local variable



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	return a * g;				
	3				
8.2					
	(b) Digorentiate letween static and dynamic scoping.				
	sconing				
	Static Scoping	Dynamic Scoping			
	' 0	• • •			
	In static scoping (or lenical	In dynamic scoping definition			
	scoping), defination of a	of a viriable is readined by			
	pariable in resolved by	searching its containing			
	searching its containing	block or function of that			
	block or function, if it	fails then searching its			
	bils it then searching	calling function, ignot			
	in outer block and ao	found the it searches in			
	on .	the junction that called the			
		the junction that called the calling function			
	11 1 10 1				
	Most modern programming	Dynamic Icohing is			
	languages support static acquing.	synamic acoping is outphorted by older eanguages			
	static ocquing.	languages			
	Eg Algol, C, Pascal	Eg. SNOBOL, APL			
	Instatio acobica . acon	la dunamia sa chian i sa ca a			
	Instatic acoping we can	In dynamic scoping we can			
	at compile time	associate the scope of windle			
	and any many	- Cu (cu) M·C			



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		And international or a	H		
8.2.	Cooler of stock allocation with holl of nont				
9.7	Exchain stack allocation with help of near lavelled diagram and example				
	Labelle anglas, and Danigue				
	The state of the s				
	Sp Subroutine D				
	86-	Aut to only to	1	. Wiles I'm say	
	JA	with a contract of		Arguments	
				to called	
		subroutine C		routines	
		OWN TOWNS TO	23	Carlon Control Control	
			V A	Temporaries	
			No.		
		Care to fine on		stu Jocal variables	
		Subsouting B			
	1	sensitive de la company		Miscellaneous	
		Comment of the later		Bookberping	
		Subroutine B.	-	1	
		alesson Transcription		Pelurnaddress	
		1 1 1			
	100000	Dulroutine A	100		
	O has a main has a				
	Procedure C main program				
	0, E +).				
	procedure B ig . then Belse ( Procedure A				
	then is else (				
	Procedure H				



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	for recursion static allocation is not possible
	Uniter of instances of variable is uneounded
	stack allocation allows us to use
	In stack allocation, activition records
1	an place allecation, actualism records
1	are hished onto the stack created
	for allocation of memory.  I stack frame will have the following
	Contents
1	Contents:-
	2. Scal variables
1	
	3. Jemperory variables 4. Return addresses
	5. His cellaneous book keeping.
	John Sayon 9
	The maintainence of stack is done by
	The maintainence of stack is done by subrocitine calling sequence prolonge
	and epilogue.
	Calling Olquence: Lode executed by the call immediately after the call.
	caller immediately after the call.
	Prologue: Lode exceeded at legring a solocates
	Daves called - Daves registers used any where
ł	Daves callo - Priver registers used any where
	insido calloe
	Epilougo: - codo excecuted at the end and
	puts return values into registers, restores
	de allocate & viane.
	de allocall & crane.