

.Subarray Sum Equals K

560. Subarray Sum Equals K

Medium

18872

551

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Given an array of integers `nums` and an integer `k`, return the total number of subarrays whose sum equals to `k`.

A subarray is a contiguous **non-empty** sequence of elements within an array.

Example 1:

Input: `nums = [1,1,1]`, `k = 2`

Output: 2

Example 2:

Input: `nums = [1,2,3]`, `k = 3`

Output: 2

Brute:-

`[1, 1, 1]` $k=2$
i ↑ j ↑
sum += arr[j] // sum = 0 + 1
= 1
if (sum == k)
cnt++

`[1, 1, 1]` j++
i ↑ j ↑
sum = 1 + 1
= 2

`[1, 1, 1]` cnt++ True // 1
i ↑ j ↑
j++
sum = 3

`[1, 1, 1]` i++
i ↑ j ↑
sum = 0 + 1
= 1

`[1, 1, 1]` j++
i ↑ j ↑
sum = 1 + 1 = 2
cnt++ // 2

`[1, 1, 1]`
i ↑ j ↑
sum = 0 + 1
= 1

return cnt // 2

T.C - $O(n^2)$

S.C - $O(1)$

Optimal solution:-

Prefix sum / Hashing



[2, 5, 7, 8, 9]

↑
standing right here

having a Hashmap which have sum of previous numbers

[0+2, $\underbrace{2+5}_a$, $\underbrace{7+a}_b$, $\underbrace{8+b}_c$, $9+c$]

[.]

sum = x

k

x - k



If we found (x-k) before x

[.]

↑
x
If any found (x-k)

then we have an subarray
== k

Day 2 un:-

arr[] = [1, 2, 3, -3, 1, 1, 1, 4, 2, -3] k=3

insert(0) = 1

prefix sum:- ϕ ~~3~~ ~~6~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~10~~ 12 9

Find = 1-3 = -2 3-3 = 0 3-6 = 3 3-3 = 0 4-3 = 1 3-5 = 2 6-3 = 3 | 10-3 = 7 | 12-3 = 9 | 9-3 = 6

cnt = ~~0~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~ ~~8~~

\downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow
 0 3 -3 1 1 1 4 2
 3rd 4th 5th 7th 7th 10th 10th
 3:2 6:2

12:1
10:1
5:1
4:1
6:2
3:2
1:1
0:1

map

Pre / freq

return cnt, 1st 2nd 3rd 4th 5th

arr[] = [3, -3, 1, 1, 1] k=3

insert(0) = 1

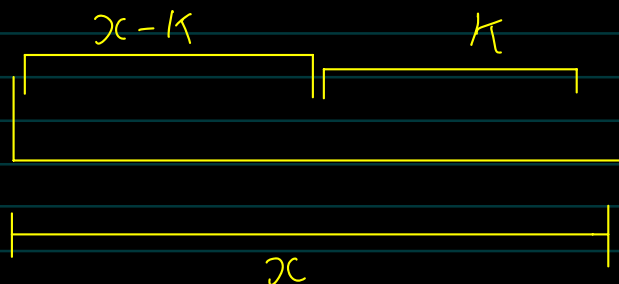
prefix = ϕ ~~3~~ ~~0~~ ~~2~~ ~~3~~

Find: 3-3 = 0 | 0-3 = -3 | 1-3 = -2 | 2-3 = -1 | 3-3 = 0

cnt = ~~0~~ ~~2~~ ~~3~~

\downarrow \downarrow \downarrow
 1st 5th 5th
 0:2

1:1
3:2
0:2



```

class Solution {
public:
    int subarraySum(vector<int>& nums, int k) {
        int cnt=0, prefix=0;
        unordered_map<int, int> mp;
        mp[0]=1;
        for(int i=0; i<nums.size(); i++){
            prefix+=nums[i];
            int findX=prefix-k;
            if(mp.find(findX)!=mp.end()){
                cnt+=mp[findX];
            }
            mp[prefix]++;
        }
        return cnt;
    }
};

```

T.C - $O(n)$
 $+ O(\log n)$

Avg case $O(1)$
 worst case $O(n)$

S.C = $O(1)$