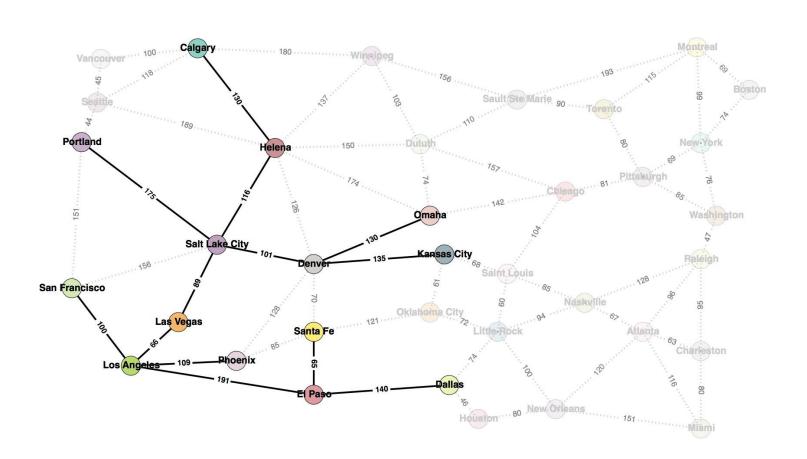
# Artificial Intelligence Search Agents Uninformed search



Use no domain knowledge!

## **Strategies:**

1.Breadth-first search (BFS): Expand shallowest node

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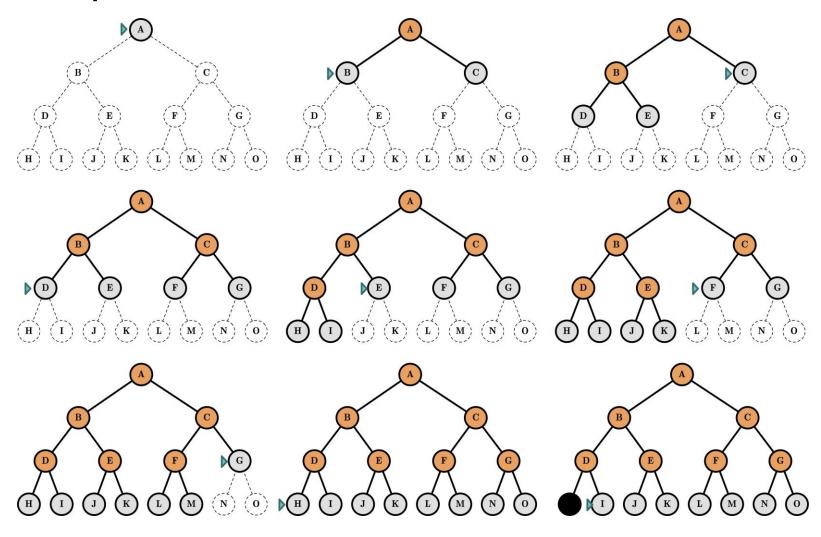
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- 4. Iterative-deepening search (IDS): DLS with increasing limit
- 5.Uniform-cost search (UCS): Expand least cost node

# Breadth-first search (BFS)

**BFS: Expandshallowestfirst.** 



# BFS search

return FAILURE

```
function Breadth-First-Search(initialState, goalTest)
     returns Success or Failure:
     frontier = Queue.new(initialState)
     explored = Set.new()
     while not frontier.isEmpty():
          state = frontier.dequeue()
          explored.add(state)
          if goalTest(state):
               return Success(state)
          for neighbor in state.neighbors():
               if neighbor not in frontier \cup explored:
                    frontier.enqueue(neighbor)
```

# **BFS** Criteria

BFS criteria?

- Complete Yes (if bis finite)
- Time  $1 + b + b^2 + b^3 + ... + b^d = O(b^d)$
- Space  $O(b^d)$ Note: If the *goal test* is applied at expansion rather than generation then  $O(b^{d+1})$
- Optimal Yes (if cost = 1 perstep).
- implementation: fringe: FIFO (Queue)

Question: If time and space complexities are exponential, why use BFS?

How bad is BFS?

How bad is BFS?

Depth	Nodes	Time	Memory
2	110	.11 milliseconds	107 kilobytes
4	11,110	11 milliseconds	10.6 megabytes
6	10 <sup>6</sup>	1.1 seconds	1 gigabyte
8	10 <sup>8</sup>	2 minutes	103 gigabytes
10	10 <sup>10</sup>	3 hours	10 terabytes
12	10 <sup>12</sup>	13 days	1 petabyte
14	10 <sup>14</sup>	3.5 years	99 petabytes
16	10 <sup>16</sup>	350 years	10 exabytes

Time and Memory requirements for breadth-first search for a branching factor b=10; 1 million nodes per second; 1,000 bytes per node.

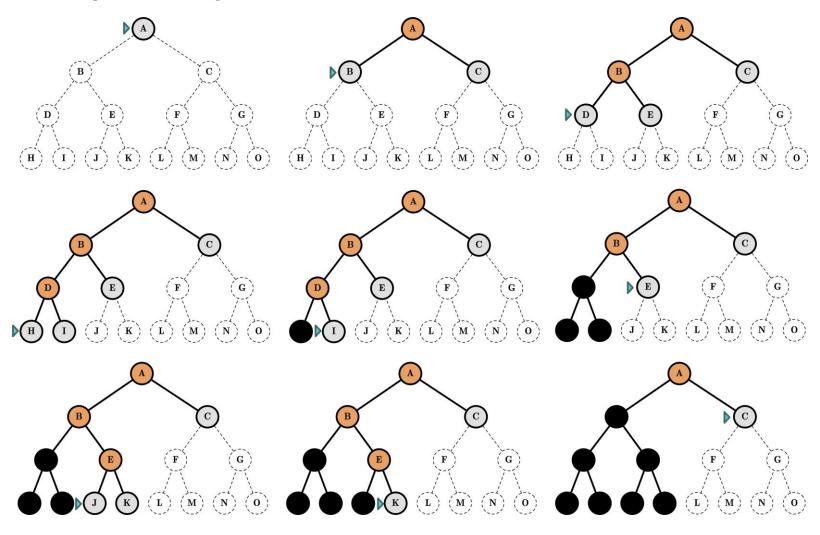
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Time and Memory requirements for breadth-first search for a branching factor b=10; 1 million nodes per second; 1,000 bytes per node.

Memory requirement + exponential time complexity are the biggest handicaps of BFS!

DFS: Expanddeepestfirst.



## DFS search

return FAILURE

```
function Depth-First-Search(initialState, goalTest)
     returns Success or Failure:
     frontier = Stack.new(initialState)
     explored = Set.new()
     while not frontier.isEmpty():
          state = frontier.pop()
          explored.add(state)
          if goalTest(state):
               return Success(state)
          for neighbor in state.neighbors():
               if neighbor not in frontier \cup explored:
                     frontier.push(neighbor)
```

DFS criteria?

- Complete No: fails in infinite-depth spaces, spaces with loops Modify to avoid repeated states along path.
  - ⇒ complete in finite spaces
- Time  $O(b^m)$ :  $1 + b + b^2 + b^3 + ... + b^m = O(b^m)$  bad if m is much larger than d but if solutions are dense, may be much faster than BFS.
- Space O(bm)linear space complexity! (needs to store only a single path from the root to a leaf node, along with the remaining unexpanded sibling nodes for each node on the path, hence the m factor.)
- Optimal No
- Implementation: fringe: LIFO (Stack)

#### How bad is DFS?

#### Recall for BFS...

Depth	Nodes	Time	Memory
2	110	.11 milliseconds	107 kilobytes
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Depth = 16.

We go down from 10 exabytes in BFS to . . . in DFS?

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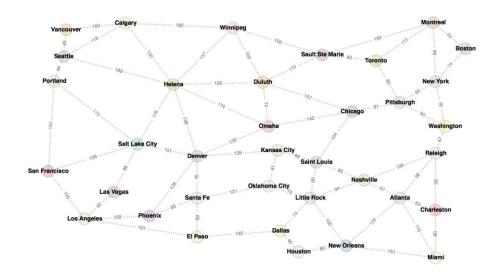
We go down from 10 exabytes in BFS to 156 kilobytes in DFS!

# Depth-limited search

- DFS with depth limit / (nodes at level / has no successors).
- Select some limit L in depth to explore with DFS
- Iterative deepening: increasing the limit /

# Depth-limited search

 If we know some knowledge about the problem, may be we don't need to go to a full depth.

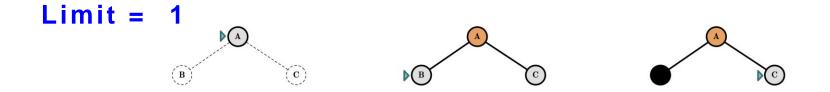


Idea: any city can be reached from another city in at most L steps with L < 36.

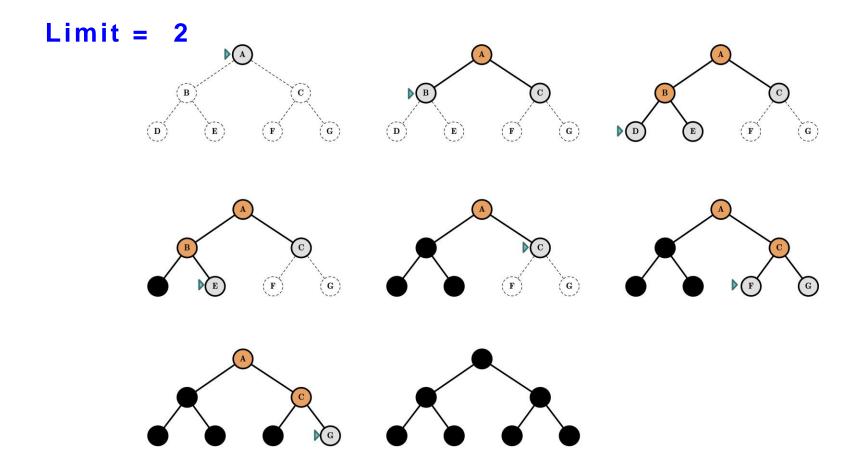
- Combines the benefits of BFS and DFS.
- Idea: Iteratively increase the search limit until the depth of the shallowest solution *d* is reached.
- AppliesDLS with increasing limits.
- The algorithm will stop if a solution is found or if DLS returns a failure (no solution).
- Because most of the nodes are on the bottom of the search tree, it not a big waste to iteratively re-generate the top
- Let's take an example with a depth limit between 0 and 3.

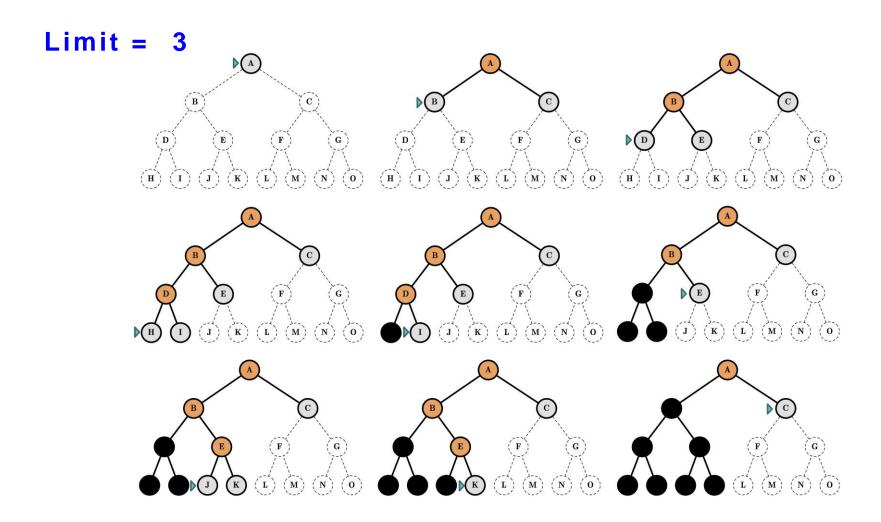
Limit = 0

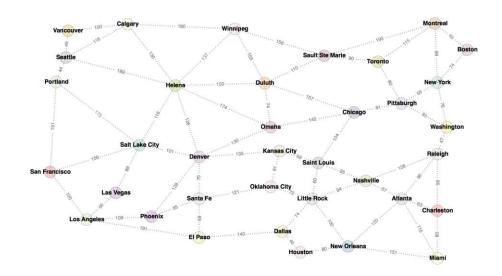
▶△



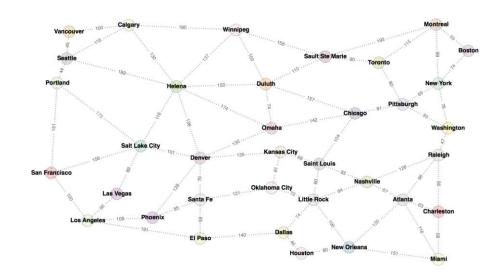




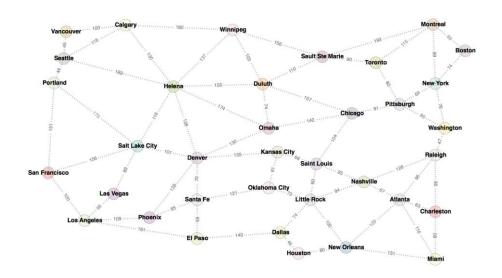




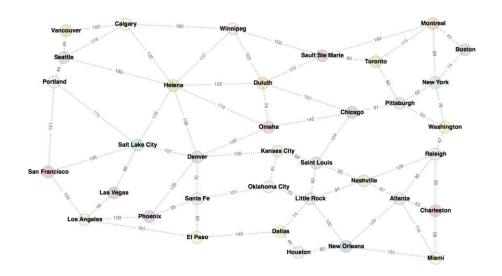
• The arcs in the search graph may have weights (different cost attached). How to leverage this information?



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- We want thecheapestnot shallowest solution.
- Modify BFS: Prioritize by cost not depth → Expand node n
   with the lowest path cost g(n)

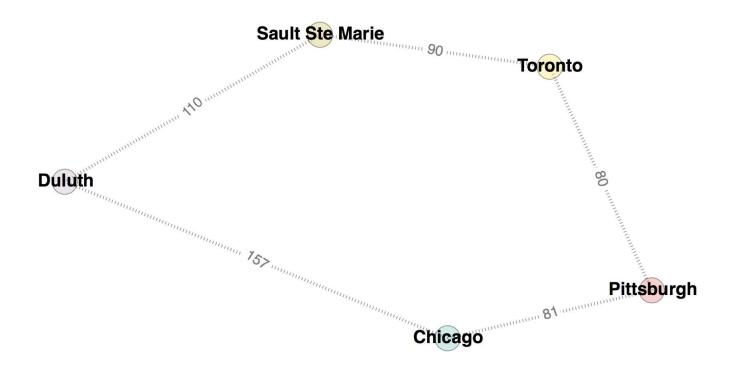


- The arcs in the search graph may have weights (different cost attached). How to leverage this information?
- BFS will find the shortest path which may be costly.
- We want thecheapestnot shallowest solution.
- Modify BFS: Prioritize by cost not depth → Expand node n
   with the lowest path cost g(n)
- Explores increasing costs.

# UCS algorithm

```
function Uniform-Cost-Search(initialState, goalTest)
     returns Success or Failure: /* Cost f(n) = g(n) */
     frontier = Heap.new(initialState)
     explored = Set.new()
     while not frontier.isEmpty():
          state = frontier.deleteMin()
          explored.add(state)
          if goalTest(state):
               return Success(state)
          for neighbor in state.neighbors():
               if neighbor not in frontier \cup explored:
                     frontier.insert(neighbor)
               else if neighbor in frontier:
                     frontier.decreaseKey(neighbor)
```

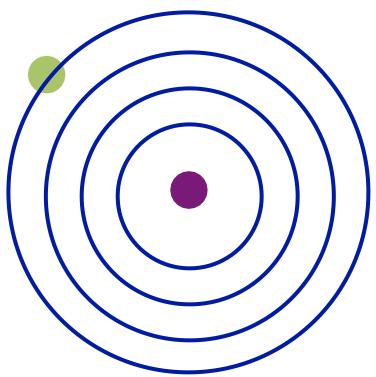
return FAILURE



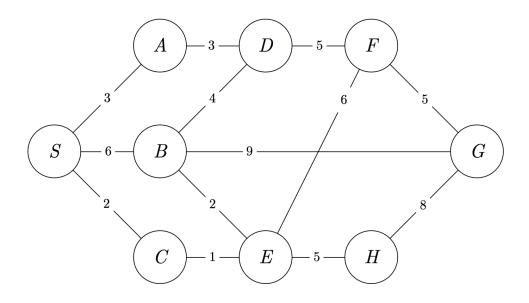
Go from Chicago to Sault Ste Marie. Using BFS, we would find Chicago-Duluth-Sault Ste Marie. However, using UCS, we would find Chicago-Pittsburgh-Toronto-Sault Ste Marie, which is actually the shortest path!

- Complete Yes, if solution has a finite cost.
- Time
  - Suppose C\*: cost of the optimal solution
  - Every action costs at least s(bound on the cost)
  - The effective depth is roughly  $C^*/s$  (how deep the *cheapest* solution could be).
  - $O(b^{C^*/s})$
- Space # of nodes with  $g \le cost of optimal solution, <math>O(b^{C^*/s})$
- Optimal Yes
- Implementation: fringe = queue ordered by path cost g(n), lowest first = Heap!

While complete and optimal, UCS explores the space in every direction because no information is provided about the goal!

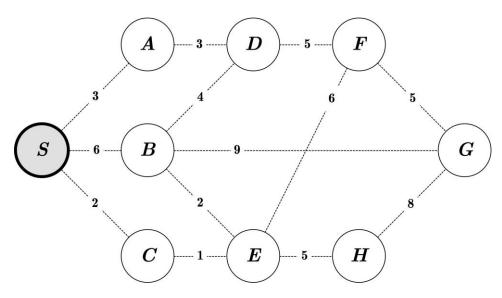


# **Exercise**



Question: What is the order of visits of the nodes and the path returned by BFS, DFS and UCS?

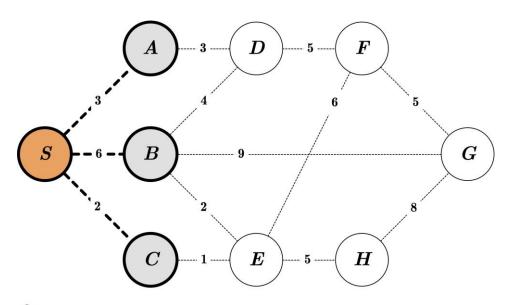
# **Exercise: BFS**



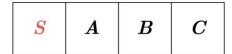
Queue:

 $\boldsymbol{S}$ 

Order of Visit:

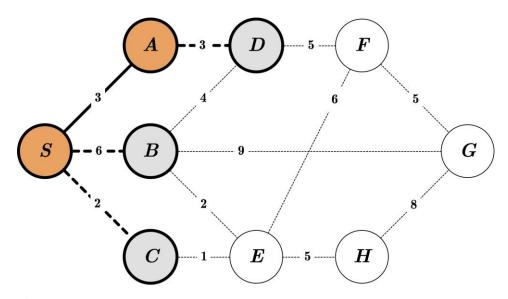


Queue:

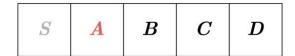


Order of Visit:

 $\boldsymbol{S}$ 

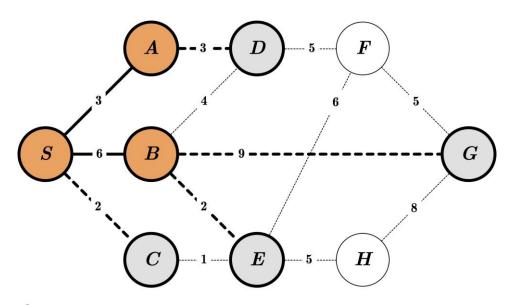


Queue:

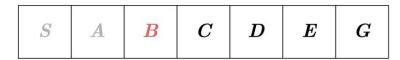


Order of Visit:

S A

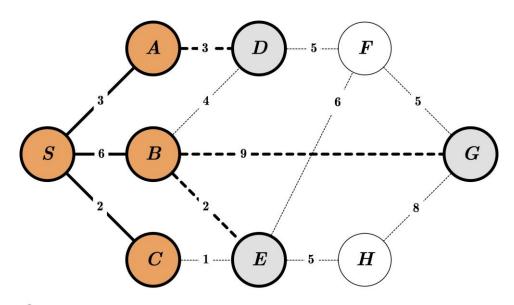


Queue:

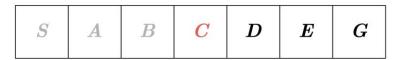


Order of Visit:

 $S \quad A \quad B$ 

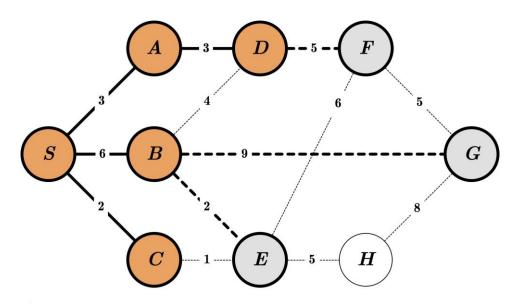


Queue:



Order of Visit:

S A B C

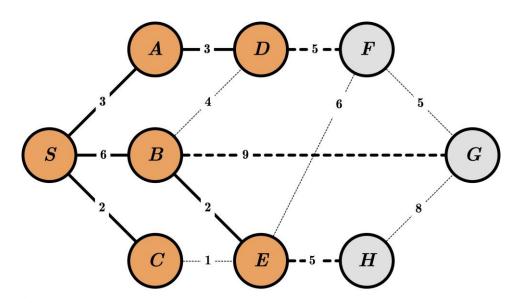


Queue:



Order of Visit:

S A B C D

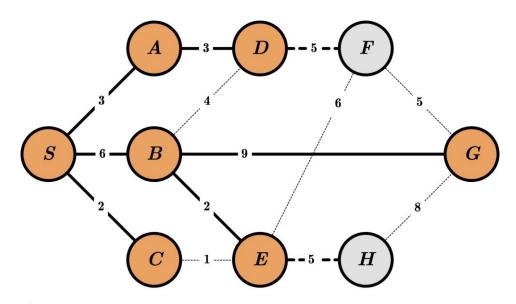


Queue:



Order of Visit:

S A B C D E

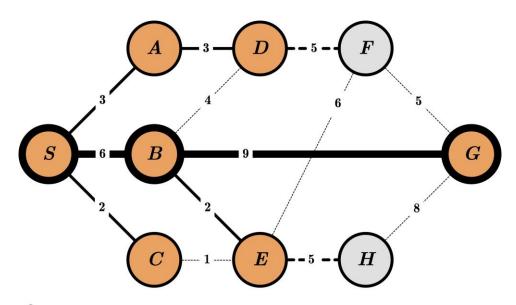


Queue:



Order of Visit:

S A B C D E G

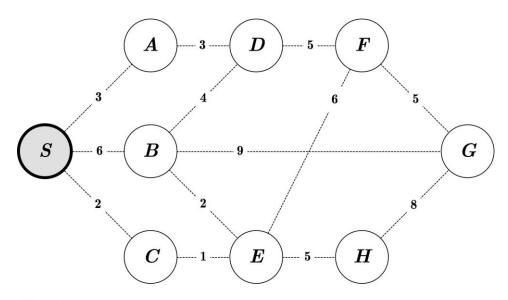


Queue:



Order of Visit:

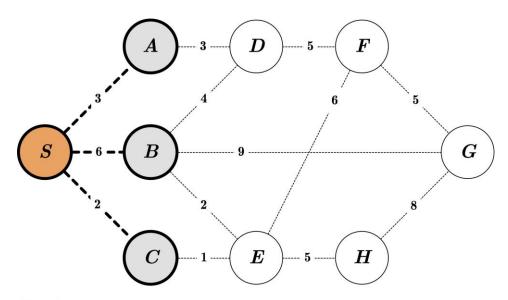
S A B C D E G



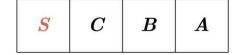
Stack:

 $\boldsymbol{S}$ 

Order of Visit:

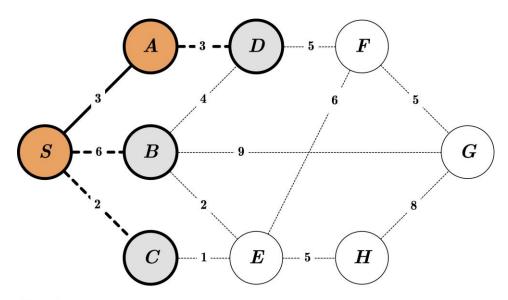


Stack:

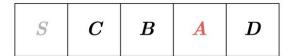


Order of Visit:

 $\boldsymbol{S}$ 

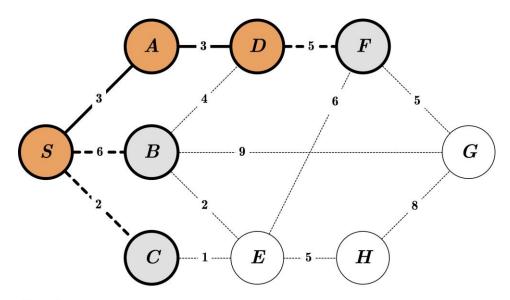


Stack:



Order of Visit:

S A

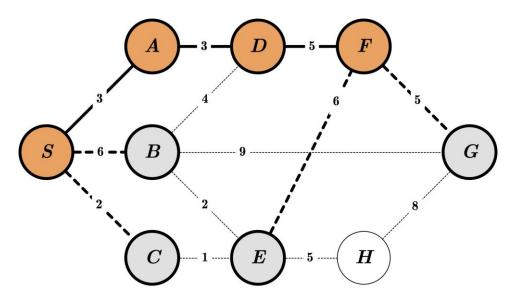


Stack:



Order of Visit:

 $S \quad A \quad D$ 

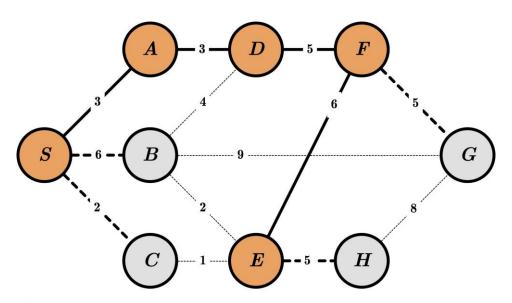


Stack:



Order of Visit:

S A D F

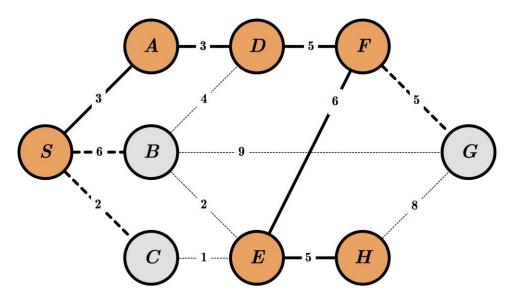


Stack:



Order of Visit:

S A D F E

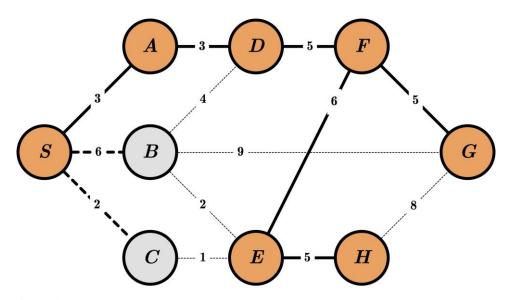


Stack:



Order of Visit:

S A D F E H

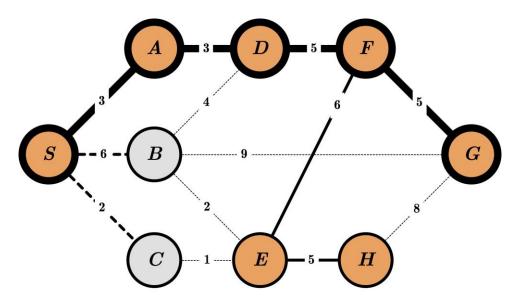


Stack:



Order of Visit:

S A D F E H G

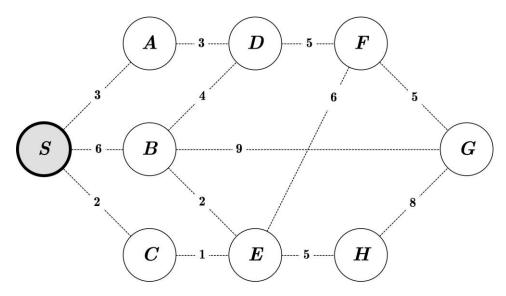


Stack:



Order of Visit:

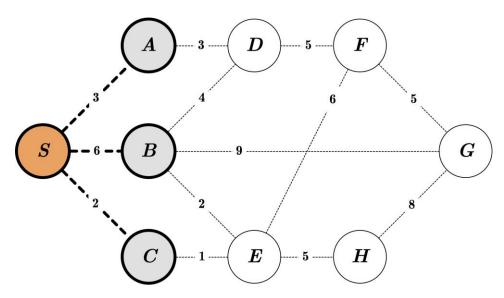
S A D F E H G



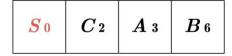
**Priority Queue:** 

 $oldsymbol{S}$  0

Order of Visit:

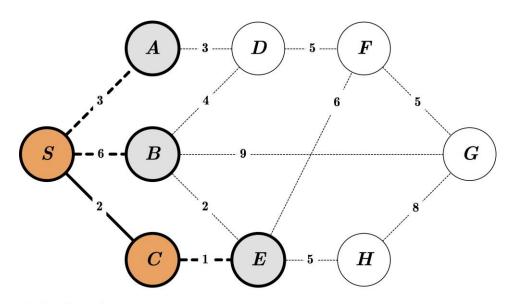


**Priority Queue:** 

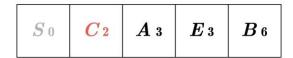


Order of Visit:

 $\boldsymbol{S}$ 

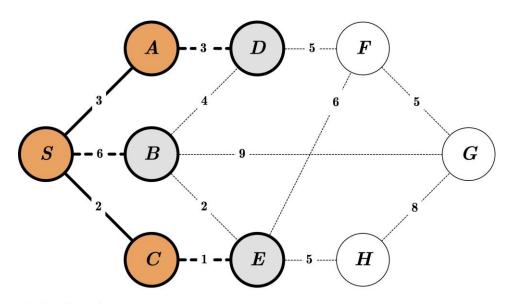


**Priority Queue:** 

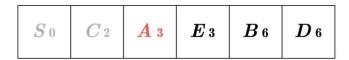


Order of Visit:

S C

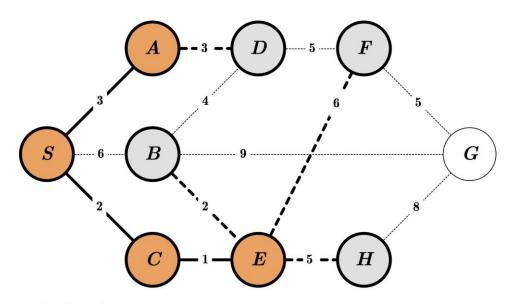


**Priority Queue:** 



Order of Visit:

 $S \quad C \quad A$ 

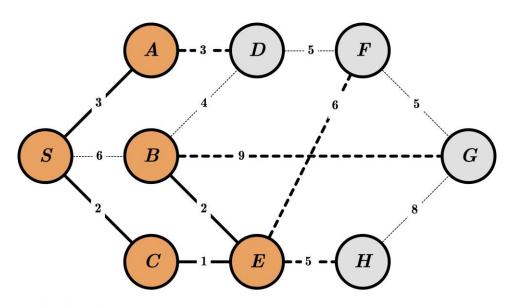


**Priority Queue:** 



Order of Visit:

S C A E

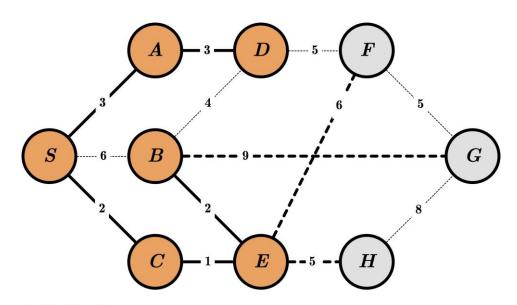


**Priority Queue:** 



Order of Visit:

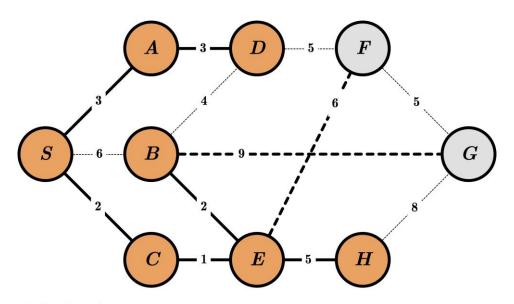
S C A E B



**Priority Queue:** 



Order of Visit:

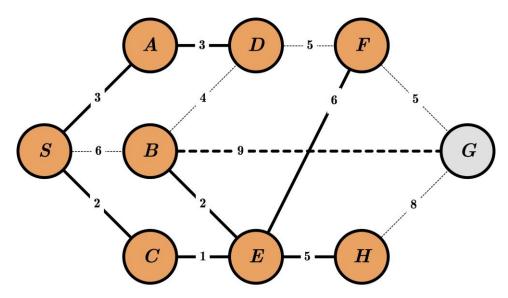


**Priority Queue:** 



Order of Visit:

S C A E B D H

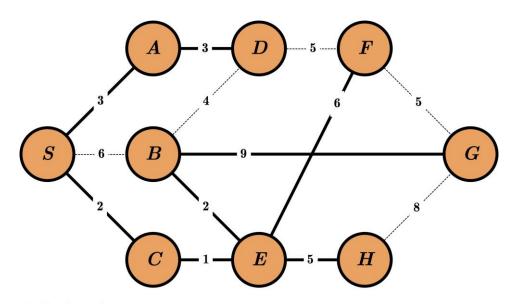


**Priority Queue:** 



Order of Visit:

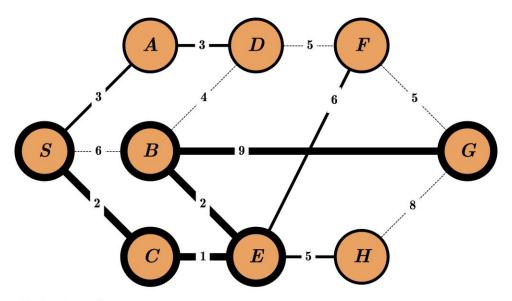
S C A E B D H F



**Priority Queue:** 



Order of Visit:



**Priority Queue:** 

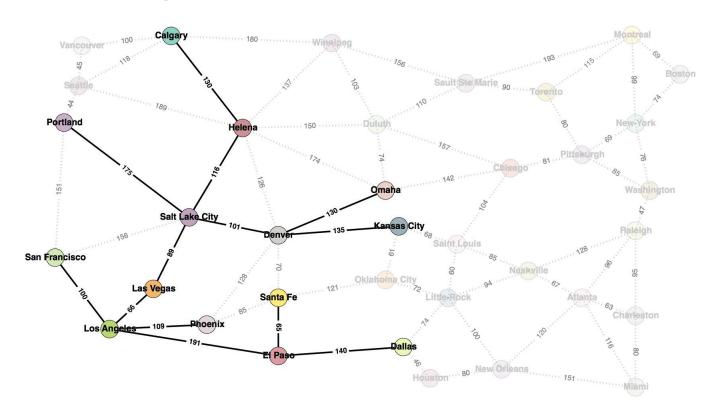


Order of Visit:

## Examples using the map

Start: Las Vegas

Goal: Calgary



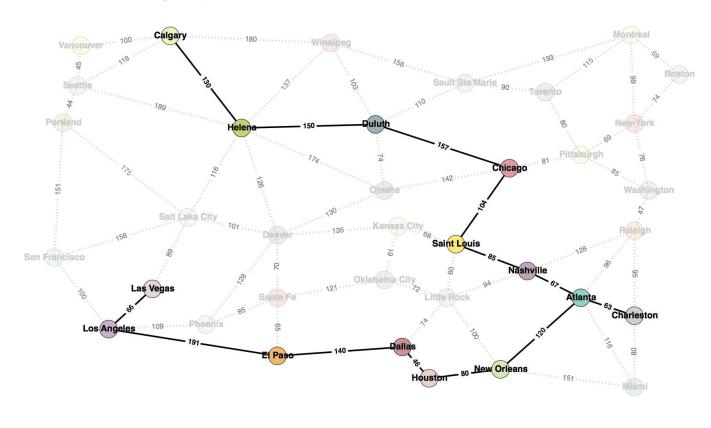
**BFS** 

Order of Visit: Las Vegas, Los Angeles, Salt Lake City, El Paso, Phoenix, San Francisco, Denver, Helena, Portland, Dallas, Santa Fe, Kansas City, Omaha, Calgary.

## Examples using the map

Start: Las Vegas

Goal: Calgary



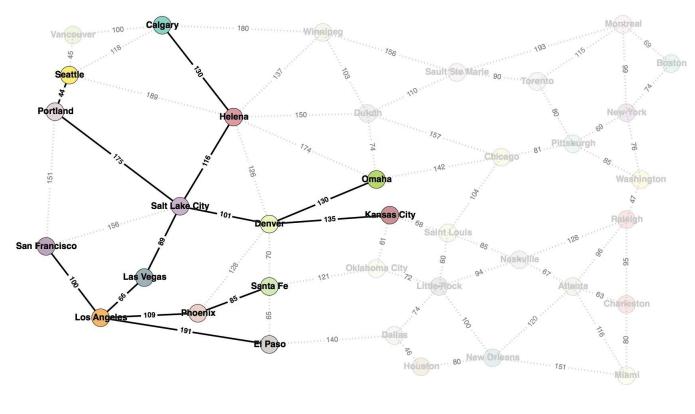
DFS

Order of Visit: Las Vegas, Los Angeles, El Paso, Dallas, Houston, New Orleans, Atlanta, Charleston, Nashville, Saint Louis, Chicago, Duluth, Helena, Calgary.

# Examples using the map

Start: Las Vegas

Goal: Calgary



UCS

Order of Visit: Las Vegas, Los Angeles, Salt Lake City, San Francisco, Phoenix, Denver, Helena, El Paso, Santa Fe, Portland, Seattle, Omaha, Kansas City, Calgary.

#### Credit

• Artificial Intelligence, A Modern Approach. Stuart Russell and Peter Norvig. Third Edition. Pearson Education.