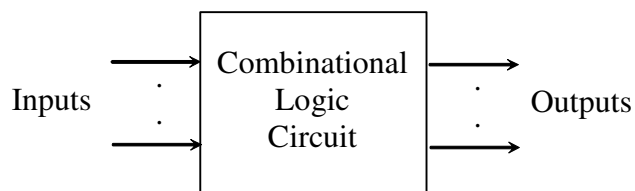


Ch. 2: Combinational Systems

2.1: Design Process for Combinational Systems



To design a logic circuit:

- Describe input/output relationships
- Translate these relationships into a digital logic circuit

2.2: Switching Algebra

An algebraic system to analyze and design *combinational logic circuits*.

Consists of *binary variables* and a set of *logical operations*:

- Binary variables: A, B, C, w, x, y, z , etc (as in algebra), taking values 0 or 1
- Three basic logical operations: AND, OR, NOT

AND $F = x \cdot y$ (F is equal to x AND y) x, y are inputs; F is output
Definition: output is 1 if and only if **both** inputs are 1.

Symbol

Truth Table

Timing Diagram

x	y	
0	0	
0	1	
1	0	
1	1	

OR $F = x + y$ (F is equal to x OR y) x, y are inputs; F is output
Definition: output is 1 if and only if either input is 1.

Symbol

Truth Table

Timing Diagram

x	y	
0	0	
0	1	
1	0	
1	1	

NOT $F = x'$ (F is equal to x prime) x is input; F is output
Definition: output is opposite (inverse) of input. (This takes one input only)

Symbol

Truth Table

Timing Diagram

x	
0	
1	

The **AND** and **OR** operators can be extended to more than two inputs:

- Definition of **AND**: output is 1 if and only if **all** inputs are 1.
- Definition of **OR** : output is 1 if and only if **any** input is 1.

3-input **AND**

$$F = x \cdot y \cdot z = x y z$$

(*dot* is often omitted)

4-input **OR**

$$G = A + B + C + D$$

Precedence of Operations (p. 46)

Order of operations - First: **NOT**; Second: **AND**; Last: **OR**

Parentheses are used to indicate or modify order of logical operations.

$a \cdot (b')$ ← Evaluate b' first, then **AND** the result with a

$(x \cdot y)'$ ← Evaluate $x \cdot y$ first, then take the **NOT** of the result

$(a \cdot b) + c$ ← Evaluate $a \cdot b$ first, then **OR** the result with c

$a \cdot (b + c)$ ← Evaluate $b + c$ first, then **AND** the result with a

Not the same!

$(X \cdot Y')' + Z$ ← Evaluate Y' first, then $(X \cdot Y')$, then $(X \cdot Y')'$, finally the **OR**

Properties (Postulates or Theorems) of Switching Algebra

Summarized in textbook's inside front cover.

P1a.	$a + b = b + a$	P1b.	$a \cdot b = b \cdot a$	Commutative
P2a.	$a + (b + c) = (a + b) + c$	P2b.	$a \cdot (b \cdot c) = (a \cdot b) \cdot c$	Associative
P3a.	$a + 0 = a$	P3b.	$a \cdot 1 = a$	Identity
P3aa.	$0 + a = a$	P3bb.	$1 \cdot a = a$	
P4a.	$a + 1 = 1$	P4b.	$a \cdot 0 = 0$	Null
P4aa.	$1 + a = 1$	P4bb.	$0 \cdot a = 0$	
P5a.	$a + a' = 1$	P5b.	$a \cdot a' = 0$	Complement
P5aa.	$a' + a = 1$	P5bb.	$a' \cdot a = 0$	
P6a.	$a + a = a$	P6b.	$a \cdot a = a$	Idempotency
P7.	$(a')' = a$			Involution
P8a.	$a \cdot (b + c) = a \cdot b + a \cdot c$	P8b.	$a + b \cdot c = (a + b) \cdot (a + c)$	Distributive
P9a.	$a \cdot b + a \cdot b' = a$	P9b.	$(a + b) \cdot (a + b') = a$	Adjacency
P9aa.	$a' \cdot b' + a' \cdot b + a \cdot b + a \cdot b' = 1$	P9bb.	$(a' + b') \cdot (a' + b) \cdot (a + b) \cdot (a + b') = 0$	
P10a.	$a + a' \cdot b = a + b$	P10b.	$a \cdot (a' + b) = a \cdot b$	Simplification
P11a.	$(a + b)' = a' \cdot b'$	P11b.	$(a \cdot b)' = a' + b'$	DeMorgan
P11aa.	$(a + b + c \dots)' = a' \cdot b' \cdot c' \dots$	P11bb.	$(a \cdot b \cdot c \dots)' = a' + b' + c' \dots$	
P12a.	$a + a \cdot b = a$	P12b.	$a \cdot (a + b) = a$	Absorption
P13a.	$a \cdot t_1 + a' \cdot t_2 + t_1 \cdot t_2 = a \cdot t_1 + a' \cdot t_2$	P13b.	$(a + t_1) \cdot (a' + t_2) \cdot (t_1 + t_2)$ $= (a + t_1) \cdot (a' + t_2)$	Consensus
P14a.	$a \cdot b + a' \cdot c = (a + c) \cdot (a' + b)$			

P1 thru P8, and P11 are basic properties. The rest can be derived from them.

Some points worth noting:

P6a means that $xyz + xyz$ can be replaced with xyz . (Eliminate duplicates)
It also means that xyz can be replaced with $xyz + xyz$. (Duplicate a term)

P8a is the distributive property over **OR** function, similar in regular algebra.

P8b means that the distributive property also applies over **AND**. (There is no analogous property in regular algebra)

Proving Basic Properties of Switching Algebra

One way: use a *truth table*. A truth table shows all possible combinations of the input variables, where each input can take on value of 0 or 1. Show if the left-hand side (LHS) is equal to the right-hand side (RHS) for every combination.

Ex. Prove $a \cdot (b + c) = a \cdot b + a \cdot c$ [**P8a**] (See book p. 48 for **P8b** proof)

a	b	c	$b + c$	$a \cdot (b + c)$	$a \cdot b$	$a \cdot c$	$a \cdot b + a \cdot c$

In a truth table: n inputs $\rightarrow 2^n$ rows

Another Example: Determine whether $(a+b)(a'+b') \stackrel{?}{=} a'b + ab'$

a	b	$(a+b)$	$(a'+b')$	LHS	$a'b$	ab'	RHS

Determine whether $(A + B)' \stackrel{?}{=} A' + B'$ [Incorrect application of P11]

A	B	$A+B$	$(A+B)'$	A'	B'	$(A' + B')$
0	0					
0	1					
1	0					
1	1					

Prove Property P11aa $(X + Y + Z)' \stackrel{?}{=} X' \cdot Y' \cdot Z'$ [DeMorgan's Property]

X	Y	Z	$X+Y+Z$	$(X+Y+Z)'$	X'	Y'	Z'	$(X' \cdot Y' \cdot Z')$
0	0	0						
0	0	1						
0	1	0						
0	1	1						
1	0	0						
1	0	1						
1	1	0						
1	1	1						

Is the equality $((A' + B)' \cdot C')' + A = B$ true?

A	B	C		
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

Determine whether $a(a + b)(a b')' = 1$ is true

a	b		
0	0		
0	1		
1	0		
1	1		

A	B	C	D		

a	b	c
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

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