

Unit 2

AES

Introduction

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- In Oct 2000 Rijndael was released.
- In Nov 2001 became U.S government standard(FIPS 197).
- *Features of AES*
 1. *Flexible*
 2. *Adapted to modern processor*
 3. *Suited to Smart Card*
 4. *Protection against cryptanalysis attacks.*

Algorithm Steps

1. Do the following one-time initialization:

- a) *Expand the 16-byte key to get the actual key block to be used.*
- b) *Do one time initialization of the 16-byte PT block (called as state).*
- c) *XOR the state with key block.*

2. For each round do the following:

- a) *Apply S-box to each of the PT bytes.*
- b) *Rotate row k of the PT block(i.e. state) by k bytes.*
- c) *Perform Mix columns operation.*
- d) *XOR the state with key block.*

Expand the 16-byte key to get the actual key block to be used.

- 16 byte
- Array size $4*4$
- 11 such array
- 1 for initialization & remain 10 for 1 round each.
- Original key copied as it is
- *Key Expansion* = $11*4*4 = 176$ bytes
- In the context of AES a word means 4 bytes
- So initial 16-byte key ($16/4 = 4$ word key)
- Will be expanded into 176 bytes key ($176/4 = 44$ words)

Key Expansion /Add round Key Algorithm

- *If the word in the W array is multiple of four.*

- $TMP = W[i-1]$ previous word
 $= W[4-1]$ word 4 place earlier
 $= W[3]$

Since $i=4$,

$I \bmod 4$ is 0 this is multiple of 4.

$TMP = S\text{-box} (\text{Rotate}(TMP)) \text{ XOR } W[i-4] \text{ XOR } Rcon$

- *Otherwise*

$TMP = W[i-1]$ previous word XOR $W[i-4]$ word 4 place earlier

Process in each Round

- a) Confusion
- b) Diffusion
- c) Matrix Multiplication using Galois Field
- d) O/P of Step (c) **XOR** Add round key

- *For step (C):*
- O/P value from step (B) & a constant matrix is used.