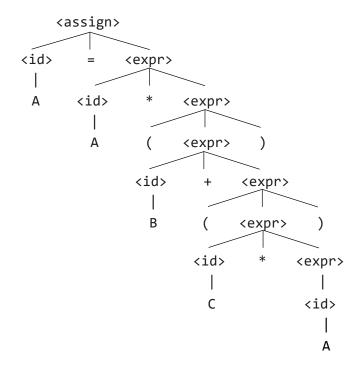
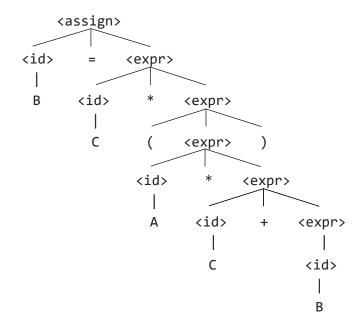
CHAPTER 3 SOLUTIONS

```
3.
\langle assign \rangle \rightarrow \langle id \rangle = \langle expr \rangle \langle id \rangle \rightarrow A \mid B \mid C
<expr> → <expr> * <term> | <term>
<term> → <factor> + <term> | <factor>
<factor> → ( <expr> )
            | <id>>
4.
<assign> -> <id> = <expr>
<id> -> A | B | C
<expr> -> <expr> + <term>
              | <term>
<term> -> <term> * <factor>
              | <factor>
<factor> -> ( <expr> )
                | <id>
                 | <id> ++
                 | <id>> - -
6.
a) A = A * (B + (C * A))
<assign> => <id> = <expr>
              => A = <expr>
              => A = <id> * <expr>
              \Rightarrow A = A * <expr>
              \Rightarrow A = A * ( <expr> )
              \Rightarrow A = A * ( \langle id \rangle + \langle expr \rangle )
              \Rightarrow A = A * (B + <expr>)
              => A = A * (B + ( <expr> ) )
              \Rightarrow A = A * ( B + ( \langle id \rangle * \langle expr \rangle ) )
              => A = A * (B + (C * < expr>))
              \Rightarrow A = A * ( B + ( C * <id> ) )
              \Rightarrow A = A * ( B + ( C * A ) )
```





Parse Tree

$$b.A = B + C + A$$

Parse Tree

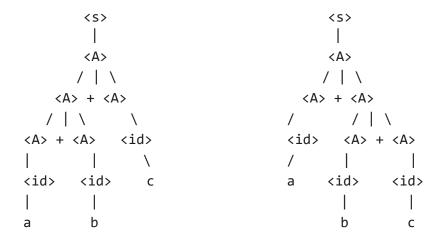
$$c.A = A * (B + C)$$

```
=> A = A * (factor>
=> A = A * (<expr>)
=> A = A * (<expr> + <term>)
=> A = A * (<term> + <term>)
=> A = A * (<factor> + <term>)
=> A = A * (<id> + <term>)
=> A = A * (B + <term>)
=> A = A * (B + <factor>)
=> A = A * (B + <id>)
```

Parse Tree

8.
$$\langle S \rangle \rightarrow \langle A \rangle$$
 $\langle A \rangle \rightarrow \langle A \rangle + \langle A \rangle \mid \langle id \rangle$ $\langle id \rangle \rightarrow a \mid b \mid c$

Take "a+b+c" as example, it has two possible parse trees:



which will cause the ambiguity.

11.
$$\langle S \rangle \rightarrow \langle A \rangle \ a \ \langle B \rangle \ b$$
 $\langle A \rangle \rightarrow \langle A \rangle \ b \ | \ b$ $\langle B \rangle \rightarrow \ a \ \langle B \rangle \ | \ a$ This BNF can be represented in EBNF: $S = "b", \{"b"\}, "a", "a", \{"a"\}, "b";$

Thus, "baab" and "bbaab" are generated by this grammar; "bbbab" and "bbaaaaa" are not.

The answer is (a) (d).

12.
$$\langle S \rangle \rightarrow a \langle S \rangle \ c \ \langle B \rangle \ | \ \langle A \rangle \ | \ b$$
 $\langle A \rangle \rightarrow c \ \langle A \rangle \ | \ c$ $\langle B \rangle \rightarrow d \ | \ \langle A \rangle$ This BNF can be represented in EBNF: $S = "a", S, "c", ("d" \ | "c", \{"c"\}) \ | \ "c", \{"c"\} \ | \ "b";$

Thus, "abcd" and "accc" are generated by this grammar; the others are not. The answer is (a) (e).

14.

For sentence aabb:

Derivation -

=> a a b b

Parse Tree



For sentence aaaabbbb:

Derivation -

=> a a a a b b b b

Parse Tree

