CSCE 221 Cover Page

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Please list all sources in the table below including web pages which you used to solve or implement the current homework. If you fail to cite sources you can get a lower number of points or even zero, read more Aggie Honor System Office $\frac{1}{2} \frac{1}{2} \frac{1}{2}$						
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I certify that I have listed all the sources that I used to develop the solutions/codes to the submitted work.

"On my honor as an Aggie, I have neither given nor received any unauthorized help on this academic work."

Your Name Rong Xu Date 2020/10/10

Homework 2

due October 16 at 11:59 pm to eCampus

- 1. (20 points) Given two sorted lists, L1 and L2, write an efficient C++ code to compute L1 ∩ L2 using only the basic STL list operations.
 - (a) Provide evidence of testing: submit your code

```
#include <iostream>
#include <list>
using namespace std;
list<int> inter_list(list<int> L1,
list < int > L2) {
         if(L1.empty() || L2.empty())
                  std::exit(-1);
         L1.merge(L2);
         list <int> return_list;
         for(list < int > :: iterator it = L1.begin(); it != L1.end(); it++) \ \{
                  if(*it == *--it)
                           return_list.push_back(*it);
                  ++it:
         return return_list;
7
int main() {
int A1[]={1,2,3,4,5,6};
int A2[]={2,4,6,8,9,10};
list<int> iL1(A1, A1+6);
list<int> iL2(A2, A2+6);
list<int> iL3 = inter_list(iL1, iL2);
list<int>::iterator it = iL3.begin();
while(it != iL3.end()) {
         cout << *it++ << "";
system("pause");
return 0;
 h: pause: command not found
4 6
```

- (b) What is the running time of your algorithm?
 - i. From the official website, the running time of function std::merge() is O(n).
 - A. https://en.cppreference.com/w/cpp/algorithm/merge
 - B. n is the number of elements of List 1.
 - ii. The big-o of for loop is O(n+k).
 - A. n is the number of elements of List 1
 - B. k is the number of elements of List 2
 - C. In this case, the list L1 was merged by List 2. Thus the List 1 now has (n + k) element.
 - iii. Base of the statements above, the running time is O(n) + O(n + k) = O(n + k)

2. (20 points) Write a C++ recursive function that counts the number of nodes in a singly linked list.

(a) Test your function using different singly linked lists. Include your code.

```
#include<iostream> using namespace std;
class Node {
        public:
                 int data:
                 Node *next;
                 Node(int da = 0, Node *p = NULL) {
                          this -> data = da;
                          this -> next = p;
                 }
};
class List{
        private:
                 Node *head, *tail; public:
        List(){head = tail = NULL;};
         ~List(){delete head; delete tail;};
        void print() {
                 Node *p = head;
                 while (p != NULL) {
                          cout << p -> data << "_{\sqcup}\a";
                          p = p->next;
                 cout << endl;
        };
         void Insert(int da) {
                 if (head == NULL) {
                          head = tail = new Node(da);
                          head -> next = NULL;
                          tail -> next = NULL;
                 } else {
                          Node *p = new Node(da);
                          tail -> next = p;
                          tail = p;
                          tail -> next = NULL;
                 }
        Node* first() {return head;}
};
int count_node(Node* n) {
        if(n == NULL)
                 return 0:
         else
                 return 1 + count node(n -> next):
}
int main() {
        List 11:
        11.Insert(1); 11.Insert(2); 11.Insert(3); 11.Insert(4);
        11.Insert(5); 11.Insert(6); 11.Insert(7);
        11.print();
        cout << "theunumberuofunodes:u" << count_node(l1.first()) << endl;
         system("pause");
        return 0;
7
   /Users/abby/Desktop/2020_fall/CSCE221_511/hw2/Q2 ; exit;
   3 4 5 6 7
number of nodes: 7
pause: command not found
```

- (b) Write a recurrence relation that represents your algorithm.
 - i. Base case: When n is a null pointer, return 0.
 - A. T(0) = c1 for some constant c1
 - ii. Recursive case: When n is not a null pointer.
 - A. T(n) = c2 + T(n 1) for some constant c2
- (c) Solve the recurrence relation using the iterating or recursive tree method to obtain the running time of the algorithm in Big-O notation.
 - i. If we knew T(n 1), we could solve T(n).
 - ii. $T(n) = T(n-1) + c2 = T(n-2) + c2 + c2 = T(n-2) + 2c2 = T(n-3) + 3c2 = \dots = T(n-k) + kc2$
 - iii. So we have T(n) = T(n-k) + k * c2 for all k
 - iv. If we set k = n, we have $T(n) = T(n-n) + nc^2 = T(0) + nc^2 = c^2 + nc^2 = O(n)$

- 3. (20 points) Write a C++ recursive function that finds the maximum value in an array (or vector) of integers without using any loops.
 - (a) Test your function using different input arrays. Include the code.

- (b) Write a recurrence relation that represents your algorithm.
 - i. Base case: When n reach to the first element in the array/vector.
 - A. T(0) = c1 for some constant c1
 - ii. Recursive case: Fund the max element when not reach to the first element..

A.
$$T(n) = c2 + T(n - 1)$$
 for some constant c2

- (c) Solve the recurrence relation and obtain the running time of the algorithm in Big-O notation.
 - i. If we knew T(n 1), we could solve T(n).
 - ii. $T(n) = T(n-1) + c2 = T(n-2) + c2 + c2 = T(n-2) + 2c2 = T(n-3) + 3c2 = \dots = T(n-k) + kc2$
 - iii. So we have T(n) = T(n-k) + k * c2 for all k
 - iv. If we set k = n, we have T(n) = T(n n) + nc2 = T(0) + nc2 = c1 + nc2 = O(n)

- 4. (20 points) What is the best, worst and average running time of quick sort algorithm?
 - (a) Provide recurrence relations and their solutions.
 - i. Best case:

[Process completed]

- A. Recurrence Relation: T(n) = T(n/2) + T(n/2) + O(n) and T(1) = 0
- B. Solve it by iteration method:

$$\begin{split} &\mathbf{T}(\mathbf{n}) = & 2\mathbf{T}(\mathbf{n}/2) + \mathbf{O}(\mathbf{n}) \\ &= & 2(2\mathbf{T}(\mathbf{n}/4) + \mathbf{O}(\mathbf{n}/2)) + \mathbf{O}(\mathbf{n}/2)) + \mathbf{O}(\mathbf{n}) \\ &= & 4\mathbf{T}(\mathbf{n}/4) + 2^*\mathbf{O}(\mathbf{n}) \\ &= & \dots \\ &= & 2^kT(n/2^k) + K*O(n) \\ &= & O(nlog_2n) \end{split}$$

ii. Worst case:

A. Recurrence Relation:
$$T(n) = T(n-1) + T(1) + n = T(n-1) + n$$
 and $T(1) = 0$

B. Solve it by iteration method:

$$\begin{split} T(n) &= T(n\text{-}1) + n \\ &= T(n\text{-}2) + (n\text{-}1) + n \\ &= T(n\text{-}3) + (n\text{-}2) + (n\text{-}1) + n \\ &= \dots \\ &= T(n\text{-}k\text{+}!) + (n\text{-}k\text{+}2) + \dots + (n\text{-}1) + n \\ &= \dots \\ &= T(1) + 2 + 3 + \dots + (n\text{-}1) + n \\ &= \operatorname{fracn}(n+1)2 - 1 \\ &= O(n^2) \end{split}$$

- iii. Average case:
 - A. Recureence Relation: T(n) = T(cn) + T((1-c)n) + n and T(1) = 0
 - B. Solve it by interation method: Let hl be the height of left subtree and hr be the hright of right subtree. Notice that hr > hl.
 - C. solve for hl and hr:

$$c^{h_L} = 1/n$$

$$h_L = -log_c h$$

$$h_L = log_{1/c} h$$

$$(1 - c)^{h_R} = 1/n$$

$$h_R = -log_{(1-c)} h$$

$$h_R = log_{1/(1-c)} h$$

- D. The big-o is $O(nlog_2n)$
- (b) Provide arrangement of the input and the selection of the pivot point for each case.
 - i. For the best case: the input has already been sorted and the pivot is just the middlest element in the list.
 - ii. For the average case: we usually choose the pivot in the middle or just choose the random index of the pivot.
 - iii. For the worst case: the input list is reversed and we start to choose the pivot at the beginning of the list until the last one.

5. (20 points) Write a C++ function that counts the total number of nodes with two children in a binary tree (do not count nodes with one or none child). You can use a STL container if you need to use an additional data structure to solve this problem. Use the big-O notation to classify your algorithm. Include your code.

```
#include <iostream>
#include <vector>
using namespace std;
std::vector < int > nodes:
struct BiTNode {
         int data:
         struct BiTNode* lchild;
         struct BiTNode* rchild;
}:
void create_tree(BiTNode* &tree) {
         int data; cin >> data;
if (data != '\n') {
                 if (data == -1) \{
                          tree = nullptr;
                  } else {
                           tree = new BiTNode;
                           tree->data = data;
                           create_tree(tree->lchild);
                           create_tree(tree->rchild);
                 }
         }
}
void pre_order_traverse(BiTNode* &tree) {
        if (tree) {
                  cout << tree->data << "_{\sqcup}";
                  pre_order_traverse(tree->lchild);
                  pre_order_traverse(tree->rchild);
         }
void count_node(BiTNode* &tree) {
         cout << "data: " << tree -> data;
         if(tree -> lchild != nullptr && tree -> rchild != nullptr)
                 nodes.push_back(tree -> data);
         if(tree -> lchild == nullptr && tree -> rchild == nullptr)
                 return;
         count_node(tree->lchild);
         count_node(tree->rchild);
int main() {
    BiTNode* T;
         create_tree(T);
                input here is:
                                     2 3 5 -1 -1 6 -1 -1 4 -1 -1
                ١
           3
          11
         5 6
pre_order_traverse(T); //2, 3, 5, 6, 4
count_node(T);
cout << "\ntheunodesuare:u" <<endl;
cout << "\ntheunumberuofunodeuwithutwouchildrenuis:u" << nodes.size() << endl; ,
\mathtt{system} \, (\, \tt{"pause"} \, ) \, ;_{\, \tt{UUUUUUUU}} \,
return<sub>□</sub>0;
the number of node with two children is: 2
sh: pause: command not found
```

(a) I think the the big-o if my algorithm is $O(\log_2 n)$