## **Topic 7: Control Flow Instructions**



CSE 30: Computer Organization and Systems Programming
Summer Session II

Dr. Ali Irturk

Dept. of Computer Science and Engineering

University of California, San Diego

#### So Far...

- \*All instructions have allowed us to manipulate data
- So we've built a calculator
- \*In order to build a computer, we need ability to make decisions...



#### Labels

- Any instruction can be associated with a label
- Example:

```
start ADD r0, r1, r2; a = b+c next SUB r1, r1, #1; b--
```

- ❖ In fact, every instruction has a label regardless if the programmer explicitly names it
  - The label is the address of the instruction
  - ❖ A label is a pointer to the instruction in memory
  - Therefore, the text label doesn't exist in binary code



#### C Decisions: if Statements

```
❖ i f statements in C

  ❖if (condition) clause
   ❖if (condition) clause1 else clause2
*Rearrange 2nd if into following:
   if (condition) goto L1;
      clause2;
       goto L2;
    L1: clause1;
    L2:
   Not as elegant as if-else, but same meaning
```



## **ARM** goto Instruction

- The simplest control instruction is equivalent to a C goto statement
- goto label (in C) is the same as:
- ❖ B label (in ARM)
- ❖ B is shorthand for "branch". This is called an unconditional branch meaning that the branch is done regardless of any conditions.
- \* There are also conditional branches



#### **ARM Decision Instructions**

❖ ARM also has variants of the branch instruction that only goto the label if a certain condition is TRUE

#### Examples:

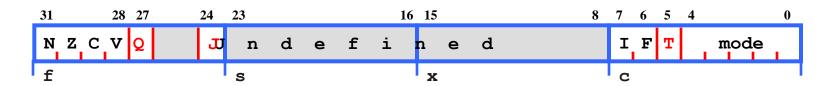
```
BEQ label; BRANCH EQUAL
BNE label; BRANCH NOT EQUAL
BLE label; BRANCH LESS THAN EQUAL
BLT label; BRANCH LESS THAN
BGE label; BRANCH GREATER THAN EQUAL
```

- Plus more ...
- The condition is T/F based upon the fields in the Program Status Register

❖ BGT label ; BRANCH GREATER THAN



## **Program Status Registers**



- Condition code flags
  - N = Negative result from ALU
  - $\star$  Z = Zero result from ALU
  - ❖ C = ALU operation Carried out
  - $\star$  V = ALU operation oVerflowed
- Sticky Overflow flag Q flag
  - Architecture 5TE/J only
  - Indicates if saturation has occurred
- \* J bit
  - Architecture 5TEJ only
  - $\star$  J = 1: Processor in Jazelle state

- Interrupt Disable bits.
  - $\bullet$  I = 1: Disables the IRQ.
  - $\star$  F = 1: Disables the FIQ.
- \* T Bit
  - Architecture xT only
  - $\star$  T = 0: Processor in ARM state
  - $\star$  T = 1: Processor in Thumb state
- Mode bits
  - Specify the processor mode

# Flags and Their Use

- The N flag
  - ❖ Set if the result is negative or equivalently if the MSB == '1'
- The Z flag
  - Set if the result is zero
- The C flag
  - Set if
    - $\bullet$  The result of an addition is greater than  $2^{32}$
    - \* The result of a subtraction is positive
    - Carryout from the shifter is '1'
- The V flag (oVerflow)
  - Set if there is overflow



#### **Condition Codes**

- The possible condition codes are listed below
  - Note AL is the default and does not need to be specified

Suffix	Description	Flags tested
EQ	Equal	Z=1
NE	Not equal	Z=0
CS/HS	Unsigned higher or same	C=1
CC/LO	Unsigned lower	C=0
MI	Minus	N=1
PL	Positive or Zero	N=0
VS	Overflow	V=1
VC	No overflow	V=0
HI	Unsigned higher	C=1 & Z=0
LS	Unsigned lower or same	C=0 or Z=1
GE	Greater or equal	N=V
LT	Less than	N!=V
GT	Greater than	Z=0 & N=V
LE	Less than or equal	Z=1 or N=!V
AL	Always	

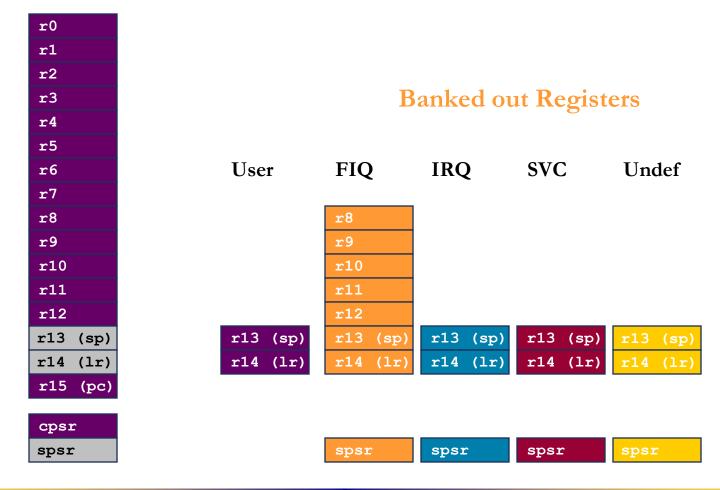


# The ARM Register Set

#### Only need to worry about cpsr (current program status register)

**Current Visible Registers** 

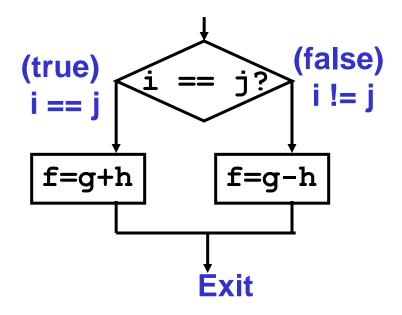
**Abort Mode** 





# Compiling C if into ARM

## Compile by hand



#### Use this mapping:



## **Comparison Instructions**

- ❖ In order to perform branch on the "==" operation we need a new instruction
- ❖ CMP Compare: subtracts a register or an immediate value from a register value and updates condition codes

#### Examples:

```
\diamond CMP r3, #0 ; set Z flag if r3 == 0
```

 $\diamond$  CMP r3, r4; set Z flag if r3 == r4

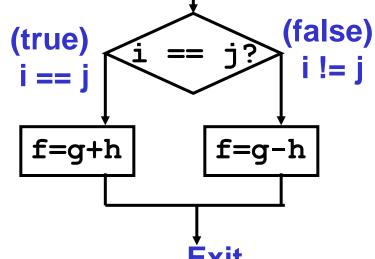
All flags are set as result of this operation, not just Z.



# Compiling C if into ARM

#### Compile by hand

```
if (i == j) f=g+h;
else f=g-h;
```



Final compiled MIPS code:

```
CMP r3, r4 ; Z = 1 if i==j

BEQ True ; goto True when i==j

SUB r0,r1,r2 ; f=g-h(false)

B Fin ; goto Fin

True ADD r0,r1,r2 ; f=g+h (true)

Fin
```

Note: Compiler automatically creates labels to handle decisions (branches) appropriately. Generally not found in C code.



## Loops in C/Assembly

#### Simple loop in C;

```
do {
    g--;
    i = i + j;}
while (i != h);
```

#### \*Rewrite this as:

```
Loop: g--;

i = i + j;

if (i != h) goto Loop;
```

#### Use this mapping:

```
g: r1, h: r2, i: r3, j: r4
```



## Loops in C/Assembly

#### Final compiled MIPS code:

```
Loop SUB r1,r1,#1 ; g--
ADD r3,r3,r4 ; i=i+j
CMP r3,r2 ; cmp i,h
BNE Loop ; goto Loop
; if i!=h
```



## Inequalities in ARM

- ❖Until now, we've only tested equalities (== and != in C). General programs need to test < and > as well.
- ♦Use CMP and BLE, BLT, BGE, BGT
- \*Examples:



## Loops in C/Assembly

- There are three types of loops in C:
  - ❖while
  - ❖do... while
  - for
- \*Each can be rewritten as either of the other two, so the method used in the previous example can be applied to while and for loops as well.
- \*Key Concept: Though there are multiple ways of writing a loop in ARM, conditional branch is key to decision making



## **Example: The C Switch Statement**

\*Choose among four alternatives depending on whether k has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {
  case 0: f=i+j; break; /* k=0*/
  case 1: f=g+h; break; /* k=1*/
  case 2: f=g-h; break; /* k=2*/
  case 3: f=i-j; break; /* k=3*/
```



## **Example: The C Switch Statement**

- This is complicated, so simplify.
- \*Rewrite it as a chain of if-else statements, which we already know how to compile:

```
if(k==0) f=i+j;
else if(k==1) f=g+h;
else if(k==2) f=g-h;
else if(k==3) f=i-j;
```

Use this mapping:

```
f: $s0, g: $s1, h: $s2, i: $s3, j: $s4, k: $s5
```



## **Example: The C Switch Statement**

```
; compare k, 0
    CMP r5, #0
                    : branch k!=0
    BNE L1
                 ; k==0 so f=i+j
    ADD r0, r3, r4
                    ; end of case so Exit
    В
       Exit
L1 CMP r5,#1
             ; compare k, -1
    BNE L2
    ADD r0, r1, r2 ; k==1 so f=g+h
                    ; end of case so Exit
       Exit
    В
L2 CMP r5,#2
                    ; compare k, 2
                    ; branch k!=2
    BNE L3
    SUB r0, r1, r2
                    ; k==2 so f=q-h
    В
       Exit
                    ; end of case so Exit
                    ; compare k, 3
L3 CMP r5,#3
                    ; branch k!=3
    BNE Exit
    SUB r0, r3, r4 ; k==3 so f=i-j
```



#### **Predicated Instructions**

\*All instructions can be executed conditionally. Simply add {EQ,NE,LT,LE,GT,GE, etc.} to end

#### C source code

# if (r0 == 0) { r1 = r1 + 1; } else { r2 = r2 + 1; }

#### **ARM** instructions

```
unconditional
```

```
CMP r0, #0
BNE else
ADD r1, r1, #1
B end
else
ADD r2, r2, #1
end
...
```

#### conditional

```
CMP r0, #0
ADDEQ r1, r1, #1
ADDNE r2, r2, #1
```

- 5 instructions
- 5 words
- 5 or 6 cycles

- 3 instructions
- 3 words
- 3 cycles



#### Conclusions

- \*A Decision allows us to decide which pieces of code to execute at run-time rather than at compile-time.
- \*C Decisions are made using conditional statements within an if, while, do while or for.
- **❖**CMP instruction sets status register bits
- \*ARM Decision making instructions are the conditional branches: BNE, BEQ, BLE, BLT, BGE, BGT.



#### Conclusion

#### **❖**Instructions so far:

```
Previously:
```

```
ADD, SUB, MUL, MULA, [U|S]MULL, [U|S]MLAL, RSB AND, ORR, EOR, BIC MOV, MVN LSL, LSR, ASR, ROR

New:

CMP, B{EQ, NE, LT, LE, GT, GE}
```

