Hint : Minimum and Maximum in the Binary Tree

1. Initialise current as root.
2. Take two variables, max and min.
3. While current is not null:
    - If the current does not have a left child:
        - Update variable max and min with current's data if required.
        - Go to the right, i.e., current = current->right.
    - Else:
        - Make current as the right child of the rightmost node in the current's left subtree.
        - Go to this left child, i.e., current = current->left.

This algorithm is based on the approach of traversing the given tree and for every node returning the maximum and minimum of 3 values: node's data, maximum in node's left subtree, and maximum in node's right subtree 1. The time complexity of this algorithm is O(N), where N is the number of nodes in the tree, as every node of the tree is processed once 2. The space complexity of this algorithm is O(1) as no extra space is used 2.

PSEUDO CODE

```
Procedure findMinMax(node, min, max)
   If node is NULL Then
      Return
   End If

   If node.data is less than min Then
      min = node.data
   End If

   If node.data is greater than max Then
      max = node.data
   End If

   Call findMinMax(node.left, min, max)
   Call findMinMax(node.right, min, max)
End Procedure
```