

Cinema Seating

Computing Potatoes

Safia Aarouss
Lalitha Boddapati
Yasmin van Dijk
Georgiana Juglan
Rebecca Ruiter



Contents

- Problem Statement
- Algorithms
- Time complexity
- Competitive ratio
- Experiments
- Conclusion



Problem Statement

- 1,5 meter distance between groups
- Groups can not be split up

Goal

- Maximize number of visitors under the given constraints



Offline / Exact Algorithms

First fit

- Small groups first
- Big groups first

Best Fit

- Small groups first
- Big groups first

Branch and bound

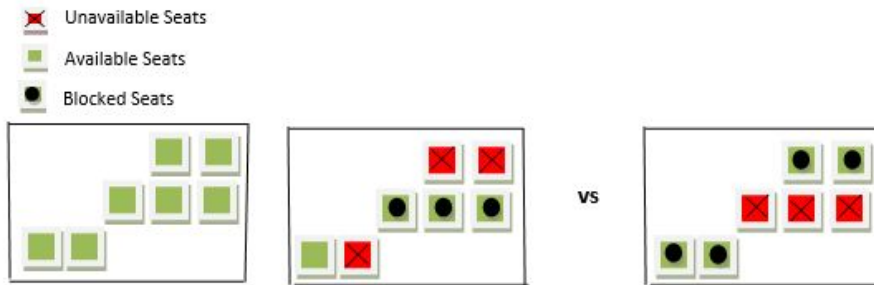


Offline First Fit

- Finds seats for groups one by one
- 2 versions:
 - Smallest groups first
 - Biggest groups first
- Places groups from top to bottom row, left to right
- First found seating given the constraints

Offline First Fit

→ Seating larger groups first does not mean maximising the total number of people



→ But neither does seating smaller groups first

Offline Best Fit

- Finds seats for groups one by one
- 2 versions:
 - Smallest groups first
 - Biggest groups first
- Minimize amount of new blocked seats



Offline Branch and Bound

- Priority queue with partial solutions
 - First partial solution: the empty cinema
 - Ordered on number of visitors placed
- Branch:
 - Find first adjacent group of available seats
 - Place groups of different sizes left-most
 - Mark one seat as unavailable

[1, 0, 1, 1, 0, 0, 0, 0]

xxxx0+111

+++++1111

new partial solutions:

[0, 0, 1, 1, 0, 0, 0, 0]

xxxx0+x++

+++++++1

[1, 0, 0, 1, 0, 0, 0, 0]

xxxx0+xxx

+++++++

[1, 0, 1, 1, 0, 0, 0, 0]

xxxx0+11

+++++1111



Offline Branch and Bound

- Priority queue with partial solutions
 - First partial solution: the empty cinema
 - Ordered on number of visitors placed
- Branch:
 - Find first adjacent group of available seats
 - Place groups of different sizes left-most
 - Mark one seat as unavailable
- Bound:
 - $\text{Nr. of visitors placed} + \text{nr. of available seats} < \text{nr. of visitors placed in best solution}$
- Stop:
 - Queue is empty
 - Time-out
 - Return best found solution

Online Algorithms

- First Fit
- Best Fit
- Order of groups is fixed



Time Complexity

D = Dictionary, G = Groups

First Fit:

→ $O(n+DG)$

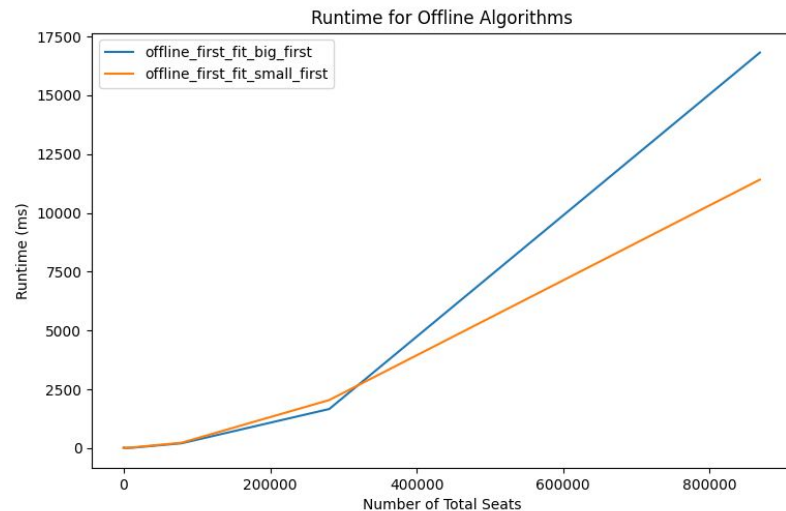
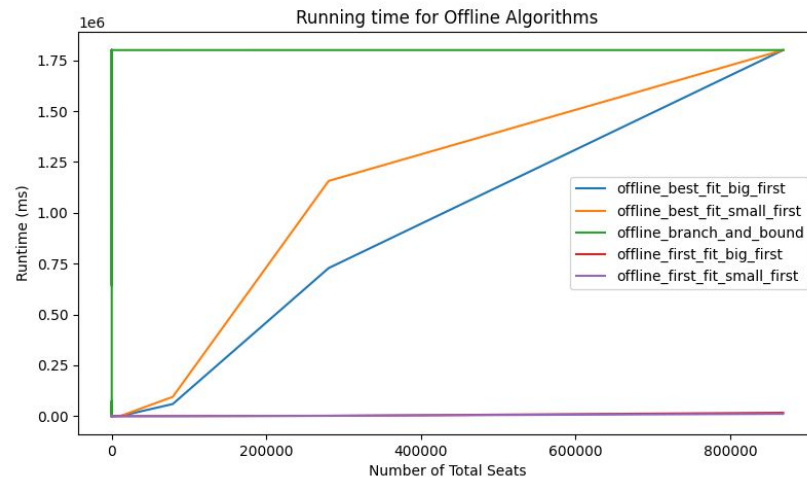
Best Fit:

→ $\Theta(n+DG)$

Branch and bound:

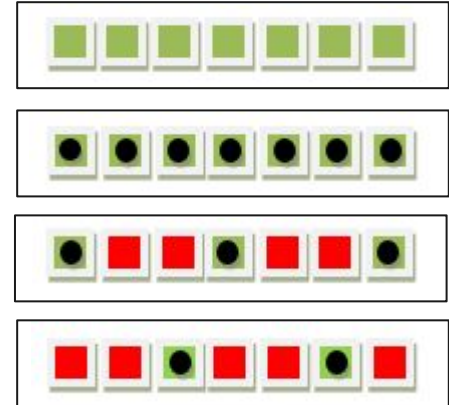
→ Likely $O(n!)$

→ Bound



Competitive Ratio

- Simple cases: row of 8
- Larger cinemas: 3 x 8 rows and beyond
- Online behavior first-fit irrelevant in all-seats cinema
- Worst-Case vs. Average-Case

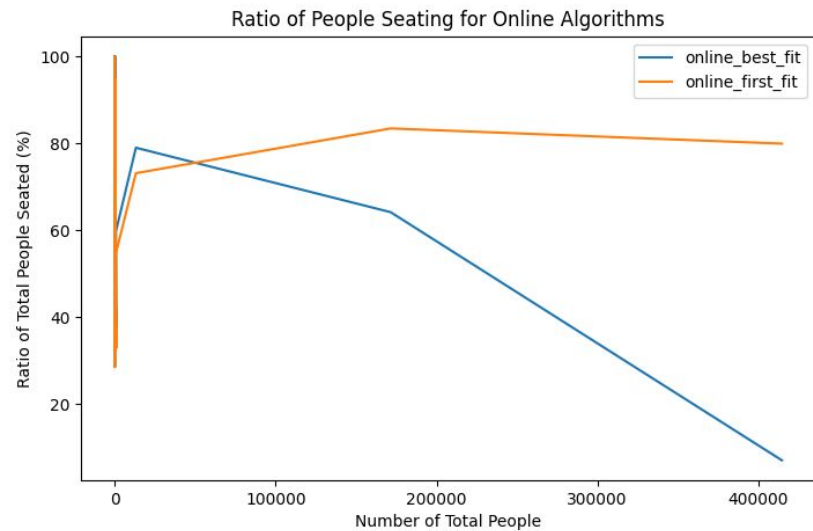
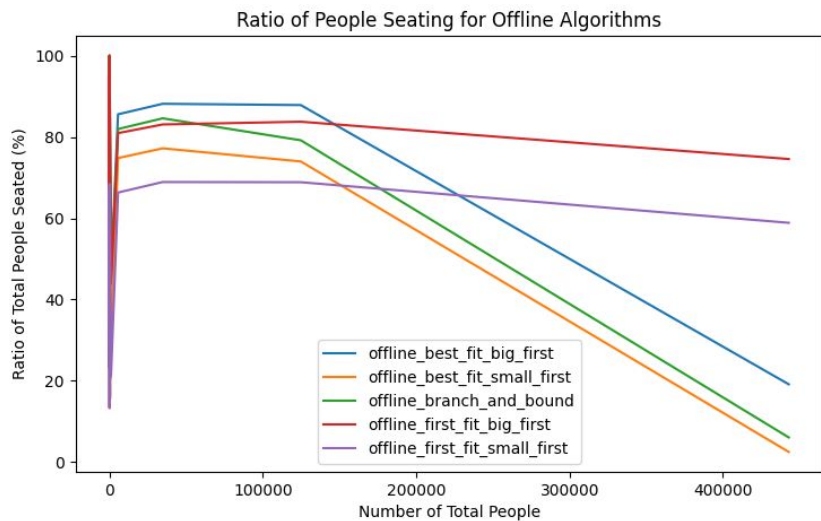




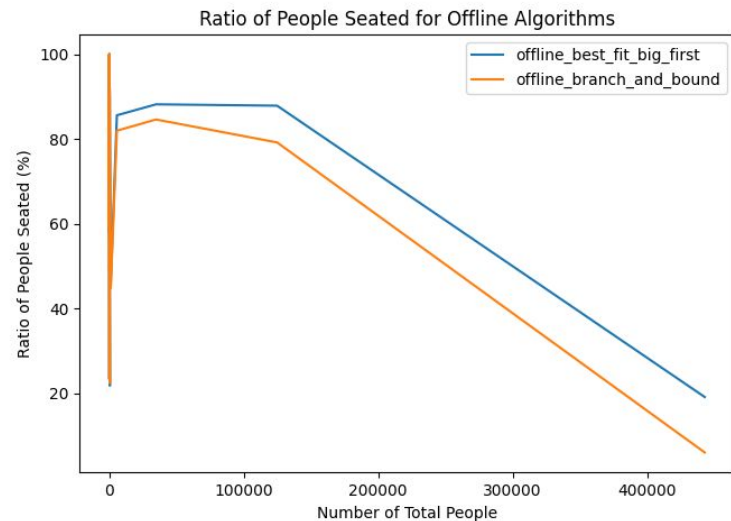
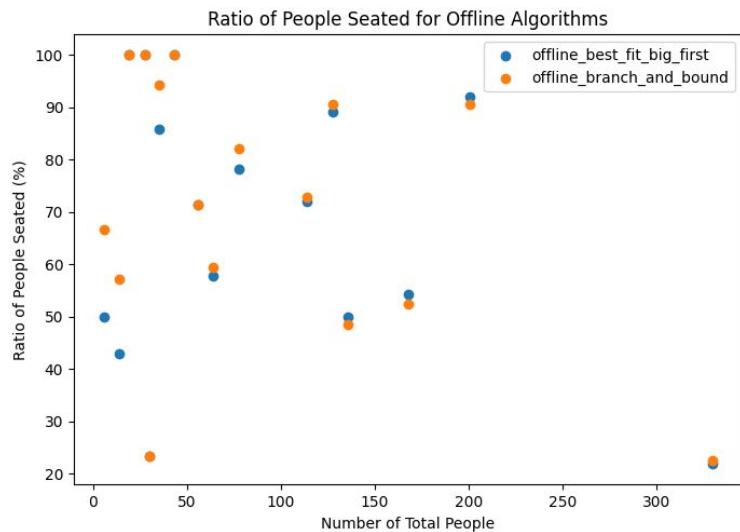
Experiments

- C#
- Time-limit: 30 minutes
- PC specifications:
 - OS: Windows 10
 - CPU: Intel Core i5-7400 3.00 GHz
 - RAM: 8 GB

Experiments

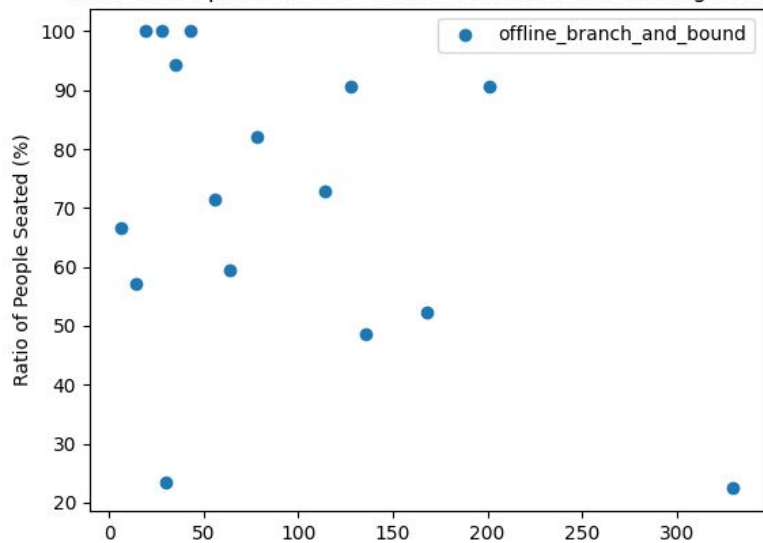


Experiments

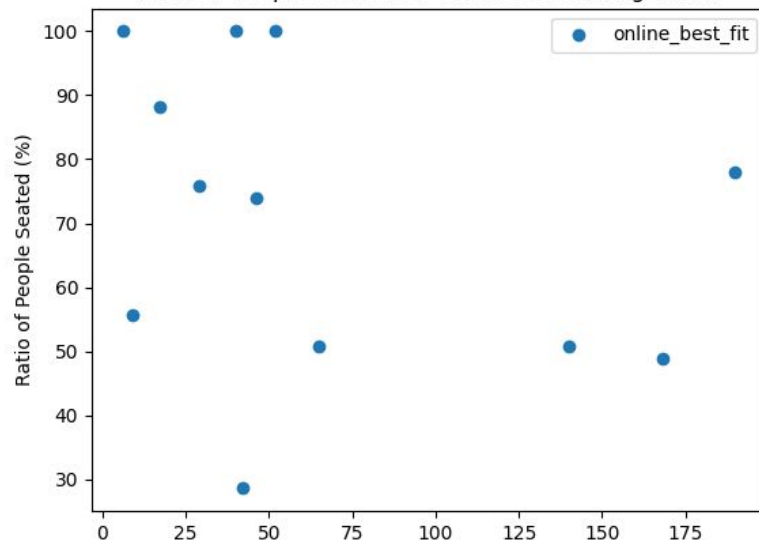


Experiments

Ratio of People Seated for Offline Branch and Bound Algorithm



Ratio of People Seated for Online Best Fit Algorithm





Conclusion

Offline / Exact algorithms:

- Smaller cinemas: branch-and-bound
- Bigger cinemas: best-fit or first-fit

Online algorithms:

- Smaller cinemas: best-fit
- Bigger cinemas: first-fit