

Automated Negotiation for Supply Chain Management

Yasser Mohammad ^{1, 2, 3}

¹ NEC Corporation, Global Innovation Unit

²National Institute of Advanced Industrial Science and Technology (AIST), Japan

³Assiut University, Egypt

December 4, 2023



ICA 2023 Tutorial Updated on December 4, 2023

Outline

- ① Theoretical Session
- ② Development Environment
- ③ Development Example

Why Now?

- ① Industries are moving online.
- ② Automation: Factory floor → The back office.
- ③ Human-Human Negotiation is cumbersome, and inefficient.
- ④ Automated Negotiation opens new possibilities:
 - Too fast for people: Repeated smart contracts.
 - Too large for people: complete supply chains



Negotiation in SCM Business

- Human negotiations lead to an estimated 17-40% *value leakage* in some estimates ¹

¹KPMG report: <https://bit.ly/3kDRy6l>

²Forrester report: <https://bit.ly/3nwXEaY>

³UN/CEFACT Project website: <https://bit.ly/38LOsLX>

⁴Y. Mohammad et al. "Supply Chain Management World: A benchmark environment for situated negotiations". In: *Proceedings of the 22nd International Conference on Principles and Practice of Multi-Agent Systems*. 2019.

Negotiation in SCM Business

- Human negotiations lead to an estimated 17-40% *value leakage* in some estimates ¹
- A recent study suggests that at least 15 companies are working in *contracting support systems* ².

¹KPMG report: <https://bit.ly/3kDRy6l>

²Forrester report: <https://bit.ly/3nwXEaY>

³UN/CEFACT Project website: <https://bit.ly/38LOsLX>

⁴Mohammad et al., "Supply Chain Management World: A benchmark environment for situated negotiations".

Negotiation in SCM Business

- Human negotiations lead to an estimated 17-40% *value leakage* in some estimates ¹
- A recent study suggests that at least 15 companies are working in *contracting support systems* ².
- A recent UNECE UN/CEFACT proposal to standardize negotiation protocols for SCM and other applications ³

¹KPMG report: <https://bit.ly/3kDRy6I>

²Forrester report: <https://bit.ly/3nwXEaY>

³UN/CEFACT Project website: <https://bit.ly/38LOsLX>

⁴Mohammad et al., "Supply Chain Management World: A benchmark environment for situated negotiations".

Negotiation in SCM Business

- Human negotiations lead to an estimated 17-40% *value leakage* in some estimates ¹
- A recent study suggests that at least 15 companies are working in *contracting support systems* ².
- A recent UNECE UN/CEFACT proposal to standardize negotiation protocols for SCM and other applications ³
- More to come⁴.

CONTRACTROOM

pactum

¹KPMG report: <https://bit.ly/3kDRy6l>

²Forrester report: <https://bit.ly/3nwXEaY>

³UN/CEFACT Project website: <https://bit.ly/38LOsLX>

⁴Mohammad et al., "Supply Chain Management World: A benchmark environment for situated negotiations".

The Automated Negotiation Challenge

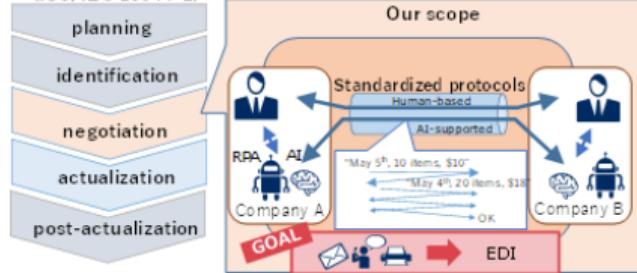
Why is it hard?

- Mechanism Design Problem:
 - Better than haggling?
- Negotiator Design Problem:
 - Generality × Effectiveness

Why is it interesting?

- Easy to state yet hard to solve.
- Multiple levels of abstraction and complexity.
- Several concrete open questions.
- Vibrant yet not saturated research space.

Five fundamental activities of a business transaction (ISO/IEC 15944-1)



attribution: UNECE eNegotiation Project



Automated Negotiating Agents Competition: 2010-

Tutorial Outline

① Theoretical Session (55min) Break (10min)

- The Negotiation Problem (35min)
- Automated Negotiating Agents Competition (ANAC) (5min)
- Supply Chain Management League 15min
- Anatomy of an Agent (10min)

② Development Environment (50min) Break (10min)

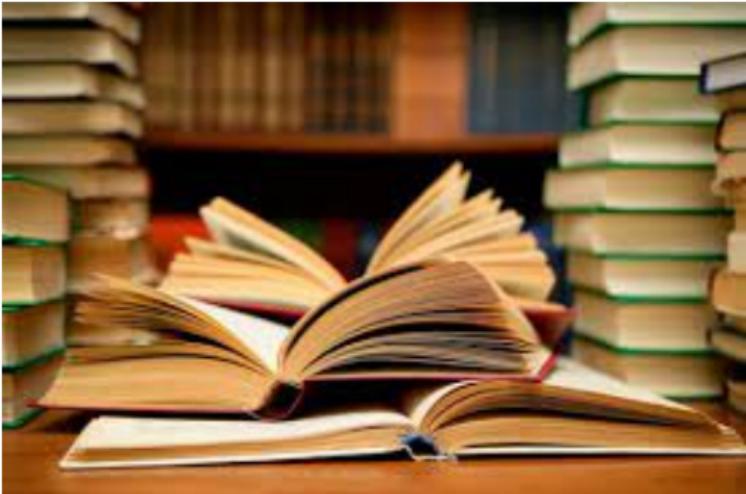
- Installing SCML (5min)
- Running a simulation (5min)
- Running Tournaments (5min)
- Visualization and Logging (15min)
- Troubleshooting and breathing time (10min)

③ Development Example (55min)

- SCML Builtin Strategies (20min)
- Example Agent Strategies (15min)
- An RL based trading strategy (20min)

④ Conclusions (5min)

Materials



- ① Tutorial Website: <http://yasserm.com/tutorial-ica2023/>
- ② Github Repository: <https://github.com/yasserfarouk/ica2023autoneg>
- ③ Handouts: <https://github.com/yasserfarouk/ica2023autoneg/raw/main/ica2023autoneg.pdf>
- ④ Negmas Documentation: <https://negmas.readthedocs.io>
- ⑤ SCML Documentation: <https://scml.readthedocs.io>
- ⑥ SCML Competition: <https://scml.cs.brown.edu>

Outline

① Theoretical Session

- The Negotiation Problem
- Automated Negotiating Agents Competition: ANAC
- Supply Chain Management League: SCML

② Development Environment

③ Development Example

Outline

① Theoretical Session

• The Negotiation Problem

- Negotiation as a Game
- Automated Negotiation Basics
- Visualizing a Negotiation

- Automated Negotiating Agents Competition: ANAC
- Supply Chain Management League: SCML

② Development Environment

③ Development Example

Types of Games

Complete Information Game

- All game parameters are **common knowledge** except players' policies.
- All information sets has a single node.

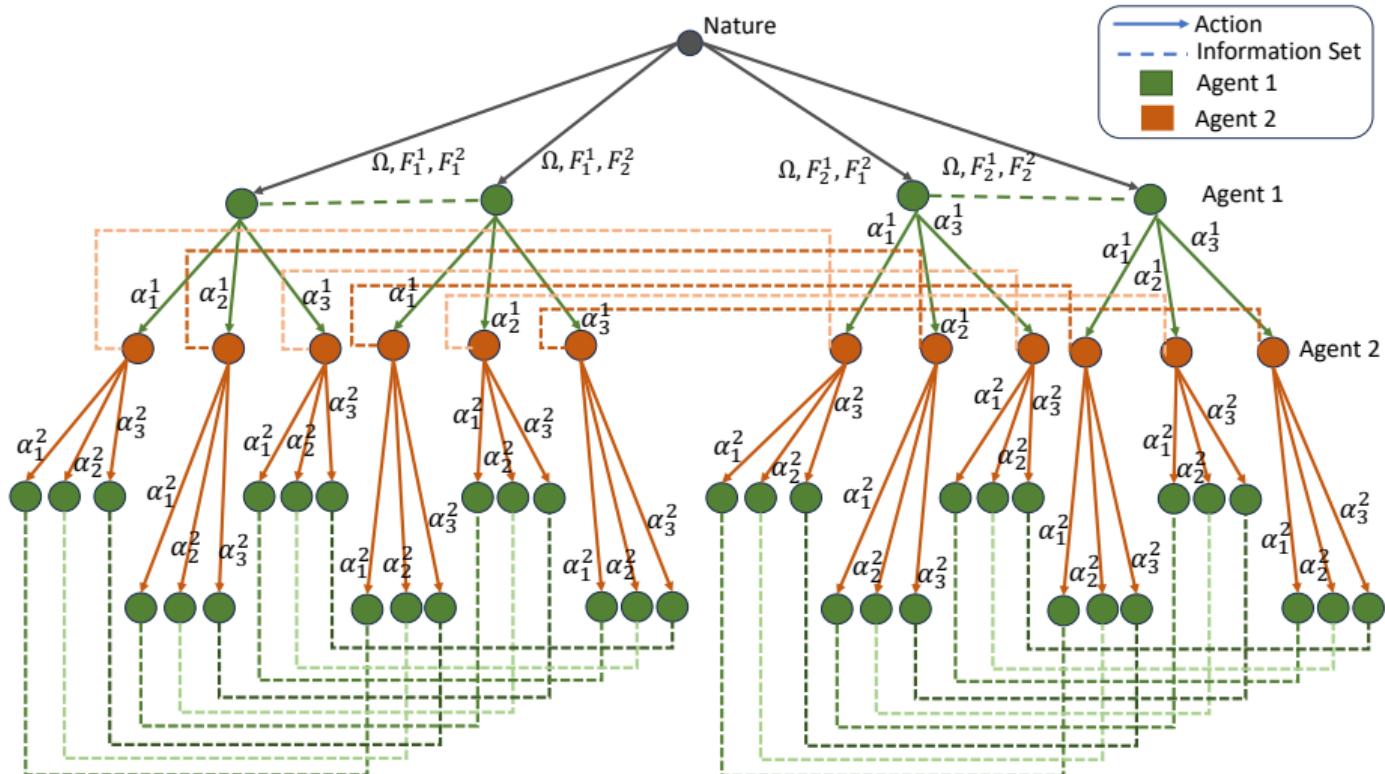
Imperfect Information Game

- Some partner infomration is unknown.
- Some information sets has multiple nodes.

Incomplete Information Game

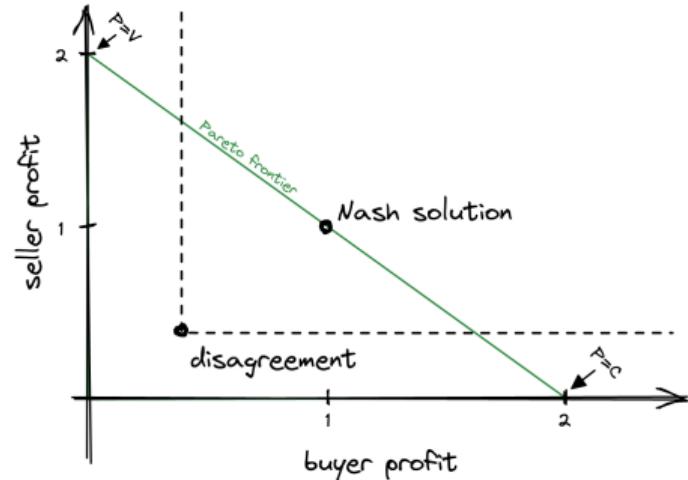
- The game is not fully specified (e.g. unknown opponent ufun)
- Can be converted to an Imperfect information game by adding **Nature** as a player.

Induced Game with Incomplete Information



A Simple Trading Problem

- A buyer values a good at V
- A seller can create the good at cost C
- If $V > C$, then there is surplus $V - C$ to be gained (**value creation**)
- Bargaining problem: how much should the buyer pay the seller for the good? (**value division**)
- We might also assume there is an outside option (e.g., eBay), if the negotiation breaks down (i.e., they do not reach an agreement):
 - The buyer (seller) can buy (sell) the good elsewhere for slightly less than V (more than C)



Sketch by Jackson de Campos

No-go Theorem

Desiderata:

- Efficient outcome.
- Individual rationality (IR).
- Incentive compatible (IC).
- Budget balance (BB).

Myerson-Satterthwaite Impossibility Result

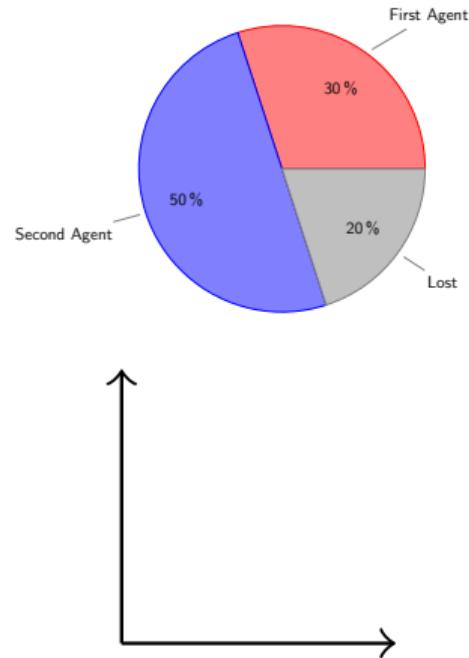
Theorem:⁵ No mechanism can achieve all four of these desiderata.

- A buyer values a good at V .
- A seller can create the good at cost C .
- V (C) is private information, known only to the buyer (seller).
- There is no IR, IC, and BB mechanism that results in agreement, for all $V > C$ values.

⁵Roger B Myerson and Mark A Satterthwaite. "Efficient mechanisms for bilateral trading". In: *Journal of economic theory* 29.2 (1983), pp. 265–281.

Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

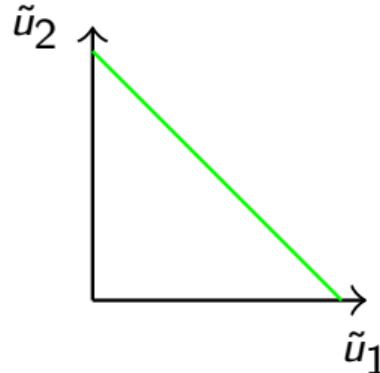
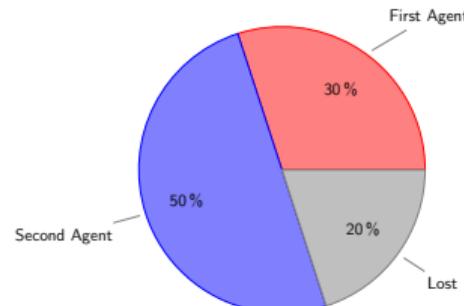


Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

- A feasible set of agreements F . A common example is to define F as all the outcomes for which the total utility received by negotiators is less than or equal to one:

$$F = \{(\omega_1, \omega_2) | \tilde{u}_2(\omega_2) + \tilde{u}_1(\omega_1) \leq 1\}.$$

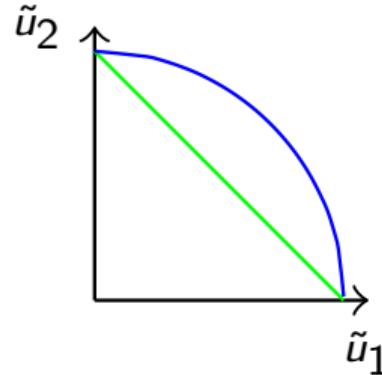
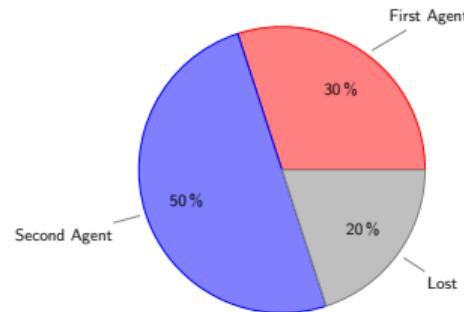


Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

- A feasible set of agreements F . A common example is to define F as all the outcomes for which the total utility received by negotiators is less than or equal to one:

$$F = \{(\omega_1, \omega_2) | \tilde{u}_2(\omega_2) + \tilde{u}_1(\omega_1) \leq 1\}.$$

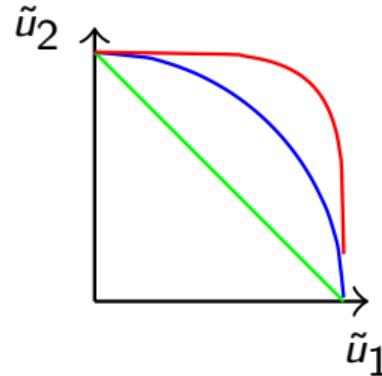
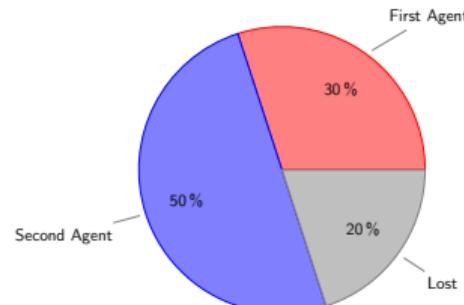


Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

- A feasible set of agreements F . A common example is to define F as all the outcomes for which the total utility received by negotiators is less than or equal to one:

$$F = \{(\omega_1, \omega_2) | \tilde{u}_2(\omega_2) + \tilde{u}_1(\omega_1) \leq 1\}.$$



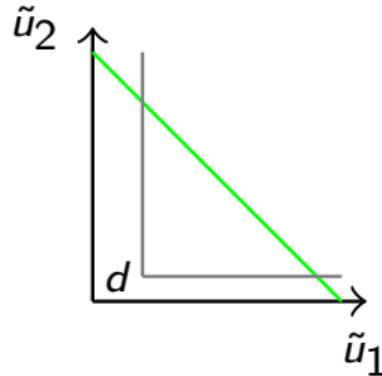
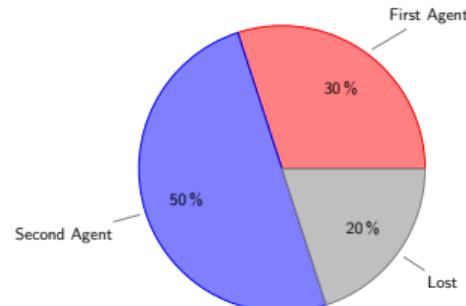
Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

- A feasible set of agreements F . A common example is to define F as all the outcomes for which the total utility received by negotiators is less than or equal to one:

$$F = \{(\omega_1, \omega_2) | \tilde{u}_2(\omega_2) + \tilde{u}_1(\omega_1) \leq 1\}.$$

- A disagreement point $d \equiv \tilde{u}_1(\phi) + \tilde{u}_2(\phi) \in \Re^2$ which is the utility value received by the two players in case of disagreement (reserved values).



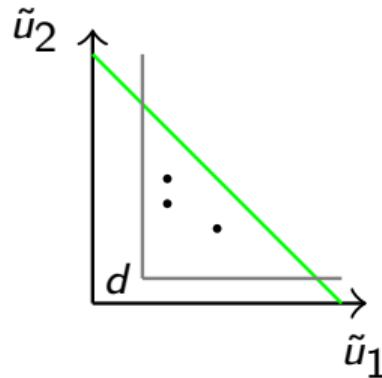
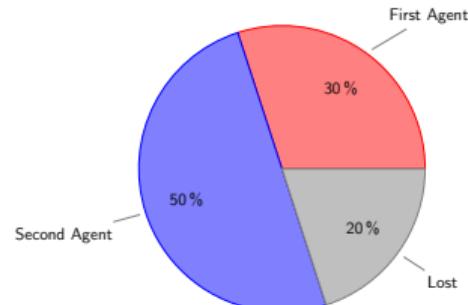
Nash Bargaining Game: Description

A single-step full-information bilateral negotiation with $\Omega = [0, 1]^2$ and two utility functions $(\tilde{u}_1, \tilde{u}_2)$ such that:

- A feasible set of agreements F . A common example is to define F as all the outcomes for which the total utility received by negotiators is less than or equal to one:

$$F = \{(\omega_1, \omega_2) | \tilde{u}_2(\omega_2) + \tilde{u}_1(\omega_1) \leq 1\}.$$

- A disagreement point $d \equiv \tilde{u}_1(\phi) + \tilde{u}_2(\phi) \in \Re^2$ which is the utility value received by the two players in case of disagreement (reserved values).



Other Bargaining Solutions

- Nash Bargaining Solution (1950): The point at which the product of surplus utility (above reservation value) of negotiators is maximized

$$\arg \max_{\omega_1, \omega_2} \prod_{i=1}^2 (u_i(\omega_i) - u_i(\phi))$$

- Kalai-Smorodinsky Bargaining Solution (1975): The Pareto outcome with equal ratios of achieved surplus utility and maximum feasible surplus utility

$$\arg \max_{\omega_1, \omega_2 \in F} (\omega_1 + \omega_2) \text{ s.t. } \left(\frac{u_1(\omega_1) - u_1(\phi)}{u_2(\omega_2) - u_2(\phi)} \right) = \left(\frac{\max_{v \in F} (u_1(v)) - u_1(\phi)}{\max_{v \in F} (u_2(v)) - u_2(\phi)} \right)$$

- Kalai Bargaining Solution (1977): The Pareto outcome maximizing the utility for the unfortunate player. Defining P as the Pareto front

$$\arg \max_{\omega_1, \omega_2 \in P} \min_{i \in \{1,2\}} (u_i(\omega_i) - u_i(\phi))$$

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:

Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.

Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:
 - Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.
 - Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$
- Agent's initial utility is the assigned part of the pie: $u_i^0 = \omega_i$.

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:
 - Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.
 - Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$
- Agent's initial utility is the assigned part of the pie: $u_i^0 = \omega_i$.
- Time pressure and utility information are common knowledge.

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:
 - Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.
 - Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$
- Agent's initial utility is the assigned part of the pie: $u_i^0 = \omega_i$.
- Time pressure and utility information are common knowledge.
- No externally imposed time-limit.

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:
 - Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.
 - Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$
- Agent's initial utility is the assigned part of the pie: $u_i^0 = \omega_i$.
- Time pressure and utility information are common knowledge.
- No externally imposed time-limit.
- Zero reservation value: $u_i^\tau(\phi) = 0 \forall \tau$.

Rubinstein's Bargaining Protocol: Description

The Game

- Two agents sharing a pie.
- Each agent is under a different time-pressure: $u_i^{t+\Delta}(\omega) < u_i^t(\omega)$. Examples of time-pressure:
 - Exponential $u_i^{t+\Delta}(\omega) = \delta_i^\Delta u_i^t(\omega)$.
 - Linear $u_i^{t+\Delta}(\omega) = u_i^t(\omega) - \Delta c_i$
- Agent's initial utility is the assigned part of the pie: $u_i^0 = \omega_i$.
- Time pressure and utility information are common knowledge.
- No externally imposed time-limit.
- Zero reservation value: $u_i^\tau(\phi) = 0 \forall \tau$.

Main Result

There is a unique *sub-game perfect equilibrium* that requires a single negotiation step in most cases.

Rubinstein's Bargaining Protocol: Equilibrium

Exponential Discounting

The negotiation ends in **one step** with the first agent proposing and the second agent accepting *for asymmetric cases*:

$$(\omega_1^*, \omega_2^*) = \left(\frac{1 - \delta_2}{1 - \delta_1 \delta_2}, \frac{\delta_2 (1 - \delta_1)}{1 - \delta_1 \delta_2} \right)$$

Rubinstein's Bargaining Protocol: Equilibrium

Exponential Discounting

The negotiation ends in **one step** with the first agent proposing and the second agent accepting *for asymmetric cases*:

$$(\omega_1^*, \omega_2^*) = \left(\frac{1 - \delta_2}{1 - \delta_1 \delta_2}, \frac{\delta_2 (1 - \delta_1)}{1 - \delta_1 \delta_2} \right)$$

Linear Discounting

The negotiation ends in **one step** with the first agent proposing and the second agent accepting:

$$(\omega_1^*, \omega_2^*) = \begin{cases} (c_2, 1 - c_2) & c_1 > c_2 \\ (x, 1 - x) \quad \forall x \in [c_1, 1] & c_1 = c_2 \\ (1, 0) & c_1 < c_2 \end{cases}$$

Negotiation with Complete Information

Hick's Paradox

Why do rational parties negotiate when they have full information?

Labor vs. Management

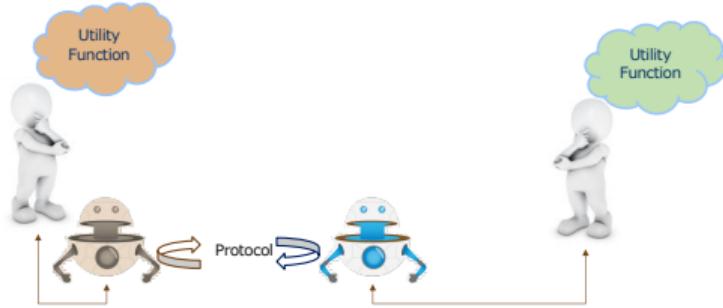
- A union negotiating (in rounds) with management about a wage raise.
- Both parties are perfectly rational and fully informed.
- Threat: the union *can* strike.

Theorem⁶ Subgame perfect equilibria exist in which there is a finite strike followed by agreement.

⁶Raquel Fernandez and Jacob Glazer. *Striking for a bargain between two completely informed agents*. Tech. rep. National Bureau of Economic Research, 1989.



Components of Negotiation



Negotiation Scenario Defines who is negotiating about what

Negotiation Protocol Defines how negotiation is to be conducted

- Alternating Offers Protocol
- Single Text Protocol
- ...

Negotiation Strategy Defines how agents behave during a negotiation

- Time-based strategies: boulware, conceder, ...
- Tit-for-tat variations
- ...

Notation

Negotiation Scenario \mathcal{S}

Agent Indices $\mathcal{A}^{\mathcal{S}}$: The set of n_A agents.

Domain $\mathcal{D}^{\mathcal{S}}$: A tuple $(K, \mathcal{F}^{\mathcal{S}})$ defining the situation:

Outcome Space K : A set of all possible agreements with $\phi \notin K$ representing disagreement. $K^+ \equiv K \cup \{\phi\}$ is called the Extended Outcome Space.

Preferences Tuple $\mathcal{F}^{\mathcal{S}}$: A tuple of n_A members that defines the preferences of each agent engaged in the negotiation.

Information Set Tuple $\mathcal{I}^{\mathcal{S}}$: A tuple of n_A information sets $(\{I_f^{\mathcal{S}}(\mathcal{F}^{\mathcal{S}}) \mid f \in \mathcal{A}^{\mathcal{S}}\})$ where $I_f^{\mathcal{S}}(\mathcal{F}^{\mathcal{S}})$ represents all the information available to agent x about the preferences $\mathcal{F}^{\mathcal{S}}$ of all agents including itself.

Protocol \mathcal{P} : The negotiation protocol.

Negotiation Protocol Execution: $\mathcal{P}(\mathcal{S})$

Protocol

Breaking Rule when the negotiation ends with a timeout.

Action Constraints what actions are valid.

Activation Rule which each agent is activated.

Evaluation Rule the next step:

Continue ▷ Goto breaking rule.

Agreement Succeed with ω .

Disagreement Fail with ϕ .

Agent (Strategy)

Selects an action from the **valid action set**.

Algorithm 1 Negotiation Mechanism Execution: $\mathcal{P}_Y(\lambda)$

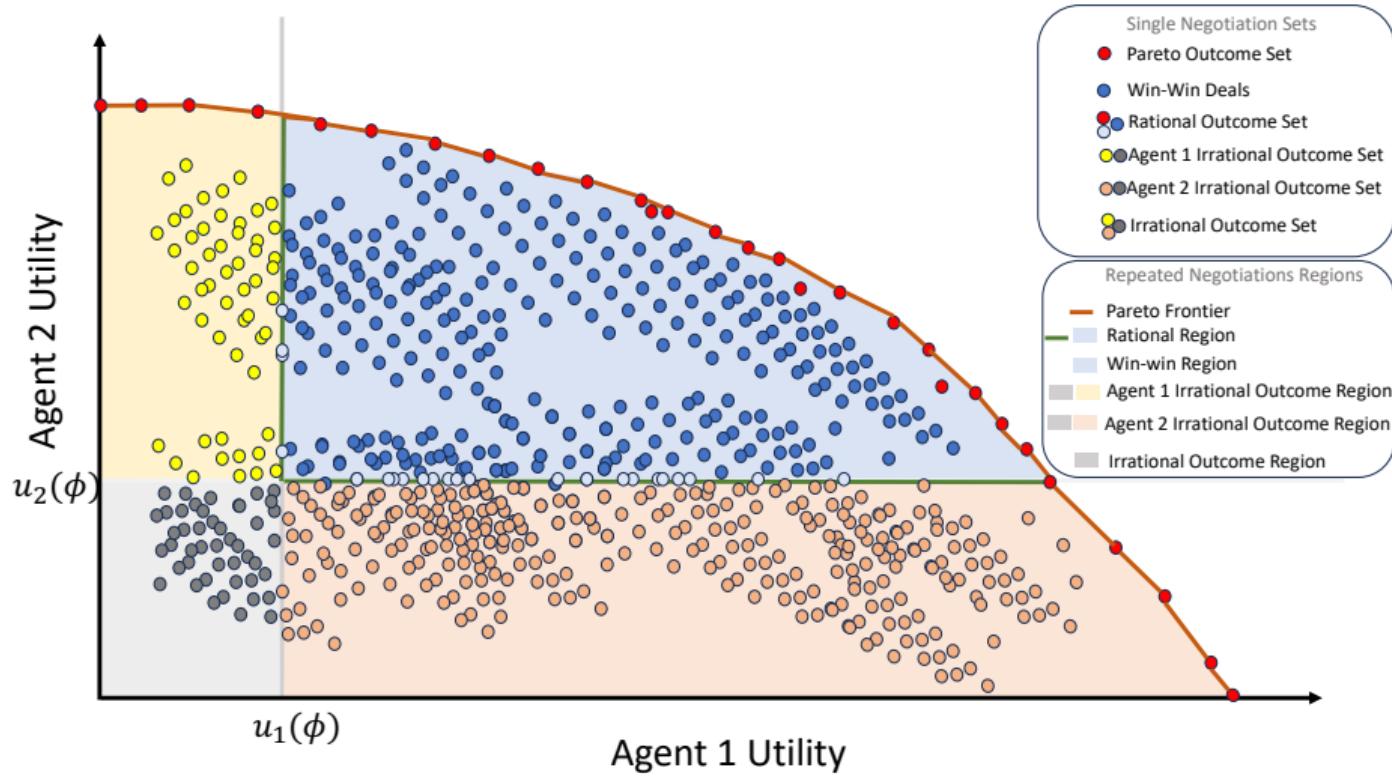
Input: Scenario $\lambda \equiv (\mathcal{N}^\lambda, (\Omega, \mathcal{F}^\lambda), \mathcal{I}^\lambda \equiv \{I_x^\lambda : x \in \mathcal{N}^\lambda\})$,
 Mechanism $\mathcal{P}_Y = (\mathcal{P}, Y) : \mathcal{P} = (\zeta, \chi, \sigma)$

Output: $\omega_* \in \Omega \cup \{\phi\}$

```

1:  $x \approx \pi_x \forall x \in \mathcal{N}^\lambda \leftarrow Y(\lambda)$                                 ▷ Assign strategies
2:  $i, t, t_0, \mathcal{T} \leftarrow 0, 0, \text{clock}(), \langle \rangle$                             ▷ Init step, time and trace
3: while  $\zeta(\mathcal{T}, i, t; \lambda) = 0$  do                                         ▷ check the breaking rule
4:    $\mathcal{N}_i^\lambda \leftarrow \chi(\mathcal{T}, i; \lambda)$                                          ▷ Next active agents
5:   for  $x \in \mathcal{N}_i^\lambda$  do                                              ▷ Activate agents
6:      $\mathbb{A}_x^v \leftarrow \gamma(\mathcal{T}, x; \lambda)$                                          ▷ Valid action set
7:      $\alpha_i^x \leftarrow \pi_x(\mathcal{T}, \mathbb{A}_x^v; I_x^\lambda)$                                ▷ Receive action
8:      $\mathcal{T} \leftarrow \mathcal{T} + \langle \alpha_i^x, x \rangle$                                      ▷ Update the trace
9:      $z \leftarrow \sigma(\mathcal{T}, i; \lambda)$                                          ▷ Evaluation rule
10:    if  $z = \phi$  then return  $\phi$  ▷ End with disagreement
11:    else if  $z \in \Omega$  then return  $\omega_i$                                          ▷ agreement
12:    end if                                         ▷ Otherwise continue to the next step
13:  end for
14:   $i, t \leftarrow i + 1, \text{clock}() - t_0$                                          ▷ Update step and time.
15: end while
16: return  $\phi$                                          ▷ Disagreement on breaking rule returning 1
  
```

Visualizing a negotiation



Outline

1 Theoretical Session

- The Negotiation Problem
- **Automated Negotiating Agents Competition: ANAC**
- Supply Chain Management League: SCML

2 Development Environment

3 Development Example

Automated Negotiating Agents Competition (ANAC)

- An international competition dedicated to automated negotiation.
- Started in conjunction with AAMAS in 2010.
- Running every year with AAMAS or IJCAI (14 iterations so far).
- Next ANAC will be in conjunction with AAMAS 2024 (not official yet).



Outline

1 Theoretical Session

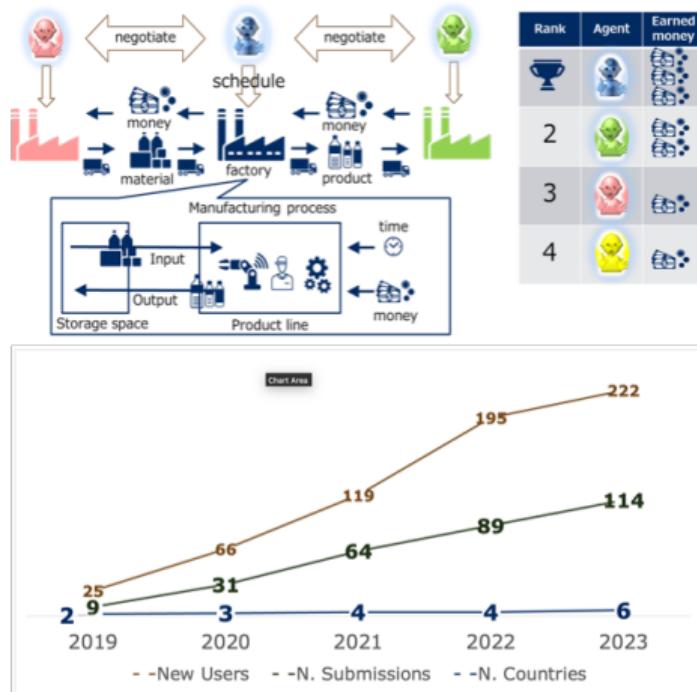
- The Negotiation Problem
- Automated Negotiating Agents Competition: ANAC
- Supply Chain Management League: SCML
 - SCML-OneShot
 - Simulation Steps

2 Development Environment

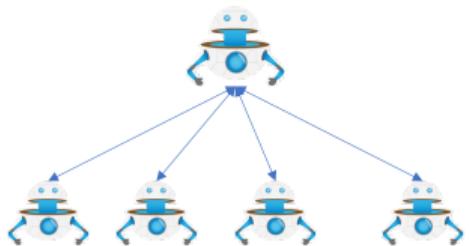
3 Development Example

Supply Chain Management League (SCML)

- Running as part of ANAC since 2019.
- **Scenario** A market in which all trade is done through **concurrent closed bilateral negotiations**:
 - Represents several applications in the real world.
- Agent preferences are endogenous to the simulation.
- **Goal** Maximize your profit.
- **Main Challenges:**
 - Situated Negotiations.
 - Sequential Negotiation (learning between negotiations).
 - Concurrent Negotiation (outside-options).
 - Trust management.



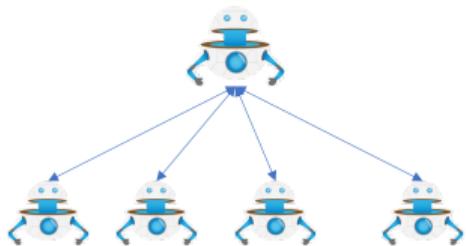
Concurrent Negotiation



Concurrent Negotiation

Generality

- Specific scenario (buyer-seller).
- General domain



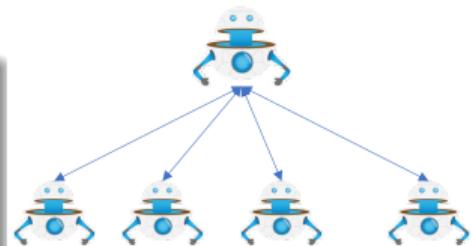
Concurrent Negotiation

Generality

- Specific scenario (buyer-seller).
- General domain

Decommitment

- Symmetric de-commitment.
- Asymmetric de-commitment.
- No de-commitment.



Concurrent Negotiation

Generality

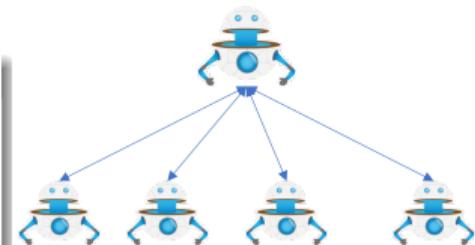
- Specific scenario (buyer-seller).
- General domain

Decommitment

- Symmetric de-commitment.
- Asymmetric de-commitment.
- No de-commitment.

Timing

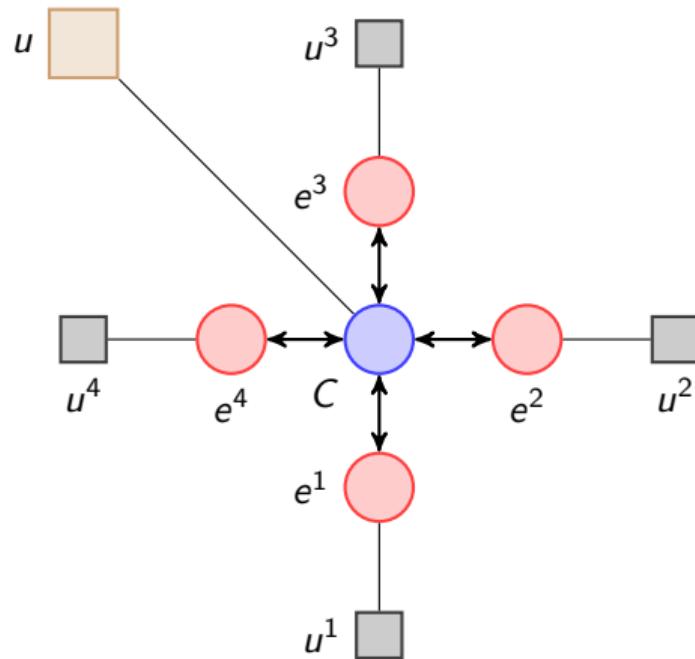
- Synchronous.
- Any-time.



Concurrent Negotiations

These negotiations are dependent, as evidenced by the global ufun u , which depends on a global outcome: i.e., the outcomes of all negotiations.

The problem of how to negotiate in this setting is reminiscent of how to bid in simultaneous auctions, when bidders' utilities are combinatorial.⁷

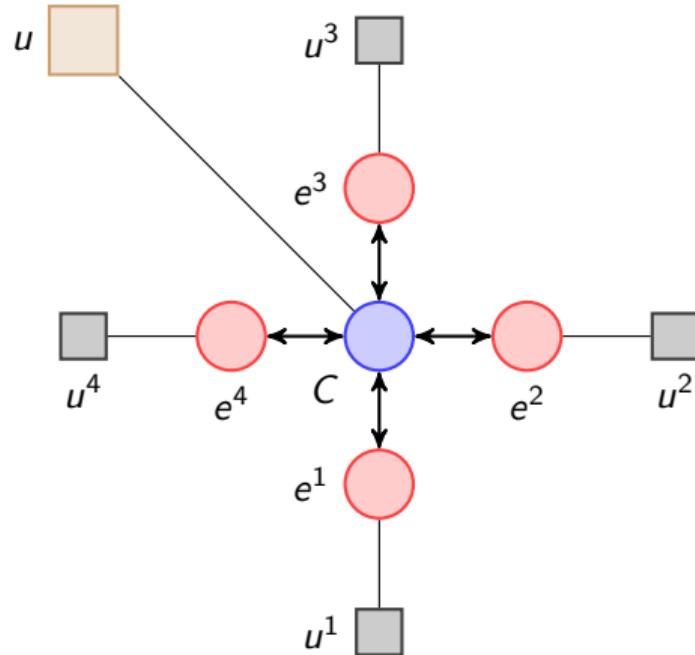


⁷ Michael P. Wellman, Eric Sodomka, and Amy Greenwald. "Self-confirming price-prediction strategies for simultaneous one-shot auctions". In: *Games and Economic Behavior* 201 (2017), pp. 339–372.

Concurrent Negotiations

In ANAC SCML, an agent negotiates with multiple agents simultaneously. Moreover, the agent's utility depends on the outcomes of all the negotiations.

An agent represents a factory, whose production depends on the inputs it acquires from all its trading partners, and so, in turn, does its sales.



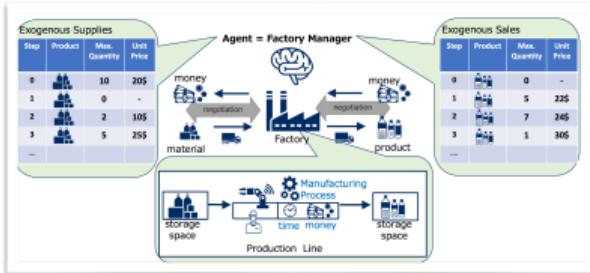
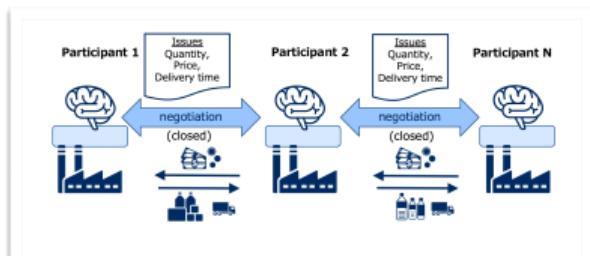
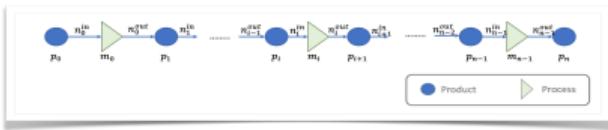
SCML World

Challenge

- Negotiation game with imperfect information
- Concurrent negotiations.
- Repeated negotiations → OneShot.
- Sequential negotiations → Standard.

Information

- **Website** <https://scml.cs.brown.edu/>
- **Code** <https://www.github.com/yasserfarouk/scml>



SCML Competition

Competition Details

- Runs as part of ANAC IJCAI.
- You control one or more factories.
 - **Oneshot track** → one factory (predefined ufun).
 - **Standard track** → one factory (you define your own ufun).
 - **Collusion track** → multiple factories (3).

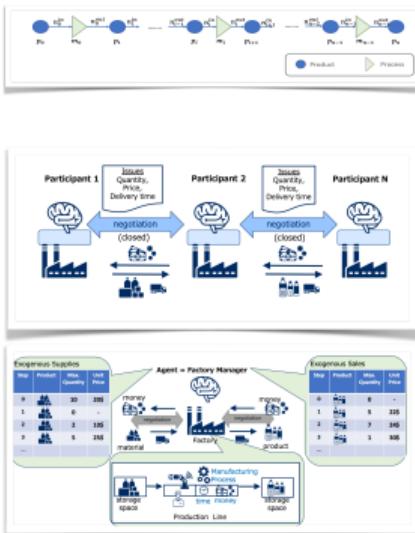
SCML Competition

Competition Details

- Runs as part of ANAC IJCAI.
- You control one or more factories.
 - Oneshot track** → one factory (predefined ufun).
 - Standard track** → one factory (you define your own ufun).
 - Collusion track** → multiple factories (3).

Flavors

- Online competition at <https://scml.cs.brown.edu>
- Official competition as part of ANAC.



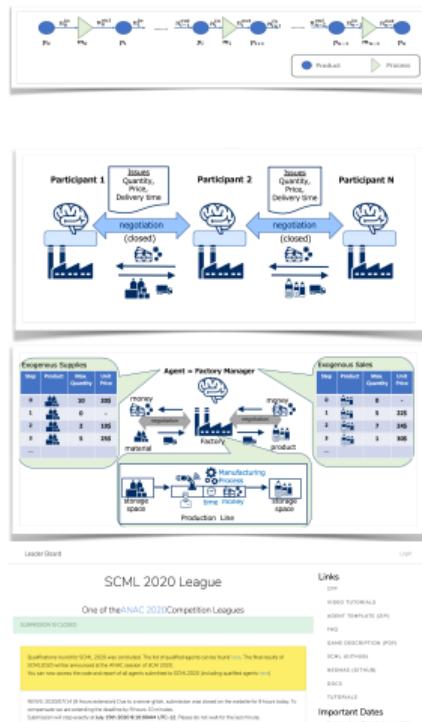
SCML Competition

Competition Details

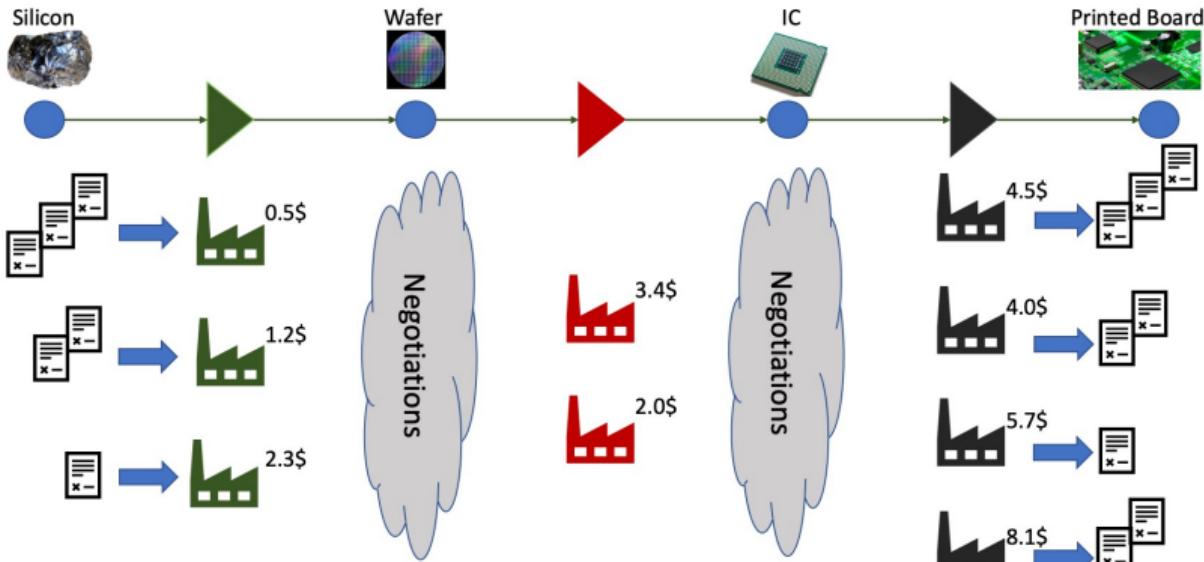
- Runs as part of ANAC IJCAI.
- You control one or more factories.
 - Oneshot track** → one factory (predefined ufun).
 - Standard track** → one factory (you define your own ufun).
 - Collusion track** → multiple factories (3).

Flavors

- Online competition at <https://scml.cs.brown.edu>
- Official competition as part of ANAC.



Example Configuration



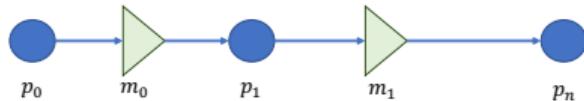
An example of an SCM world showing four products (circles), three processes (triangles) and few factories.

Each process consumes one item of its input and generates one output of its output in one day.

Each factory requires a different cost to run its process (shown in its top right).

Factories in the first level have exogenous contracts to buy raw material (silicon) and factories at the last level have exogenous contracts to sell the final product (printed boards). These contracts drive the market.

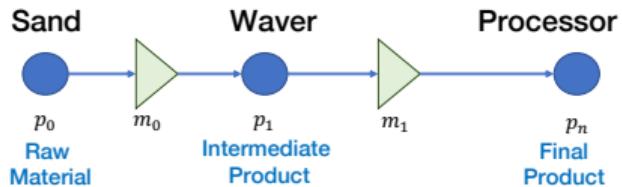
Overview



- A **production-graph** defines what can be produced and how.



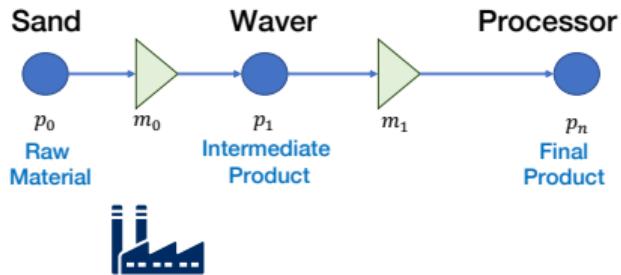
Overview



- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.

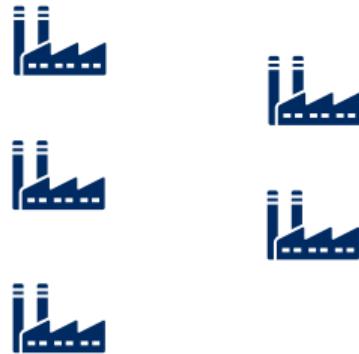
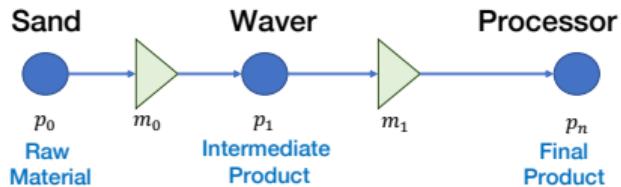


Overview



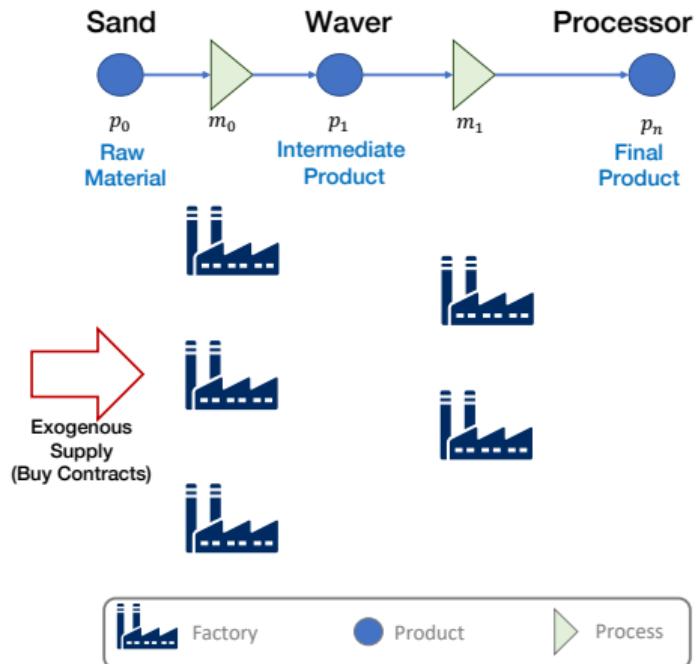
- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.
- Factories can run **manufacturing processes** converting input products into output products on their **production lines**.

Overview



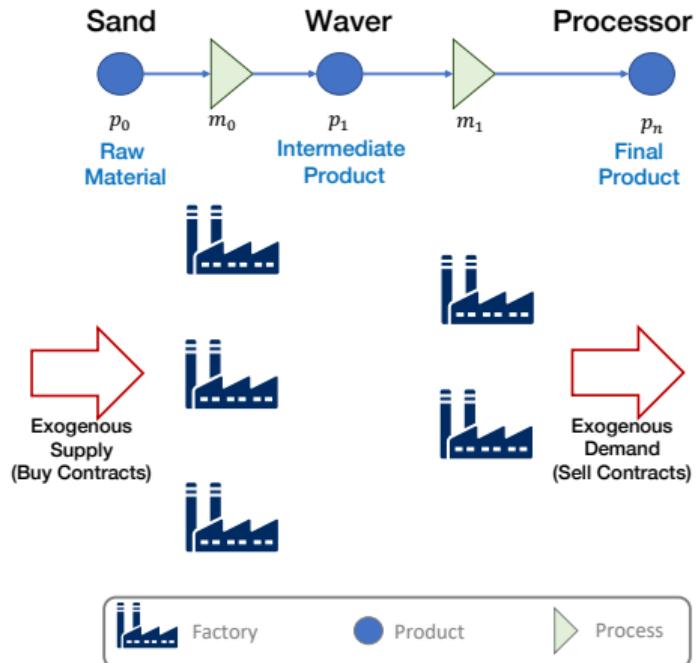
- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.
- Factories can run **manufacturing processes** converting input products into output products on their **production lines**.
- We have two layers of factories/agents.

Overview



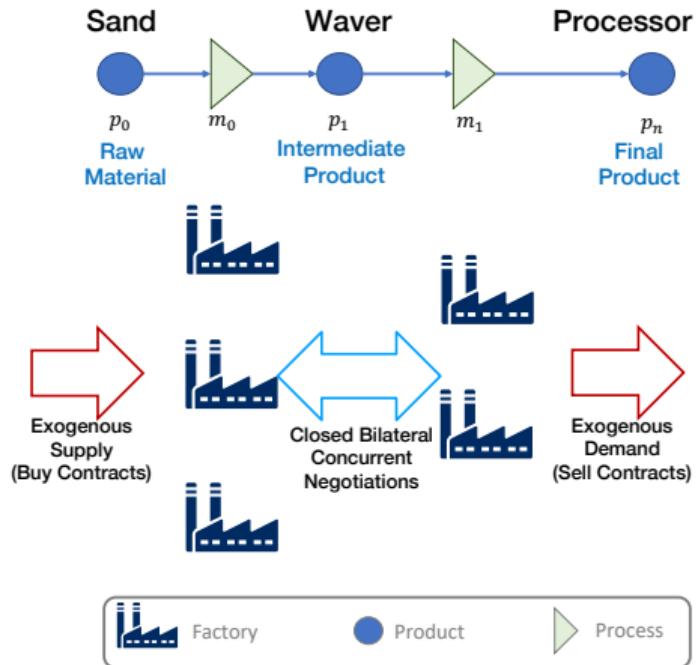
- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.
- Factories can run **manufacturing processes** converting input products into output products on their **production lines**.
- We have two layers of factories/agents.
- L_0 factories/agents receive exogenous supplies of **raw material**.

Overview



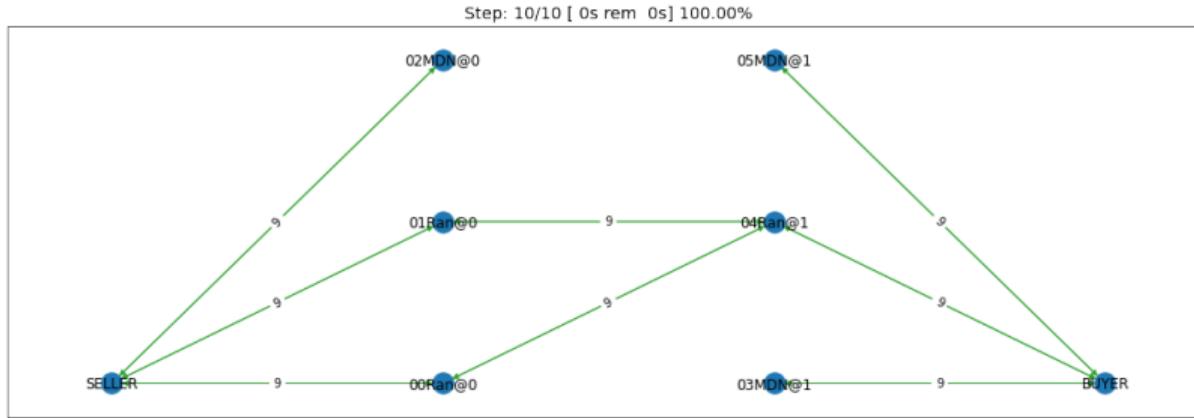
- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.
- Factories can run **manufacturing processes** converting input products into output products on their **production lines**.
- We have two layers of factories/agents.
- L_0 factories/agents receive exogenous supplies of **raw material**.
- L_1 factories/agents receive exogenous sales of **final product**.

Overview



- A **production-graph** defines what can be produced and how.
- We have 3 products, 2 processes.
- Factories can run **manufacturing processes** converting input products into output products on their **production lines**.
- We have two layers of factories/agents.
- L_0 factories/agents receive exogenous supplies of **raw material**.
- L_1 factories/agents receive exogenous sales of **final product**.
- L_0 negotiate with L_1 agents to exchange **intermediate product**

SCML-OneShot Track



Main Idea

- Agents arranged in two production levels (3 products, 2 processes)
- Every day you get a **fresh set** of exogenous contracts.
- All products perish in one day (no inventory accumulation).

Utility Function

General Form

Utility = Profit = Sales - Supply cost - Production cost - Disposal cost - Delivery Penalty

Sales unit price \times quantity \forall feasible sales.

Supply cost unit price \times quantity \forall supplies.

Production cost unit production cost \times quantity produced.

Disposal cost unit disposal cost \times quantity bought but not produced.

Shortfall penalty unit shortfall penalty \times infeasible sales.

Information about self

Static Information

- Number of production lines.
- Production cost.
- Mean and variance of disposal cost and shortfall penalty.
- Input/output product, consumers/suppliers, n. input/output negotiations.

Dynamic Information

- current input/output negotiation issues.
- current input/output exogenous contracts (quantity, unit price).
- current disposal cost and shortfall penalty.
- Current balance (money in wallet).

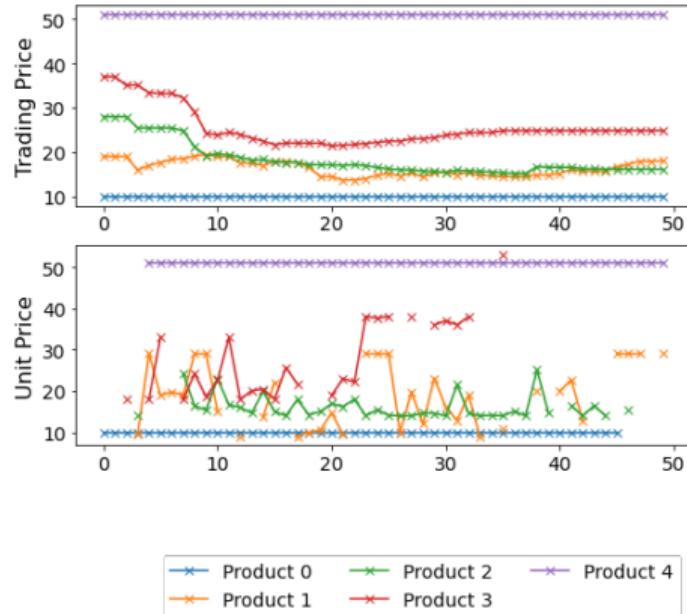
Market Information

Trading prices

- Trading prices represent a weighted running average of different product prices.
- Available to the agent through the AWI in all tracks.

Exogenous Contract Summary

- The total quantity and average prices of all exogenous contracts are now available through the AWI.
 - Exogenous contracts for individual agents are private information.



Product 0 Product 2 Product 4
Product 1 Product 3

Other Agents' Information

Financial Reports

For each agent, a financial report is published every m days (e.g. 5) with the following information:

- Current balance (money in wallet).
- breach probability (fraction of sale contracts not satisfied).
- breach level (average fraction of sales not satisfied).

Simulation Steps

Once



Initialize all agents
• init()

Simulation Steps

Once



Initialize all agents
• init()



Update trading prices

Simulation Steps

Once



Simulation Steps

Once



Every

Simulation Steps

Once



Simulation Steps

Once



Simulation Steps

Once



Every
Day

Simulation Steps

Once



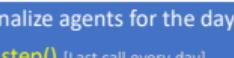
Every Day



Create exogenous contracts and sample agent's disposal cost, and shortfall penalty

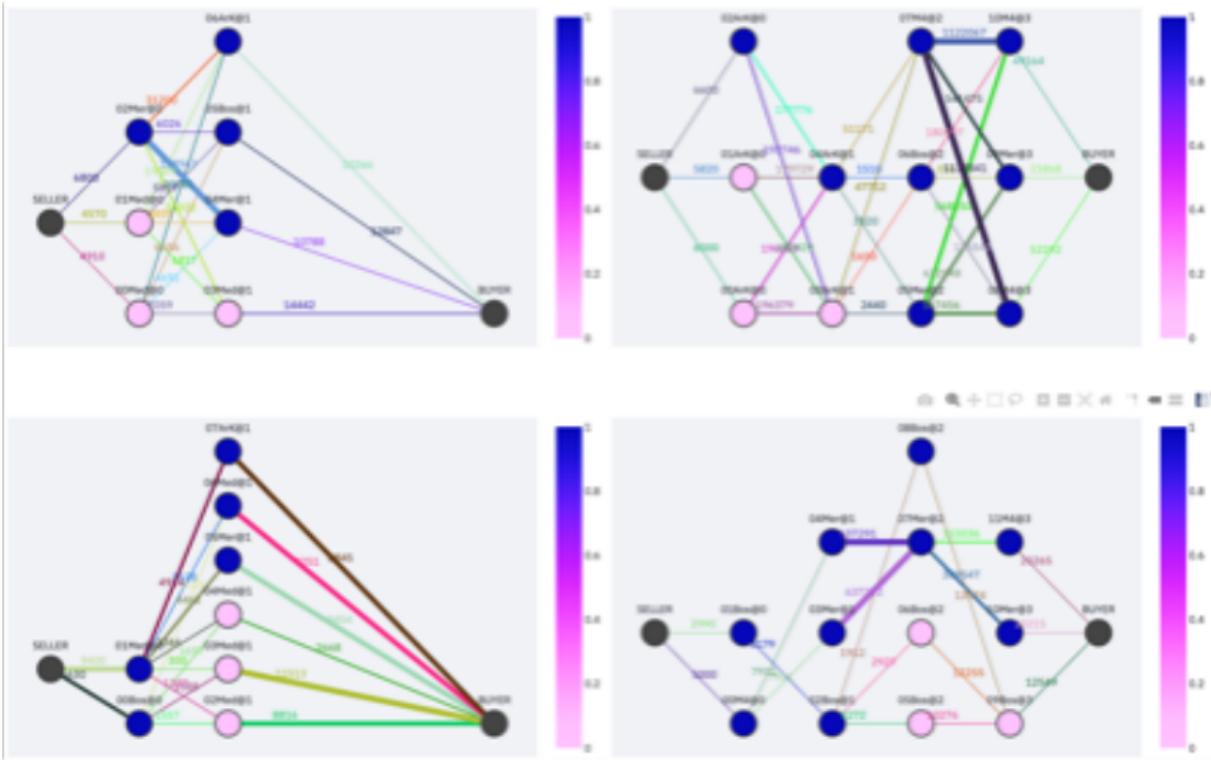


Calculate profits for all agents by simulating contract execution.



Publish Financial Reports

SCML Environment



Outline

1 Theoretical Session

2 Development Environment

- Installing SCML
- Running a simulation
- Running a tournament
- Visualization and logging
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

3 Development Example

Outline

1 Theoretical Session

2 Development Environment

- **Installing SCML**
- Running a simulation
- Running a tournament
- Visualization and logging
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

3 Development Example

Installing SCML

Pre-Installation

```
> python3 -m venv .venv  
> source .venv/bin/activate  
> .venv\\Scripts\\activate
```

Basic Installation

```
> pip install scml
```

Optional Installations

```
> pip install scml-agents  
> pip install scml-vis
```

Outline

1 Theoretical Session

2 Development Environment

- Installing SCML
- **Running a simulation**
- Running a tournament
- Visualization and logging
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

3 Development Example

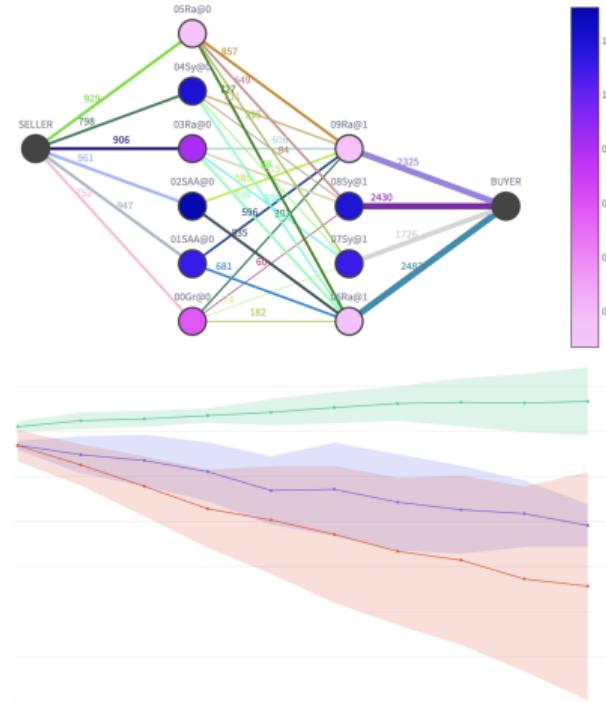
Running a single simulation (oneshot)

Command Line Tool

```
> scml run2024 —oneshot
```

Python

```
SCML2024OneShotWorld(  
**SCML2024OneShotWorld.generate(  
agent_types=DefaultAgentsOneShot2024  
)).run()
```



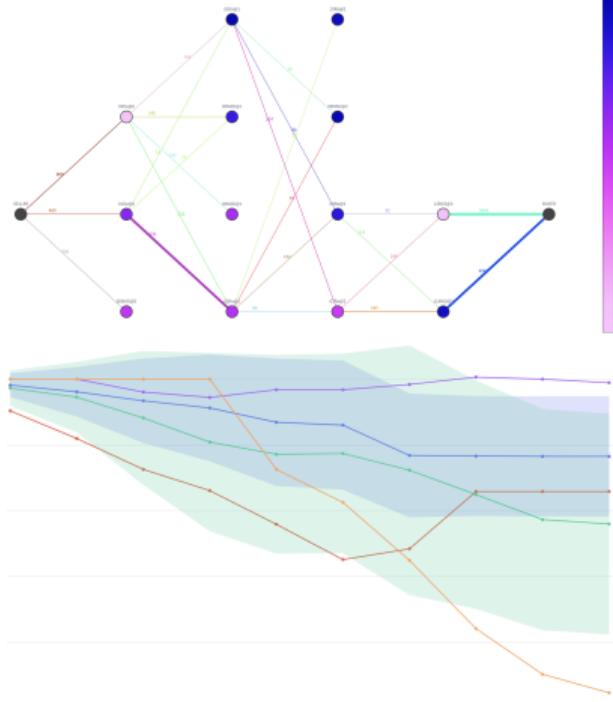
Running a single simulation (std)

Command Line Tool

```
> scml run2024
```

Python

```
SCML2024StdWorld(  
    **SCML2024StdWorld.generate(  
        agent_types=DefaultAgents2024  
    )).run()
```



Outline

1 Theoretical Session

2 Development Environment

- Installing SCML
- Running a simulation
- Running a tournament**
- Visualization and logging
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

3 Development Example

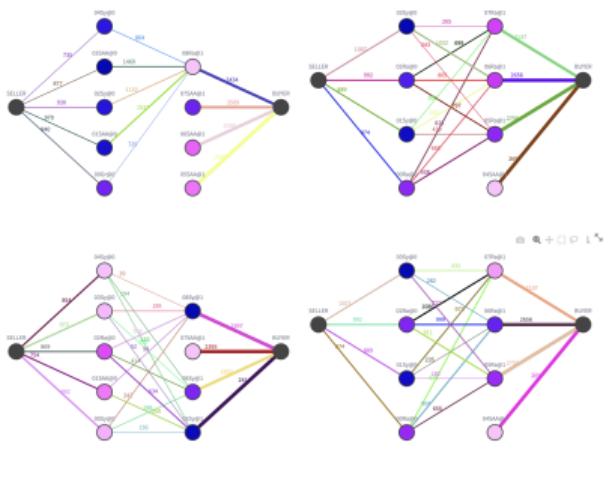
Running a complete tournament (oneshot)

Command Line Tool

```
> scml tournament2024 --type=oneshot
```

Python

```
anac2024_oneshot(
    competitors=[
        RandomOneShotAgent,
        SyncRandomOneShotAgent,
        GreedyOneShotAgent
])
```



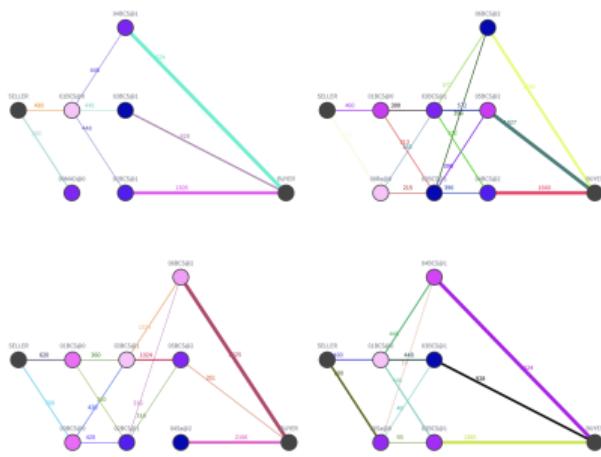
Running a complete tournament (std)

Command Line Tool

```
> scml tournament2024 --type=std
```

Python

```
anac2024_std(
    competitors=[
        RandomOneShotAgent,
        SyncRandomOneShotAgent,
        GreedyOneShotAgent
])
```



Outline

1 Theoretical Session

2 Development Environment

- Installing SCML
- Running a simulation
- Running a tournament
- Visualization and logging**
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

3 Development Example

Visualization CLI

Command Line Tool

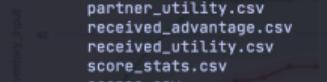
```
> scmlv show
```

```

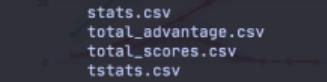
20231129H194614733893Ps7-stage- 2.2 K 0000S0_0.n_A-RCFR.0008a55a
20231129H194614733893Ps7-stage- 2.2 K 0001S0_1.H_A-R4oJ.016ed8bf
20231129H194846359611Uby-stage-08 1 K 0002S0_2.B_A-RUup.0242aa54
20231129H19495076256MB4-stage- 2.2 K 0003S0_3.5_A-Rw8m.03165be0
20231129H195118577752sfJ-stage- 2.7 K 0004S1_0.q_A-RNHn.0499dc63
20231129H195216244186qsC-stage- 2.7 K 0005S1_1.p_A-R3ND.05a5aabb
20231129H19485262934603k-stage-08 1 K 0006S1_2.A_A-RnFK.068292d7
20231130H080023219868Phm-stage- 960 B 0007S1_3.l_A-RSG8.077b18fc
agent_stats.csv
agg_stats.csv
assigned_configs.json
assigned_configs.pickle
attempts
base_configs.json
configs
kstats.csv
params.json
partner_advantage.csv
partner_utility.csv
received_advantage.csv
received_utility.csv
score_stats.csv
scores.csv
stats.csv
total_advantage.csv
total_scores.csv
tstats.csv
type_stats.csv

```

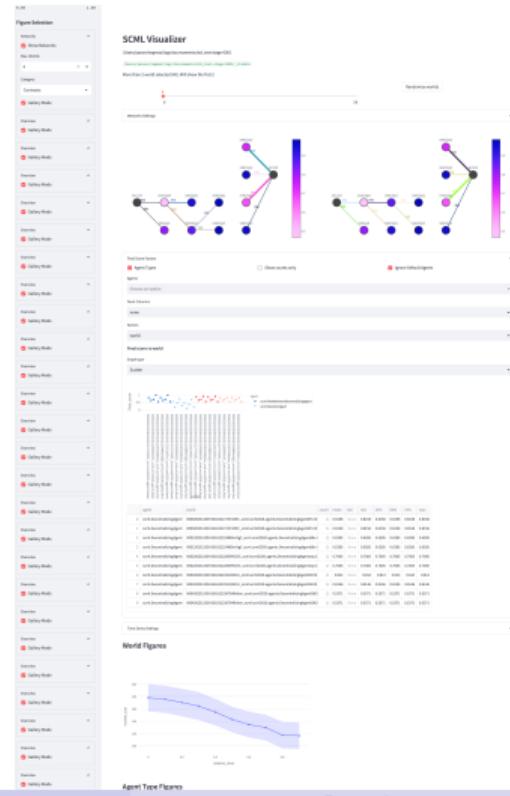
Inventory Input



Final Score Factors



Variables



Outline

1 Theoretical Session

2 Development Environment

- Installing SCML
- Running a simulation
- Running a tournament
- Visualization and logging
- **Troubleshooting and breathing time**
- NegMAS: The Underlying Platform

3 Development Example

Troubleshooting and breathing time



Outline

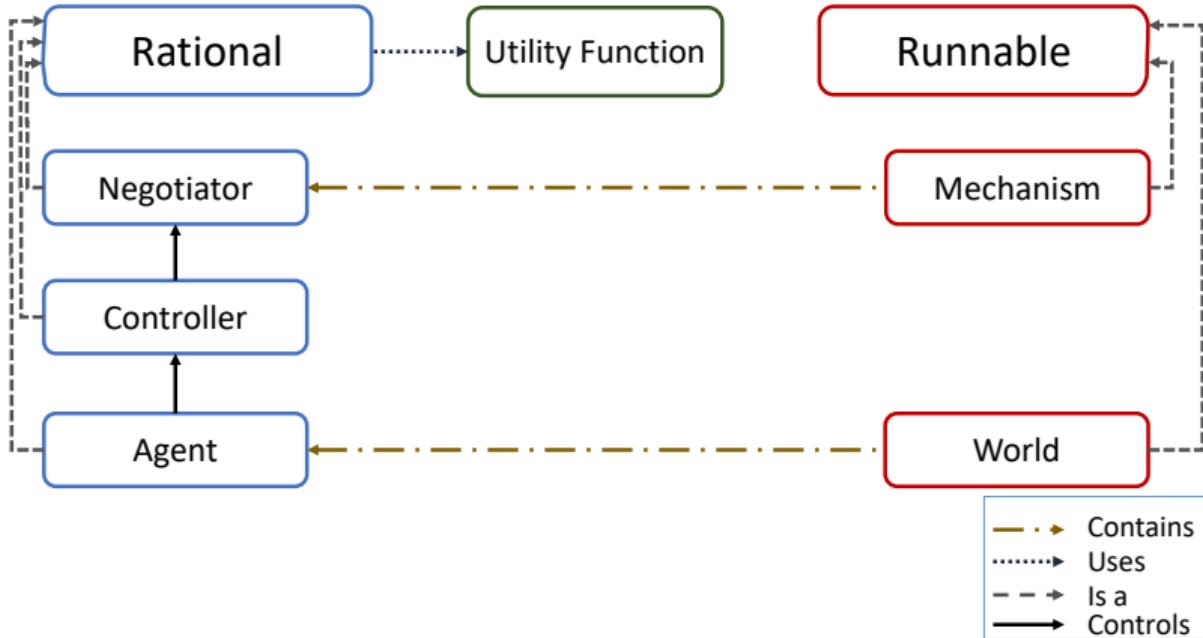
1 Theoretical Session

2 Development Environment

- Installing SCML
- Running a simulation
- Running a tournament
- Visualization and logging
- Troubleshooting and breathing time
- NegMAS: The Underlying Platform

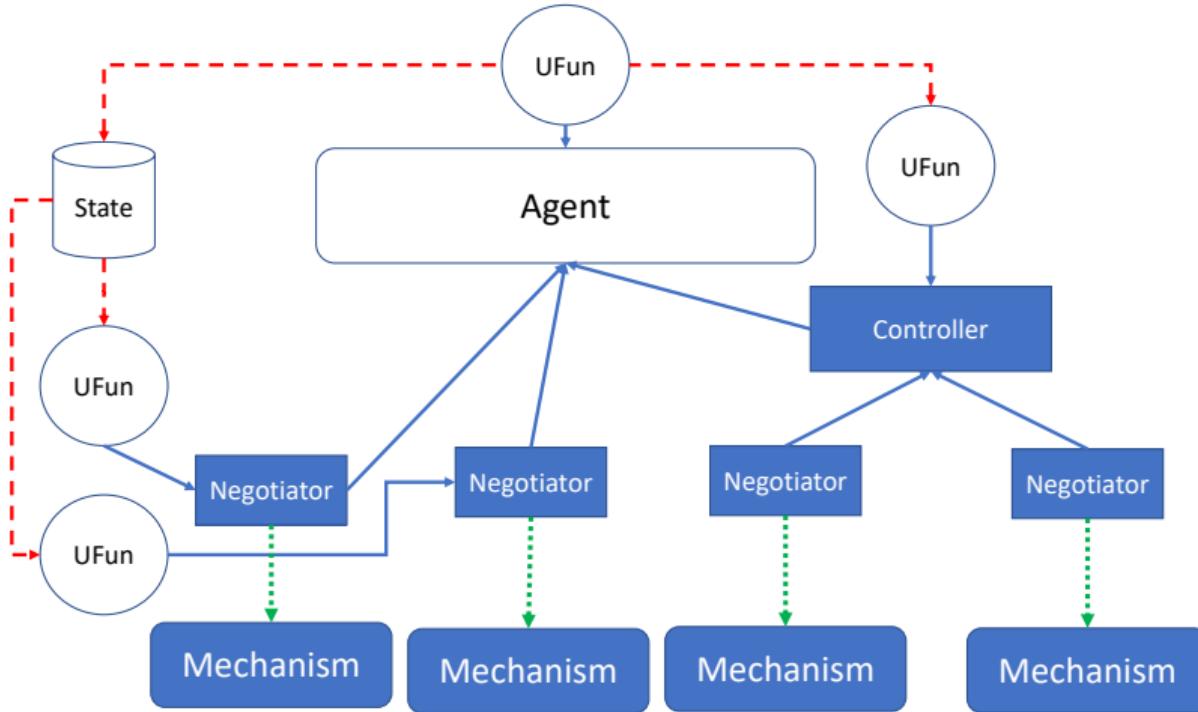
3 Development Example

NegMAS⁷ in two slides

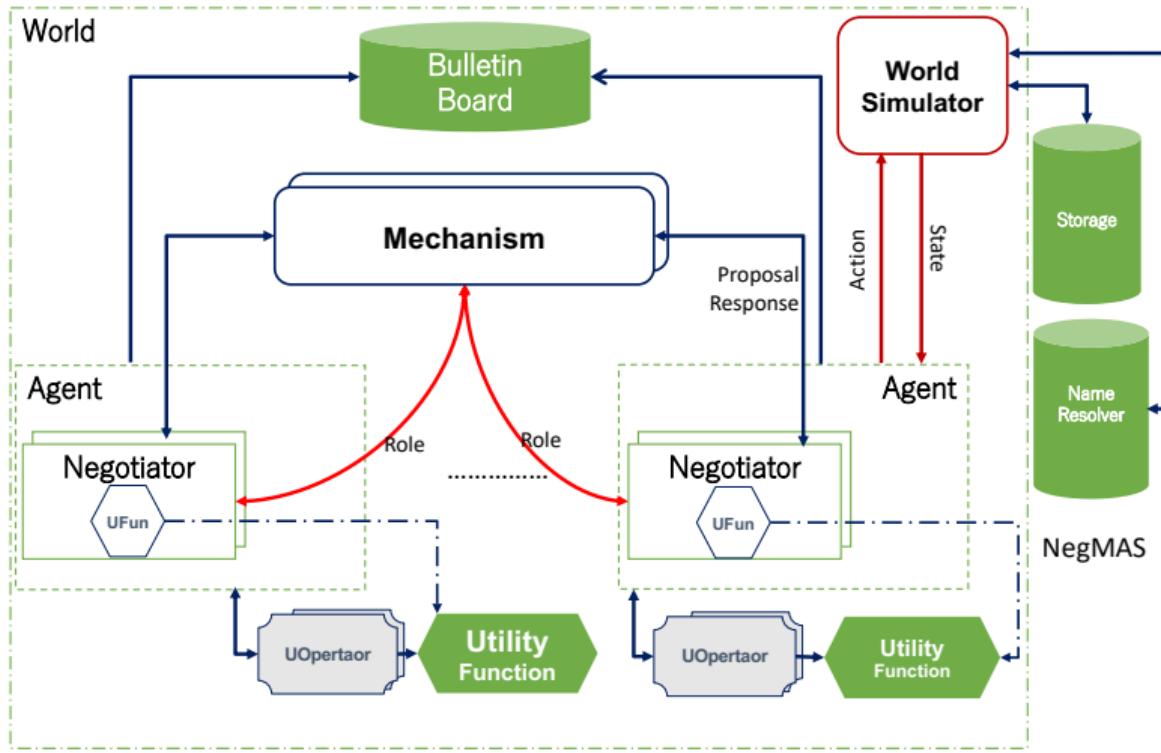


⁷<https://www.github.com/yasserfarouk/negmas>

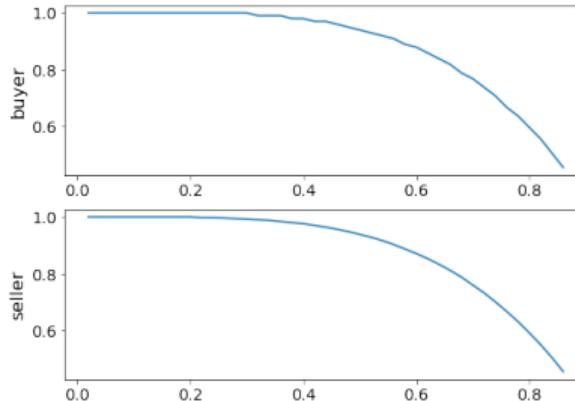
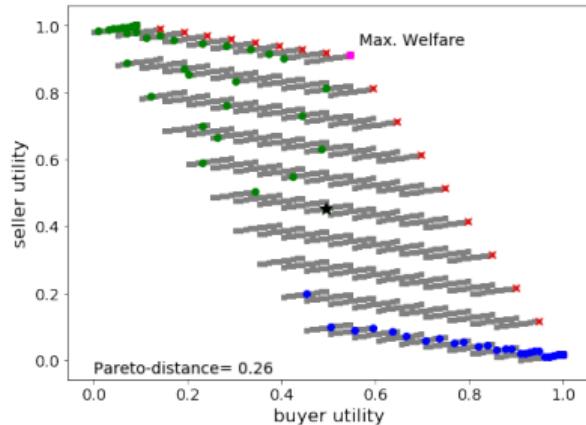
NegMAS in two slides



NegMAS in two slides (... OK 3)



NegMAS in two slides (... really!!!)



- An Example negotiation.
- Can you spot a problem?

Outline

- ① Theoretical Session
- ② Development Environment
- ③ Development Example
 - Anatomy of An Agent
 - SCML Builtin Strategies
 - Example Agent Strategies

Outline

- 1 Theoretical Session
- 2 Development Environment
- 3 Development Example
 - Anatomy of An Agent
 - SCML Builtin Strategies
 - Example Agent Strategies

SCML-OneShot Agent: Core

```
class OneShotSyncAgent(SAOSyncController, OneShotAgent, ABC):
    """
    An agent that automatically accumulate offers from opponents and allows
    you to control all negotiations centrally in the `counter_all` method.

    Args:
    +--- 4 lines: owner: The adapter owning the agent. You do not need to directly deal...
    """

    def __init__(self, *args, **kwargs):
        +--- 2 lines: kwargs["global_ufun"] = True.....
        |
        @abstractmethod
        def counter_all(
            self, offers: dict[str, Outcome], states: dict[str, SAOState]
        ) -> dict[str, SAOResponse]:
        +--- 20 lines: """Calculate a response to all offers from all negotiators.....
        |
        @abstractmethod
        def first_proposals(self) -> dict[str, Outcome]:
        +--- 10 lines: """.....
```

SCML-OneShot Agent: Optional Timing Callbacks

```
def init(self):
--- 6 lines: """"....."
|
def before_step(self):
--- 8 lines: """"....."
|
def step(self):
--- 8 lines: """"....."
|
```

SCML-OneShot Agent: Optional Events

```
def on_negotiation_failure(  
--- 5 lines: self,.....  
) -> None:  
--- 12 lines: """.....  
|  
def on_negotiation_success(  
|   self, contract: Contract, mechanism: NegotiatorMechanismInterface  
) -> None:  
--- 8 lines: """.....  
|
```

Outline

- 1 Theoretical Session
- 2 Development Environment
- 3 Development Example
 - Anatomy of An Agent
 - SCML Builtin Strategies
 - Example Agent Strategies

Greedy Random Agent



- Link to builtin agent examples (OneShot)
- Link to builtin agent examples (Std)

Outline

- 1 Theoretical Session
- 2 Development Environment
- 3 Development Example
 - Anatomy of An Agent
 - SCML Builtin Strategies
 - Example Agent Strategies

Example Winner Agents



- Link to 2023 winner agent descriptions
- Link to source code and description of every agent