15CSE302 Database Management Systems Lecture 18 **Second Normal Form**

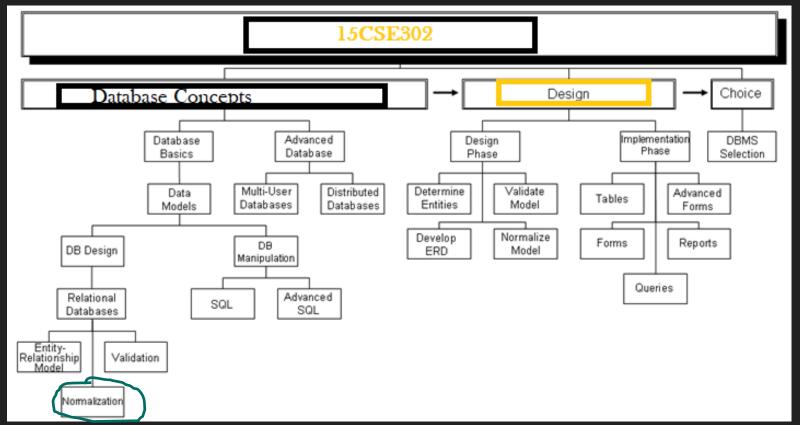
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Syllabus



Brief Recap of Previous Lecture

- Functional Dependency
- Armstrong Axioms –Rules for Functional Dependency



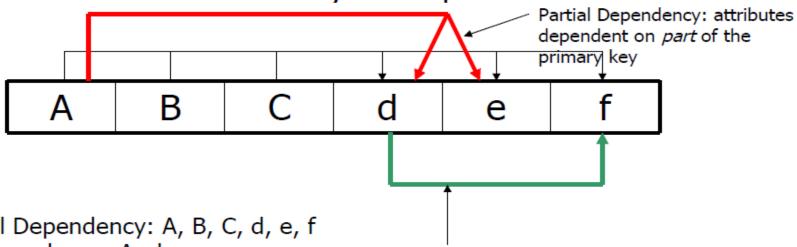
Today we'll discuss

Partial Dependency

Second Normal Form

Dependency Diagram

First Normal Form: identify **all** dependencies



Functional Dependency: A, B, C, d, e, f Partial Dependency: A, d, e Transitive Dependency: d, f

Transitive Dependency (non-key attribute depends on another non-key attribute)

Full Dependency

- \blacksquare R(S,T,U,V)
- \blacksquare S \rightarrow T T \rightarrow U U \rightarrow V V \rightarrow S
- S,T,U ,V are candidate keys



Partial Dependency

- **Determinant FD** is part of Candidate key
- **Dependent FD is a Non Prime Attribute**

R(A,B,C,D)

 $A \rightarrow C, B \rightarrow D$

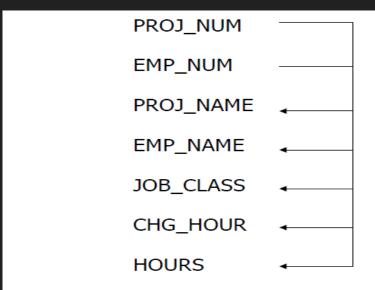
AB is candidate key

A->C A is part of Candidate key,

C is Non Prime Attribute



Dependency Diagram



Functional Dependency: each attribute is uniquely identified by, or is dependent on the primary key.

Primary key is a composite primary key and is made up of PROJ_NUM and EMP_NUM

EMPLOYEE_PROJECT (PROJ_NUM (pk), EMP_NUM (pk), PROJ_NAME, EMP_NAME, JOB_CLASS, CHG_HOUR, HOURS)

TA	BLE_PRODUC	CT
Product ID	Color	Price
1	red, green	15.99
2	yellow	23.99
3	green	17.50
4	yellow, blue	9.99
5	red	29.99

- A database is in first normal form if it satisfies the following conditions:
 - > Contains only atomic values
 - > There are no repeating groups
- **An atomic value is a value that cannot be divided.**
- For example, in the table shown below, the values in the [Color] column in the first row can be divided into "red" and "green", hence [TABLE_PRODUCT] is not in 1NF.

Product ID

3

5

Color

red, green

vellow

green vellow, blue

red

Price

15.99

23.99

17.50

9.99

29.99

- This table is not in first normal form because the [Color] column can contain multiple values. For example, the first row includes values "red" and "green."
- To bring this table to first normal form, we split the table into two tables and now we have the resulting tables:

TA	BLE_PRODUC	CT
Product ID	Color	Price
1	red, green	15.99
2	yellow	23.99
3	green	17.50
4	yellow, blue	9.99
5	red	29.99

TABLE_PRODUCT_PRICE

Product ID	Price
1	15.99
2	23.99
3	17.50
4	9.99
5	29.99

TABLE_PRODUCT_COLOR

Product ID	Color
1	red
1	green
2	yellow
3	green
4	yellow
4	blue
5	red

A database is in **second normal form** if it satisfies the following conditions:

- It is in first normal form
- All non-key attributes are fully functional dependent on the primary key

- In a table, if attribute B is functionally dependent on A,
 - but is not functionally dependent on a proper subset of A, then B is considered fully functional dependent on A.
- If a table is in 2NF, all non-key attributes cannot be dependent on a subset of the primary key.

- Note that if the primary key is not a composite key, all non-key attributes are always fully functional dependent on the primary key.
- A table that is in 1st normal form and contains only a single key as the primary key is automatically in 2nd normal form.

Second Normal Form Example

Consider the following example:

	TABLE_PU	RCHASE_DETAIL
CustomerID	Store ID	Purchase Location
1	1	Los Angeles
1	3	San Francisco
2	1	Los Angeles
3	2	New York
4	3	San Francisco

TABLE_PURCHASE_DETAIL Customer ID Store ID Purchase Location 1 1 Los Angeles 1 3 San Francisco 2 1 Los Angeles 3 2 New York 4 3 San Francisco

2nd Normal Form ExampleConsider the following example:

- This table has a composite primary key [Customer ID, Store ID].
- The non-key attribute is [Purchase Location].
- In this case, [Purchase Location] only depends on [Store ID], which is only part of the primary key.
- Therefore, this table does not satisfy second normal form.
- To bring this table to second normal form, we decompose the table into two tables, and now we have the following:

Second Normal Form Example

TABLE_P	JRCHASE	TA	BLE_STORE
Customer ID	Store ID	Store ID	Purchase Location
1	1	1	Los Angeles
1	3	2	New York
2	1	3	San Francisco
3	2		
4	3		

Second Normal Form Example

- What we have done is to remove the partial functional dependency
- Now, in the table [TABLE_STORE], the column [Purchase Location] is fully dependent on the primary key of that table, which is [Store ID].

First normal form (1NF)

- First normal form: A relation is in first normal form if every attribute in every row can contain only one single (atomic) value.
- **A** university uses the following relation:

Student(Surname, Name, Skills)

First normal form (1NF)

A university uses the following relation:

Student(Surname, Name, Skills)

The attribute Skills can contain multiple values and therefore the relation is not in the first normal form.

But the attributes Name and Surname are atomic attributes that can contain only one value.

FirstName	LastName	Knowledge	
Thomas	Mueller	Java, C++, P	\mathbf{HP}
Ursula	Meier	PHP, Java	
Igor	Mueller	C++, Java	
	Result	after Normalis	-+:
	- Result	arter reormans	auc
Students	- Result	arter ryomans	ащ
Students FirstName	LastName	Knowledge	апс
	1		au
FirstName	LastName	Knowledge	
FirstName Thomas	LastName Mueller	Knowledge C++	
FirstName Thomas Thomas	LastName Mueller Mueller	Knowledge C++ PHP	
FirstName Thomas Thomas Thomas	LastName Mueller Mueller Mueller	Knowledge C++ PHP Java	
FirstName Thomas Thomas Thomas Ursula	LastName Mueller Mueller Mueller Meier	Knowledge C++ PHP Java Java	atic

References

- Hillyer Mike, MySQL AB. <u>An Introduction to Database Normalization</u>, http://dev.mysql.com/tech-resources/articles/intro-to-normalization.html, accessed October 17, 2006.
- Microsoft. Description of the database normalization basics, http://support.microsoft.com/kb/283878, accessed October 17, 2006.
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- https://www.db-book.com/db6/index.html
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Summary

- Partial Dependency
- Second Normal Form

Next Lecture Transitive dependency

Third Normal Form

Thank You

Happy to answer any questions!!!