# Non-monotone Continuous DR-submodular Maximization: Structure and Algorithms



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DR-submodularity captures a subclass of non-convex/non-concave functions that enables exact minimization and approximate maximization in poly. time.

Investigate geometric properties that underlie such objectives, e.g., a strong relation between stationary points & global optimum is proved.

Devise two guaranteed algorithms: i) A "two-phase" algorithm with 1/4 approximation guarantee. ii) A non-monotone Frank-Wolfe variant with 1/e approximation guarantee

Extend to a much broader class of submodular functions on "conic" lattices.

## DR-submodular (Diminishing Returns) **Maximization & Its Applications**



 $f \colon \mathcal{X} \to \mathbb{R}$  is continuous DR-submodular.  $\mathcal{X}$  is a hypercube. Wlog, let  $\mathcal{X} = [\mathbf{0}, \overline{\mathbf{u}}]$ .  $\mathcal{P} \subseteq \mathcal{X}$  is convex and down-closed:  $x \in \mathcal{P} \& 0 \le y \le x \text{ implies } y \in \mathcal{P}.$ 

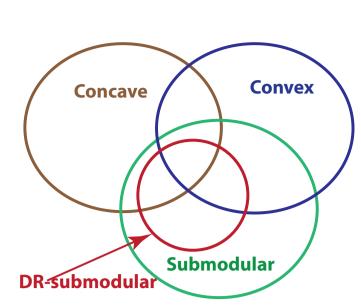
**DR-submodular (DR property)** [Bian et al '17]:  $\forall a \leq b \in \mathcal{X}, \forall i, \forall k \in \mathbb{R}_+$ , it holds,

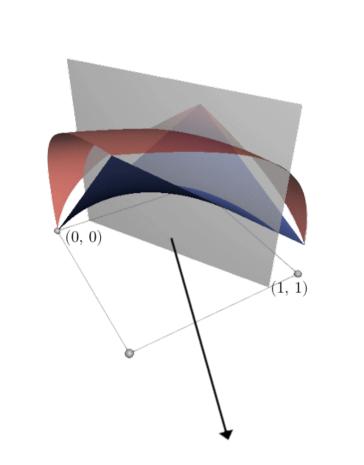
$$f(k\boldsymbol{e}_i + \boldsymbol{a}) - f(\boldsymbol{a}) \ge f(k\boldsymbol{e}_i + \boldsymbol{b}) - f(\boldsymbol{b}).$$

- If f differentiable,  $\nabla f()$  is an antitone mapping  $(\forall a \leq b)$ , it holds  $\nabla f(a) \geq \nabla f(b)$
- If f twice differentiable,  $\nabla_{ij}^2 f(x) \leq 0$ ,  $\forall x$

# Applications

- Softmax extension for determinantal point processes (DPPs) [Gillenwater et al '12]
- Mean-field inference for log-submodular models [Djolonga et al '14]
- DR-submodular quadratic programming
- (Generalized submodularity over conic lattices) e.g., logistic regression with a non-convex separable regularizer [Antoniadis et al '11]
- **Etc...** (more see paper)





Softmax (red) & multilinear (blue) extensions, & concave cross-sections Fig. from [Gillenwater et al '12]

## **Underlying Properties of DR-submodular Maximization**

**Concavity Along Non-negative Directions:** 

Quadratic Lower Bound. With a L-Lipschitz gradient, for all x and  $v \in \pm \mathbb{R}^n_+$ , it holds,  $f(\mathbf{x} + \mathbf{v}) \ge f(\mathbf{x}) + \langle \nabla f(\mathbf{x}), \mathbf{v} \rangle - \frac{L}{2} ||\mathbf{v}||^2$ 

**Strongly DR-submodular & Quadratic Upper Bound**. f is  $\mu$ -strongly DR-submodular if for all x and  $v \in \pm \mathbb{R}^n_+$ , it holds,

$$f(x + v) \le f(x) + \langle \nabla f(x), v \rangle - \frac{\mu}{2} ||v||^2$$

#### **Approximately Stationary Points & Global Optimum:**

**Lemma.** For any x, y,  $\langle y - x, \nabla f(x) \rangle \ge f(x \vee y) + f(x \wedge y) - 2f(x) + \frac{\mu}{2} ||x - y||^2$ 

If  $\nabla f(x) = 0$ , then  $2f(x) \ge f(x \lor y) + f(x \land y) + \frac{\mu}{2} ||x - y||^2 \rightarrow \text{ implicit relation between } x \& y$ . (finding an exact stationary point is difficult (2))

**Non-stationarity Measure** [Lacoste-Julien '16]. For any  $Q \subseteq \mathcal{X}$ , the non-stationarity of  $x \in Q$  is,  $g_{\mathcal{Q}}(\mathbf{x}) \coloneqq \max_{\mathbf{v} \in \mathcal{Q}} \langle \mathbf{v} - \mathbf{x}, \nabla f(\mathbf{x}) \rangle$ 

(Local-Global Relation). Let  $x \in \mathcal{P}$  with non-stationarity  $g_{\mathcal{P}}(x)$ . Define  $Q := \mathcal{P} \cap \{y \mid y \leq \overline{u} - x\}$ . Let  $z \in Q$  with non-stationarity  $g_Q(z)$ . Then,  $\max\{f(\mathbf{x}), f(\mathbf{z})\} \ge \frac{1}{4} [f(\mathbf{x}^*) - g_{\mathcal{P}}(\mathbf{x}) - g_{\mathcal{Q}}(\mathbf{z})] + \frac{\mu}{8} (\|\mathbf{x} - \mathbf{x}^*\|^2 + \|\mathbf{z} - \mathbf{z}^*\|^2),$ where  $z^* \coloneqq x \vee x^* - x$ .

Proof using the essential **DR** property on carefully constructed auxiliary points

- Good empirical performance for the Two-Phase algorithm: if x is away from  $x^*$ ,  $||x - x^*||^2$  will augment the bound; if x is close to  $x^*$ , by the smoothness of f, should be near optimal.

# **Two Guaranteed Algorithms**

#### NON-MONOTONE FRANK-WOLFE VARIANT

**Input**: step size  $\gamma \in (0,1]$ 

 $x^{(0)} \leftarrow 0, k \leftarrow 0, t^{(0)} \leftarrow 0$ // t: cumulative step size

While  $t^{(k)} < 1$  do:

 $v^{(k)} \leftarrow \operatorname{argmax}_{v \in \mathcal{P}, v \leq \overline{u} - x^{(k)}} \langle v, \nabla f(x^{(k)}) \rangle$ // shrunken LMO

 $\gamma_k \leftarrow \min\{\gamma, 1 - t^{(k)}\}$ 

 $\mathbf{x}^{(k+1)} \leftarrow \mathbf{x}^{(k)} + \gamma_k \mathbf{v}^{(k)}, \ t^{(k+1)} \leftarrow t^{(k)} + \gamma_k, \ k + +$ 

key difference from the monotone Frank-Wolfe variant [Bian et al '17]

Output:  $x^{(K)}$ 

Guarantee of Non-Monotone Frank-Wolfe Variant.

$$f(\mathbf{x}^{(K)}) \ge e^{-1} f(\mathbf{x}^*) - O(\frac{1}{K^2}) f(\mathbf{x}^*) - \frac{D^2 L}{2K}$$

D: diameter of  $\mathcal{P}$ *L*: smooth gradient Based on Local-Global Relation, can use any solver for finding an approximately stationary point as the subroutine, e.g., the Nonconvex Frank-Wolfe solver in [Lascote-Julien '16]

## **TWO-PHASE ALGORITHM**

**Input**: stopping tolerances  $\epsilon_1$ ,  $\epsilon_2$ , #iterations  $K_1$ ,  $K_2$ 

 $x \leftarrow \text{Non-convex Frank-Wolfe}(f, \mathcal{P}, K_1, \epsilon_1)$ // Phase I on  $\mathcal P$ 

 $Q \leftarrow \mathcal{P} \cap \{y \mid y \leq \overline{u} - x\}$ 

 $z \leftarrow \text{Non-convex Frank-Wolfe}(f, Q, K_2, \epsilon_2)$ // Phase II on Q

Output:  $argmax{f(x), f(z)}$ 

#### **Guarantee of Two-Phase Algorithm.**

$$\max\{f(x), f(z)\} \ge \frac{\mu}{8} (\|x - x^*\|^2 + \|z - z^*\|^2) + \frac{1}{4} \left[ f(x^*) - \min\left\{\frac{C_1}{\sqrt{K_1 + 1}}, \epsilon_1\right\} - \min\left\{\frac{C_2}{\sqrt{K_2 + 1}}, \epsilon_2\right\} \right],$$

where  $z^* := x \vee x^* - x$ 

# Experimental Results (more see paper)

#### Baselines:

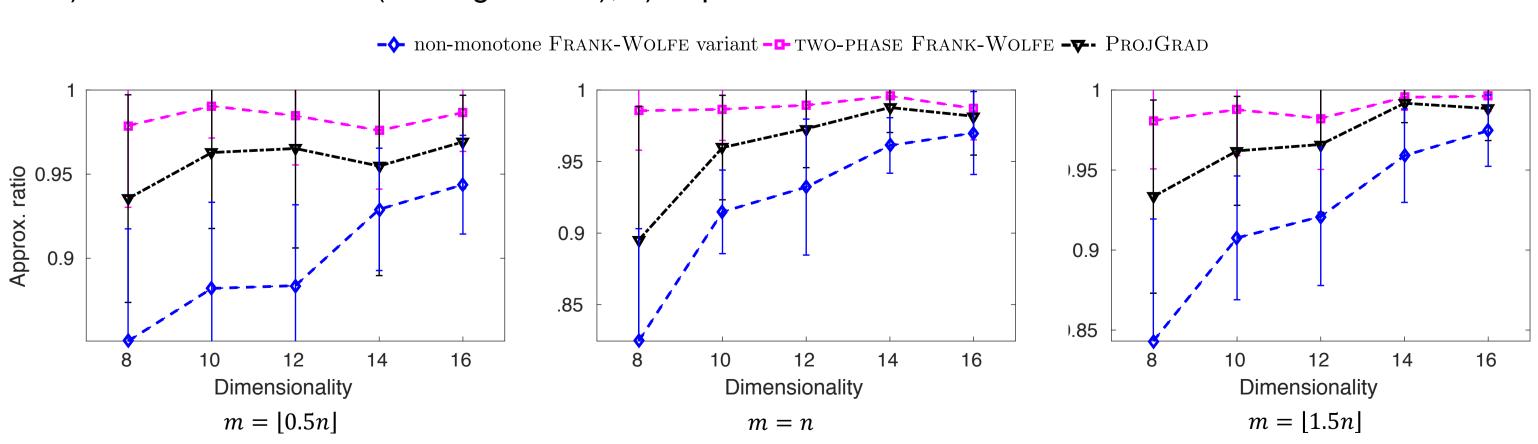
- QUADPROGIP: global solver for non-convex quadratic programming (possibly in exponential time)
- Projected Gradient Ascent (PROJGRAD) with diminishing step sizes (1/k+1)

# DR-submodular Quadratic Programming.

Synthetic problem instances  $f(x) = \frac{1}{2}x^THx + h^Tx + c$ ,  $\mathcal{P} = \{x \in \mathbb{R}^n_+ \mid Ax \leq b, x \leq b\}$  $\overline{u}$ ,  $\mathbf{A} \in \mathbb{R}_{++}^{m \times n}$ ,  $\mathbf{b} \in \mathbb{R}_{+}^{m}$ } has m linear constraints.

Randomly generated in two manners:

1) Uniform distribution (see Figs below); 2) Exponential distribution



Feldman, Naor, and Schwartz. A unified continuous greedy algorithm for submodular maximization. FOCS 2011

Gillenwater, Kulesza, and Taskar. Near-optimal map inference for determinantal point processes. NIPS 2012.

Bach. Submodular functions: from discrete to continous domains. arXiv:1511.00394, 2015.

Lacoste-Julien. Convergence rate of frank-wolfe for non-convex objectives. arXiv:1607.00345, 2016.

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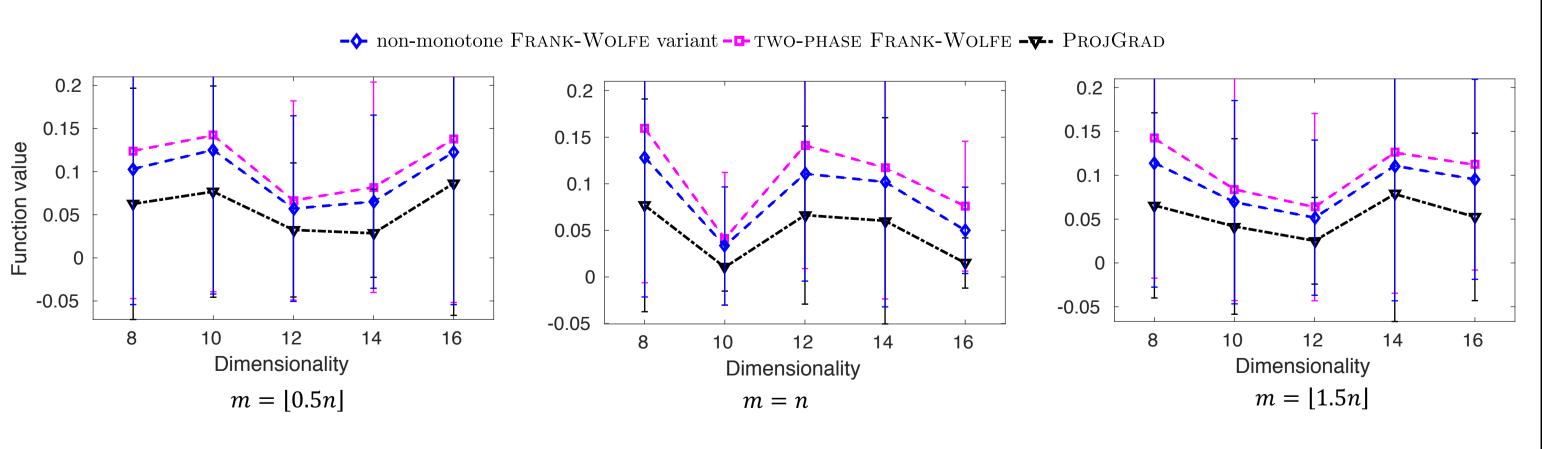
# Maximizing Softmax Extensions for MAP inference of DPPs.

 $f(\mathbf{x}) = \log \det(\operatorname{diag}(\mathbf{x}) (\mathbf{L} - \mathbf{I}) + \mathbf{I}), \mathbf{x} \in [0,1]^n$ 

L: kernel/similarity matrix.  $\mathcal{P}$  is a matching polytope for matched summarization.

#### Synthetic problem instances:

- Softmax objectives: generate  $\mathbf{L}$  with n random eigenvalues
- Generate polytope constraints similarly as that for quadratic programming



#### Real-world results on matched summarization:

Select a set of document pairs out of a corpus of documents, such that the two documents within a pair are similar, and the overall set of pairs is as diverse as possible. Setting similar to [Gillenwater et al '12], experimented on the 2012 US Republican detates data.

