Memory Overhead Analysis of Container-based Android sandboxes

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Objective

Adapt VPBox for Emulators

Collect Memory Usage Traces of instances of Virtual Phones

Analyse physical page sharing patterns

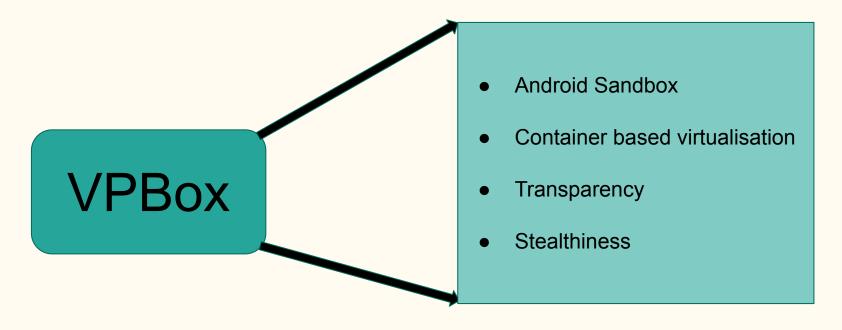
Container-Based Virtualization

Hardware resource multiplexing for all devices

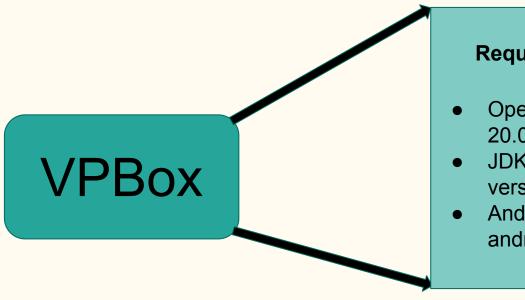
Incomplete device virtualization

Lack of transparency & stealthiness

Introduction



Introduction



Requirements for building*:

- Operating system : Ubuntu 20.04 LTS
- JDK version : OpenJDK version 9
- Android Source Code : android-10.0.0 r33

2 Levels of Virtualization

KERNEL LEVEL

- 1. Device namespace
- 2. Binder, GPS multiplexing
- 3. Efficient hardware resource multiplexing

USER LEVEL

- 1. Achieve no in-guest virtualization component
- 2. Binder service sharing

Memory Usage Bottleneck: Proposed Solutions



Extend Android

- 1. KSM
- 2. Low memory killer

Extra Features

- 1. AUFS
- 2. Screen off

Picture Credits: https://dl.acm.org/doi/pdf/10.1145/3460120.3484544

Objective

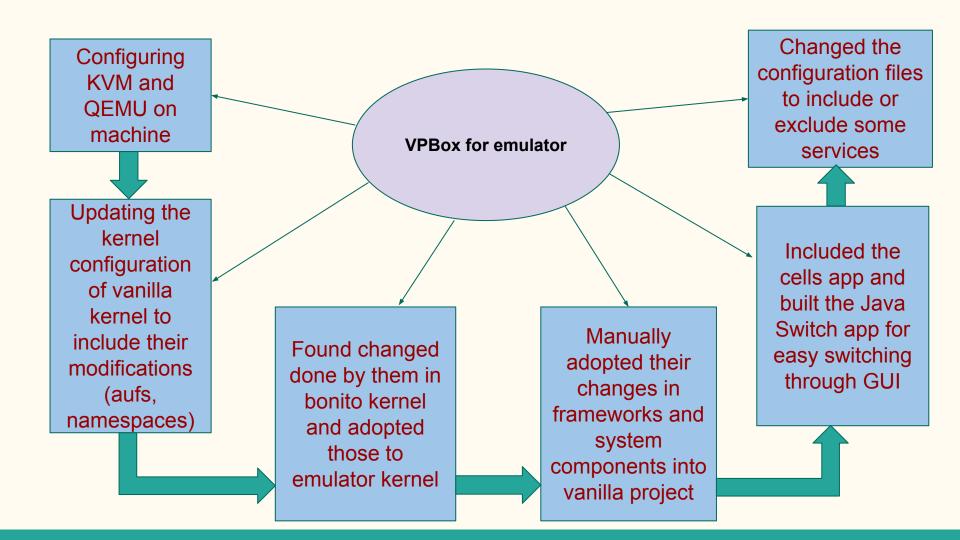
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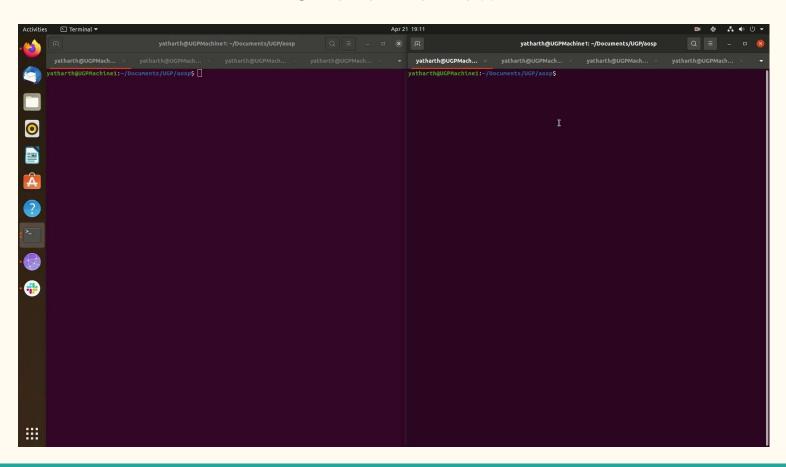
Analyse physical page sharing patterns

Device Specifications

- Emulated Physical Device: Pixel 3a XL
- Android version: Android 10.0.0_r33
- Linux kernel version: goldfish-kernel-4.14.112
- Android Emulator API Version 29
- Memory: 8GB RAM



Short Demo..



Objective

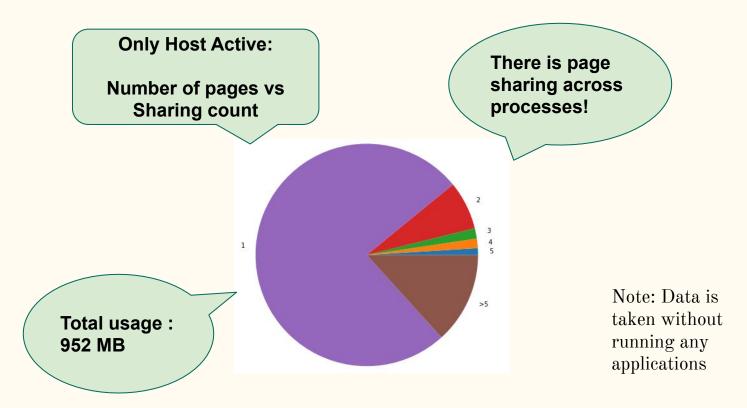
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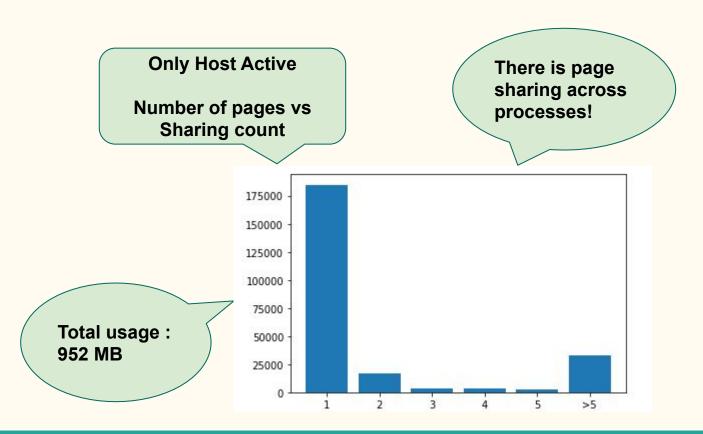
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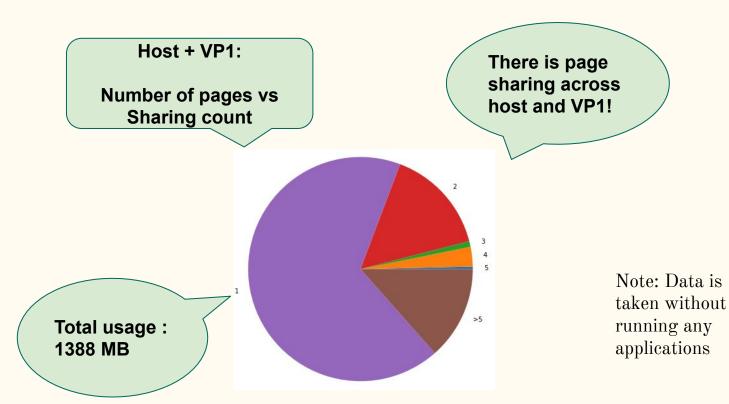
Analyse physical page sharing patterns

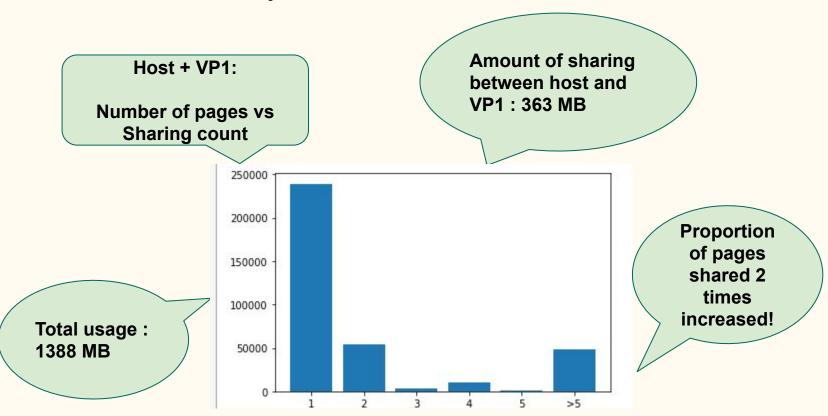
Collecting memory usage data

- 1. Implemented a BFS inside kernel which on being given a start pid, walks over VM areas of all the processes in the subtree of the given process.
- 2. Identified physical mappings for the processes by following page table entries.
- 3. Collected the following information about each virtual address
 - a. Physical mapping
 - b. VM area permissions + PTE permissions
 - c. File path
- 4. Used the pseudo sysfs filesystem in linux kernel to get the above information for init processes of all the virtual phones as well as host by setting up callbacks for the same inside kernel.
- 5. We have removed the physical memory mappings corresponding to device /goldfish_pipe (related to emulator) so as to compare the results with the paper.









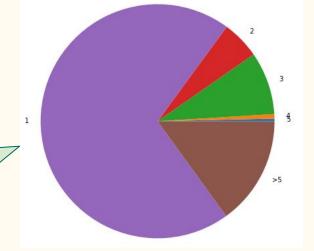


Number of pages vs Sharing count

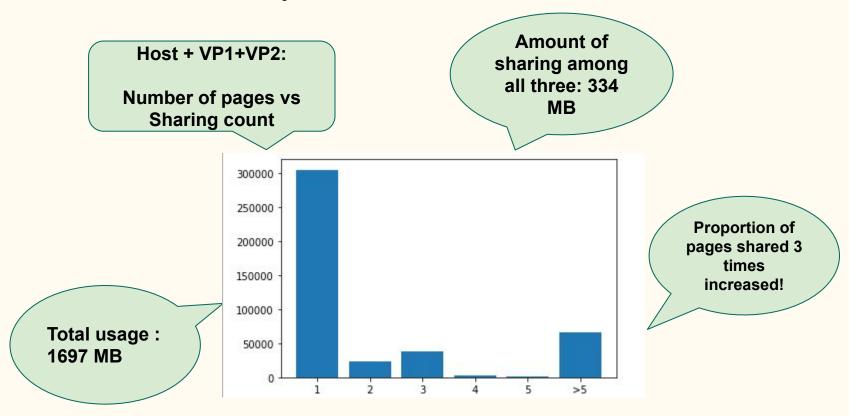
Total usage:

1697 MB

There is page sharing across host, VP1 and VP2!



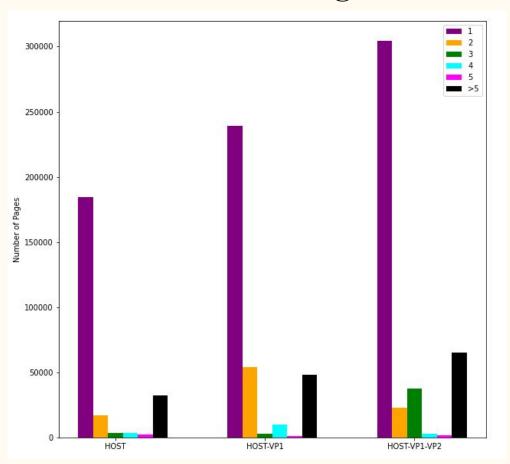
Note: Data is taken without running any applications



More Information

```
Host pages: 896.00390625 MB
VP1 pages: 744.94921875 MB
VP2 pages: 763.66796875 MB
TOTAL pages: 1697.078125 MB
Unique VP1-VP2 1147.94140625 MB
Unique VP1-HOST 1299.53515625 MB
Unique VP2-HOST 1319.80078125 MB
VP1 VP2 common 360.67578125 MB
VP1 HOST common 341.41796875 MB
VP2 HOST common 339.87109375 MB
TOTAL Common 334.421875 MB
```

Summarizing it..



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Observations

Ineffectiveness in sharing shared object files (libraries):

- We found the top VMAs which have unshared physical pages and their corresponding files.
- Analysed the library files(.so) present in the above list to see their association with namespaces.
- Found that the unshared pages corresponding to these files are equally distributed among HOST, VP1 and VP2.
- Example: '/system/lib64/libcrypto.so' has 399 physical pages unique to host and 304 pages unique to both VP1 and VP2

Observations

Memory overhead of Virtual Phone

- We observe that overhead of each of the virtual phone is around 350-400 MB. This is somewhat consistent with what are results shown in the paper.
- The amount of pages which are unique across VP phones increase by around 250-300 MB, which is also consistent with what is shown in the paper.

Observations

KSM analysis

- We took the pte flags as well while translating a virtual address.
- Next, we found the addresses which have the VM area permissions as 'rw' but pte_flags as only 'r'.
- These might be the pages that are being handled by KSM.
- Most of these pages corresponded to ANONYMOUS mappings or same namespace.

Future Work

Root Cause of the memory overheads

Aggressiveness of KSM

Security related analysis and network overheads

AUFS => OverlayFS?

QUESTIONS?

THANK YOU!