

Review Week 08

2022-11-14

Hamming and simplex codes + useful facts

Hamming code $\text{Ham}(r, q)$ for prime power q

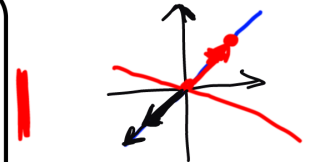
Reminder: COURSEWORK

Week 10 (timed online test)
Material up to and incl. Week 09

M.J.E. Golay (1949)

Def Projective space, $\mathbb{P}_{n-1}(\mathbb{F}_q)$:

- a line in \mathbb{F}_q^n is a 1-dimensional subspace of \mathbb{F}_q^n .
- a representative vector of a line is any non-zero vector from that line (it spans the line)
- $\mathbb{P}_{n-1}(\mathbb{F}_q) = \{\text{all lines in } \mathbb{F}_q^n\}$



Check matrix for $\text{Ham}(r, q)$: - take one representative column vector from each line in \mathbb{F}_q^r - this guarantees $d = 3$

Construct a check matrix H for $\text{Ham}(3, 2)$

$\begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$ represents $\left\{ \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \right\}$ ($\mathbb{F}_2 = \{0, 1\}$)

$$\left[\begin{array}{cccccc} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{array} \right] \quad r=3$$

Construct a check matrix H for $\text{Ham}(2, 3)$

$$\left\{ \begin{array}{l} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \rightarrow \begin{bmatrix} 1 \\ 0 \end{bmatrix} \rightarrow \begin{bmatrix} 2 \\ 0 \end{bmatrix} \\ \begin{bmatrix} 0 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 0 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 0 \\ 2 \end{bmatrix} \\ \begin{bmatrix} 1 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 2 \\ 2 \end{bmatrix} \\ \begin{bmatrix} 2 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 2 \\ 1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 \\ 2 \end{bmatrix} \end{array} \right\} \quad 4 \text{ lines}$$

$$r=2 \quad \mathbb{F}_3^2 = \{ \dots \}$$

$$\Rightarrow H = \begin{bmatrix} 1 & 2 & 1 & 0 \\ 2 & 2 & 0 & 1 \end{bmatrix} \quad 4 \text{ columns}$$

Parameters of $\text{Ham}(r, q)$

- n
- k
- d

length $\# \mathbb{P}_{r-1}(\mathbb{F}_q) = \frac{q^r - 1}{q - 1} = q^{r-1} + q^{r-2} + \dots + q + 1$

$d=3 \Rightarrow t = \left\lfloor \frac{d-1}{2} \right\rfloor = 1$ $n - \dim \text{Ham}(r, q)^\perp = n - r$

The decoder for $\text{Ham}(r, q)$

$$H = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{bmatrix}$$

$\underline{a} = 0000000$
 $\underline{a} = 1000000$
 $\underline{a} = 0100000$
 $\underline{a} = 0010000$ $\left. \begin{array}{l} \text{vectors of} \\ \text{wt}=1 \end{array} \right\} n$

$S(\underline{a}) = 000$
 $S(\underline{a}) = 100$
 $S(\underline{a}) = 010$
 $S(\underline{a}) = 110$ $\left. \begin{array}{l} \text{vectors of} \\ \text{wt}=1 \end{array} \right\} 2^r - 1$

$\# \text{ columns} = 2^r - 1$

$\text{Ham}(r, q)$ is a perfect code

$\underline{y} \in \mathbb{F}_2^n$ received

$$S(\underline{y}) = \underline{\begin{bmatrix} * \\ * \end{bmatrix}} = S(\underline{e}_i)$$

$$\underline{e}_i = 00 \dots 0 \underset{i}{1} 0 \dots 0$$

$$S(\underline{y}) = \begin{bmatrix} 1 \\ 1 \end{bmatrix}; \quad \underline{h}_2 = 101 \Rightarrow \text{Decode}(\underline{y}) = \underline{y} - \underline{e}_2$$

$$q \neq 2: \quad S(\underline{y}) = \lambda \underline{h}_i, \quad \text{Decode}(\underline{y}) = \underline{y} - \lambda \underline{e}_i$$

The MacWilliams identity

$C \neq D$ but $w_C(x, y) = w_D(x, y)$

THM

$$w_{C^\perp}(x, y) = \frac{1}{\#C} w_C(x + (q-1)y, xy)$$

\Downarrow

$$w_{C^\perp}(x, y) = w_{D^\perp}(x, y)$$

The Average Weight Equation

$$\frac{1}{\#C} \sum_{c \in C} w(c) = (n - z)(1 - \frac{1}{q})$$

The Plotkin bound

average weight = $n(1 - q^{-1})$ if G has no zero columns

e.g. if we know

w_{C^\perp} for $C = \text{Ham}(2, 3)$

we can calculate $w_C(x, y)$

$$C \subseteq \mathbb{F}_2^n \text{ linear, } d = d(C) > \frac{n}{2} \Rightarrow$$

$$\#C \leq \frac{d}{d - n/2}$$

The simplex code $\Sigma(r, q)$

Def The simplex code: $\Sigma(r, q) = \text{Ham}(r, q)^\perp$



THM $\Sigma(r, q)$, of length $n = (q^r - 1)/(q - 1)$ and dimension r , has the property that the Hamming distance between each pair of codewords is q^{r-1} .

Find the average weight of a codeword of $\Sigma(3, 2)$ in three ways.

Bonus question: does $\Sigma(3, 2)$ attain the Plotkin bound?

$$c \in \Sigma(r, q) \setminus \{0\}, w(c) = q^{r-1}$$