

Translations between bases in Base-extension Semantics

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Goals for this talk

- Introduce a more generalised framework of atomic rules and bases.
- Introduce two notions of atomic derivability.
- Show how one may relate these notions of atomic derivability.
- Recall support relations for IPL and ILL.
- Show how one may thus relate these notions of support.

Presentation root directory

- 1 Bases and atomic rules
- 2 Atomic derivability
- 3 Comparing relations
- 4 BeS for IPL and ILL
- 5 Translating between the semantics

Notation

- \mathbb{A} represents a fixed, countably infinite set of propositional atoms.
- Lower case latin letters represent propositional atoms.
- Upper case latin letters represent finite multisets of propositional atoms.
- The sum of two multisets P and Q is denoted $P \uplus Q$.

Atomic rules

Atomic rules take the form:

$$(P_1 \Rightarrow p_1), \dots, (P_n \Rightarrow p_n) \Rightarrow p$$

Atomic rules

Pictorially, we can represent this as:

$$\frac{\begin{array}{ccc} [P_1] & & [P_n] \\ p_1 & \dots & p_n \end{array}}{p}$$

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or linearly, like this:

$$(P_1 \Rightarrow p_1), \dots, (P_n \Rightarrow p_n) \Rightarrow p$$

Contextual atomic rules

Contextual atomic rules are rules which may have contextual brackets distributed across them.

Contextual atomic rules

Pictorially, we can represent such rules as:

$$\frac{\left\{ \begin{array}{ccc} [P_{1_1}] & & [P_{1_{l_1}}] \\ p_{1_1} & \dots & p_{1_{l_1}} \end{array} \right\} \quad \dots \quad \left\{ \begin{array}{ccc} [P_{n_1}] & & [P_{n_{l_n}}] \\ p_{n_1} & \dots & p_{n_{l_n}} \end{array} \right\}}{p}$$

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or linearly, like this:

$$\{(P_{1_i} \Rightarrow p_{1_i})\}_{i=1}^{l_1}, \dots, \{(P_{n_i} \Rightarrow p_{n_i})\}_{i=1}^{l_n} \Rightarrow p$$

Examples of atomic rules

- $\{\Rightarrow a\} \Rightarrow c$
- $\{\Rightarrow a\}, \{(c, d \Rightarrow e), (f \Rightarrow g)\} \Rightarrow q$
- $\{a \Rightarrow b\}, \{\Rightarrow e\}, (f \Rightarrow g) \Rightarrow q$

- $(\Rightarrow a) \Rightarrow c$
- $(\Rightarrow b) \Rightarrow c$
- $(\Rightarrow c), (a \Rightarrow p), (b \Rightarrow p) \Rightarrow p$

Bases

Definition (Base)

A base is a set of atomic rules.

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Else it is said to be context-free.

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- $\{((\Rightarrow a) \Rightarrow c), ((\Rightarrow b) \Rightarrow c), ((\Rightarrow c), (a \Rightarrow p), (b \Rightarrow p) \Rightarrow p)\}$
- $\{(\{ \Rightarrow a \} \Rightarrow c), (\{ a \Rightarrow c \} \Rightarrow b), (\{ \Rightarrow a \}, \{ \Rightarrow a \} \Rightarrow a)\}$
- $\{(\{ \Rightarrow a \} \Rightarrow c), (\{ a \Rightarrow c \} \Rightarrow b), (\{ \Rightarrow d \}, \{ \Rightarrow d \} \Rightarrow a), (\{ \Rightarrow a \} \Rightarrow d)\}$

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Definition (Context-free atomic derivability $\triangleright_{\mathcal{B}}$)

The relation of derivability in a context-free base \mathcal{B} , is defined as so:

Ref $S \triangleright_{\mathcal{B}} p$ if $p \in S$.

App For $((P_1 \Rightarrow p_1), \dots, (P_n \Rightarrow p_n) \Rightarrow q) \in \mathcal{B}$ and $S, P_i \triangleright_{\mathcal{B}} p_i$ for each $i \in \{1, \dots, n\}$ then $S \triangleright_{\mathcal{B}} q$.

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Definition (Contextual atomic derivability $\vdash_{\mathcal{B}}$)

The relation of derivability in a contextual base \mathcal{B} , is defined as so:

Ref $p \vdash_{\mathcal{B}} p$.

App For $(\{(P_{1_i} \Rightarrow p_{1_i})\}_{i=1}^{l_1}, \dots, \{(P_{n_i} \Rightarrow p_{n_i})\}_{i=1}^{l_n} \Rightarrow q) \in \mathcal{B}$ and $C_i, P_{ij} \vdash_{\mathcal{B}} p_{ij}$ for each $i \in \{1, \dots, n\}$ and $j \in \{1, \dots, l_i\}$ then $C_1, \dots, C_n \vdash_{\mathcal{B}} q$.

Example derivations

Example

Let $\mathcal{B} = \{(\Rightarrow a), ((\Rightarrow a), (\Rightarrow b) \Rightarrow c)\}$

$$\frac{\frac{}{b \triangleright_{\mathcal{B}} b} \text{Ref} \quad \frac{}{b \triangleright_{\mathcal{B}} a} \Rightarrow a}{b \triangleright_{\mathcal{B}} c} (\Rightarrow a), (\Rightarrow b) \Rightarrow c$$

Example derivations

Example

Let $\mathcal{B} = \{(\Rightarrow a), (\{\Rightarrow a\}, \{\Rightarrow b\} \Rightarrow c)\}$

$$\frac{\frac{}{b \vdash_{\mathcal{B}} b} \text{Ref} \quad \frac{}{\vdash_{\mathcal{B}} a} \Rightarrow a}{b \vdash_{\mathcal{B}} c} \{\Rightarrow a\}, \{\Rightarrow b\} \Rightarrow c$$

Example derivations

Example

By the definition of $\vdash_{\mathcal{B}}$, deriving a from a, a in the empty base is not possible, i.e. $a, a \vdash_{\emptyset} a$ is not possible. But $a, a \triangleright_{\emptyset} a$ is possible.

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Example

Let $\mathcal{B} = \{(\{ \Rightarrow a \} \Rightarrow c), (\{ a \Rightarrow c \} \Rightarrow b)\}$

$$\frac{\frac{\frac{\times}{a, a \vdash_{\mathcal{B}} a} \quad \{ \Rightarrow a \} \Rightarrow c}{a, a \vdash_{\mathcal{B}} c} \quad \{ a \Rightarrow c \} \Rightarrow b}{a \vdash_{\mathcal{B}} b}$$

We see that in this base, a is not derivable from a, a .

Example derivations

Example (A possible fix)

Let $(\mathcal{B})_{\star} = \{((\Rightarrow a) \Rightarrow c), ((a \Rightarrow c) \Rightarrow b)\}$

$$\begin{array}{c}
 \frac{}{a, a \triangleright_{(\mathcal{B})_{\star}} a} \text{Ref} \\
 \frac{}{a, a \triangleright_{(\mathcal{B})_{\star}} c} (\Rightarrow a) \Rightarrow c \\
 \frac{}{a \triangleright_{(\mathcal{B})_{\star}} b} (a \Rightarrow c) \Rightarrow b
 \end{array}$$

Example derivations

Example (Another possible fix)

Let $\mathcal{C} = \{(\{ \Rightarrow a \} \Rightarrow c), (\{ a \Rightarrow c \} \Rightarrow b), (\{ \Rightarrow a \}, \{ \Rightarrow a \} \Rightarrow a)\}$

$$\begin{array}{c}
 \frac{}{a \vdash_{\mathcal{C}} a} \text{Ref} \quad \frac{}{a \vdash_{\mathcal{C}} a} \text{Ref} \\
 \frac{}{a \vdash_{\mathcal{C}} a} \text{Ref} \quad \frac{}{a \vdash_{\mathcal{C}} a} \text{Ref} \\
 \frac{a, a \vdash_{\mathcal{C}} a}{\{ \Rightarrow a \}, \{ \Rightarrow a \} \Rightarrow a} \\
 \frac{a, a \vdash_{\mathcal{C}} a}{\{ \Rightarrow a \} \Rightarrow c} \\
 \frac{a, a \vdash_{\mathcal{C}} c}{\{ a \Rightarrow c \} \Rightarrow b} \\
 \frac{}{a \vdash_{\mathcal{C}} b}
 \end{array}$$

Example derivations

Example (A very interesting derivation)

Let \mathcal{B} be the following set of rules:

$\{(\{ \Rightarrow b \} \Rightarrow c), (\{ \Rightarrow a \}, \{ b \Rightarrow c \} \Rightarrow d), (\{ \Rightarrow d, \Rightarrow a \} \Rightarrow e), (\{ a \Rightarrow e \} \Rightarrow f)\}$

$$\begin{array}{c}
 \frac{}{a \vdash_{\mathcal{B}} a} \text{Ref} \quad \frac{\frac{}{b \vdash_{\mathcal{B}} b} \text{Ref}}{b \vdash_{\mathcal{B}} c} \{ \Rightarrow b \} \Rightarrow c}{\frac{}{a \vdash_{\mathcal{B}} d} \{ \Rightarrow a \}, \{ b \Rightarrow c \} \Rightarrow d} \text{Ref} \quad \frac{}{a \vdash_{\mathcal{B}} a} \text{Ref} \\
 \frac{\frac{}{a \vdash_{\mathcal{B}} d} \{ \Rightarrow a \}, \{ b \Rightarrow c \} \Rightarrow d \quad \frac{}{a \vdash_{\mathcal{B}} a} \{ \Rightarrow d, \Rightarrow a \} \Rightarrow e}{\frac{a \vdash_{\mathcal{B}} e}{\vdash_{\mathcal{B}} f} \{ a \Rightarrow e \} \Rightarrow f} \text{Ref}
 \end{array}$$

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Comparing relations

Definition (Structural rules)

We define two rules:

- $\text{Wk}_q^p = (\{\Rightarrow p\}, \{\Rightarrow q\} \Rightarrow q)$
- $\text{Ctn}_q^p = (\{\Rightarrow p\}, \{p, p \Rightarrow q\} \Rightarrow q)$

Mappings between bases

Mappings between bases

Definition

Let \mathcal{B} be a context-free base. We define structural contextualisation of that base $(\mathcal{B})^*$ as follows:

$$(\mathcal{B})^* = \{ \{P_1 \Rightarrow p_1\}, \dots, \{P_n \Rightarrow p_n\} \Rightarrow q \mid ((P_1 \Rightarrow p_1), \dots, (P_n \Rightarrow p_n) \Rightarrow q) \in \mathcal{B} \} \\ \cup \{ \text{Wk}_q^p, \text{Ctn}_q^p \mid \forall p, q \in \mathbb{A} \}$$

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We define the decontextualisation of a contextual base as:

$$(\mathcal{B})_\star = \{ (P_{1_1} \Rightarrow p_{1_1}), \dots, (P_{n_{l_n}} \Rightarrow p_{n_{l_n}}) \Rightarrow q \\ \mid (\{P_{1_i} \Rightarrow p_{1_i}\}_{i=1}^{l_1}, \dots, \{P_{n_i} \Rightarrow p_{n_i}\}_{i=1}^{l_n} \Rightarrow q) \in \mathcal{B} \}$$

Properties of these mappings

Let \mathcal{B} be a context-free base. Then the following hold:

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- For all $\mathcal{C} \supseteq \mathcal{B}$ we have that $(\mathcal{C})^* \supseteq (\mathcal{B})^*$
- For all $\mathcal{C} \supseteq (\mathcal{B})^*$ there exists an extension $\mathcal{X} \supseteq \mathcal{B}$ such that $(\mathcal{X})^* = \mathcal{C}$

Properties of these mappings

Let \mathcal{B} be a context-free base. Then the following hold:

- $((\mathcal{B})^*)_{\star} \supseteq \mathcal{B}$
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- For all $\mathcal{C} \supseteq \mathcal{B}$ we have that $(\mathcal{C})^* \supseteq (\mathcal{B})^*$
- For all $\mathcal{C} \supseteq (\mathcal{B})^*$ there exists an extension $\mathcal{X} \supseteq \mathcal{B}$ such that $(\mathcal{X})^* = \mathcal{C}$
- For all $\mathcal{C} \supseteq (\mathcal{B})^*$, it follows immediately that $\mathcal{C} = ((\mathcal{C})_{\star})^*$.

Structural admissibility in a structurally contextualised base

Lemma

Suppose $L \vdash_{(\mathcal{B})^} p$ holds. Then $R, L \vdash_{(\mathcal{B})^*} p$ also holds for any atomic multiset R .*

Proof.

Let $R = \{r_1, \dots, r_n\}$ for some n . Then we have that we can effectively weaken S away as follows:

$$\begin{array}{c}
 \frac{\frac{\frac{}{r_1 \vdash_{(\mathcal{B})^*} r_1} \text{Ref} \quad \frac{\frac{\frac{}{r_n \vdash_{(\mathcal{B})^*} r_n} \text{Ref} \quad \frac{L \vdash_{(\mathcal{B})^*} p}{r_n, L \vdash_{(\mathcal{B})^*} p} \text{Wk}_p^{r_n}}{r_n, L \vdash_{(\mathcal{B})^*} p} \text{Wk}_p^{r_{n-1}}}{r_{n-1} \vdash_{(\mathcal{B})^*} r_{n-1}} \text{Ref} \quad \vdots \quad \frac{\frac{}{r_1 \vdash_{(\mathcal{B})^*} r_1} \text{Ref} \quad \frac{r_2, \dots, r_n, L \vdash_{(\mathcal{B})^*} p}{R, L \vdash_{(\mathcal{B})^*} p} \text{Wk}_p^{r_1}}{R, L \vdash_{(\mathcal{B})^*} p}
 \end{array}$$

□

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Structural admissibility in a structurally contextualised base

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Corollary

For arbitrary $m \geq 1$, if $R^m \vdash_{(\mathcal{B})^} p$ then $R \vdash_{(\mathcal{B})^*} p$.*

Proof.

Let $R = \{r_1, \dots, r_n\}$ for some n . Then we have that we can effectively contract on R as follows:

$$\begin{array}{c}
 \frac{\frac{\frac{}{r_1 \vdash_{(\mathcal{B})^*} r_1} \text{Ref} \quad \vdots \quad \frac{\frac{\frac{}{r_n \vdash_{(\mathcal{B})^*} r_n} \text{Ref} \quad \frac{R, R \vdash_{(\mathcal{B})^*} p}{r_1, \dots, r_{n-1}, R \vdash_{(\mathcal{B})^*} p} \text{Ctn}_p^{r_n}}{r_1, \dots, r_{n-1}, R \vdash_{(\mathcal{B})^*} p} \text{Ctn}_p^{r_{n-1}}}{R \vdash_{(\mathcal{B})^*} p} \text{Ctn}_p^{r_1}
 \end{array}
 \quad \square$$

Key results under this base translation

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- $R \vdash_{(\mathcal{B})^*} p$ iff for all bases $\mathcal{X} \supseteq (\mathcal{B})^*$ where for each $r \in R$ we have $\vdash_{\mathcal{X}} r$ then it follows that $\vdash_{\mathcal{X}} p$.

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- $R \vdash_{(\mathcal{B})^*} p$ iff $!R \Vdash_{(\mathcal{B})^*}^{\emptyset} p$
- $\Vdash_{(\mathcal{B})^*}^R p$ iff $!R \Vdash_{(\mathcal{B})^*}^{\emptyset} p$

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- $R \vdash_{(\mathcal{B})^*} p$ iff $!R \Vdash_{(\mathcal{B})^*}^{\emptyset} p$
- $\Vdash_{(\mathcal{B})^*}^R p$ iff $!R \Vdash_{(\mathcal{B})^*}^{\emptyset} p$
- If $R \Vdash_{\mathcal{B}} p$ then $!R \Vdash_{(\mathcal{B})^*}^{\emptyset} p$

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BeS for IPL: Summary

(At) $\Vdash_{\mathcal{B}} p$ iff $\triangleright_{\mathcal{B}} p$.

(\wedge) $\Vdash_{\mathcal{B}} \varphi \wedge \psi$ iff $\Vdash_{\mathcal{B}} \varphi$ and $\Vdash_{\mathcal{B}} \psi$.

(\vee) $\Vdash_{\mathcal{B}} \varphi \vee \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, p , if $\varphi \Vdash_{\mathcal{C}} p$ and $\psi \Vdash_{\mathcal{C}} p$ then $\Vdash_{\mathcal{C}} p$.

(\rightarrow) $\Vdash_{\mathcal{B}} \varphi \rightarrow \psi$ iff $\varphi \Vdash_{\mathcal{B}} \psi$.

(\perp) $\Vdash_{\mathcal{B}} \perp$ iff $\Vdash_{\mathcal{B}} p$ for all p .

(\top) $\Vdash_{\mathcal{B}} \top$ iff always.

(Inf) $\Gamma \Vdash_{\mathcal{B}} \varphi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$ if $\Vdash_{\mathcal{C}} \gamma$ for all $\gamma \in \Gamma$ then $\Vdash_{\mathcal{C}} \varphi$.

(Val) $\Gamma \Vdash \varphi$ iff $\Gamma \Vdash_{\mathcal{B}} \varphi$ for all \mathcal{B} .

BeS for ILL: Structure and the Exponential

(At) $\Vdash_{\mathcal{B}}^L p$ iff $L \vdash_{\mathcal{B}} p$.

(!) $\Vdash_{\mathcal{B}}^L !\varphi$ iff
 $\forall \mathcal{C} \supseteq \mathcal{B}, K, p$, if $(\forall \mathcal{D} \supseteq \mathcal{C} (\Vdash_{\mathcal{D}}^{\emptyset} \varphi \text{ implies } \Vdash_{\mathcal{D}}^K p))$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

(Inf) $\Gamma \Vdash_{\mathcal{B}}^L \varphi$ iff
 $\forall \mathcal{C} \supseteq \mathcal{B}, K$, if $\Vdash_{\mathcal{C}}^{\emptyset} \Delta$ and $\Vdash_{\mathcal{C}}^K \Theta$ where $!\Delta, \Theta = \Gamma$ then $\Vdash_{\mathcal{C}}^{L,K} \varphi$.

(Val) $\Gamma \Vdash \varphi$ iff $\Gamma \Vdash_{\mathcal{B}}^{\emptyset} \varphi$ for all \mathcal{B} .

BeS for ILL: Multiplicatives

(\otimes) $\Vdash_{\mathcal{B}}^L \varphi \otimes \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, K and p , if $\varphi, \psi \Vdash_{\mathcal{C}}^K p$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

(\multimap) $\Vdash_{\mathcal{B}}^L \varphi \multimap \psi$ iff $\varphi \Vdash_{\mathcal{B}}^L \psi$.

(1) $\Vdash_{\mathcal{B}}^L 1$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, K and p , if $\Vdash_{\mathcal{C}}^K p$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

BeS for ILL: Additives

($\&$) $\Vdash_{\mathcal{B}}^L \varphi \& \psi$ iff $\Vdash_{\mathcal{B}}^L \varphi$ and $\Vdash_{\mathcal{B}}^L \psi$.

(\oplus) $\Vdash_{\mathcal{B}}^L \varphi \oplus \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, K and p , if $\varphi \Vdash_{\mathcal{C}}^K p$ and $\psi \Vdash_{\mathcal{C}}^K p$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

(0) $\Vdash_{\mathcal{B}}^L 0$ iff $\Vdash_{\mathcal{B}}^L p$ for all p .

(\top) $\Vdash_{\mathcal{B}}^L \top$ iff always.

BeS for ILL: Summary

- (At) $\Vdash_{\mathcal{B}}^L p$ iff $L \vdash_{\mathcal{B}} p$.
- (\otimes) $\Vdash_{\mathcal{B}}^L \varphi \otimes \psi$ iff $\forall \mathcal{C} \supseteq \mathcal{B}, K, p (\varphi, \psi \Vdash_{\mathcal{C}}^K p \text{ implies } \Vdash_{\mathcal{C}}^{L,K} p)$.
- (\multimap) $\Vdash_{\mathcal{B}}^L \varphi \multimap \psi$ iff $\varphi \Vdash_{\mathcal{B}}^L \psi$.
- (1) $\Vdash_{\mathcal{B}}^L 1$ iff $\forall \mathcal{C} \supseteq \mathcal{B}, K, p (\Vdash_{\mathcal{C}}^K p \text{ implies } \Vdash_{\mathcal{C}}^{L,K} p)$.
- (&) $\Vdash_{\mathcal{B}}^L \varphi \& \psi$ iff $\Vdash_{\mathcal{B}}^L \varphi$ and $\Vdash_{\mathcal{B}}^L \psi$.
- (\oplus) $\Vdash_{\mathcal{B}}^L \varphi \oplus \psi$ iff $\forall \mathcal{C} \supseteq \mathcal{B}, K, p (\varphi \Vdash_{\mathcal{C}}^K p \text{ and } \psi \Vdash_{\mathcal{C}}^K p \text{ implies } \Vdash_{\mathcal{C}}^{L,K} p)$.
- (0) $\Vdash_{\mathcal{B}}^L 0$ iff $\Vdash_{\mathcal{B}}^L p$ for all p .
- (\top) $\Vdash_{\mathcal{B}}^L \top$ iff always.
- (!) $\Vdash_{\mathcal{B}}^L !\varphi$ iff
 $\forall \mathcal{C} \supseteq \mathcal{B}, K, p (\forall \mathcal{D} \supseteq \mathcal{C} (\Vdash_{\mathcal{D}}^{\emptyset} \varphi \text{ implies } \Vdash_{\mathcal{D}}^K p) \text{ implies } \Vdash_{\mathcal{C}}^{L,K} p)$.
- (Inf) $\Gamma \Vdash_{\mathcal{B}}^L \varphi$ iff
 $\forall \mathcal{C} \supseteq \mathcal{B}, K, \text{ if } \Vdash_{\mathcal{C}}^{\emptyset} \Delta \text{ and } \Vdash_{\mathcal{C}}^K \Theta \text{ where } !\Delta, \Theta = \Gamma \text{ then } \Vdash_{\mathcal{C}}^{L,K} \varphi$
- (Val) $\Gamma \Vdash \varphi$ iff $\Gamma \Vdash_{\mathcal{B}}^{\emptyset} \varphi$ for all \mathcal{B} .

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Comparing clauses

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(IPL) $\Vdash_{\mathcal{B}} \varphi \vee \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, p , if $\varphi \Vdash_{\mathcal{C}} p$ and $\psi \Vdash_{\mathcal{C}} p$ then $\Vdash_{\mathcal{C}} p$.

(ILL) $\Vdash_{\mathcal{B}}^L \varphi \oplus \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, K , p , if $\varphi \Vdash_{\mathcal{C}}^K p$ and $\psi \Vdash_{\mathcal{C}}^K p$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

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(ILL) $\Vdash_{\mathcal{B}}^L \varphi \oplus \psi$ iff for all $\mathcal{C} \supseteq \mathcal{B}$, K , p , if $\varphi \Vdash_{\mathcal{C}}^K p$ and $\psi \Vdash_{\mathcal{C}}^K p$ then $\Vdash_{\mathcal{C}}^{L,K} p$.

(IPL) $\Vdash_{\mathcal{B}} \varphi \rightarrow \psi$ iff $\varphi \Vdash_{\mathcal{B}} \psi$.

(ILL) $\Vdash_{\mathcal{B}}^L \varphi \multimap \psi$ iff $\varphi \Vdash_{\mathcal{B}}^L \psi$.

Comparing clauses

(IPL) $\Vdash_{\mathcal{B}} p \text{ iff } \triangleright_{\mathcal{B}} p.$

(ILL) $\Vdash_{\mathcal{B}}^L p \text{ iff } L \vdash_{\mathcal{B}} p.$

(IPL) $\Vdash_{\mathcal{B}} \varphi \wedge \psi \text{ iff } \Vdash_{\mathcal{B}} \varphi \text{ and } \Vdash_{\mathcal{B}} \psi.$

(ILL) $\Vdash_{\mathcal{B}}^L \varphi \& \psi \text{ iff } \Vdash_{\mathcal{B}}^L \varphi \text{ and } \Vdash_{\mathcal{B}}^L \psi.$

(IPL) $\Vdash_{\mathcal{B}} \varphi \vee \psi \text{ iff for all } \mathcal{C} \supseteq \mathcal{B}, p, \text{ if } \varphi \Vdash_{\mathcal{C}} p \text{ and } \psi \Vdash_{\mathcal{C}} p \text{ then } \Vdash_{\mathcal{C}} p.$

(ILL) $\Vdash_{\mathcal{B}}^L \varphi \oplus \psi \text{ iff for all } \mathcal{C} \supseteq \mathcal{B}, K, p, \text{ if } \varphi \Vdash_{\mathcal{C}}^K p \text{ and } \psi \Vdash_{\mathcal{C}}^K p \text{ then } \Vdash_{\mathcal{C}}^{L,K} p.$

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(IPL) $\Vdash_{\mathcal{B}} \perp \text{ iff } \Vdash_{\mathcal{B}} p \text{ for all } p.$

(ILL) $\Vdash_{\mathcal{B}}^L 0 \text{ iff } \Vdash_{\mathcal{B}}^L p \text{ for all } p.$

Translation of formulas

We seem to have a correspondence between the clauses for the following pairs of connectives:

- \wedge and $\&$
- \vee and \oplus
- \rightarrow and \multimap
- \perp and 0

Translation of formulas

One possible translation of formulas from IPL to ILL is:

- $\llbracket p \rrbracket \equiv p$, where p is an atom.
- $\llbracket \varphi \wedge \psi \rrbracket \equiv \llbracket \varphi \rrbracket \& \llbracket \psi \rrbracket$
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Translation of formulas

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Sequents are translated as follows:

- If $\Gamma \vdash_{\text{IPL}} \varphi$ then $!\llbracket \Gamma \rrbracket \vdash_{\text{ILL}} \llbracket \varphi \rrbracket$

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- If $\Gamma \vdash_{\text{IPL}} \varphi$ then $!\llbracket \Gamma \rrbracket \vdash_{\text{ILL}} \llbracket \varphi \rrbracket$

Semantically we should therefore have:

- If $\Gamma \Vdash \varphi$ then $!\llbracket \Gamma \rrbracket \Vdash \llbracket \varphi \rrbracket$

Immediate correspondence between bases

$$\begin{array}{lll}
 \Gamma \vdash_{\text{IPL}} \varphi & \text{iff } \Gamma \Vdash \varphi & \text{iff for all } \mathcal{B}, p, \text{ if } \Vdash_{\mathcal{B}} \gamma \text{ then } \Vdash_{\mathcal{B}} \varphi \\
 ![\Gamma] \vdash_{\text{ILL}} [\varphi] & \text{iff } ![\Gamma] \Vdash [\varphi] & \text{iff for all } \mathcal{C}, p, \text{ if } \Vdash_{\mathcal{C}}^{\emptyset} [\gamma] \text{ then } \Vdash_{\mathcal{C}}^{\emptyset} [\varphi]
 \end{array}$$

Immediate correspondence between bases

$$\begin{aligned} \Gamma \vdash_{\text{IPL}} \varphi & \quad \text{iff } \Gamma \Vdash \varphi & \quad \text{iff for all } \mathcal{B}, p, \text{ if } \Vdash_{\mathcal{B}} \gamma \text{ then } \Vdash_{\mathcal{B}} \varphi \\ ![\Gamma] \vdash_{\text{ILL}} [\varphi] & \quad \text{iff } ![\Gamma] \Vdash [\varphi] & \quad \text{iff for all } \mathcal{C}, p, \text{ if } \Vdash_{\mathcal{C}}^{\emptyset} [\gamma] \text{ then } \Vdash_{\mathcal{C}}^{\emptyset} [\varphi] \end{aligned}$$

Since we have that $\Gamma \vdash_{\text{IPL}} \varphi$ implies $![\Gamma] \vdash_{\text{ILL}} [\varphi]$, then we have a mapping between bases \mathcal{B} and bases \mathcal{C} .

But what about if we have $\Gamma \Vdash_{\mathcal{B}} \varphi$?

But what about if we have $\Gamma \Vdash_{\mathcal{B}} \varphi$?




Theorem

If $\Gamma \Vdash_{\mathcal{B}} \varphi$ then $\llbracket \Gamma \rrbracket \Vdash_{(\mathcal{B})^}^{\emptyset} \llbracket \varphi \rrbracket$*





Thank you!

Thank you for listening!
Comments? Observations? Please ask and/or feel free to
email me at y.buzoku@ucl.ac.uk

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