Operating System Labs

Yuanbin Wu cs@ecnu

Announcement

- Project 1 due
 - 21:00 Oct. 2
- Project 2 due
 - 21:00 Oct. 23
- TA email update:
 - os_ecnu_class1@163.com
 - os_ecnu_class2@163.com

Operating System Labs

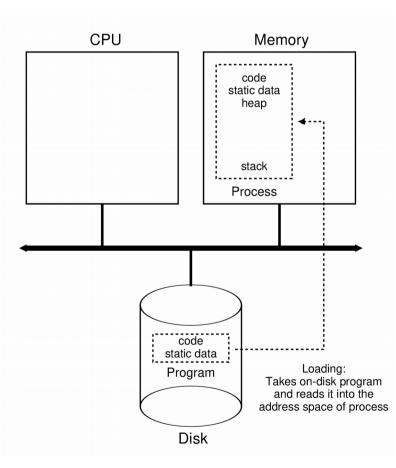
- The abstraction of process
- CPU virtualization
 - Low level and high level mechanisms
- Process API
- Project 2

- The abstraction: Process
 - Running programs
- What does a process consist of?
 - CPU
 - Program Counter (PC)
 - Stack Pointer / Frame Pointer
 - Memory
 - Address space
 - Disk
 - Set of file descriptors

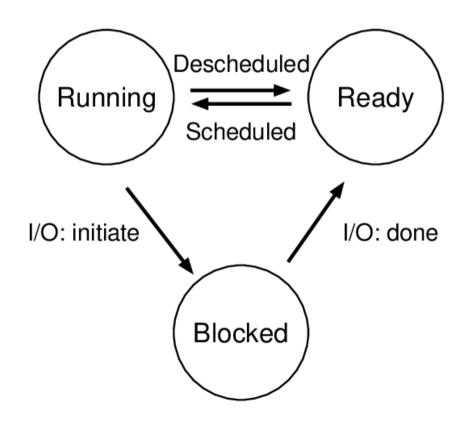
- What does a process consist of?
 - proc file system
 - Example
 - cat /proc/<PID>/status
 - cat /proc/<PID>/maps
 - cat /proc/<PID>/fd
 - cat /proc/<PID>/io
 - Provide a method of communication between kernel space and user space
 - ps command

- Process API
 - Create
 - Destroy
 - Wait
 - Miscellaneous Control
 - Get status

- Details of process creation
 - Load code and static data
 - Establish stack
 - local variables, function calls
 - Init heap
 - malloc, free
 - Allocate file descriptors
 - STDIN_FILENO
 - STDOUT_FILENO
 - STDERR_FILENO



Process States



Process States

Time	$\mathbf{Process}_0$	$\mathbf{Process}_1$	Notes
1	Running	Ready	
2	Running	Ready	
3	Running	Ready	
4	Running	Ready	Process ₀ now done
5	_	Running	
6	_	Running	
7	_	Running	
8	_	Running	Process ₁ now done

Time	$\mathbf{Process}_0$	$\mathbf{Process}_1$	Notes
1	Running	Ready	
2	Running	Ready	
3	Running	Ready	Process ₀ initiates I/O
4	Blocked	Running	Process ₀ is blocked,
5	Blocked	Running	so $Process_1$ runs
6	Blocked	Running	
7	Ready	Running	I/O done
8	Ready	Running	Process ₁ now done
9	Running	_	
10	Running	_	Process ₀ now done

```
// the registers xv6 will save and restore
// to stop and subsequently restart a process
struct context {
 int eip;
 int esp;
 int ebx;
 int ecx;
 int edx;
 int esi;
 int edi;
 int ebp;
};
// the different states a process can be in
enum proc_state { UNUSED, EMBRYO, SLEEPING,
                 RUNNABLE, RUNNING, ZOMBIE };
// the information xv6 tracks about each process
// including its register context and state
struct proc {
                             // Start of process memory
 char *mem;
                             // Size of process memory
 uint sz;
 char *kstack;
                             // Bottom of kernel stack
                             // for this process
 enum proc_state state;
                            // Process state
 int pid;
                             // Process ID
 struct proc *parent;
                             // Parent process
 void *chan;
                             // If non-zero, sleeping on chan
 int killed;
                             // If non-zero, have been killed
 struct file *ofile[NOFILE]; // Open files
 struct inode *cwd;
                    // Current directory
 struct context; // Switch here to run process
 struct trapframe *tf; // Trap frame for the
                             // current interrupt
```

CPU Virtualization

- What
 - Provide the illusion of many CPUs
- Why
 - Multi-task
- How
 - Time sharing

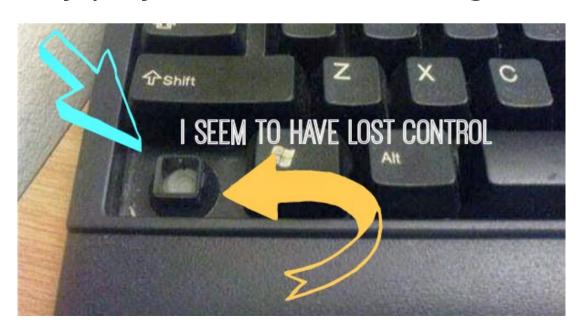
CPU Virtualization

- Mechanisms
 - Low level mechanisms
 - Context switch
 - High level intelligence
 - Scheduling policy

- Low level mechanisms
 - Direct Execution
 - Just run a program on CPU directly

OS	Program
Create entry for process list	
Allocate memory for program	
Load program into memory	
Set up stack with argc/argv	
Clear registers	
Execute call main()	
· ·	Run main()
	Execute return from main
Free memory of process	
Remove from process list	

- Problems of direct execution
 - No control
 - Visit any memory address
 - Open any file
 - Directly play with hardwares (e.g. I/O)



- Limited Direct Execution
 - Kernel model and user model
 - "restricted operations"
 - By the kernel
 - When a thread need do restricted operations
 - System call

- User mode
 - Code is restricted in what it can do
- Kernel mode
 - Code can do what it likes
 - Issue I/O, executing all types of instructions
- How to switch?
 - System call

- Hardware supports on system call
 - A bit in CPU identifies kernel/user mode
 - "trap" instruction
 - "return-from-trap" instruction
 - Save the registers before do the restricted operation (kernel stack)

OS @ run (kernel mode) Hardware

Program (user mode)

Run main()

... Call system call **trap** into OS

- Switching between processes
 - Cooperative approach
 - OS trusts the process to yield CPU properly
 - Non-cooperative approach
 - OS revokes the control of CPU periodically
 - Time interrupt
 - Scheduler

OS @ boot (kernel mode)	Hardware	
initialize trap table start interrupt timer	remember addresses of syscall handler timer handler start timer interrupt CPU in X ms	
OS @ run (kernel mode)	Hardware	Program (user mode)
Handle the trap Call switch() routine save regs(A) to proc-struct(A) restore regs(B) from proc-struct(B) switch to k-stack(B)	timer interrupt save regs(A) to k-stack(A) move to kernel mode jump to trap handler	Process A
return-from-trap (into B)	restore regs(B) from k-stack(B) move to user mode jump to B's PC	Process B

- Low-level mechanisims: summary
 - Direct execution
 - Limited direct execution
 - Switch between processes

Scheduling Policy

- High level intelligence
 - Scheduling policy
 - First In, First Out
 - Shortest job first
 - Shortest time to complete first
 - Round Roubin

CPU virtualization

- Summary of CPU virtualization
 - Low level mechanisms
 - A little hardware support goes a long way
 - High level mechanisms

- Process API
 - fork(), exec(), wait(), exit()
 - Create, execute, wait and terminate a process
 - May be the strangest API you've ever met

- fork()
 - Create a new process
 - Exactly copy the calling process
- The return code of fork() is different
 - In parent: fork() return the pid of the child
 - In child: fork() return 0
- Who will run first is not determined

- wait()
 - Wait for child to finish his job
 - The parent will not proceed until wait() return.
- waitpid()

- exec()
 - Execute a different program in child process
- A group of system calls:
 - execl, execv, execle, execve, execlp, execvp, fexecv

- Some Coding
 - fork
 - fork, wait
 - fork, wait, execvp

- What's happening behind fork()?
 - The child get a "copy" of parent's data space, stack, heap
 - the system call: clone()
 - "Copy-on-write"
 - Not really copy the data, but share the data with "read only" flag
 - If parent or child writes on a shared address, the kernel make a copy of that piece of memory only (usually a page)

- What's happening behind fork()?
 - File sharing
 - fd
 - File offsets

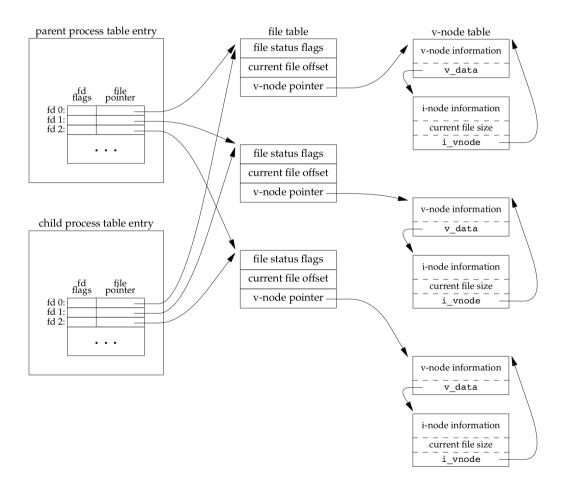


Figure 8.2 Sharing of open files between parent and child after fork

- What's happening behind fork()?
 - Other shared data:
 - User ID, group ID...
 - Current working directory
 - Environment
 - Memory mapping
 - Resources limits
 - ...

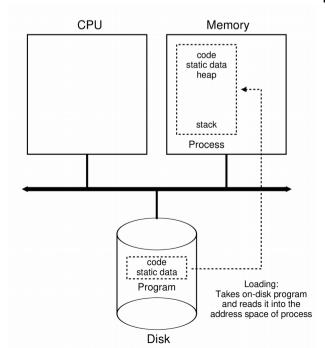
- What's happening behind exit()?
 - Close all fds, release all memory, ...
 - Inform the exit status to the parent process, which can be captured by wait()

- What's happening behind wait()?
 - The parent terminates first?
 - The init process (PID=0)
 - The child terminates first?
 - The kernel keeps a small amount of information for every terminating process
 - Available when the parent calls wait()
 - PID, termination status, the amount of CPU time
 - zombies



- What's happen behind wait()/waitpid()
 - wait(): block the caller until a child process terminates
 - waitpid(): wait which child, and some other options

- What's happening behind exec()?
 - Replace the current process with a new program from disk
 - Text, data, heap, stack
 - Start from the main() of that program



- Process API summary
 - fork(): create a new process
 - wait(): wait for a child
 - exit(): destroy a process
 - exec(): execute a program in child

Project2

- Implement your own shell
 - Use fork, wait, execvp
 - Also open, close, dup2

- Basic shell
 - Run your shell by: ./mysh
 - It will print a prompt:

```
mysh>
```

- You can type some commands

```
mysh> Is
```

Hit ENTER, the command will be executed

- Build-in Commands
 - When "mysh" execute a command, it will check weather it is a **build-in** or not.
 - For build-in commands, you should involve your implementation.
 - They are:
 - exit
 - wait
 - cd
 - pwd

- Redirection
 - Your shell should support redirection:

mysh> Is -I > output

- The file "output" contain the result of "Is -I"

- Background Jobs
 - Your shell should be able to run jobs in the background

mysh> Is &

Your shell will continue to work rather than wait.

- Batch mode
 - Your shell should be able to run in batch mode

```
./mysh batch_file
```

- Your shell will run the commands in batch_file
- E.g, "batch_file" containsls -lcat batch file

- Bonus: Pipe
 - The pipe connect the input/output of different commands

mysh> grep "hello" FILE | wc -l

How many lines have "hello"