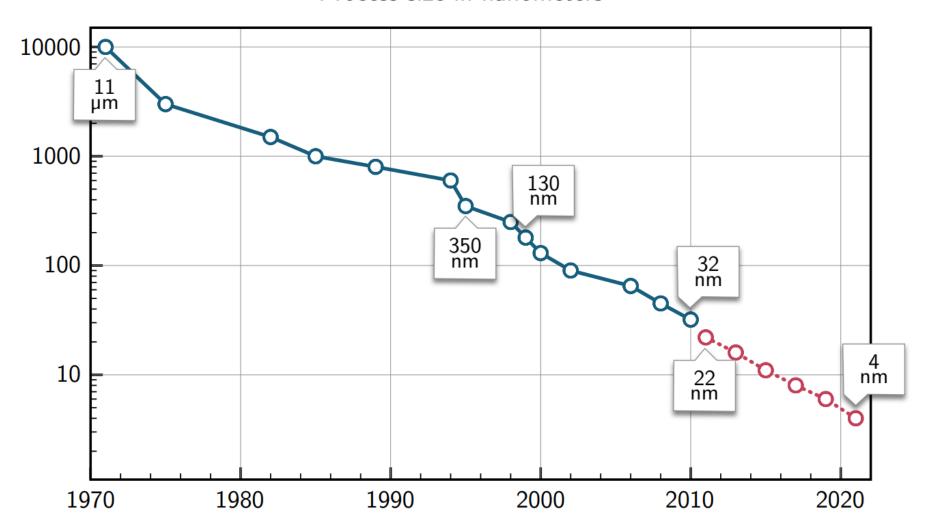
First Quantum Programming

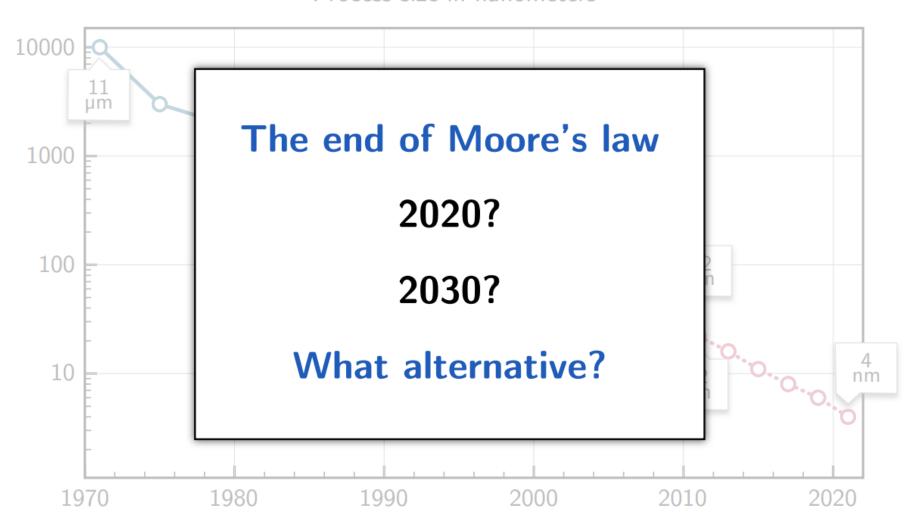
Cheng Lin Hong

Moore's Law

Process size in nanometers



Process size in nanometers



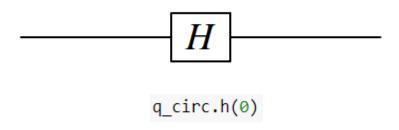
Qubits, Operators and Measurement

- A quantum computation is a collection of three elements
 - ① A quantum register or a set of quantum register.
 - ② A unitary matrix, which is used to execute a given quantum algorithm.
 - ③ Measurement to extract information we need.

- Quantum circuit model
 - Universal quantum computer

The superposition

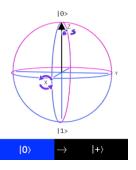
Hadamard can be used to create Quantum superposition



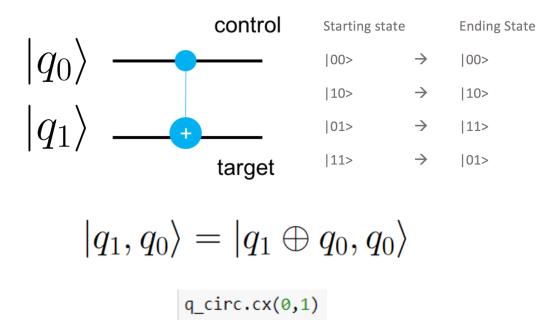
Hadamard

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \frac{|0\rangle + |1\rangle}{\sqrt{2}}$$

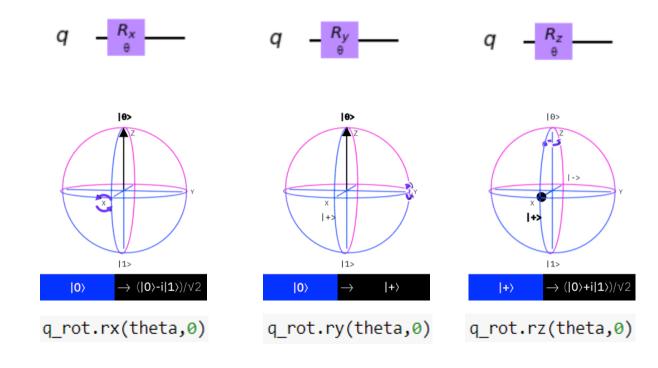
$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$



Controlled-NOT

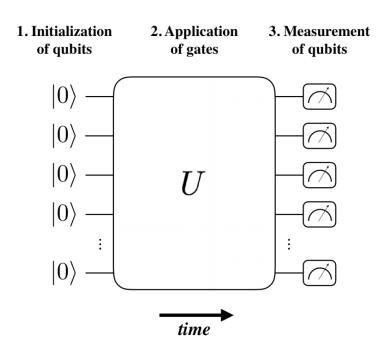


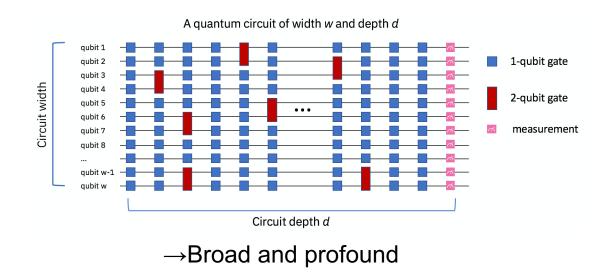
Single Qubits Rotations



Reference: https://quantum-computing.ibm.com/docs/circ-comp/q-gates

Quantum Circuit



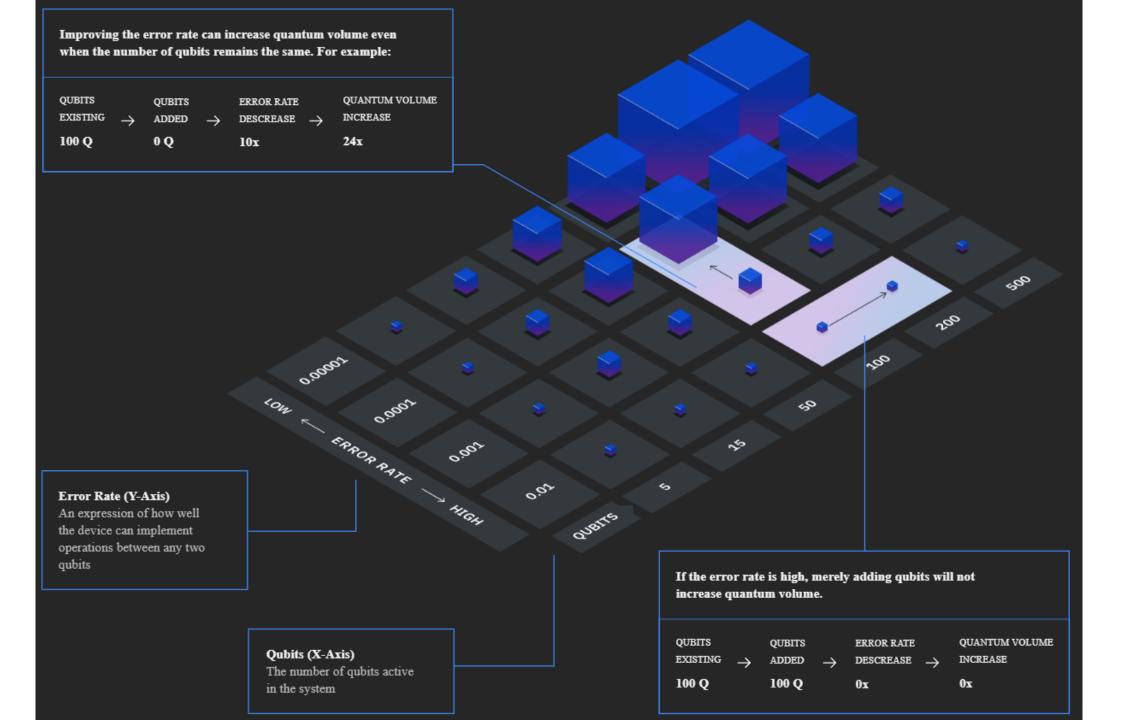


Source: https://qiskit.org/textbook/ch-states/introduction.html

iPhone 6 Bigger than bigger

Learn more > Watch the film (•)
Experience the keynote (•)





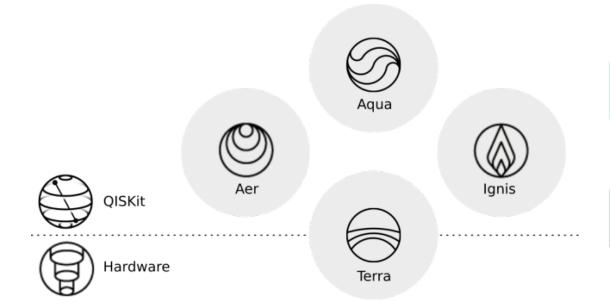
Qiskit Overview

Institution	IBM	
First Release	0.1 on March 7, 2017	
Open Source	Yes	
License	Apache-2.0	
HomePage	https://qiskit.org/	
Github	https://github.com/Qiskit	
Documentation	https://qiskit.org/documentation/	
OS	Mac, Windows, Linux	
Language	Python	
Quantum Language	OpenQASM	

Version Information

Qiskit Software	Version
Qiskit	0.17.0
Terra	0.12.0
Aer	0.4.1
Ignis	0.2.0
Aqua	0.6.5
IBM Q Provider	0.6.0

The Qiskit Elements



Terra, the 'earth' element, is the foundation on which the rest of the software lies.

Aer, the 'air' element, permeates all Qiskit elements. For accelerating development via simulators, emulators and debuggers

Aqua, the 'water' element, is the element of life. For building algorithms and applications.

Ignis, the 'fire' element, is dedicated to fighting noise and errors and to forging a new path

Source: https://qiskit.org/documentation/the_elements.html

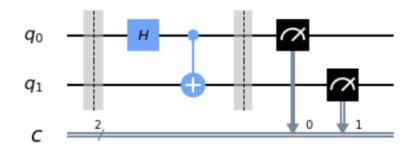
Qiskit Code Example

```
In [1]: from qiskit import QuantumCircuit

q_bell = QuantumCircuit(2, 2)
q_bell.barrier()
q_bell.h(0)
q_bell.cx(0, 1)
q_bell.barrier()
q_bell.measure([0, 1], [0, 1])

q_bell.draw(output='mpl',plot_barriers=True)
```

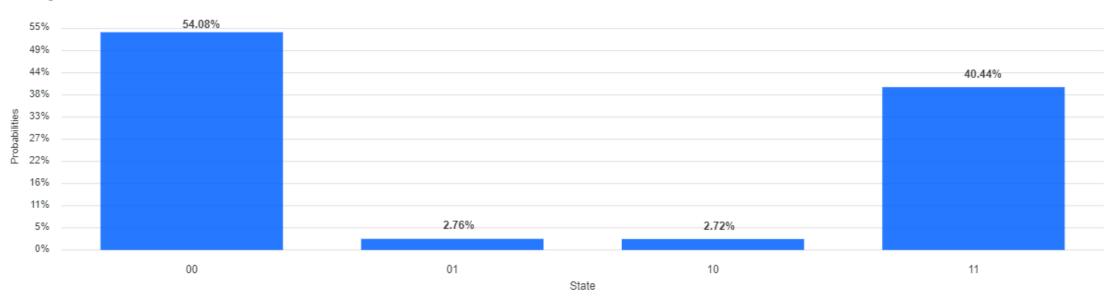
Out[1]:



```
In [2]: from qiskit import Aer, execute
        from qiskit.visualization import plot histogram
        backend = Aer.get_backend('qasm_simulator')
        job sim = execute(q bell, backend, shots=100000)
         sim result = job sim.result()
        print(sim result.get counts(q bell))
        plot histogram(sim result.get counts(q bell))
         {'11': 50254, '00': 49746}
Out[2]:
            0.60
                                                          0.503
                          0.497
            0.45
         Probabilities
00
00
            0.15
            0.00
                           8
```

Real QC

Histogram



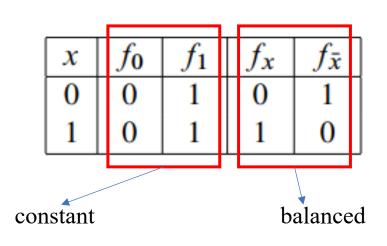
QA times Quantum Algorithms 101

The Deutsch Algorithm

The first to demonstrate quantum over classical computing.

Deutsch Problem:

Given a black box that implement some Boolean function $f: \{0,1\} \rightarrow \{0,1\}$. We are promised that the function is either constant or balanced.

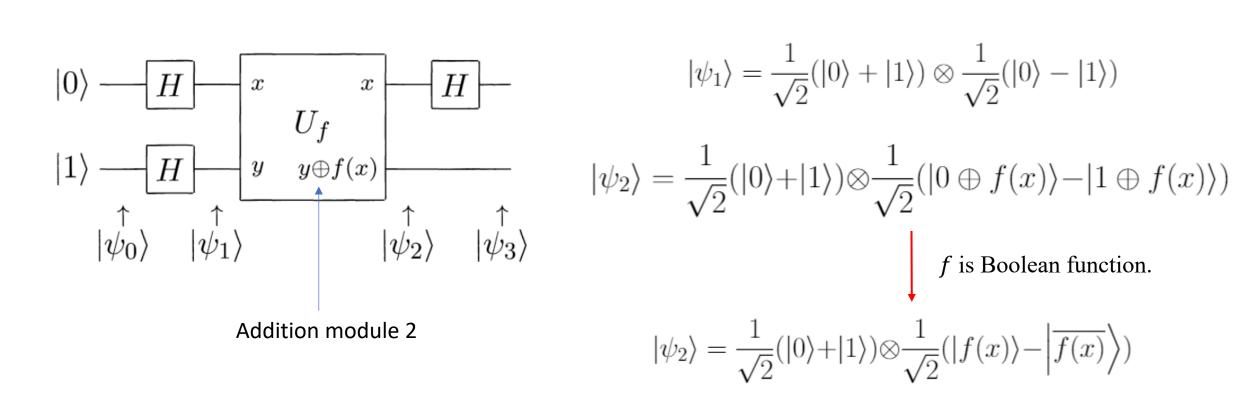


```
if f(0) = 0:
    if f(1) = 0:
        print("Constant")
    else:
        print("Balanced")
else:
    if f(1) = 0:
        print("Balanced")
    else:
        print("Constant")
```

require two function evaluations to figure out the answer.

The Deutsch Algorithm

We need only one query on a quantum computer!



 $|\psi_0\rangle = |0\rangle \otimes |1\rangle$

Phase Kick Back

• Useful trick in many quantum algorithm.

Consider Quantum Black box function f

$$|x\rangle - x \qquad x - |x\rangle$$

$$|y\rangle - y \qquad y \oplus f(x) - |y \oplus f(x)\rangle$$

$$|0\rangle$$
 — x x — U_f $y = y \oplus f(x)$ —

$$U_f\left(|x\rangle\otimes\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)\right)=|x\rangle\otimes\frac{1}{\sqrt{2}}(|0\oplus f(x)\rangle-|1\oplus f(x)\rangle)$$

$$f(x) = 0: \quad \frac{|0 \oplus f(x)\rangle - |1 \oplus f(x)\rangle}{\sqrt{2}} = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

$$f(x) = 1: \quad \frac{|0 \oplus f(x)\rangle - |1 \oplus f(x)\rangle}{\sqrt{2}} = \frac{|1\rangle - |0\rangle}{\sqrt{2}} = -\left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)$$

$$\frac{|0 \oplus f(x)\rangle - |1 \oplus f(x)\rangle}{\sqrt{2}} = (-1)^{f(x)} \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)$$

$$U_f\left(|x\rangle\otimes\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)\right)=(-1)^{f(x)}\left(|x\rangle\otimes\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)\right)$$

$$U_f\left((\alpha_0|0\rangle+\alpha_1|1\rangle)\otimes\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)\right)=\left((-1)^{f(0)}\alpha_0|0\rangle+(-1)^{f(1)}\alpha_1|1\rangle\right)\otimes\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)$$

The Deutsch Algorithm

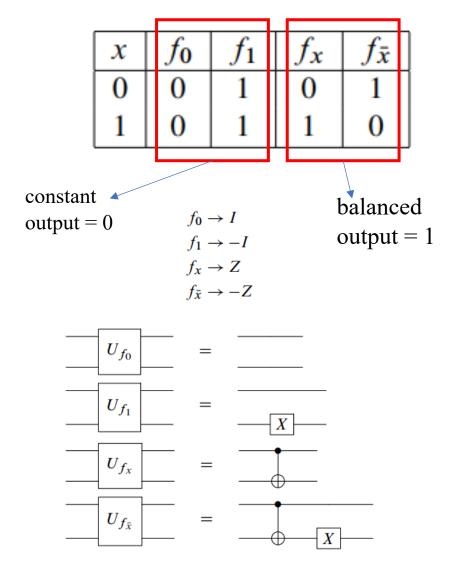
$$\begin{aligned} |\psi_0\rangle &= |0\rangle \otimes |1\rangle \\ |\psi_1\rangle &= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \\ |\psi_2\rangle &= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \otimes \frac{1}{\sqrt{2}}(|f(x)\rangle - \left|\overline{f(x)}\right\rangle) \\ U_f\left((\alpha_0|0\rangle + \alpha_1|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)\right) &= \left((-1)^{f(0)}\alpha_0|0\rangle + (-1)^{f(1)}\alpha_1|1\rangle\right) \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \end{aligned}$$

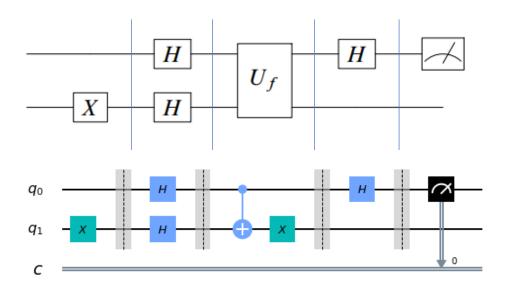
$$\begin{aligned} |\psi_{2}\rangle &= \frac{(-1)^{f(0)}}{\sqrt{2}} |0\rangle \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) + \frac{(-1)^{f(1)}}{\sqrt{2}} |1\rangle \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) \\ &= \left(\frac{(-1)^{f(0)} |0\rangle + (-1)^{f(1)} |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) \\ &= (-1)^{f(0)} \left(\frac{|0\rangle + (-1)^{f(0) \oplus f(1)} |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) \end{aligned}$$

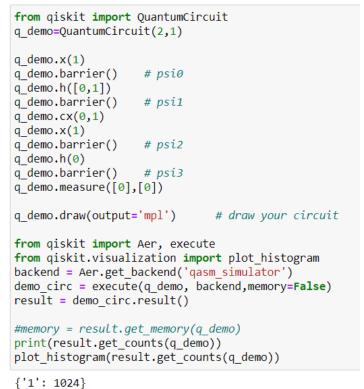
$$|\psi_{2}\rangle = \begin{cases} \pm \left[\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right] \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right] & \text{if } f(0) = f(1) \\ \pm \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right] \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right] & \text{if } f(0) \neq f(1). \end{cases}$$

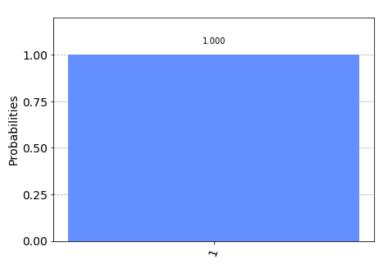
$$|\psi_3\rangle = \begin{cases} \pm |0\rangle \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right] & \text{if } f(0) = f(1) \\ \pm |1\rangle \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right] & \text{if } f(0) \neq f(1). \end{cases}$$

Qiskit Example









Extension: The Deutsch-Jozsa algorithm

• This time the function f is a function from n bits string to a bit.

The Deutsch–Jozsa Problem

Input: A black-box for computing an unknown function $f: \{0,1\}^n \to \{0,1\}$.

Promise: f is either a constant or a balanced function.

Problem: Determine whether f is constant or balanced by making queries to

f.

The Deutsch-Jozsa algorithm

$$|\psi_0\rangle = |0\rangle^{\otimes n} \otimes |1\rangle$$

The Deutsch-Jozsa Problem

Input: A black-box for computing an unknown function $f: \{0,1\}^n \to \{0,1\}$.

Promise: f is either a constant or a balanced function.

Problem: Determine whether f is constant or balanced by making queries to

f.

$$|\psi_1\rangle = \frac{1}{\sqrt{2^n}} \sum_x |x\rangle \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$|0\rangle \xrightarrow{f} H^{\otimes n} - \begin{bmatrix} x & x \\ U_f \\ y & y \oplus f(x) \end{bmatrix} - \begin{bmatrix} \uparrow \\ \psi_0 \end{pmatrix} + \begin{bmatrix} \uparrow \\ \psi_1 \end{pmatrix} + \begin{bmatrix} \uparrow \\ \psi_2 \end{pmatrix} + \begin{bmatrix} \uparrow \\ \psi_3 \end{bmatrix}$$

$$|\psi_2\rangle = \frac{1}{\sqrt{2^n}} \sum_x (-1)^{f(x)} |x\rangle \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$|\psi_3\rangle \qquad |\psi_3\rangle = \frac{1}{\sqrt{2^n}} \sum_{y} \sum_{x} (-1)^{f(x)+x\cdot y} |y\rangle \otimes \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

Summary

Algorithm: Deutsch-Jozsa

Inputs: (1) A black box U_f which performs the transformation $|x\rangle|y\rangle \to |x\rangle|y\oplus f(x)\rangle$, for $x\in\{0,\ldots,2^n-1\}$ and $f(x)\in\{0,1\}$. It is promised that f(x) is either *constant* for all values of x, or else f(x) is *balanced*, that is, equal to 1 for exactly half of all the possible x, and 0 for the other half.

Outputs: 0 if and only if f is constant.

Runtime: One evaluation of U_f . Always succeeds.

Procedure:

1.
$$|0\rangle^{\otimes n}|1\rangle$$
 initialize state

2. $\rightarrow \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right]$ create superposition using Hadamard gates

3. $\rightarrow \sum_{x} (-1)^{f(x)}|x\rangle \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right]$ calculate function f using U_f

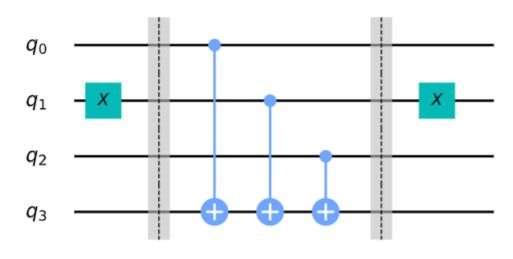
4. $\rightarrow \sum_{z} \sum_{x} \frac{(-1)^{x \cdot z + f(x)}|z\rangle}{\sqrt{2^n}} \left[\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right]$ perform Hadamard transform

5. $\rightarrow z$ measure to obtain final output z

Qiskit Example

Balanced Oracle

Outputs 0	Outputs 1
001	000
010	011
100	101
111	110



Now, your turn

Improvement?

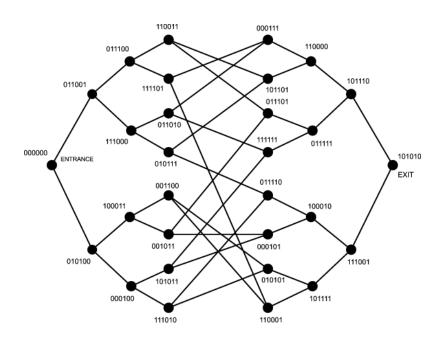
- The worst case in classical take us $\frac{2^n}{2} + 1$
- In quantum we need only one query But, if there exist error? How about randomized algorithm?

• That's why we have Bernstein-Vazirani, Simon's algorithm.

BQP class

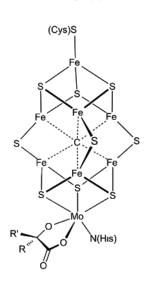
- BQP: Bounded-error quantum polynomial time
 - What we focus on for QC

Quantum Walks

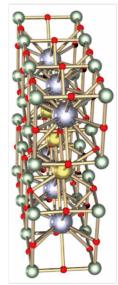


Quantum Simulation

Promising applications of quantum simulators



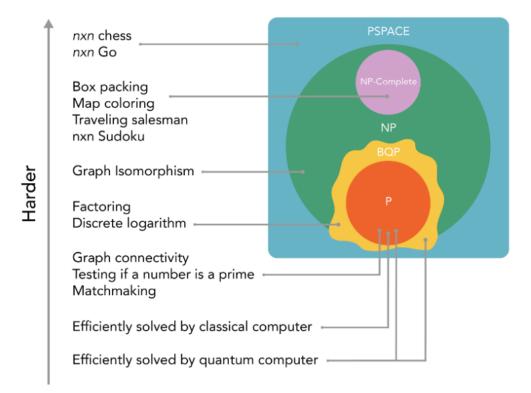
Complex molecules applications in nitrogen fixation, fuel cells, drug developement



Strongly correlated materials explaining high-temperature superconductors and engineering quantum materials

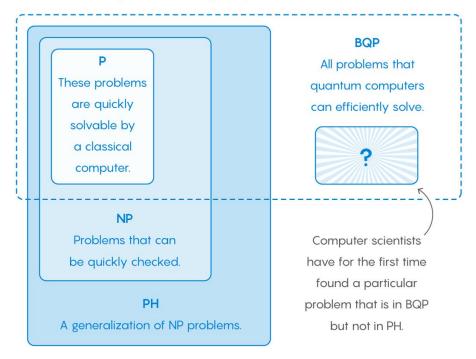
Where Quantum Computers Fits In

Example Problems



A New Island on the Complexity Map

What can a quantum computer do that any possible classical computer cannot? Computer scientists have finally found a way to separate two fundamental computational complexity classes.



The Zoo of Quantum Algorithms

Algebraic and Number Theoretic Algorithms

- Algorithm: Factoring (Super-polynomial)
- Algorithm: Discrete-log (Super-polynomial)
- Algorithm: Verifying Matrix Products (Polynomial)

Oracular Algorithms

- Algorithm: Searching (Polynomial)
- **Algorithm:** Hidden Shift (Super-polynomial)
- Algorithm: Counterfeit
 Coins

 (Polynomial)

Approximation and Simulation Algorithms

- Algorithm: Quantum Simulation (Super-polynomial)
- Algorithm: Quantum Approximate Optimization (Super-polynomial)

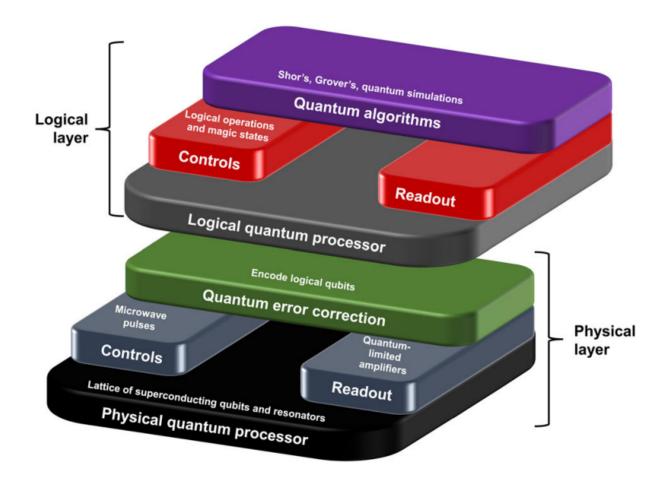
Optimization, Numerics, and Machine Learning

- Algorithm: Machine Learning (Varies)
- Algorithm: Adiabatic Algorithms (Unknown)
- Algorithm: Quantum Dynamic Programming (Polynomial)

Why have we not shown that?

- All of above QA require a quantum circuit implemented with near zero error rates
 - → Need Quantum Error Correction

Quantum Computing Stack



Source: Gambetta, J.M., Chow, J.M. & Steffen, M. Building logical qubits in a superconducting quantum computing system. npj Quantum Inf 3, 2 (2017).

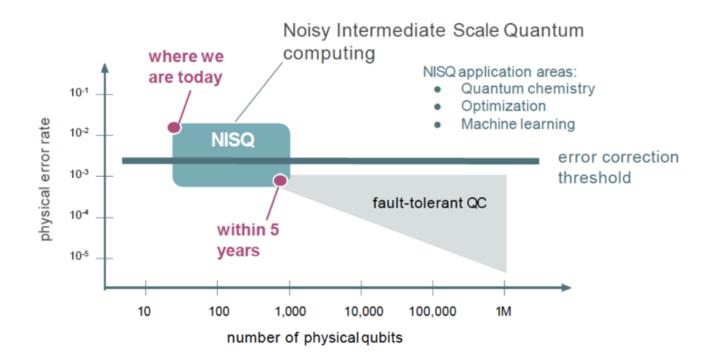
Where are we now?

Resource	Current Capability	Estimate for 2048-bit Shor's factorization*
Operational Fidelity	~99%	99.99%
# of physical gate operations	100s	1.5×10^{21}
# of physical qubits	53	98,000,000
# of FT logical operations (perfect operations)	0	450,000,000,000
# of FT logical qubits (ideal qubits)	0	12000

^{*}N. C. Jones, R. Van Meter, A. G. Fowler, P. L. McMahon, J. Kim, T. D. Ladd, and Y. Yamamoto, Layered Architecture for Quantum Computing, Phys. Rev. X 2, 031007 (2012).

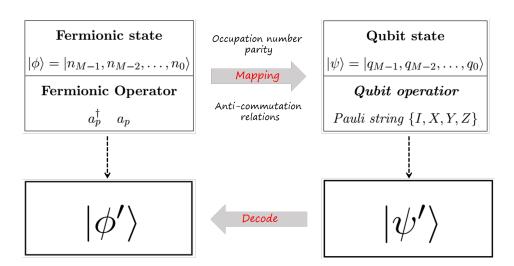
Near term

- Noisy Intermediate-Scale Quantum (NISQ)
 - we don't expect to be able to execute a circuit that contains many more than about 1000 gates



Application: Quantum meets CS

• When simulating Fermionic Problem



PHYSICAL REVIEW A 95, 032332 (2017)

Operator locality in the quantum simulation of fermionic models

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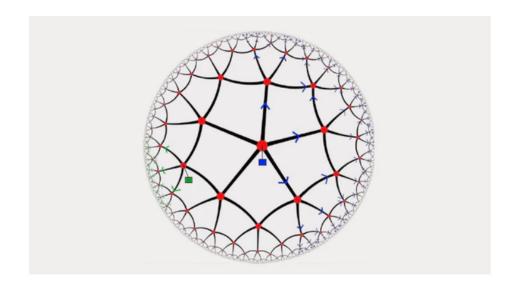
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Department of Physics and Astronomy, Dartmouth College, 6127 Wilder Laboratory, Hanover, New Hampshire 03755, USA (Received 6 February 2017; published 29 March 2017)

Simulating fermionic lattice models with qubits requires mapping fermionic degrees of freedom to qubits. The simplest method for this task, the Jordan-Wigner transformation, yields strings of Pauli operators acting on an extensive number of qubits. This overhead can be a hindrance to implementation of qubit-based quantum simulators, especially in the analog context. Here we thus review and analyze alternative fermion-to-qubit mappings, including the two approaches by Bravyi and Kitaev and the auxiliary fermion transformation. The Bravyi-Kitaev transform is reformulated in terms of a classical data structure and generalized to achieve a further locality improvement for local fermionic models on a rectangular lattice. We conclude that the most compact encoding of the fermionic operators can be done using ancilla qubits with the auxiliary fermion scheme. Without introducing ancillas, a variant of the Bravyi-Kitaev transform provides the most compact fermion-to-qubit mapping for Hubbard-like models.

Application: Doing Physics with Quantum Computers

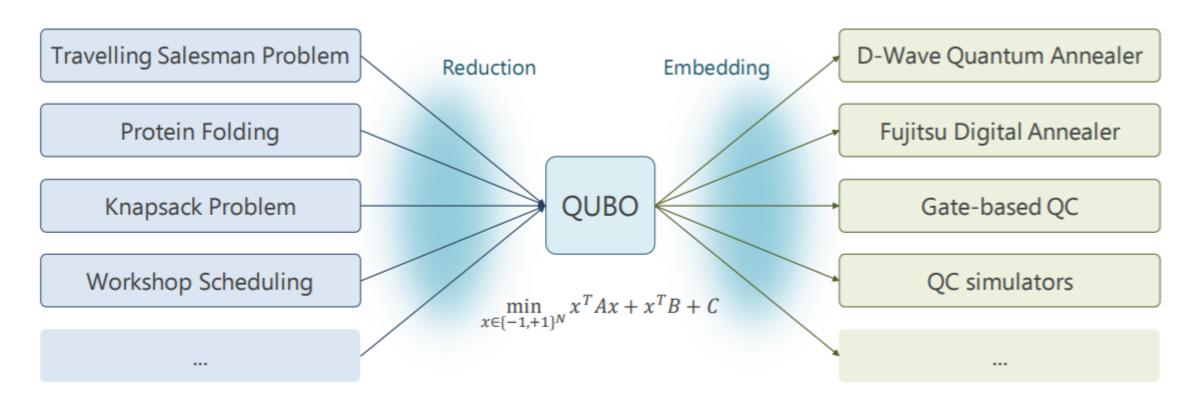
• It from Qubit: Projects



Developments over the past ten years have shown that major advances in our understanding of quantum gravity, quantum field theory and other aspects of fundamental physics can be achieved by bringing to bear insights and techniques from quantum information theory. Nonetheless, fundamental physics and quantum information theory remain distinct disciplines and communities, separated by significant barriers to communication and collaboration. Funded by a grant from the Simons Foundation, "It from Qubit" is a large-scale effort by some of the leading researchers in both communities to foster communication, education and collaboration between them, thereby advancing both fields and ultimately solving some of the deepest problems in physics. The overarching scientific questions motivating the collaboration include:

- Does spacetime emerge from entanglement?
- Do black holes have interiors?
- Does the universe exist outside our horizon?
- What is the information-theoretic structure of quantum field theories?
- Can quantum computers simulate all physical phenomena?
- How does quantum information flow in time?

Application: Quantum + NP = NO problems?



QUBO reduction from NP-hard problems