CS350: Data Structures Doubly Linked Lists

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Doubly Linked Lists

- Adds an additional pointer to a the list nodes that points to the previous node in the list
- Traversable in either the forward or backward direction
- Resolves the issue of removing the last node in the linked list
 - Becomes a O(1) operation as opposed to O(N)



Doubly Linked List Operations

Basic operations include:

```
-insert / add
-remove / removeFirst / removeLast
```

Additional operations may include:

```
-getFirst / getLast
```

- -find
- isEmpty
- makeEmpty

 Basic implementation uses head and tail pointers that points to the first node and the last node in the list

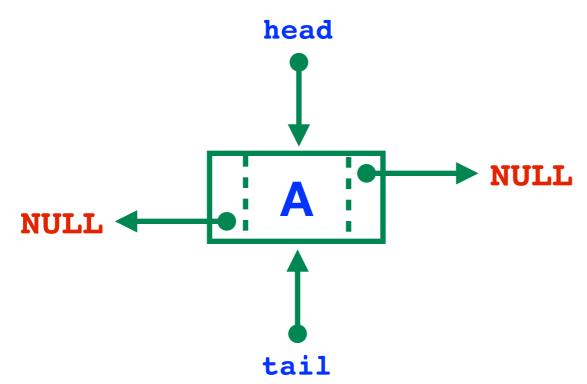
- Both pointers points to **NULL** upon initialization when no nodes exist in the list



 Depending on implementation, insertion may take place at the head of the list, at the tail of the list, or at some other specified node

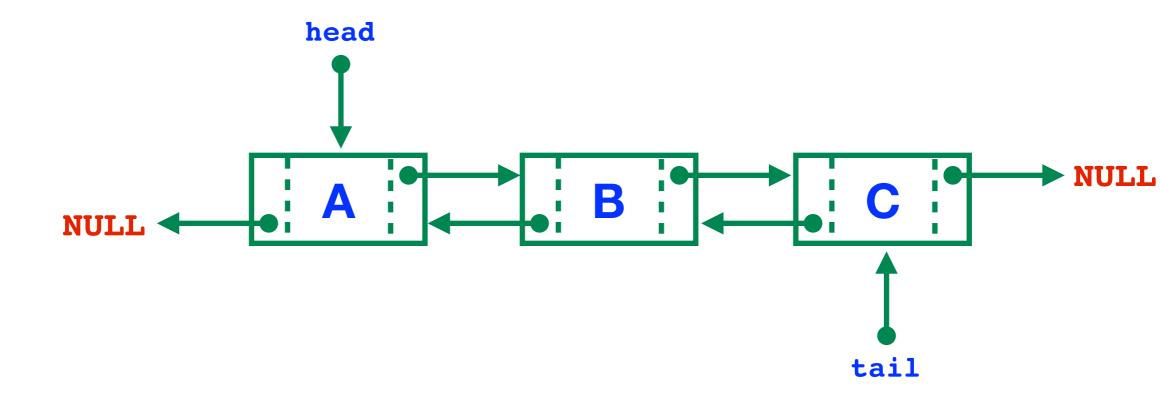
After Inserting a Node

- In this illustration, nodes are inserted at the tail end (nodes can be inserted at either the head or the tail)
 - After inserting the first node, A
 - The **head** and the **tail** pointers are reassigned to point to the first node



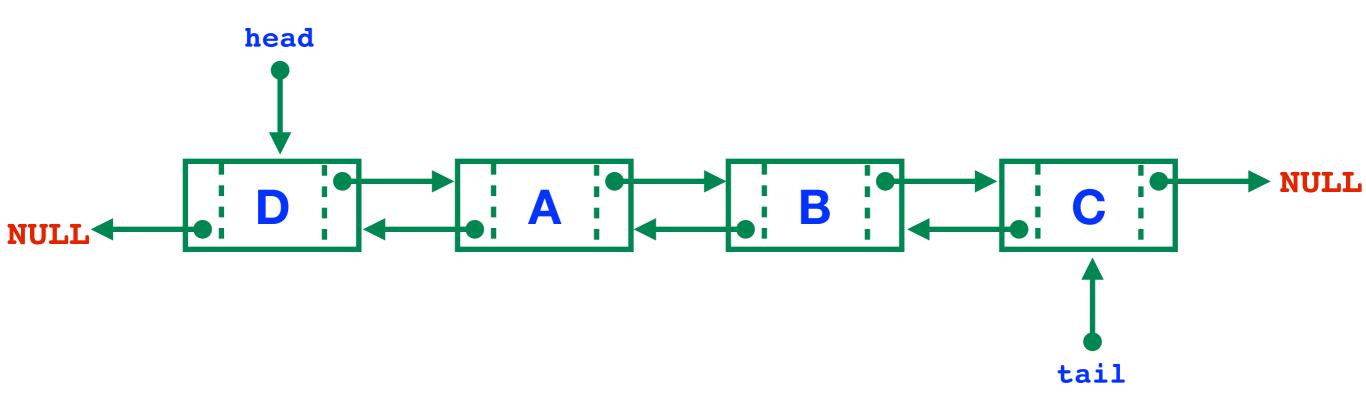
After Inserting Additional Nodes

- In this illustration, nodes are inserted at the tail end (nodes can be inserted at either the head or the tail)
 - After inserting nodes in the sequence A, B, C
 - The tail pointer advances with each insertion at the tail end



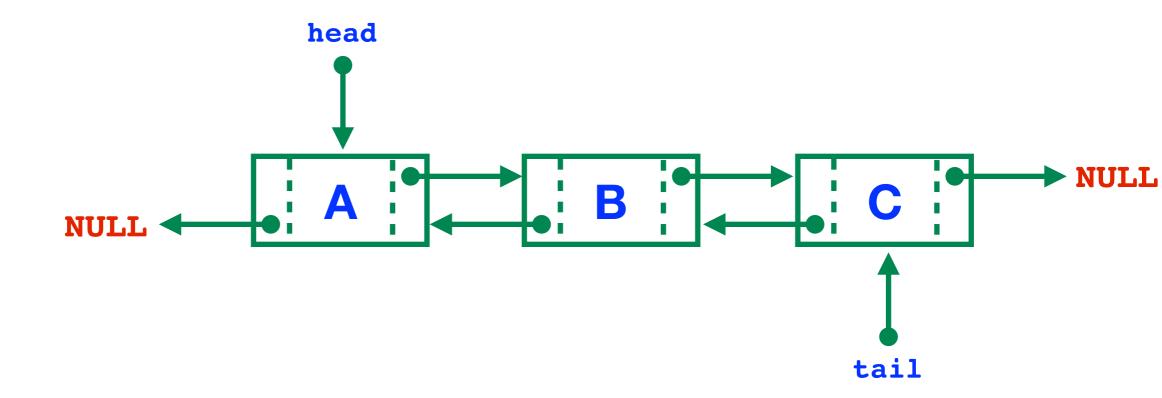
After Inserting Additional Nodes

- In this illustration, nodes are inserted at the head
 - After inserting the node D
 - The head pointer retreats with each insertion at the head



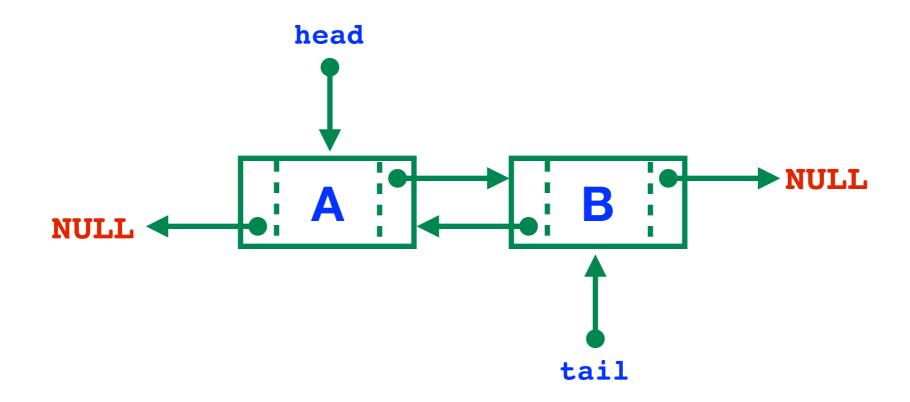
After Removing the First Node

- Nodes can be removed from the head of the list or the tail
 - After removing a single node from the head of the list
 - The **head** pointer advances with a removal from the **head** of the list



After Removing the Last Node

- Nodes can be removed from the head of the list or the tail
 - After removing a single node from the tail of the list
 - The tail pointer retreats with a removal from the tail of the list



```
public class DLinkedListNode<E> {
   public E data;
   public DLinkedListNode<E> next;
   public DLinkedListNode<E> prev;
}
```

```
// Inserts at the tail of the list
public void append (E data) {
   DLinkedListNode<E> newNode = new DLinkedListNode<E>();
   newNode.data = data; // assign data to newNode
   newNode.prev = tail;
   tail.next = newNode;
   tail = newNode;
}
```

This method is oversimplified, what happens if this is called when the list is empty?

Fixed append method

```
// Inserts at the tail of the list
public void append (E data) {
   DLinkedListNode<E> newNode = new DLinkedListNode<E>();
   newNode.data = data; // assign data to newNode
   if (isEmpty()) {
     head = tail = newNode;
   } else {
     newNode.prev = tail;
     tail.next = newNode;
     tail = newNode;
   }
}
```

```
// Inserts at the head of the list
public void prepend (E data) {
   DLinkedListNode<E> newNode = new DLinkedListNode<E>();
   newNode.data = data; // assign data to newNode
   newNode.next = head;
   head.prev = newNode;
   head = newNode;
}
```

This method is oversimplified, what happens if this is called when the list is empty?

```
// Removes node from head of list and returns its value
public E removeFirst() {
   if (head != null) {
      E nodeData = head.data;
      head.next.prev = null;
      head = head.next;
      return nodeData;
   } else {
      return null;
   }
}
```

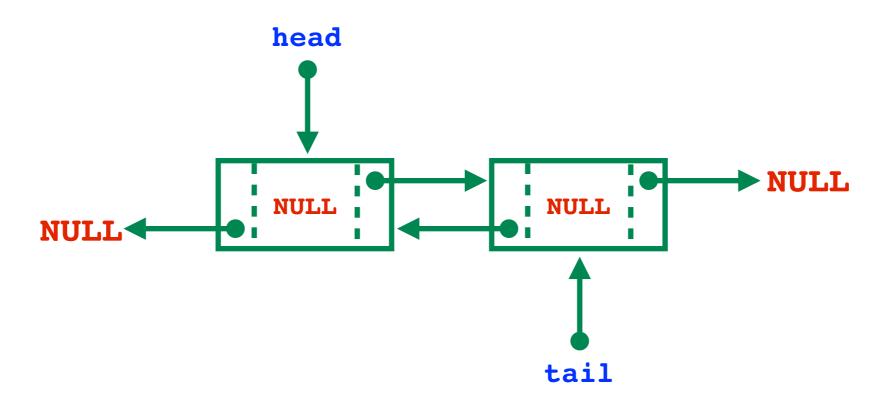
Considerations for Linked List Implementation

- Implementation as previously shown requires error checking in the insertion and removal methods to check for edge cases (i.e. checking for an empty list)
- To improve the speed of operations, it is possible to remove these tests
 - Tradeoff: speedup comes at the expense of one/two additional 'dummy' nodes in the Linked List
- Idea: create one or two dummy nodes (sentinel nodes) that exists in the linked list at ALL times
 - Eliminates the need to always check for NULL
 - Generalized the insertion and removal methods

Doubly Linked List with Two Sentinel Nodes

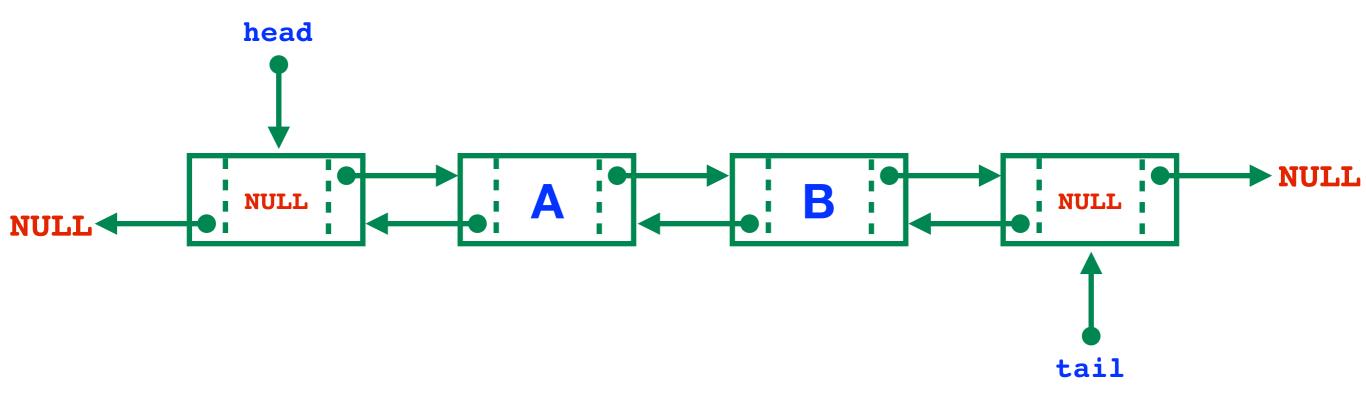
- To check for empty: (head.next == tail);
- When traversing the list check to see if current position points to either the head or the tail to determine if at the end of the list

Empty list



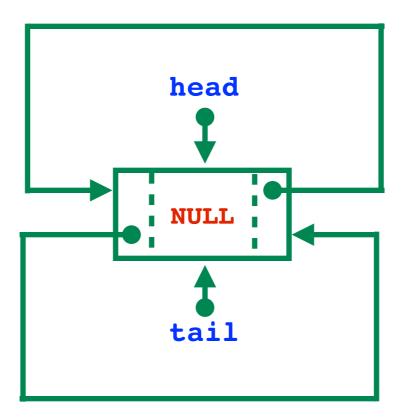
Doubly Linked List with Two Sentinel Nodes

List with two nodes



Circular Doubly Linked List with a Sentinel Node

Empty list



How should is Empty () be implemented?

Circular Doubly Linked List with a Sentinel Node

List with two nodes

