CS420: Operating Systems I/O Systems

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Overview

- I/O management is a major component of operating system design and operation
 - Important aspect of computer operation
 - I/O devices vary greatly
 - Various methods to control them
 - Performance management
 - New types of devices frequent

Ports, buses, device controllers connect to various devices

- Device drivers encapsulate device details
 - Present uniform device-access interface to I/O subsystem

I/O Hardware

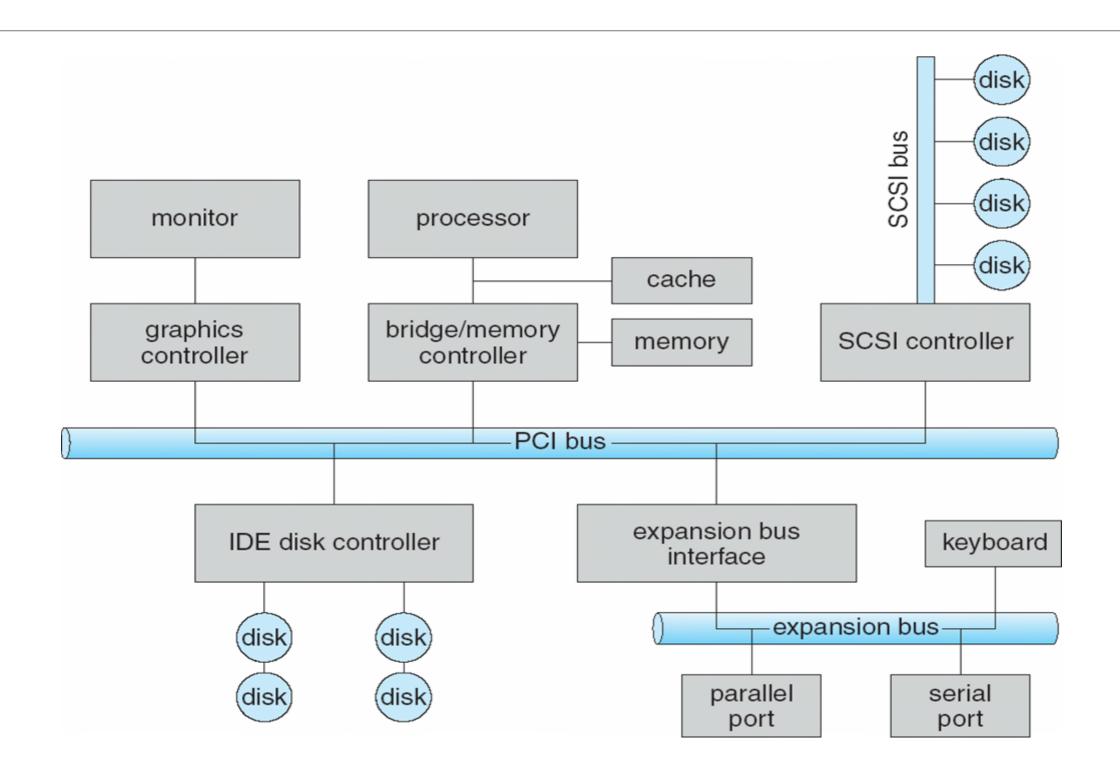
Incredible variety of I/O devices

- Storage (hard disk drives, flash drives, CDs, DVD, etc.)
- Transmission (wireless network cards, Ethernet cards, modems ;-)
- Human-interface (mouse, keyboard, trackpad, etc.)

Common concepts – signals from I/O devices interface with computer

- Port connection point for device
- Bus daisy chain or shared direct access
- Controller (host adapter) electronics that operate port, bus, device
 - Sometimes integrated into device
 - Sometimes separate circuit board
 - · Contains processor, microcode, private memory, bus controller, etc

A Common PC Bus Structure



I/O Hardware (Cont.)

- I/O instructions control devices
- Devices usually have registers where device driver places commands, addresses, and data to write, or read data from registers after command execution
 - Data-in register, data-out register, status register, control register
 - Registers are typically 1-4 bytes, or FIFO buffer
- Devices have addresses, used by
 - Direct I/O instructions
 - Memory-mapped I/O
 - Device data and command registers mapped to processor address space
 - Especially useful for large address spaces (graphics)

Device I/O Port Locations on PCs (partial list)

I/O address range (hexadecimal)	device	
000-00F	DMA controller	
020–021	interrupt controller	
040–043	timer	
200–20F	game controller	
2F8–2FF	serial port (secondary)	
320–32F	hard-disk controller	
378–37F	parallel port	
3D0-3DF	graphics controller	
3F0-3F7	diskette-drive controller	
3F8-3FF	serial port (primary)	

Polling

- One option when communicating between the host and a controller is polling
 - For each byte of I/O
 - Read busy bit from status register until 0 (busy-waiting)
 - Host sets read or write bit and if write copies data into data-out register
 - Host sets command-ready bit
 - Controller sets busy bit, executes transfer
 - Controller clears busy bit, error bit, command-ready bit when transfer done
 - Step 1 is busy-wait cycle to wait for I/O from device
 - Reasonable if device is fast
 - But inefficient if device slow
 - If CPU switches to another task, data may be lost (buffer overflows)

Interrupts

Polling can happen in 3 instruction cycles

- Read status, logical-and to extract status bit, branch if not zero
- How to be more efficient if non-zero infrequently?

CPU Interrupt-request line triggered by I/O device

- Checked by processor after each instruction

Interrupt handler receives interrupts

- Maskable to ignore or delay some interrupts -- important when kernel is modifying shared data structures

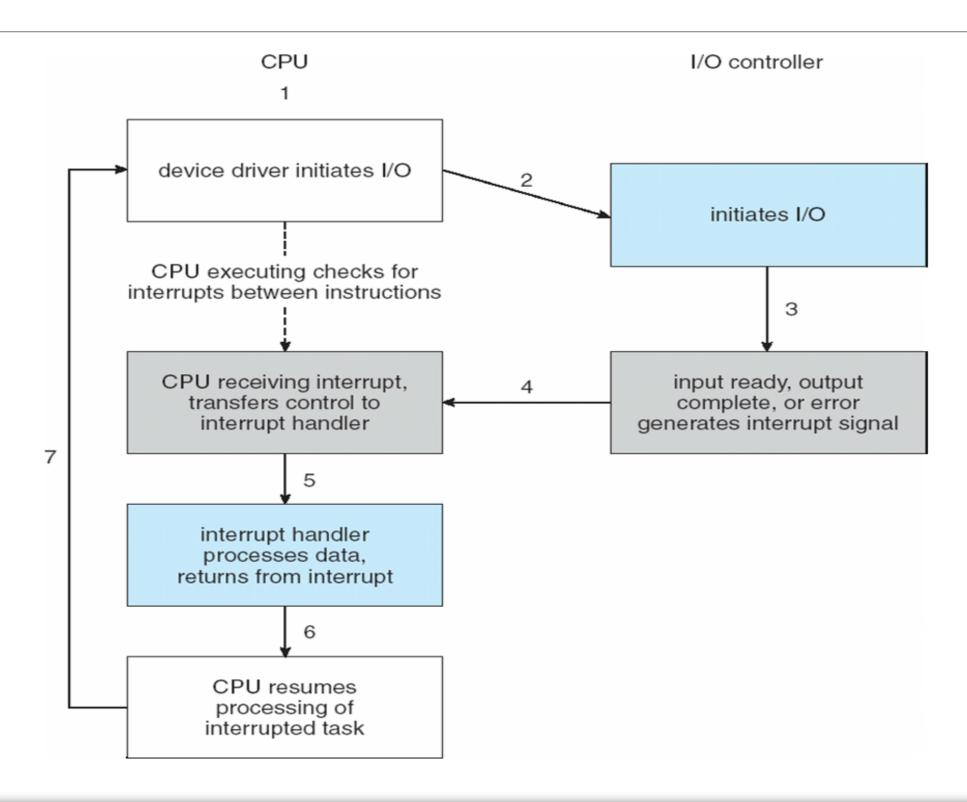
Interrupt vector is used to dispatch interrupt to correct interrupt handler

- Context switch at start and end
- Priority based
- Some interrupts in vector are non-maskable
- May need to chain interrupts if more than one device at same interrupt number

Intel Pentium Processor Event-Vector Table

vector number	description	
0	divide error	
1	debug exception	
2	null interrupt	
3	breakpoint	
4	INTO-detected overflow	
5	bound range exception	
6	invalid opcode	
7	device not available	
8	double fault	
9	coprocessor segment overrun (reserved)	
10	invalid task state segment	
11	segment not present	
12	stack fault	
13	general protection	
14	page fault	
15	(Intel reserved, do not use)	
16	floating-point error	
17	alignment check	
18	machine check	
19–31	(Intel reserved, do not use)	
32–255	maskable interrupts	

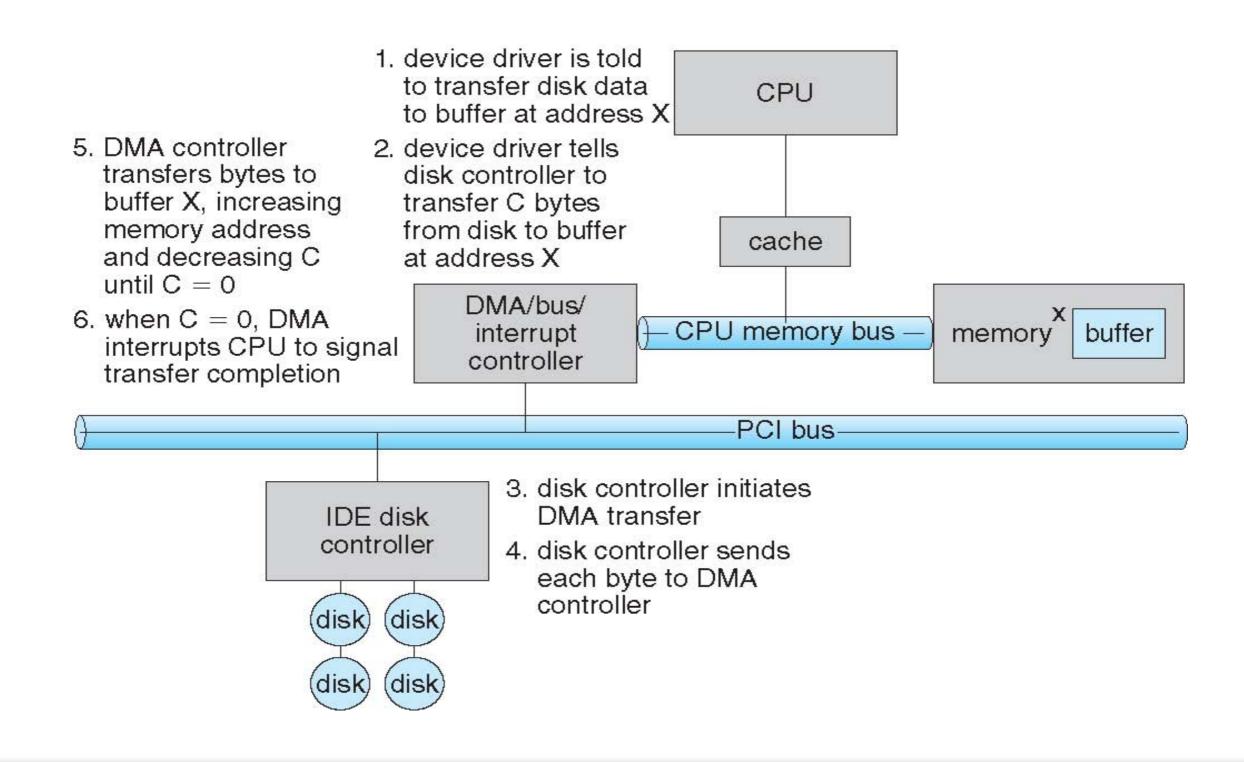
Interrupt-Driven I/O Cycle



Direct Memory Access

- Used to avoid programmed I/O (one byte at a time) for large data movement
- Requires DMA controller
- Bypasses CPU to transfer data directly between I/O device and memory
 - Don't want to waste CPU cycles to do something as simple as transferring data
- OS writes DMA command block into memory
 - Source and destination addresses
 - Read or write mode
 - Count of bytes
 - Writes location of command block to DMA controller
 - Bus mastering of DMA controller grabs memory bus from CPU (CPU CANNOT use memory bus during this time)
 - When done, interrupts to signal completion

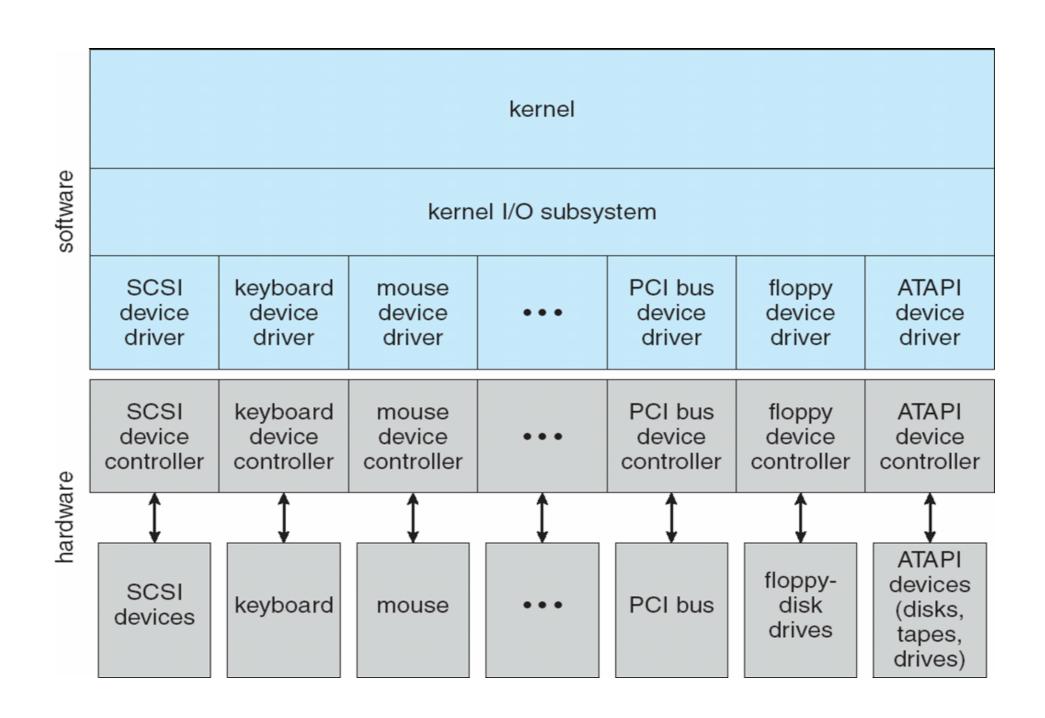
Six Step Process to Perform DMA Transfer



Application I/O Interface

- · It is useful for the operating system to be able to treat devices in a standard, uniform way
 - Too many different types of devices for OS designers to handle each individually
 - Devices types are grouped and the OS provides a standard interface to those groups
- Device drivers specific to each device implement one of the standard interfaces provided by the OS
- Device-driver layer hides differences among I/O controllers from kernel
- New devices using already-implemented protocols to communicate need no extra work
- Each OS has its own I/O subsystem structures and device driver frameworks
- Devices vary in many dimensions
 - Character-stream or block transfer
 - Sequential or random-access
 - Synchronous or asynchronous transfers (or both)
 - Sharable or dedicated (can it be used concurrently by multiples threads/processes?)
 - Speed of operation
 - read-write, read only, or write only

A Kernel I/O Structure



Characteristics of I/O Devices

aspect	variation	example
data-transfer mode	character block	terminal disk
access method	sequential random	modem CD-ROM
transfer schedule	synchronous asynchronous	tape keyboard
sharing	dedicated sharable	tape keyboard
device speed	latency seek time transfer rate delay between operations	
I/O direction	read only write only read–write	CD-ROM graphics controller disk

Characteristics of I/O Devices (Cont.)

Subtleties of devices handled by device drivers

- Broadly I/O devices can be grouped by the OS into:
 - Block I/O
 - Character I/O (Stream)
 - Memory-mapped file access
 - Network sockets

Block and Character Devices

Block devices include disk drives

- Commands include read, write, seek
- Raw I/O, direct I/O, or file-system access
- Memory-mapped file access possible
- DMA

- Character devices include keyboards, mice, serial ports
 - Commands include get, put
 - Libraries layered on top allow additional functionality

Network Devices

- Different enough from block and character devices to have its own interface
- Unix and Windows NT/9x/2000 include socket interface
 - Commands include open, close, send, receive
 - Separates network protocol from network operation

Clocks and Timers

Provide current time, elapsed time, timer

- Programmable interval timer used for timings, periodic interrupts
 - Determining when a time-slice for a thread is complete
 - Periodically flushing disk buffers