CS420: Operating Systems

Threads

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Threads

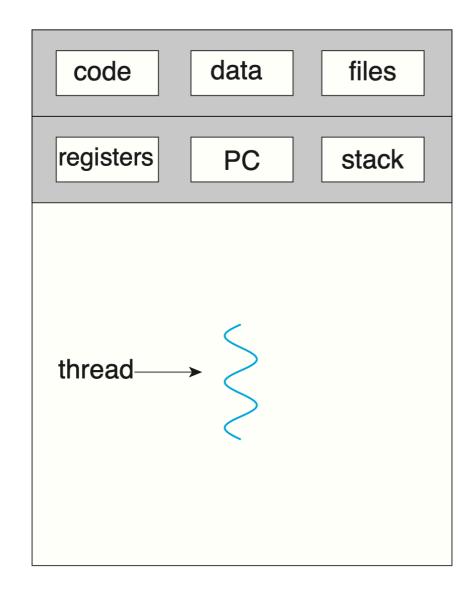
- A thread is a basic unit of processing
 - Has the following components:
 - Thread ID
 - Program counter
 - Register set
 - Stack
 - Shares some resources with other threads in same process
 - Code section
 - Data section
 - OS Resources (e.g. open files, signals, etc.)

Threads

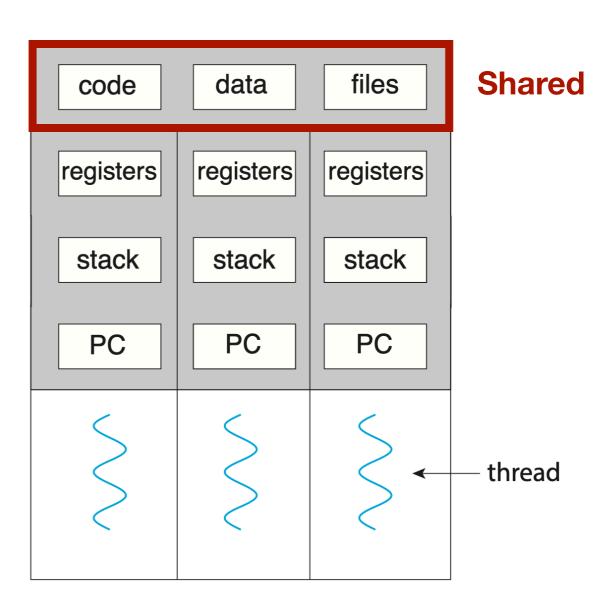
- A heavyweight process is a process that has a single thread of control
 - Can only perform a single task at a time

- A multi-threaded process is a process that has multiple threads of control
 - Can perform more than one task at a time
 - Render images (e.g. thumbnails for an image library)
 - · Fetch data
 - Update display
 - Check spelling

Single and Multithreaded Processes



single-threaded process



multithreaded process

Thread vs Process

Processes

- Independent units of execution
- Each contains its own state information
- Each contains its own address space
- Interact with each other through various IPC mechanisms

Threads (within the same process)

- Share the same state
- Share the same memory space
- Share the same variables
- Can communicate directly through shared variables
- Share signal handling

Benefits of Multithreaded Programming

Responsiveness

- Interactive applications are more responsive when multithreaded (e.g. a thread for the GUI, another for socket, a third for rendering, etc.)

Resource Sharing

- Unlike processes, threads share memory (code and data sections) and resources

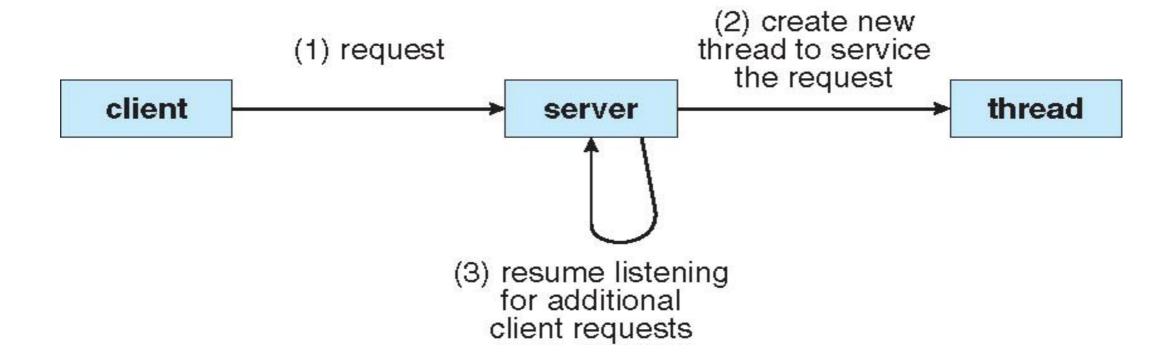
Economy

- Since threads share resources, creating threads and switching between them is more efficient than processes

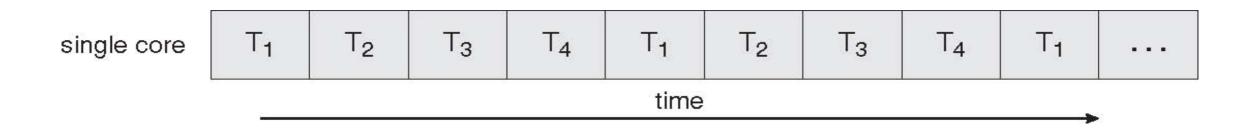
Scalability

 Multithreading allows for increased parallelism on multicore systems as each thread can run on a different CPU core (kernel threads only)

Multithreaded Server Architecture

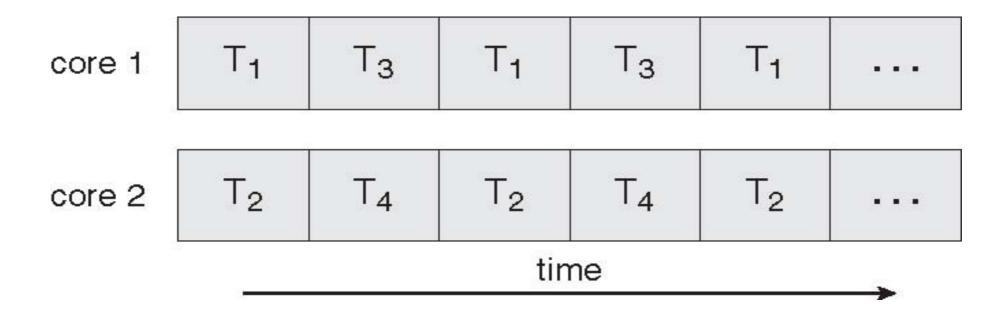


Concurrent Execution on a Single-core System



- Only a single thread can execute at a time
- Threads are interleaved so each gets time on the processor

Parallel Execution on a Multicore System



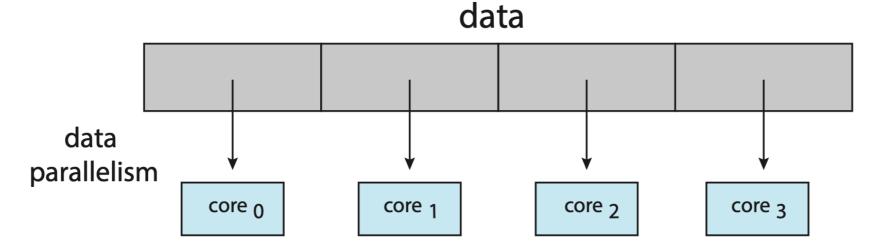
- With multiple cores, threads can be divided over the cores and run in parallel
- May still interleave threads if not enough cores are available for all of the threads

Multicore Programming

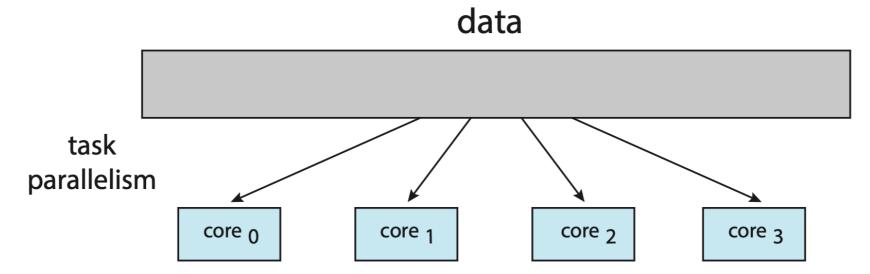
- The performance of a process can be increased by properly threading the process to take advantage of modern multicore CPUs
- Multicore systems have challenges not faced in a single-core/single-threaded environment
 - Dividing activities How can an application be divided into separate, concurrent tasks?
 - Balance How can those tasks be divided in such a way that each does an equal amount of work?
 - Data splitting Can the data for those task be divided for processing on separate CPU cores?
 - Data dependency Are there data dependencies between different tasks?
 - Testing and debugging What is the best way to debug a multithreaded program with many different execution paths?

Data Parallelism and Task Parallelism

Data Parallelism - distribute subsets of the same data across multiple cores



 Task Parallelism - distribute threads across multiple cores (each performing a unique operation)



Thread Support

Threads may be supported at different levels of the OS

- User threads

- Supported above the kernel
- Managed and scheduled without kernel support
- Main user thread libraries currently in use POSIX PThreads / Java threads

- Kernel threads

- Supported by the kernel/operating system
- Managed and scheduled by the kernel/operating system
- Most modern operating systems support kernel threads (e.g. Windows 2000/XP/.../10/11, Solaris, Linux, macOS)

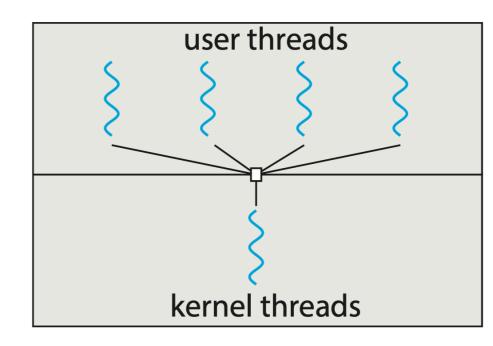
Multithreading Models

There must be a relationship between user level threads and kernel threads

- Different models of threading exist to define this relationship
 - Many-to-One
 - One-to-One
 - Many-to-Many

Many-to-One

- Many user-level threads mapped to single kernel thread
 - Examples: Solaris Green Threads, GNU Portable Threads
 - Not many systems use this model
- Thread management is done in user space
- Entire process will block if any single thread blocks (no other threads will run)
- Unable to run multiple user-level threads in parallel on a multiprocessor system
 - Not very common anymore

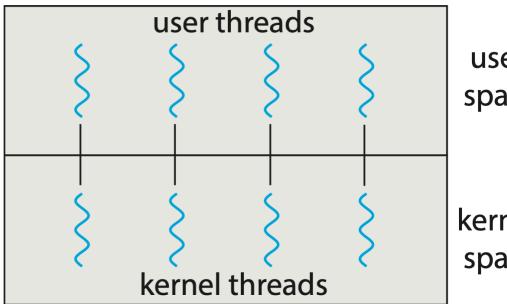


user space

kernel space

One-to-One

- Each user-level thread maps to kernel thread
 - Examples: Windows NT/XP/2000, Linux, Solaris 9 and later
- Allows more concurrency
 - A thread can run when another thread has made a blocking system call
 - Multiple user-level threads can run in parallel on multiprocessor systems
- Downside for each thread created, a corresponding kernel thread must also be created
 - Large number of kernel threads may degrade system performance



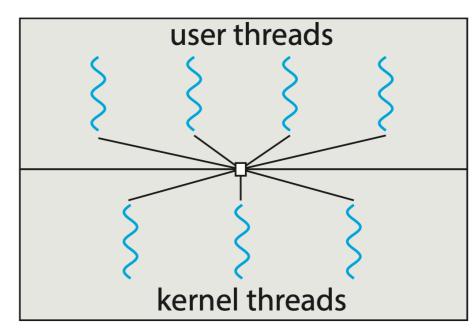
user space

kernel space

Many-to-Many Model

- Many user-level threads are mapped to a smaller number of kernel threads
 - Avoids blocking of threads when a single thread makes a blocking system call

- Allows the operating system to create a sufficient number of kernel threads
 - OS may allocate more kernel threads on a machine with more CPU cores
- Reduces the overhead associated with too many kernel threads as was present in the one-to-one model



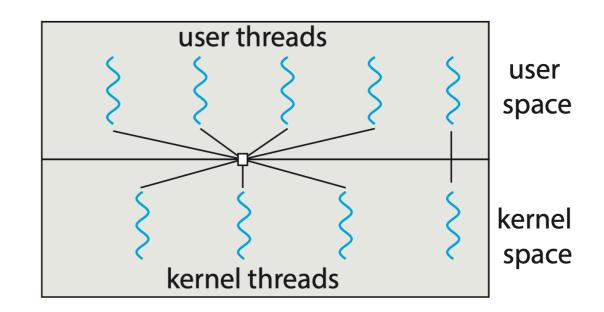
user space

kernel space

Two-level Model

 Similar to the many-to-many model except that it allows a user thread to be bound to a specific kernel thread

- Examples include
 - IRIX
 - HP-UX
 - Tru64 UNIX
 - Solaris 8 and earlier



Thread Libraries

 A thread library provides programmer with API for creating and managing threads

- Two primary ways of implementing
 - Library entirely in user space (no kernel support)
 - Kernel-level library supported by the OS
- Three main thread libraries currently in use
 - (1) POSIX Pthreads user-level or kernel-level threads for POSIX-compliant systems
 - (2) Windows thread library kernel-level threads for Windows systems
 - (3) Java threads threads created and managed in Java programs (typically mapped to thread library of host system)

POSIX Pthreads

May be provided either as user-level or kernel-level

• A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization

 API specifies behavior of the thread library, implementation is up to development of the library

Common in UNIX operating systems (Solaris, Linux, macOS)

Pthreads Example

```
#include <pthread.h>
#include <stdio.h>
#include <stdlib.h>
int sum; /* this data is shared by the thread(s) */
void *runner(void *param); /* threads call this function */
int main(int argc, char *argv[])
  pthread_t tid; /* the thread identifier */
  pthread_attr_t attr; /* set of thread attributes */
  /* set the default attributes of the thread */
  pthread_attr_init(&attr);
  /* create the thread */
  pthread_create(&tid, &attr, runner, argv[1]);
  /* wait for the thread to exit */
  pthread_join(tid,NULL);
  printf("sum = %d\n",sum);
```

Pthreads Example (Cont.)

```
/* The thread will execute in this function */
void *runner(void *param)
{
  int i, upper = atoi(param);
  sum = 0;

  for (i = 1; i <= upper; i++)
     sum += i;

  pthread_exit(0);
}</pre>
```

Pthreads Code for Joining 10 Threads

```
#define NUM_THREADS 10

/* an array of threads to be joined upon */
pthread_t workers[NUM_THREADS];

for (int i = 0; i < NUM_THREADS; i++)
   pthread_join(workers[i], NULL);</pre>
```

Windows API - Multithreaded C Program

```
#include <windows.h>
#include <stdio.h>
DWORD Sum; /* data is shared by the thread(s) */
int main(int argc, char *argv[])
  DWORD ThreadId;
  HANDLE ThreadHandle;
  int Param:
  Param = atoi(argv[1]);
  /* create the thread */
  ThreadHandle = CreateThread(
     NULL, /* default security attributes */
     0, /* default stack size */
     Summation, /* thread function */
     &Param, /* parameter to thread function */
     0, /* default creation flags */
     &ThreadId); /* returns the thread identifier */
   /* now wait for the thread to finish */
  WaitForSingleObject(ThreadHandle,INFINITE);
  /* close the thread handle */
  CloseHandle(ThreadHandle);
  printf("sum = %d\n",Sum);
```

Windows API - Multithreaded C Program (Cont.)

```
/* The thread will execute in this function */
DWORD WINAPI Summation(LPVOID Param)
{
    DWORD Upper = *(DWORD*)Param;
    for (DWORD i = 1; i <= Upper; i++)
        Sum += i;
    return 0;
}</pre>
```

Java Threads

Java threads are managed by the JVM (Java Virtual Machine)

Typically implemented using the threads model provided by underlying OS

- Java threads may be created in two different ways:
 - Extend the Thread class and override run () method
 - Implement the Runnable interface

Java Multithreaded Program

```
import java.util.concurrent.*;
class Summation implements Callable<Integer>
  private int upper;
  public Summation(int upper) {
    this.upper = upper;
  /* The thread will execute in this method */
  public Integer call() {
     int sum = 0;
    for (int i = 1; i <= upper; i++)
       sum += i;
    return new Integer(sum);
public class Driver
 public static void main(String[] args) {
   int upper = Integer.parseInt(args[0]);
   ExecutorService pool = Executors.newSingleThreadExecutor();
   Future<Integer> result = pool.submit(new Summation(upper));
   try {
       System.out.println("sum = " + result.get());
   } catch (InterruptedException | ExecutionException ie) { }
```

Implicit Threading

- Creation and management of threads done by compilers and run-time libraries rather than programmers
 - Growing in popularity as numbers of threads increase, program correctness more difficult with explicit threads
- Three methods explored
 - Thread Pools
 - OpenMP
 - Grand Central Dispatch
- Other methods include Intel Threading Building Blocks (TBB), java.util.concurrent package

Thread Pools

- In applications where threads are repeatedly being created/destroyed thread pools might provide a performance benefit
 - Example: A server that spawns a new thread each time a client connects to the system and discards that thread when the client disconnects

- A thread pool is a group of threads that have been pre-created and are available to do work as needed
 - Threads may be created when the process starts
 - A thread may be kept in a queue until it is needed
 - After a thread finishes, it is placed back into a queue until it is needed again
 - Avoids the extra time needed to spawn new threads when they're needed

Thread Pools

Advantages of thread pools:

- Typically faster to service a request with an existing thread than create a new thread (performance benefit)
- Bounds the number of threads in a process
 - The only threads available are those in the thread pool
 - If the thread pool is empty, then the process must wait for a thread to re-enter the pool before it can assign work to a thread
 - Without a bound on the number of threads in a process, it is possible for a process to create so many threads that all of the system resources are exhausted

Implicit Threading with OpenMP

 Set of compiler directives and an API that provides support for parallel programming in shared-memory environments

- Identifies parallel regions as blocks of code that may run in parallel
 - Developers insert compiler directives into their code to define parallel region

Implicit Threading with OpenMP (Cont.)

```
#include <omp.h>
#include <stdio.h>
int main(int argc, char *argv[])
{
    /* standard sequential code here */
    /* the next bit is automatically parallelized */
    #pragma omp parallel
       printf("I am a parallel region.");
    #pragma omp parallel for
    for (int i = 0; i < 1000; i++) {
        c[i] = a[i] + b[i];
    /* more standard sequential code can go here */
    return 0;
```

Grand Central Dispatch

- Developed for macOS / iOS (and other Apple operating systems)
- Tasks are placed into dispatch queues as program runs
- Utilizes multiple dispatch queues
 - Serial dispatch queue
 - Tasks are removed in FIFO order
 - Once a task is removed, it must be completed prior to next dequeue
 - Concurrent dispatch queue
 - Tasks are removed in FIFO order
 - Multiple tasks may be removed at a time and run in parallel