CS420: Operating Systems Interprocess Communication

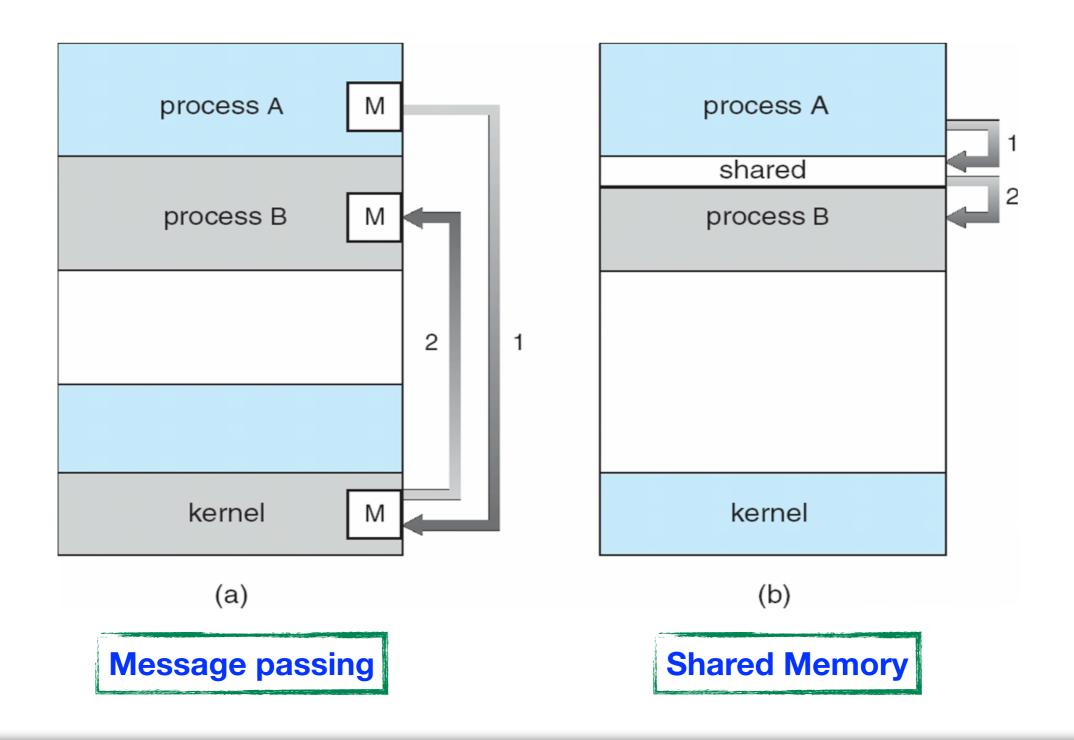
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Interprocess Communication

- Processes within a system may be independent or cooperating
 - An independent process cannot affect or be affected by the execution of another process
 - A cooperating process can affect or be affected by the execution of another process
- Reasons for cooperating processes:
 - Information sharing
 - Computation speedup
 - Modularity
 - Convenience
- Cooperating processes need interprocess communication (IPC)
- Two models of IPC
 - Message passing
 - Shared memory

Communication Models



Producer-Consumer Problem

Producer-consumer problem is a common paradigm for cooperating processes

- Producer process produces information that is consumed by a consumer process
 - One solution is to use shared memory for the two processes to communicate
 - Useful to have a buffer that can be filled by the producer and emptied by the consumer if they are to run concurrently
 - Unbounded-buffer places no practical limit on the size of the buffer
 - Bounded-buffer assumes that there is a fixed buffer size

Bounded-Buffer - Shared-Memory Approach

 The following information is in shared memory and is available to both the producer and the consumer

```
#define BUFFER_SIZE 10

typedef struct {
    /* info to be passed */
} item;

item buffer[BUFFER_SIZE]; /* circular buffer */
int in = 0; /* updated by producer */
int out = 0; /* updated by consumer */
```

• This implementation can only use BUFFER SIZE-1 elements

Bounded-Buffer – Producer

```
while (true) {
    /* Produce an item */
    while ( ((in + 1) % BUFFER SIZE) == out)
        ; /* do nothing -- no free buffers */
    buffer[in] = item; /* buffer not full, add item */
    in = (in + 1) % BUFFER SIZE;
```

Bounded Buffer – Consumer

```
while (true) {
    while (in == out)
        ; /* do nothing -- nothing to consume */
    // remove an item from the buffer
    item = buffer[out];
    out = (out + 1) % BUFFER SIZE;
    /* consume/process item retrieved from buffer */
```

Interprocess Communication – Message Passing

 Message passing – processes communicate with each other without resorting to shared variables

- IPC facility provides two operations:
 - send(message) message size fixed or variable
 - receive(message)

- If two processes want to communicate, they need to:
 - Establish a communication link between them
 - Exchange messages via send/receive

Interprocess Communication – Message Passing

 Communication link can be implemented in variety of ways (including shared memory)

- There are several choices when implementing the communication link
 - Direct or indirect communication
 - Synchronous or asynchronous communication
 - Automatic or explicit buffering

Direct Communication

Processes must name each other explicitly:

- send(P, message) send a message to process P
- receive(Q, message) receive a message from process Q

Properties of direct communication link

- Links are established automatically between the two processes
- A link is associated with exactly one pair of communicating processes
- Between each pair there exists exactly one link
- The link may be unidirectional, but is usually bidirectional

Indirect Communication

- Messages are directed and received from mailboxes (also referred to as ports)
 - Each mailbox has a unique id
 - Processes can communicate only if they share a mailbox

- Properties of indirect communication link
 - Link established only if processes share a common mailbox
 - A link may be associated with more than two processes
 - Each pair of processes may share several communication links (multiple mailboxes)
 - Link may be unidirectional or bidirectional

Indirect Communication (Cont.)

Operations

- Create a new mailbox
- Send and receive messages through mailbox
- Destroy a mailbox

Primitives are defined as:

- send(A, message) send a message to mailbox A
- receive(A, message) receive a message from mailbox A

Indirect Communication (Cont.)

- Mailbox sharing consider the following ...
 - P_1 , P_2 , and P_3 share mailbox A
 - P_1 , sends; P_2 and P_3 receive
 - Who gets the message?

- Possible solutions to avoid unpredictable behavior
 - Allow a link to be associated with at most two processes
 - Allow only one process at a time to execute a receive operation
 - Allow the system to arbitrarily select the receiver. Sender is notified who the receiver was.

Synchronization

Message passing may be either blocking or non-blocking

- Blocking is considered synchronous
 - Blocking send has the sender block until the message is received
 - Blocking receive has the receiver block until a message is available

- Non-blocking is considered asynchronous
 - Non-blocking send has the sender send the message and continue
 - Non-blocking receive has the receiver receive a valid message or null

Buffering

- Regardless of how messages are exchanged between processes, messages are queued
- Queueing can be implemented in one of three ways
 - (1) Zero capacity queue has maximum length of 0 Sender must wait (or block) until the receiver gets the message
 - (2) **Bounded capacity** queue has finite length of *n* messages Sender must wait if link full
 - (3) Unbounded capacity queue has 'infinite' length Sender never waits

Examples of IPC Systems - POSIX

POSIX Shared Memory

- Process first creates shared memory segment

```
segment_id = shmget(IPC_PRIVATE, size, S_IRUSR | S_IWUSR);
```

- Any process wanting access to that shared memory must attach to it

```
shared memory = (char *) shmat(segment id, NULL, 0);
```

- Now the process could write to the shared memory

```
sprintf(shared memory, "Writing to shared memory");
```

- When done a process can detach the shared memory from its address space

```
shmdt(shared memory);
```

- When the shared memory space is no longer needed, free it

```
shmctl(segment id, IPC RMID, NULL);
```