

ECE260: Fundamentals of Computer Engineering

MIPS Branching Instructions

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Conditional Operations & Simple Jumps

- Not all programs involve a simple set of sequential instructions
 - MOST programs have conditionally executed code
 - If some condition is true, execute a specific block of code, otherwise execute some other code
 - If-statements / If-else statements / If-elseif-statements
- MIPS has **branch** and **jump** instructions to handle conditional code
 - Branch to a labeled instruction if a condition is true, otherwise, continue sequentially

Conditional Operations & Simple Jumps

- Branch-on-equal (beq)

- I-type instruction with 16-bit immediate that stores distance (in words) to label
- if (rs == rt) branch to label L1

```
beq rs, rt, L1
```

- Branch-on-not-equal (bne)

- I-type instruction with 16-bit immediate that stores distance (in words) to label
- if (rs != rt) branch to label L1

```
bne rs, rt, L1
```

- Jump (j)

- J-type instruction with 26-bit immediate for address of label
- unconditional jump to label L1

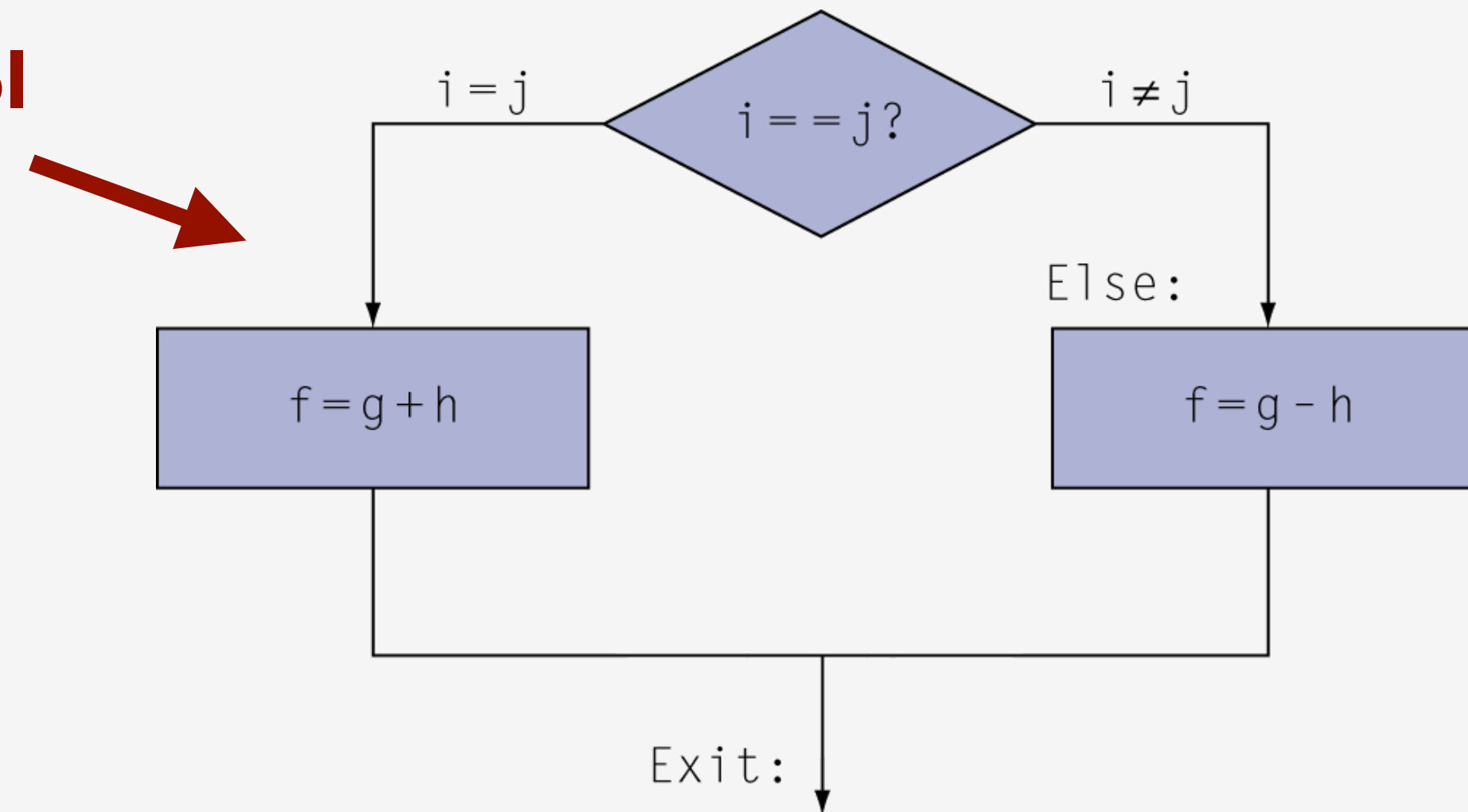
```
j L1
```

Compiling If Statements

- Example C code

```
if (i == j)
    f = g + h;
else
    f = g - h;
```

**Flow Control
of C code**



- Oftentimes, the above code structure will result in a **branch** and a **jump** instruction

```
        bne $s3, $s4, Else  # check i==j, branch if they differ
        add $s0, $s1, $s2  # if i==j, do (g + h)
        j    Exit          # then jump to the end
Else:    sub $s0, $s1, $s2  # if i!=j, do (g - h)
Exit:
```

Assume:

**i in \$s3
j in \$s4
g in \$s1
h in \$s2**

Compiling Loop Statements

- Example C code

```
while (save[i] == k) {  
    i += 1;  
}
```

Assume:

i in \$s3

k in \$s5

save baseAddr in \$s6

- Looping structures continue to execute until some condition is met

```
Loop: sll    $t1, $s3, 2      # convert i to byte-addressing (i.e. i*4)  
      add    $t1, $t1, $s6    # add offset of index i to baseAddr of save  
      lw     $t0, 0($t1)      # load save[i]  
      bne    $t0, $s5, Exit   # check save[i]==k, branch if they differ  
      addi   $s3, $s3, 1      # increment i  
      j      Loop            # jump back to beginning of loop
```

Exit:

More Conditional Operations

- Set result to 1 if a condition is true, otherwise, set to 0
- Set-on-less-than (slt)
 - if ($rs < rt$) $rd = 1$; else $rd = 0$;
- Set-on-less-than-immediate (slti)
 - if ($rs < \text{constant}$) $rt = 1$; else $rt = 0$;

```
slt rd, rs, rt
```

```
slti rt, rs, constant
```

- Set-on-less-than-immediate-unsigned (sltiu)
 - if ($rs < \text{constant}$) $rt = 1$; else $rt = 0$;
- Use slt instructions in combination with beq, and bne instructions

```
sltiu rt, rs, constant
```

```
slt $t0, $s1, $s2    # if ($s1 < $s2)  
bne $t0, $zero, L     # branch to L
```

Branch Instruction Design

- Only beq and bne branch instructions
 - No native instructions for branch-on-less than (blt), branch-on-greater-or-equal (bge), etc.
- Hardware for $<$, $>$, \leq , \geq , is slower than hardware for $=$, \neq
 - Combining comparison with branch involves more work per instruction
 - At hardware level, additional levels of hardware logic result in slower hardware
 - Would require a slower clock
 - All instructions would be penalized and run at slower clock
- beq and bne are the common case, so they are implemented in hardware
- Pseudo-instructions exist for blt, bgt, ble, bge
 - Converted into some combination of slt, bne, beq

Signed vs. Unsigned Comparisons

- Signed comparison: slt, slti
- Unsigned comparison: sltu, sltui

- Example

- Assume:

$\$s0 = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111$ (-1 or +4,294,967,295)

$\$s1 = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001$

- Signed comparison

```
slt $t0, $s0, $s1 # -1 < +1 therefore, sets $t0=1
```

- Unsigned comparison

```
sltu $t0, $s0, $s1 # +4,294,967,295 > +1 therefore, sets $t0=0
```