

ECE260: Fundamentals of Computer Engineering

Data Representation & 2's Complement

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Data Representation

- Internally, computers represent all data as binary
 - Provides a simple method to design and build hardware
 - A switch (i.e. a transistor) is either "on" or "off"
 - A wire is either charged or not charged
- All information is encoded as 1's and 0's
 - Characters
 - Integers (positive and negative numbers)
 - Non-integers (fixed-point and floating point numbers)
- Standards ensure interoperability between computers
 - 2's Complement, ASCII, Unicode, IEEE Floating Point

American Standard Code for Information Interchange (ASCII)

- Common character encoding standard
- Available in 7-bit and extended 8-bit
 - 7-bit version encodes 2^7 (128) characters
 - 8-bit version encodes 2^8 (256) characters
- Not so great for languages based on non-English alphabets
- Unicode has replaced ASCII in many contexts
 - Backwards compatible with ASCII
 - Supports a much wider range of characters

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Binary	Character	Binary	Character	Binary	Character	Binary	Character
00000000	NUL	00100000	SP	01000000	@	01100000	`
00000001	SOH	00100001	!	01000001	A	01100001	a
00000010	STX	00100010	"	01000010	B	01100010	b
00000011	ETX	00100011	#	01000011	C	01100011	c
00000100	EOT	00100100	\$	01000100	D	01100100	d
00000101	ENQ	00100101	%	01000101	E	01100101	e
00000110	ACK	00100110	&	01000110	F	01100110	f
00000111	BEL	00100111	'	01000111	G	01100111	g
00001000	BS	00101000	(01001000	H	01101000	h
00001001	HT	00101001)	01001001	I	01101001	i
00001010	LF	00101010	*	01001010	J	01101010	j
00001011	VT	00101011	+	01001011	K	01101011	k
00001100	FF	00101100	,	01001100	L	01101100	l
00001101	CR	00101101	-	01001101	M	01101101	m
00001110	SO	00101110	.	01001110	N	01101110	n
00001111	SI	00101111	/	01001111	O	01101111	o
00010000	DLE	00110000	0	01010000	P	01110000	p
00010001	DC1	00110001	1	01010001	Q	01110001	q
00010010	DC2	00110010	2	01010010	R	01110010	r
00010011	DC3	00110011	3	01010011	S	01110011	s
00010100	DC4	00110100	4	01010100	T	01110100	t
00010101	NAK	00110101	5	01010101	U	01110101	u
00010110	SYN	00110110	6	01010110	V	01110110	v
00010111	ETB	00110111	7	01010111	W	01110111	w
00011000	CAN	00111000	8	01011000	X	01111000	x
00011001	EM	00111001	9	01011001	Y	01111001	y
00011010	SUB	00111010	:	01011010	Z	01111010	z
00011011	ESC	00111011	;	01011011	[01111011	{
00011100	FS	00111100	<	01011100	\	01111100	
00011101	GS	00111101	=	01011101]	01111101	}
00011110	RS	00111110	>	01011110	^	01111110	~
00011111	US	00111111	?	01011111	_	01111111	DEL

Encoding Numbers

- Numbers can be represented in any base
 - Humans typically use base 10 (**decimal**) so we can count on our fingers
 - Programmers often represent data in base 16 (**hexadecimal**)
 - Computers represent all data using base 2 (**binary**)
- In any base, the decimal value of the i^{th} digit d can be calculated as:
 - Sum the decimal values of each digit to determine the total value:

$$d \times \text{Base}^i$$

$$\sum_{i=0}^{\text{\#digits}} (d \times \text{Base}^i)$$

base 10

5	3	8
2	1	0

digit index

= **538_{ten}**

(5×10^2)	+	(3×10^1)	+	(8×10^0)
(5×100)	+	(3×10)	+	(8×1)
500	+	30	+	8

base 2

1	0	1	1
3	2	1	0

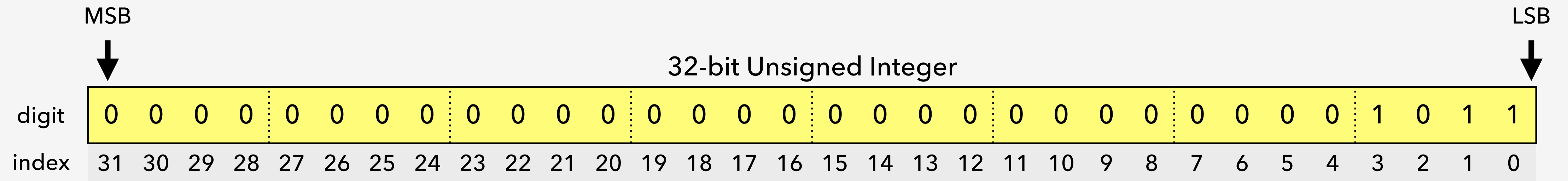
digit index

= **11_{ten}**

(1×2^3)	+	(0×2^2)	+	(1×2^1)	+	(1×2^0)
(1×8)	+	(0×4)	+	(1×2)	+	(1×1)
8	+	0	+	2	+	1

Unsigned Integers

- Represented as binary value (a bit string)
 - Leftmost bit is called the **most significant bit (MSB)**
 - Rightmost bit is called the **least significant bit (LSB)**

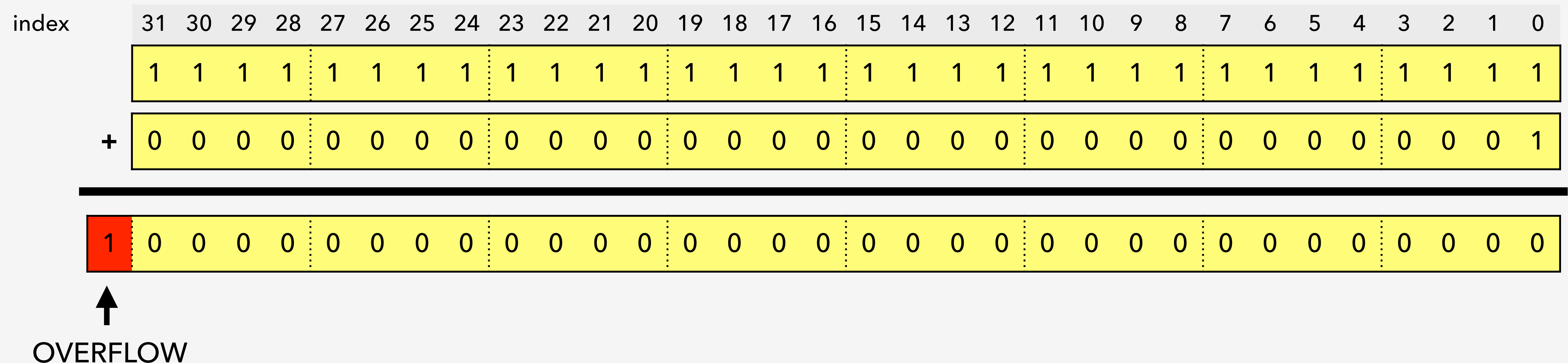


- The number of bits is determined by the **bitness** of the computer architecture
 - A 32-bit computer can represent unsigned integers from 0 to $2^{32} - 1$
 - 0 to 4,294,967,295 (4.29 Billion)
 - A 64-bit computer can represent unsigned integers from 0 to $2^{64} - 1$
 - 0 to 18,446,744,073,709,551,615 (18.4 Quintillion)



Arithmetic Limitations & Overflow

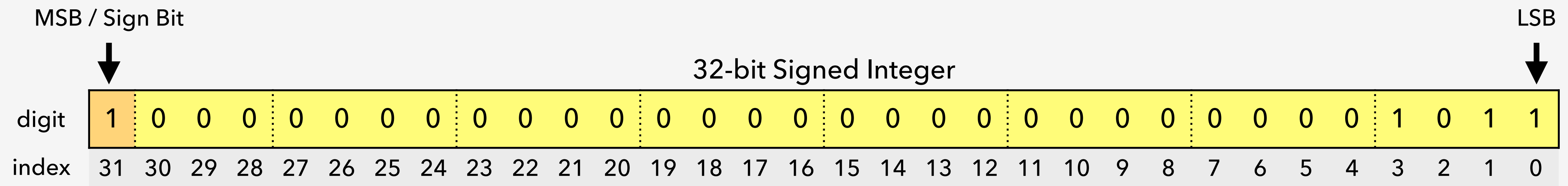
- The largest unsigned value a 32-bit computer can represent is **4,294,967,295**
- What happens when user requests **4,294,967,295 + 1**?



- CPUs have a special **status register** with an **overflow flag** to indicate when this happens

2's Complement Signed Integers

- Must be able to represent both positive AND negative numbers
- Represented as binary value where the MSB is now a **sign bit**
 - When sign bit is 1, the signed integer is a negative number
 - When sign bit is 0, the signed integer is a positive number



- Maximum value less than that of unsigned integers since 1 bit is being used to represent the sign
 - A 32-bit computer can represent signed integers from -2^{31} to $2^{31} - 1$
 - $-2,147,483,648$ to $2,147,483,647$

2's Complement Signed Integer Examples

- Positive Number Examples

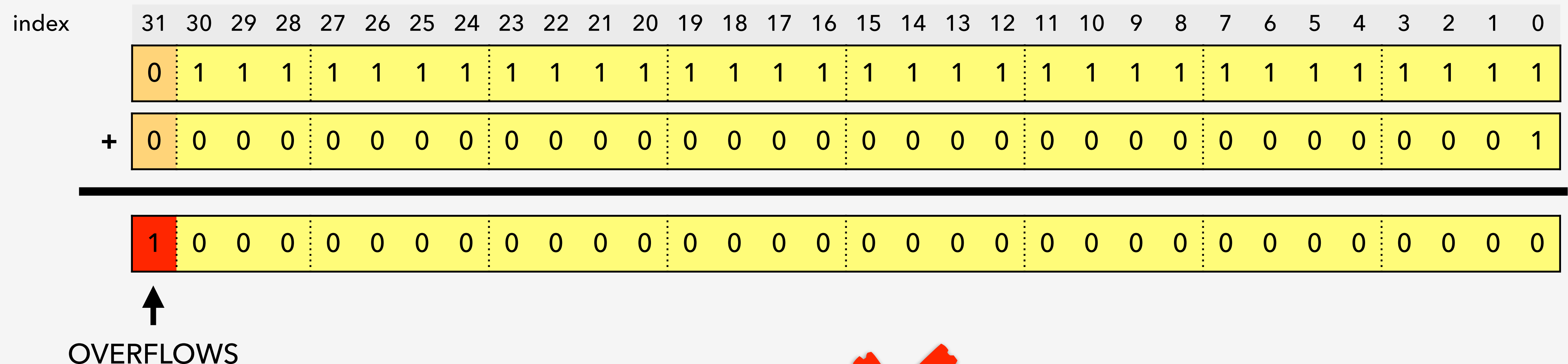
2,147,483,647	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
247	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Negative Number Examples

-1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
-2	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	
-3	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	0	1		
-2,147,483,648	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	

Arithmetic Overflow with Signed Integers

- The largest signed value a 32-bit computer can represent is **2,147,483,647**
- What happens when user requests **2,147,483,647 + 1**?



$$2,147,483,647 + 1 = -2,147,483,648$$

- Similar badness happens if you subtract 1 from $-2,147,483,648$

Binary to Decimal Conversion for Signed Integers

- Similar to conversion with unsigned integers, but sign bit is multiplied by -2^{31} instead of 2^{31}
 - If the signed integer is positive $\Rightarrow (0 \times -2^{31}) = 0$
 - If the signed integer is negative $\Rightarrow (1 \times -2^{31}) = -2,147,483,648$

For a 32-bit 2's complement integer, compute decimal value From binary as follows:

$$(d \times -2^{31}) + \sum_{i=0}^{i=30} (d \times 2^i)$$

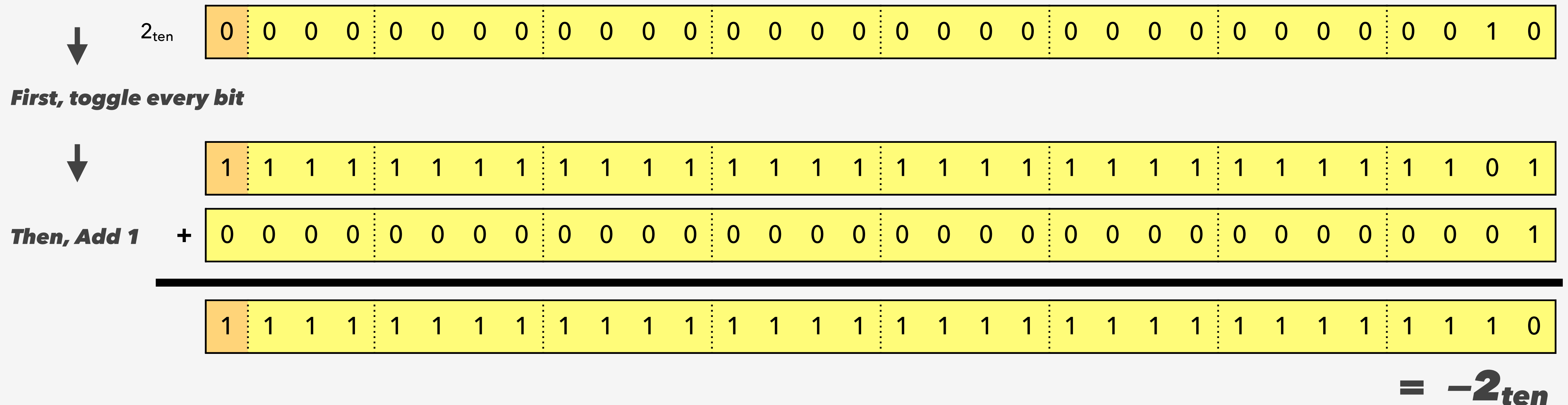
MSB / Sign Bit																																LSB			
↓																32-bit Signed Integer																↓			
digit	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	1	1	0	1			
index	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			

(1×-2^{31})	+	...	+	(1×2^5)	+	(0×2^4)	+	(1×2^3)	+	(1×2^2)	+	(0×2^1)	+	(1×2^0)
	+	...	+	(1×32)	+	(0×16)	+	(1×8)	+	(1×4)	+	(0×2)	+	(1×1)
-2,147,483,648	+	0	+	32	+	0	+	8	+	4	+	0	+	1

= **-2,147,483,603**

Negating 2's Complement Signed Integers

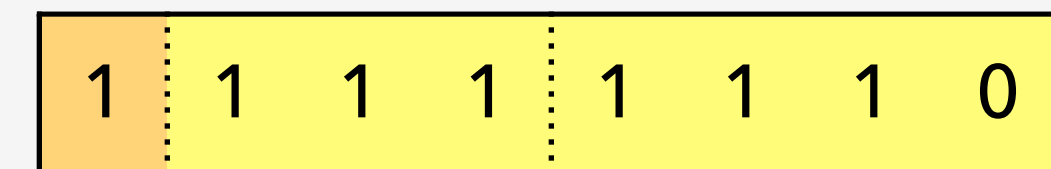
- **Step #1:** Toggle every bit, including the sign bit (i.e. all 1's become 0's and all 0's become 1's)
- **Step #2:** Add 1 to result of step #1
- Example: Negate the integer +2



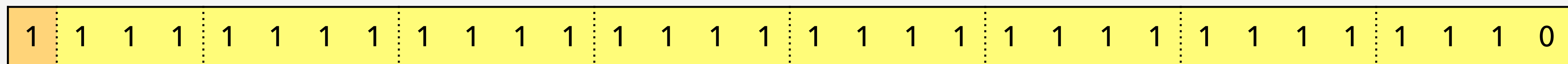
Sign Extension for 2's Complement Signed Integers

- Must sign extend when increasing the number of bits used to represent a **signed integer**
 - Must **extend sign bit** to preserve the numeric value (i.e. fill all leading zeros with sign bit)
- Example: Loading a signed 8-bit value into a 32-bit register
 - If bits [31:9] of register are left as 0's then sign of 8-bit value may be lost

Signed 8-bit value holds -2_{ten}



without sign extension 32-bit register holds value 254_{ten}



WITH sign extension 32-bit register holds correct value -2_{ten}