

## ECE335 Summer 2019 - Assignment 4

### Section 2.3 - More Operations on Sets

2.3.1 . Analyze the logical forms of the following statements

a.  $\mathcal{F} \subseteq \mathcal{P}(A)$

Using the definition of subset we first expand

$$\forall x(x \in \mathcal{F} \rightarrow x \in \mathcal{P}(A))$$

Then applying the definition of powerset, i.e. that an element of the powerset is a *subset* gives

$$\forall x(x \in \mathcal{F} \rightarrow \forall y(y \in x \rightarrow y \in A))$$

b.  $A \subseteq \{2n + 1 | n \in \mathbb{N}\}$

Using the definition of subset we first expand

$$\forall x(x \in A \rightarrow x \in \{2n + 1 | n \in \mathbb{N}\})$$

Then rewriting the elementhood test gives

$$\forall x(x \in A \rightarrow \exists n \in \mathbb{N}(x = 2n + 1))$$

c.  $\{n^2 + n + 1 | n \in \mathbb{N}\} \subseteq \{2n + 1 | n \in \mathbb{N}\}$

Using the definition of subset we first expand

$$\forall x(x \in \{n^2 + n + 1 | n \in \mathbb{N}\} \rightarrow x \in \{2n + 1 | n \in \mathbb{N}\})$$

Since the  $n$ 's in both equations may be different for the same  $x$  we can more consisely write the statement by using  $m$  and  $n$  as

$$\forall n \in \mathbb{N} \exists m \in \mathbb{N}(n^2 + n + 1 = 2m + 1)$$

d.  $\mathcal{P}(\cup_{i \in I} A_i) \not\subseteq \cup_{i \in I} \mathcal{P}(A_i)$

Using the definition of not a subset (which implies *there exists* an element not in the superset) we first expand

$$\exists x(x \in \mathcal{P}(\cup_{i \in I} A_i) \wedge x \notin \cup_{i \in I} \mathcal{P}(A_i))$$

Now expanding the first term using the definition of powerset gives

$$\exists x(\forall y(y \in x \rightarrow y \in (\cup_{i \in I} A_i)) \wedge \neg(x \in \cup_{i \in I} \mathcal{P}(A_i)))$$

Using the definition of indexed set union gives

$$\exists x(\forall y(y \in x \rightarrow \exists i \in I(y \in A_i)) \wedge \neg \exists i \in I(x \in \mathcal{P}(A_i)))$$

Expanding the second powerset term gives

$$\exists x(\forall y(y \in x \rightarrow \exists i \in I(y \in A_i)) \wedge \neg \exists i \in I(\forall y(y \in x \rightarrow y \in A_i)))$$

Negating the quantifier in the second term and using DeMorgan's law (after converting the conditional to  $\vee$ ) gives

$$\exists x(\forall y(y \in x \rightarrow \exists i \in I(y \in A_i)) \wedge \forall i \in I(\exists y(y \in x \wedge y \notin A_i)))$$

2.3.8 . Let  $I = \{2, 3\}$  and for each  $i \in I$  let  $A_i = \{i, 2i\}$  and  $B_i = \{i, i + 1\}$

a. List the elements of the sets  $A_i$  and  $B_i$  for  $i \in I$ .

Substituting in the two values for  $i$  produces 4 sets

$$A_2 = \{2, 4\}, A_3 = \{3, 6\}, B_2 = \{2, 3\}, B_3 = \{3, 4\}$$

b. Find  $\cap_{i \in I}(A_i \cup B_i)$  and  $(\cap_{i \in I} A_i) \cup (\cap_{i \in I} B_i)$

First, to find  $\cap_{i \in I}(A_i \cup B_i)$  we compute the unions for both indices giving

$$(A_2 \cup B_2) = \{2, 3, 4\} \text{ and } (A_3 \cup B_3) = \{3, 4, 6\}$$

Taking the intersection of these two sets gives

$$\cap_{i \in I}(A_i \cup B_i) = \{3, 4\}$$

Second, to find  $(\cap_{i \in I} A_i) \cup (\cap_{i \in I} B_i)$  we compute the two family intersections giving

$$(\cap_{i \in I} A_i) = \emptyset \text{ and } (\cap_{i \in I} B_i) = \{3\}$$

Taking the union of these two sets gives

$$(\cap_{i \in I} A_i) \cup (\cap_{i \in I} B_i) = \{3\}$$

which is clearly different than the first result.

c. What can you conclude about whether or not these statements are equivalent?

Since part (b) gives two sets that provide a *counterexample*

$$\cap_{i \in I}(A_i \cup B_i) \neq (\cap_{i \in I} A_i) \cup (\cap_{i \in I} B_i)$$

i.e. family intersection *does not* distribute over indexed set union.

2.3.10 . Show that for any sets  $A$  and  $B$  that  $\mathcal{P}(A \cap B) = \mathcal{P}(A) \cap \mathcal{P}(B)$ .

We shall show this equivalency by showing that  $x \in \mathcal{P}(A \cap B)$  is equivalent to  $x \in \mathcal{P}(A) \cap \mathcal{P}(B)$ .

$x \in \mathcal{P}(A \cap B) = x \subseteq (A \cap B)$	definition of powerset
$= \forall y (y \in x \rightarrow y \in (A \cap B))$	definition of subset
$= \forall y (y \in x \rightarrow (y \in A \wedge y \in B))$	definition of intersection
$= \forall y (\neg(y \in x) \vee (y \in A \wedge y \in B))$	conditional law
$= \forall y ((\neg(y \in x) \vee (y \in A)) \wedge (\neg(y \in x) \vee (y \in B)))$	distributive law
$= \forall y ((\neg(y \in x) \vee (y \in A)) \wedge \forall y (\neg(y \in x) \vee (y \in B)))$	quantifier distribution
$= \forall y (y \in x \rightarrow y \in A) \wedge \forall y (y \in x \rightarrow y \in B)$	conditional law
$= (x \subseteq A) \wedge (x \subseteq B)$	definition of subset
$= (x \in \mathcal{P}(A)) \wedge (x \in \mathcal{P}(B))$	definition of powerset
$= x \in \mathcal{P}(A) \cap \mathcal{P}(B)$	definition of intersection

Thus  $\mathcal{P}(A \cap B) = \mathcal{P}(A) \cap \mathcal{P}(B)$

### Section 3.1 - Proof Strategies

3.1.8 . Suppose  $A \setminus B \subseteq C \cap D$  and  $x \in A$ . Prove that if  $x \notin D$  then  $x \in B$ .

<i>Givens</i>	<i>Goal</i>
$A \setminus B \subseteq C \cap D$	$x \notin D \rightarrow x \in B$
$x \in A$	

We can modify the goal by assuming  $x \notin D$  as a given thus transforming the goal into

<i>Givens</i>	<i>Goal</i>
$A \setminus B \subseteq C \cap D$	$x \in B$
$x \in A$	
$x \notin D$	

At this point we begin by noting that since  $x \notin D$  that  $x \notin C \cap D$  and therefore since  $A \setminus B \subseteq C \cap D \Rightarrow x \notin A \setminus B$ . But since  $x \in A$ , for  $x \notin A \setminus B$  we must have  $x \in B$ .

Formally,

**Theorem.** Suppose  $A \setminus B \subseteq C \cap D$  and  $x \in A$ . If  $x \notin D$  then  $x \in B$ .

*Proof.* Suppose  $x \notin D$ . Then this implies that  $x \notin (C \cap D)$  and therefore  $x \notin (A \setminus B)$ . But since  $x \in A$  for  $x \notin (A \setminus B)$  means that  $x \in B$ . Thus if  $x \notin D$  then  $x \in B$ .

3.1.11 . Suppose  $a, b, c$ , and  $d$  are real numbers,  $0 < a < b$  and  $d > 0$ . Prove that if  $ac \geq bd$  then  $c > d$ .

<i>Givens</i>	<i>Goal</i>
$a, b, c, d \in \mathbb{R}$	$ac \geq bd \rightarrow c > d$
$0 < a < b$	
$d > 0$	

Consider proving the contrapositive of the goal giving

<i>Givens</i>	<i>Goal</i>
$a, b, c, d \in \mathbb{R}$	$c \leq d \rightarrow ac < bd$
$0 < a < b$	
$d > 0$	

Then using the conditional proof technique, we move the first part of the goal to the givens

<i>Givens</i>	<i>Goal</i>
$a, b, c, d \in \mathbb{R}$	$ac < bd$
$0 < a < b$	
$d > 0$	
$c \leq d$	

Since  $a > 0$  we can multiply both sides of the last given without reversing the inequality giving  $ac \leq ad$ . Likewise, since  $d > 0$  we can multiply both sides of the second given without reversing the inequality giving  $ad < bd$ . Combining these two inequalities gives  $ac \leq ad < bd \Rightarrow ac < bd$  which is the desired goal.

Formally

**Theorem.** Suppose  $a, b, c$ , and  $d$  are real numbers,  $0 < a < b$  and  $d > 0$ . If  $ac \geq bd$  then  $c > d$ .

*Proof.* We will prove by contrapositive, thus suppose  $c \leq d$ . Multiplying by the positive number  $a$  gives  $ac \leq ad$ . Similarly multiplying by the positive number  $d$  gives  $ad < bd$ . Combining  $ac \leq ad$  and  $ad < bd$  we can conclude that  $ac < bd$ . Thus if  $ac \geq bd$  then  $c > d$ .

## Section 3.2 - Proofs Involving Negations and Conditionals

3.2.5. Use the method of contradiction to prove that supposing  $A \cap C \subseteq B$  and  $a \in C$ , that  $a \notin (A \setminus B)$ .

Since the initial goal is  $a \notin (A \setminus B)$  we will instead use proof by contradiction by *assuming*  $a \in (A \setminus B)$  and attempt to reach a contradiction. Thus our givens and goals are

<i>Givens</i>	<i>Goal</i>
$A \cap C \subseteq B$	<i>contradiction</i>
$a \in C$	
$a \in (A \setminus B)$	

Since  $a \in (A \setminus B) \Rightarrow (a \in A \wedge a \notin B)$ . Thus we can add those to our givens (since they both must be true)

<i>Givens</i>	<i>Goal</i>
$A \cap C \subseteq B$	<i>contradiction</i>
$a \in C$	
$a \in A$	
$a \notin B$	

Since by givens 2 and 3,  $a \in A$  and  $a \in C$ , therefore  $a \in (A \cap C)$ . By given 1, this implies that  $a \in B$  which contradicts the last given. Thus our initial assumption was incorrect, hence  $a \notin (A \setminus B)$ .

Formally,

**Theorem.** Suppose  $A \cap C \subseteq B$  and  $a \in C$ . Then  $a \notin (A \setminus B)$ .

*Proof.* Suppose not, then  $a \in (A \setminus B)$ . This implies that  $a \in A$  and  $a \notin B$ . Since  $a \in C$  we have that  $a \in (A \cap C)$ . But since  $A \cap C \subseteq B$ , this means that  $a \in B$  which contradicts the fact that  $a \notin B$ . Thus  $a \notin (A \setminus B)$ .