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# Putnam Compendium

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## Abstract

The work contains the solutions to Putnam problems.

## 1 Solutions

### Question 2012 A-1.

**Solution.** Assume without loss of generality that  $d_1 \leq d_2 \leq \dots \leq d_n$ . Suppose for sake of contradiction that there does not exist distinct indices  $i, j, k$  such that  $d_i, d_j, d_k$  are the side lengths of an acute triangle. By the property of an acute triangle we have (PATCH)

$$d_{i+2}^2 \geq d_{i+1}^2 + d_i^2 \tag{1}$$

for all  $i$  such that  $1 \leq i \leq 10$ . We first claim that

$$d_i^2 \geq F_i d_1^2$$

holds for all  $1 \leq i \leq 12$  where  $F_i$  is a  $i$ th Fibonacci number with  $F_1 = 1$  and  $F_2 = 1$ . We proceed to prove the claim by strong induction. Base case of the induction trivially holds as  $F_1 = 1$  yielding

$$d_1^2 \geq d_1^2.$$

Now assume that the statement holds true from 1 to  $i$ . From (1) we obtain

$$d_{i+1}^2 \geq d_i^2 + d_{i-1}^2.$$

With the inductive hypothesis we can lower bound the RHS as

$$d_i^2 + d_{i-1}^2 \geq F_i d_1^2 + F_{i-1} d_1^2. \tag{2}$$

Factoring and substituting the Fibonacci recurrence to RHS of (2) we obtain

$$d_i^2 + d_{i-1}^2 \geq F_{i+1} d_1^2.$$

Hence, we finally get that

$$d_{i+1}^2 \geq F_{i+1} d_1^2$$

which completes the induction. Hence the claim  $d_i^2 \leq d_1^2$  holds true for all  $i$  such that  $1 \leq i \leq 12$ . Notice that for  $i = 12$  case we have  $F_{12} = 144$  yielding

$$d_{i+1}^2 \geq 144 d_1^2.$$

As the numbers are chosen from the open interval  $(1, 12)$  we have that  $d_{i+1}^2$  is strictly less than 144, but  $144 d_1^2$  is strictly greater than 144, which is a contradiction. Therefore, we have shown given any  $d_1, d_2, \dots, d_{12}$  chosen from  $(1, 11)$ , there exists three distinct indices  $i, j, k$  such that  $d_i, d_j, d_k$  form side lengths of an acute triangle.  $\square$

**Remark.**

**Question 2012 A-2.**

**Solution.** Assume that for every  $x$  and  $y$  in  $S$  there exists  $z$  in  $S$  such that  $x * z = y$ . Assume that  $a, b, c$  are in  $S$  and  $a * c = b * c$  holds. Let  $d$  be the element such that  $d = a * c = b * c$  holds. From the assumption we can deduce that there exist  $e$  and  $f$  in  $S$  such that

$$\begin{aligned} d * e &= a \\ d * f &= b \end{aligned}$$

**Question 2008 A-1.**

**Solution.** We claim that  $g(x) = f(x, 0)$  satisfies the given properties. Substituting  $(x, y, z) = (0, 0, 0)$  into the functional equations yields

$$f(0, 0) + f(0, 0) + f(0, 0) = 3f(0, 0) = 0$$

which gives that  $f(0, 0) = 0$ . Substituting  $(x, y, z) = (0, 0, 0)$  we obtain

$$f(x, 0) + f(0, 0) + f(0, x) = 0;$$

hence,  $f(x, 0) = -f(0, x)$ . Substituting  $(x, y, z) = (x, y, 0)$  gives

$$f(x, y) + f(y, 0) + f(0, x) = 0.$$

Rearranging and substituting  $f(x, 0) = -f(0, x)$  results in

$$f(x, y) = f(x, 0) - f(y, 0)$$

Let  $g(x) = f(x, 0)$  which is a well-defined function  $g : \mathbb{R} \rightarrow \mathbb{R}$ , we see that

$$f(x, y) = g(x) - g(y)$$

as desired. Hence, we have shown that there exists a function  $g : \mathbb{R} \rightarrow \mathbb{R}$ ,  $g(x) = f(x, 0)$  in particular, that  $f(x, y) = g(x) - g(y)$  is satisfied.

**Question 2010 A-1.**

**Solution.** We claim that  $k = \lceil \frac{n}{2} \rceil$ . For  $n$  odd, the lower bound can be achieved by the partition

$$\{n\}, \{1, n-1\}, \{2, n-2\}, \dots$$

Since  $n-1$  is even, the partition is well defined. Now for the case where

$$\{n, 1\}, \{n-1, 2\}, \{n-2, 3\}, \dots$$

Now we show that such partition strategy is indeed optimal for all  $n$ . Let us denote the sum of each box of the optimal strategy as  $S$ . Since the number  $n$  itself belongs to one of the boxes we obtain a lower bound,  $S \geq n$ . As the partition strategy for  $n$  odd case yields  $S = n$  and  $n$  is indeed a lower bound on  $S$ , we have shown that the strategy is optimal for  $n$  odd case. For  $n$  even we show that  $S$  has to be lower bounded by  $n+1$ . Suppose that  $S = n$ . A partition that can achieve  $S$  must have the following configuration

**Question 2010 B-1.**

**Solution.** Suppose that such sequence exists. Suppose that  $0 \leq a_i^2 \leq 1$  for  $i$ . Then we get that

$$\sum_{i=1}^{\infty} a_i^4 \leq \sum_{i=1}^{\infty} a_i^2.$$

Since  $\sum_{i=1}^{\infty} a_i^k = k$ , by substitution we get  $4 \leq 2$ , which is a contradiction. Hence, there exists an index  $k$  such that  $a_k^2 > 1$ . Now, for large enough  $m$ , we have  $a_k^{2m} > 2m$ , which is a contradiction of the condition of the sequence. Therefore, there does not exist an infinite sequence of real numbers such that

$$a_1^m + a_2^m \dots = m$$

for every positive integer  $m$ .

**Question 2001 A-1.**

**Solution.** Consider a set  $S$  and a binary operation  $*$  and  $S$  is closed under  $*$ . Assume that  $(a*b)*a = b$  for all  $a, b \in S$ . Substituting  $(b*a)$  into the assumption, we obtain

$$((b*a)*b)*(b*a) = b.$$

From  $(b*a)*b = a$ , we can re-write the above equality as

$$a*(b*a) = b,$$

as desired.

**Question 2001 A-2.**

**Solution.** Let us denote the probability of having odd number of heads after  $i$  tosses as  $P_i$ . Then, for  $i \geq 2$  we can express  $P_i$  as

$$P_i = (1 - C_i)P_{i-1} + C_i(1 - P_{i-1}),$$

where  $C_i$  is a probability of flipping a head for the  $i$ th coin. Simplifying the above equality yields

$$P_i = (1 - 2C_i)P_{i-1} + C_i.$$

Substituting  $C_i = \frac{1}{2i+1}$  into the last equality, we obtain

$$P_i = P_{i-1} + \frac{1}{2i+1} - 2$$