## 1. The Basic Problem

Consider real-valued input functions with domain  $\mathcal{X} = [a, b]$ , where a, b be two fixed real numbers with a < b. We want to do the function recovery and let S be a solution operator. Here we consider:

$$S: f \mapsto f, \qquad S: \mathcal{W}^{2,\infty}[a,b] \to \mathcal{L}_{\infty}[a,b].$$

where for all real numbers  $\alpha, \beta$  with  $\alpha < \beta$ ,

$$\mathcal{W}^{2,\infty} := \mathcal{W}^{2,\infty}[a,b], \qquad \mathcal{W}^{2,\infty}[\alpha,\beta] := \{ f \in C[\alpha,\beta] : \left\| f'' \right\|_{\infty,[\alpha,\beta]} < \infty \},$$
 
$$\mathcal{L}_{\infty} := \mathcal{L}_{\infty}[a,b], \qquad \mathcal{L}_{\infty}[\alpha,\beta] := \mathcal{W}^{0,\infty}[\alpha,\beta],$$
 
$$\left\| f'' \right\|_{[\alpha,\beta]} := \left\| f'' \right\|_{\infty,[\alpha,\beta]}, \qquad \left\| f \right\|_{[\alpha,\beta]} := \left\| f \right\|_{\infty,[\alpha,\beta]} = \max_{\alpha \le x \le \beta} |f(x)|.$$

Here we only consider infinity norm, so we omit sub index  $\infty$ .

## 2. Solving the Problems on Partitions

For any interval  $[\alpha, \beta] \subset [a, b]$ , let let  $f_{[\alpha, \beta]}$  denote  $\mathcal{W}^{2,\infty}$  restricted to the interval  $[\alpha, \beta]$ , i.e.,

$$f_{[\alpha, \beta]}: [\alpha, \beta] \to \mathbb{R}, \qquad f_{[\alpha, \beta]}: \mathbf{x} \mapsto f(\mathbf{x}).$$

Moreover, let  $W^{2,\infty}[\alpha,\beta]$  denote the space of functions in  $W^{2,\infty}$  restricted to the set interval  $[\alpha,\beta]$ .

For function recovery, these partitions take the form of subintervals of [a, b]:

$$\mathcal{T} = \{ [\alpha, \beta] : a \le \alpha < \beta \le b \}$$

$$\mathcal{P} = \{ [t_0, t_1], [t_1, t_2], \dots, [t_{L-1}, t_L] \}, \qquad a = t_0 < t_1 < \dots < t_L = b.$$

For each  $[\alpha, \beta] \in \mathcal{T}$ , we let  $||f''||_{[\alpha,\beta]} = ||f''_{[\alpha,\beta]}||_{\infty,[\alpha,\beta]}$ . For partitions  $\mathcal{P}$ , define algorithm and error

$$\mathbf{A}(f;\mathcal{P}) = \{A(f; [\alpha, \beta])\}_{[\alpha, \beta] \in \mathcal{P}}, \qquad \boldsymbol{\varepsilon}_{\mathcal{P}} = \{\boldsymbol{\varepsilon}_{[\alpha, \beta]}\}_{[\alpha, \beta] \in \mathcal{P}}$$

$$\mathbf{A}(f; [t_0, t_1]) \qquad t_0 \leq x \leq t_1,$$

$$A(f; [t_1, t_2]) \qquad t_1 \leq x \leq t_2,$$

$$\vdots \qquad \qquad \vdots,$$

$$A(f; [t_{i-1}, t_i]) \qquad t_{i-1} \leq x \leq t_i,$$

$$\vdots \qquad \qquad \vdots,$$

$$A(f; [t_{i-1}, t_i]) \qquad t_{i-1} \leq x \leq t_i$$

Then our goal is

$$||f - \mathbf{A}(f; \mathcal{P})|| \le ||\varepsilon_{\mathcal{P}}|| < \varepsilon.$$

i.e.

## 3. Algorithms

Now we consider numerical algorithms for solving the problems on a subinterval  $[\alpha, \beta]$  of [a, b].

We use piecewise linear interpolation. The number of possible function values is  $n \in \mathbb{N}$ , then we have

$$x_j = \alpha + (\beta - \alpha)\frac{j-1}{n-1}, \qquad j = 1, \dots, n,$$
(1)

$$A_n(f; [\alpha, \beta]) := \frac{n-1}{\beta - \alpha} \left[ f(x_j)(x_{j+1} - x) + f(x_{j+1})(x - x_j) \right], \qquad x_j \le x \le x_{j+1} \quad (2)$$

The difference between f and its linear spline can be bounded in terms of an integral involving the second derivative using integration by parts. For  $x \in [\alpha, \beta]$  it follows that

$$f(x) - A_n(f; [\alpha, \beta])(x) = f(x) - \frac{n-1}{\beta - \alpha} [f(x_j)(x_{j+1} - x) + f(x_{j+1})(x - x_j)]$$
$$= \frac{n-1}{\beta - \alpha} \int_{x_j}^{x_{j+1}} v_j(t, x) f''(t) dt, \tag{3}$$

$$f'(x) - A_n(f; [\alpha, \beta])'(x) = \frac{n-1}{\beta - \alpha} \int_{x_i}^{x_{j+1}} \frac{\partial v_j}{\partial x}(t, x) f''(t) dt, \tag{4}$$

where the continuous, piecewise differentiable kernel v is defined as

$$v_j(t,x) := \begin{cases} (x_{j+1} - x)(x_j - t), & x_j \le t \le x, \\ (x - x_j)(t - x_{j+1}), & x < t \le x_{j+1}, \end{cases}$$

To derive the error bounds for  $A_n(f; [\alpha, \beta])$  we have:

$$||f - A_{n}(f; [\alpha, \beta])||_{[\alpha, \beta]} \leq \sup_{\substack{x_{j} \leq x \leq x_{j+1} \\ j=1, \dots, n_{i}-1}} |f(x) - A_{n}(f; [\alpha, \beta])(x)|$$

$$= \frac{n-1}{\beta - \alpha} \sup_{\substack{x_{j} \leq x \leq x_{j+1} \\ j=1, \dots, n-1}} \int_{x_{j}}^{x_{j+1}} |v_{j}(t, x)f''(t)| dt$$

$$\leq \frac{n-1}{\beta - \alpha} ||f''||_{[\alpha, \beta]} \sup_{\substack{x_{j} \leq x \leq x_{j+1} \\ j=1, \dots, n-1}} \int_{x_{j}}^{x_{j+1}} |v_{j}(t, x)| dt$$

$$= ||f''||_{[\alpha, \beta]} \sup_{\substack{x_{j} \leq x \leq x_{j+1} \\ j=1, \dots, n-1}} \frac{(x - x_{j})(x_{j+1} - x)}{2}$$

$$= \frac{(\beta - \alpha)^{2}}{8(n-1)^{2}} ||f''||_{[\alpha, \beta]}.$$
 (5)

## 4. Functions in Cones

The challenge in applying the error bounds for the non-adaptive function recovery algorithm is that  $\|f''\|_{[\alpha,\beta]}$  is not known a priori. Our approach is to assume that the input functions lie inside cones.

Let  $\boldsymbol{\tau}_{\mathcal{P}} = \{\tau_{[\alpha,\beta]}\}_{[\alpha,\beta]\in\mathcal{P}}$ . For each  $[\alpha,\beta]\in\mathcal{T}$ , there exists a norm  $\left\|f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha}\right\|_{[\alpha,\beta]}$  on space  $\mathcal{W}^{2,\infty}[\alpha,\beta]$  weaker than  $\|f''\|_{[\alpha,\beta]}$ . Here we define our cone condition

$$C_{\tau} := \left\{ f \in \mathcal{W}^{2,\infty} : \|f''\|_{[\alpha,\beta]} \le \frac{\tau_{[\alpha,\beta]}}{\beta - \alpha} \left\| f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right\|_{[\alpha,\beta]} \right\}. \tag{6}$$

where  $\tau_{[\alpha,\beta]}$  is a parameter depend on the length of interval  $[\alpha,\beta]$ .

Lower bound on  $||f''||_{[\alpha,\beta]}$  can be derived similarly to the previous section using a centered difference. Specifically, for  $n_i \geq 3$ ,

$$F_n(f; [\alpha, \beta]) := \frac{(n-1)^2}{(\beta - \alpha)^2} \sup_{j=1,\dots,n-2} |f(x_j) - 2f(x_{j+1}) + f(x_{j+2})|.$$
 (7)

It follows using the Hölder's inequality that

$$F_n(f; [\alpha, \beta]) = \frac{(n-1)^2}{(\beta - \alpha)^2} \sup_{j=1, \dots, n-2} \left| \int_{x_j}^{x_{j+2}} \left[ \frac{\beta - \alpha}{n-1} - |x - x_{j+1}| \right] f''(x) \, \mathrm{d}x \right|$$

$$\leq \frac{(n-1)^2}{(\beta - \alpha)^2} \sup_{j=1, \dots, n-2} ||f''||_{[\alpha, \beta]} \int_{x_j}^{x_{j+2}} \left| \frac{\beta - \alpha}{n-1} - |x - x_{j+1}| \right| \, \mathrm{d}x = ||f''||_{[\alpha, \beta]} \, .$$

Define

$$G_n(f; [\alpha, \beta]) := \|A_n(f; [\alpha, \beta])' - A_2(f; [\alpha, \beta])'\|_{\infty}$$

$$= \sup_{j=1, \dots, n-1} \left| \frac{n-1}{\beta - \alpha} [f(x_{j+1}) - f(x_j)] - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right|. \quad (8)$$

Note  $G_n(f; [\alpha, \beta])$  never overestimates  $\left\| f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right\|_{[\alpha, \beta]}$  because

$$\left\| f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right\|_{[\alpha, \beta]} = \left\| f' - A_2(f; [\alpha, \beta])' \right\|_{[\alpha, \beta]}$$

$$= \sup_{\substack{x_j \le x \le x_{j+1} \\ j=1, \dots, n-1}} |f'(x) - A_2(f; [\alpha, \beta])'(x)|$$

$$\geq \sup_{j=1, \dots, n-1} \frac{n-1}{\beta - \alpha} \int_{x_j}^{x_{j+1}} \left| f'(x) - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right| dx$$

$$\geq \sup_{j=1, \dots, n-1} \frac{n-1}{\beta - \alpha} \left| \int_{x_j}^{x_{j+1}} \left[ f'(x) - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right] dx \right|$$

$$= \sup_{j=1, \dots, n-1} \frac{n-1}{\beta - \alpha} \left| f(x_{j+1}) - f(x_j) - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right|$$

$$= G_n(f; [\alpha, \beta]).$$

Since  $(f - A_n(f; [\alpha, \beta]))(x)$  vanishes for  $x = \alpha, \beta$ . Using (4) it then follows that

$$\left\| f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right\|_{[\alpha,\beta]} - G_n(f; [\alpha,\beta])$$

$$\leq \left\| f' - A_n(f; [\alpha,\beta])' \right\|_{\infty}$$

$$= \sup_{\substack{x_j \le x \le x_{j+1} \\ j=1,\dots,n-1}} \left| f'(x) - \frac{n-1}{\beta - \alpha} [f(x_{j+1}) - f(x_j)] \right|$$

$$= \frac{n-1}{\beta - \alpha} \sup_{\substack{x_j \le x \le x_{j+1} \\ j=1,\dots,n_{i-1}}} \left| \int_{x_j}^{x_{j+1}} \frac{\partial v_j}{\partial x} (t, x) f''(t) dt \right|$$

$$\leq \frac{n-1}{\beta - \alpha} \left\| f'' \right\|_{[\alpha,\beta]} \sup_{\substack{x_j \le x \le x_{j+1} \\ j=1,\dots,n_{i-1}}} \int_{x_j}^{x_{j+1}} \left| \frac{\partial v_j}{\partial x} (t, x) \right| dt$$

$$= \frac{n-1}{\beta - \alpha} \left\| f'' \right\|_{[\alpha,\beta]} \sup_{\substack{x_j \le x \le x_{j+1} \\ j=1,\dots,n-1}} \left\{ \frac{(\beta - \alpha)^2}{2(n-1)^2} - (x - x_j)(x_{j+1} - x) \right\}$$

$$= \frac{\beta - \alpha}{2(n-1)} \left\| f'' \right\|_{[\alpha,\beta]}.$$

Thus, we obtain if  $n > 1 + \tau_{[\alpha,\beta]}/2$ 

$$G_n(f; [\alpha, \beta]) \leq \left\| f' - \frac{f(\beta) - f(\alpha)}{\beta - \alpha} \right\|_{[\alpha, \beta]} \leq \frac{G_n(f; [\alpha, \beta])}{1 - \tau_{[\alpha, \beta]} / (2n - 2)} \quad \forall f \in \mathcal{C}_{\tau}.$$

Therefore, we obtain

$$F_n(f; [\alpha, \beta]) \le ||f''||_{[\alpha, \beta]} \le \frac{\tau_{[\alpha, \beta]}}{(\beta - \alpha)} \cdot \frac{G_n(f; [\alpha, \beta])}{1 - \tau_{[\alpha, \beta]}/(2n - 2)}$$

$$\Rightarrow \tau_{\min, n} = \frac{F_n(f; [\alpha, \beta])}{G_n(f; [\alpha, \beta])/(\beta - \alpha) + F_n(f; [\alpha, \beta])/(2n_i - 2)} \le \tau_{[\alpha, \beta]}$$

And by (5), we obtain

$$||f - A_n(f; [\alpha, \beta])||_{[\alpha, \beta]} \le \frac{\tau_{[\alpha, \beta]}(\beta - \alpha)G_n(f; [\alpha, \beta])}{4(n - 1)(2n - 2 - \tau_{[\alpha, \beta]})} = \varepsilon_{[\alpha, \beta]}.$$

Let  $\mathbf{n} = \{n_{[\alpha,\beta]}\}_{[\alpha,\beta] \in \mathcal{P}}$  denote the number of points. Define

$$\begin{aligned} \boldsymbol{A_n}(f;\mathcal{P}) &= \{A_{n_{[\alpha,\beta]}}(f;[\alpha,\beta])\}_{[\alpha,\beta] \in \mathcal{P}} \\ \boldsymbol{G_n}(f;\mathcal{P}) &= \{G_{n_{[\alpha,\beta]}}(f;[\alpha,\beta])\}_{[\alpha,\beta] \in \mathcal{P}} \\ \boldsymbol{F_n}(f;\mathcal{P}) &= \{F_{n_{[\alpha,\beta]}}(f;[\alpha,\beta])\}_{[\alpha,\beta] \in \mathcal{P}} \end{aligned}$$

For simplification, we can denote

$$\tau_{[t_{l-1},t_l]} = \tau_l, \qquad n_{[t_{l-1},t_l]} = n_l, \qquad \varepsilon_{[t_{l-1},t_l]} = \varepsilon_l, \qquad l = 1,\ldots,L$$

where  $L = |\mathcal{P}|$  is the cardinality of  $\mathcal{P}$ .

**Algorithm 1** (Locally Adaptive Univariate Function Recovery). Let the sequence of algorithms  $\{A_n(\cdot;\mathcal{P})\}$ ,  $\{G_n(\cdot;\mathcal{P})\}$ ,  $\{F_n(\cdot;\mathcal{P})\}$ , and  $n,\mathcal{P}$  be as described above. Set  $L=1,\ l=1,\ldots,L$ . Choose integer  $\tau_{\text{lo}},\tau_{\text{hi}}$ , where  $\tau_{\text{hi}}\geq\tau_{\text{lo}}$  such that

$$\tau_l = \max \left\{ \left\lceil \tau_{\rm hi} \left( \frac{\tau_{\rm lo}}{\tau_{\rm hi}} \right)^{\frac{1}{1+t_i-t_{i-1}}} \right\rceil, 3 \right\}.$$

Then let  $n_l = \left\lceil \frac{\tau_{[t_{l-1},t_l]}+1}{2} \right\rceil + 1$ ,  $\varepsilon_l = \infty$ . For any error tolerance  $\varepsilon$  and input function f, do the following:

Stage 1. Find the maximum error If  $\max_{l=1,...,L} \varepsilon_l < \varepsilon$ , stop. Otherwise, find

$$k = \arg\max_{l=1,\dots,L} \varepsilon_l.$$

Stage 2. Compute  $\left\|f' - \frac{f(t_k) - f(t_{k-1})}{t_k - t_{k-1}}\right\|_{[t_{k-1}, t_k]}$  and bound  $\|f''\|_{[t_{k-1}, t_k]}$ . Compute  $G_{n_k}(f; [t_{k-1}, t_k])$  in (8) and  $F_{n_k}(f; [t_{k-1}, t_k])$  in (7).

Stage 3. Check the necessary condition for  $f \in \mathcal{C}_{\tau_k}$ . Compute

$$\tau_{\min,n_k} = \frac{F_{n_k}(f;[t_{k-1},t_k])}{G_{n_k}(f;[t_{k-1},t_k])/(t_k-t_{k-1}) + F_{n_k}(f;[t_{k-1},t_k])/(2n_k-2)}.$$

If  $\tau_k \geq \tau_{\min,n_k}$ , then go to stage 4. Otherwise, set  $\tau_k = 2\tau_k$ . If  $n_k \geq (\tau_k + 1)/2$ , then go to stage 4. Otherwise, go to Stage 5.

Stage 4. Check for convergence. Estimate  $\varepsilon_k$ .

$$\varepsilon_k = \frac{\tau_k(t_k - t_{k-1})G_{n_k}(f)}{4(n_k - 1)(2n_k - 2 - \tau_k)}.$$

If  $\varepsilon_{[t_{l-1},t_l]} < \varepsilon$ , go to Stage 1. Otherwise go to Stage 5.

Stage 5. Double the initial number of points and split the interval Let L=L+1. Then

$$\begin{split} t_l &= t_{l-1}, & n_l = n_{l-1}, & \varepsilon_l = \varepsilon_{l-1}, & \tau_l = \tau_{l-1}, & \text{when } l > k, \\ t_k &= \frac{t_{k-1} + t_k}{2}, & \varepsilon_k = \varepsilon_{k+1} = \infty, \\ \tau_k &= \tau_{k+1} = \max \left\{ \left\lceil \tau_{\text{hi}} \left( \frac{\tau_{\text{lo}}}{\tau_{\text{hi}}} \right)^{\frac{1}{1 + t_k - t_{k-1}}} \right\rceil, 3 \right\}. \end{split}$$

Go to Stage 1.

**Theorem 1.** Let  $\{A(\cdot; \mathcal{P})\}$  be the adaptive linear spline defined by Algorithm 1, and let L,  $\tau$ , and  $\varepsilon$  be the inputs and parameters described there. Let  $\mathcal{C}_{\tau}$  be the cone of functions defined in (6). Then it follows that Algorithm 1 is successful for all functions in  $\mathcal{C}_{\tau}$ , i.e.,  $\|f - A(f, \varepsilon; \mathcal{P})\|_{\infty} \leq \varepsilon$ . Moreover, the cost of this algorithm is bounded below and above as follows:

$$\operatorname{cost}(\boldsymbol{A}, f; \mathcal{P}, \varepsilon) \leq \sqrt{\frac{\|\boldsymbol{\tau}\|_{\infty} (b - a)^{2} \|f''\|_{[a, b]}}{4\varepsilon}} + \|\boldsymbol{\tau}\|_{1} + 4L.$$

The algorithm is computationally stable, meaning that maximum costs for all functions, f, with fixed  $||f''||_{\infty}$  are an  $\varepsilon$ -independent constant of each other.

*Proof.* For each subinterval  $[t_{k-1}, t_k] \subset \mathcal{P}$ , we want

$$\frac{\tau_k(t_k - t_{k-1})G_{n_k}(f; [t_{k-1}, t_k])}{4(n_k - 1)(2n_k - 2 - \tau_k)} \le \varepsilon.$$

And we know that

$$\frac{\tau_k(t_k - t_{k-1})G_{n_k}(f; [t_{k-1}, t_k])}{4(n_k - 1)(2n_k - 2 - \tau_k)} \le \frac{\tau_k(t_k - t_{k-1})^2 \|f''\|_{[t_{k-1}, t_k]}}{8(n_k - 1)(2n_k - 2 - \tau_k)}$$

$$n_k \le \sqrt{\frac{\tau_k(t_k - t_{k-1})^2 \|f''\|_{[t_{k-1}, t_k]}}{4\varepsilon}} + \tau_k + 4.$$

Hence,

$$cost(\mathbf{A}, f; \mathcal{P}, \varepsilon) = \sum_{k=1}^{L} n_{k} \leq \sum_{k=1}^{L} \left[ \sqrt{\frac{\tau_{k}(t_{k} - t_{k-1})^{2} \|f''\|_{[t_{k-1}, t_{k}]}}{4\varepsilon}} + \tau_{k} + 4 \right] \\
\leq \sqrt{\frac{\|\boldsymbol{\tau}\|_{\infty} \|f''\|_{[a, b]}}{4\varepsilon}} \sum_{k=1}^{L} (t_{k} - t_{k-1}) + \sum_{k=1}^{L} \tau_{k} + 4L \\
= \sqrt{\frac{\|\boldsymbol{\tau}\|_{\infty} (b - a)^{2} \|f''\|_{[a, b]}}{4\varepsilon}} + \|\boldsymbol{\tau}\|_{1} + 4L.$$