

# Parallel Programming SS21 Final Project

Project 02 – Gaussian Elimination

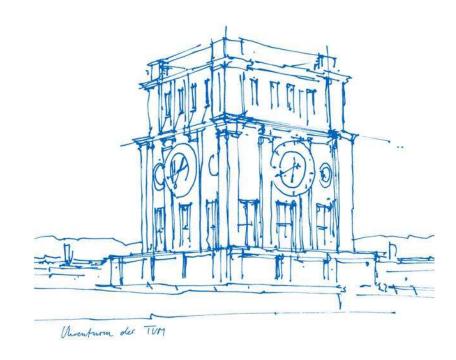
Group 215

09.07.2021

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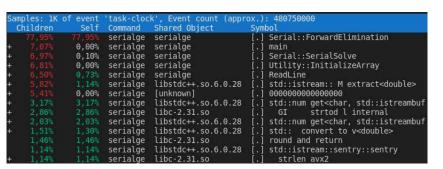
Jing Xiong



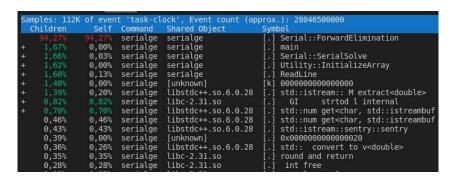
## Sequential code analysis - Profiling



size 1024x1024



size 4096x4096



size 2048x2048

	Children	Self	Command	Shared Object	Symbol
			serialge	serialge	[.] Serial::ForwardElimination
		0,00%	serialge	serialge	[.] main
	3,31%	0,07%	serialge	serialge	[.] Serial::SerialSolve
+	3,24%	0,00%	serialge	serialge	[.] Utility::InitializeArray
+	3,21%	0,28%	serialge	serialge	[.] ReadLine
	2,76%	0,41%	serialge	libstdc++.so.6.0.28	[.] std::istream:: M extract <double></double>
	1,48%	1,48%	serialge	libc-2.31.so	[.] GI strtod l internal
	1,46%	0,00%	serialge	[unknown]	[.] 0000000000000000
+	1,27%	1,27%	serialge	libstdc++.so.6.0.28	[.] std::num get <char, std::istreambuf<="" td=""></char,>
	1,12%	0,00%	serialge	[unknown]	[.] 0x00007ffc69927ad4
+	0,86%	0,86%	serialge	libstdc++.so.6.0.28	[.] std::num get <char, std::istreambuf<="" td=""></char,>
	0.79%	0.79%	serialge	libstdc++.so.6.0.28	[.] std::istream::sentry::sentry
	0.72%	0.72%	serialge	libc-2.31.so	[.] round and return

size 8192x8192

- (	Children	Self	Command	Shared Object	Symbol
			serialge	serialge	[.] Serial::ForwardElimination
H)		0,00%	serialge	serialge	[.] main
+	0,87%	0,02%	serialge	serialge	[.] Serial::SerialSolve
+	0,85%	0,00%	serialge	serialge	[.] Utility::InitializeArray
+	0,84%	0,07%	serialge	serialge	[.] ReadLine
+	0.74%	0,12%	serialge	libstdc++.so.6.0.28	[.] std::istream:: M extract <double></double>
	0,43%	0,43%	serialge	libc-2.31.so	[.] GI strtod l internal
	0,42%	0,00%	serialge	[unknown]	[.] 0000000000000000
	0,38%	0,38%	serialge	libstdc++.so.6.0.28	[.] std::num get <char, std::istreambu<="" td=""></char,>
	0,29%	0,00%	serialge	[unknown]	[.] 0x00007ffe546a5954
	0,24%	0,24%	serialge	libstdc++.so.6.0.28	[.] std::num get <char, std::istreambu<="" td=""></char,>
	0,23%	0,23%	serialge	libstdc++.so.6.0.28	[.] std::istream::sentry::sentry
	0.17%	0.13%	serialge	libstdc++.so.6.0.28	[.] std:: convert to v <double></double>

## Sequential code analysis - Profiling



```
Performance counter stats for './serialge ./ge data/size8192x8192':
                                                1,000 CPUs utilized
      216.171,10 msec task-clock
                      context-switches
                                                0,018 K/sec
           3.996
                     cpu-migrations
                                              # 0,000 K/sec
                                              # 0,607 K/sec
         131.257
                     page-faults
  893.168.658.549
                     cycles
                                                4,132 GHz
                                                   2,75 insn per cycle
                     instructions
2.459.203.457.971
  565.625.107.143
                     branches
                                              # 2616,562 M/sec
                                                   0,01% of all branches
                     branch-misses
      42.950.841
   216,187621527 seconds time elapsed
   216,013328000 seconds user
     0,160030000 seconds sys
```

### Sequential code analysis - Amdahl's law



#### Parameters:

f = fraction of parallel execution f = 99.8%(size2048x2048)

p = number of parallel tasks/threads/processes p = 4

Speedup refers to speedup of function "Solve"

Theoretical Speedup:

$$SU(p) = \frac{T}{T(p)} = \frac{T}{(1-f)*T + \frac{f*T}{p}} = \frac{1}{1-f+\frac{f}{p}} = \frac{1}{1-0.998 + \frac{0.998}{4}} = 3.976$$

## Sequential code analysis - Amdahl's law



#### Parameters:

f = fraction of parallel execution  $f \approx 100\%$ 

p = number of parallel tasks/threads/processes p = 4

Speedup refers to speedup of function "Solve"

$$SU(p) = \frac{T}{T(p)} = \frac{T}{(1-f)*T + \frac{f*T}{p}} = \frac{1}{1-f + \frac{f}{p}} = \frac{1}{1-1 + \frac{1}{4}} = 4$$

**Maximal Speedup: Around 4** 



```
void ForwardElimination(double *matrix, double *rhs, int rows, int columns){
    for(int row = 0; row < rows; row++){</pre>
       // Extract Diagonal element
       int diag idx = row*rows + row;
       double diag_elem = matrix[diag_idx];
       #pragma omp parallel for schedule(dynamic)
       for (int lower_rows=row+1; lower_rows<rows; lower_rows++){</pre>
           int below_diag_idx = lower_rows*rows + row;
           assert(diag_elem!=0);
           double elimination_factor = matrix[below_diag_idx]/diag_elem;
           int element_idx;
            for (int column=row+1; column<columns; column++){</pre>
                // set the column index of the entry to be operated
                element idx = lower rows*rows + column;
                // subtract the row
                matrix[element_idx] -= elimination_factor*matrix[row*rows+column];
            rhs[lower_rows] -= elimination_factor*rhs[row];
            // matrix[below diag idx] = 0.;
```

#### First try

-Totally 3 "for loops" in this function

-Parallel region: 2nd "for loop"

-Schedule:dynamic

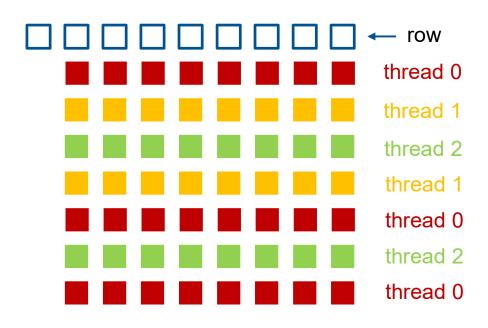
# OpenMP - Intermediate Speed-up results, profiling



size	Overhead of ForwardElimination	Overhead of vmovapd	Speedup (on my PC)	Speedup (on server)
2048	80.99%	69.15%	1.39	<b>≈</b> 10
4096	90.44%	73.9%	1.23	
8192	95.06%	75.82%	1.14	
	of first try eck: ForwardEliminationDa	PS: speedup on speedup for function compared with to optimized seque	ction "Solve" he	

<sup>-</sup>Speedup: Not bad on my PC, but on the server it is not fast enough

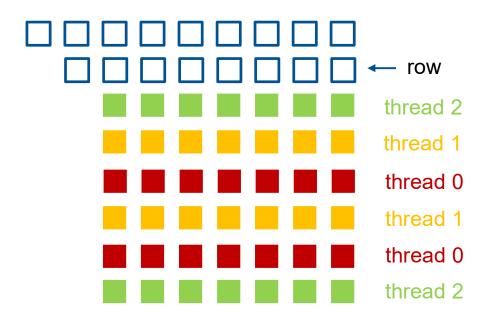




#### First try

-Distribution: The task row addresses distributed to one thread are not continuous





#### First try

-Distribution: The task row addresses distributed to one thread are not continuous and randomly

#### OpenMP - Final Implementation improvements and new speed-up



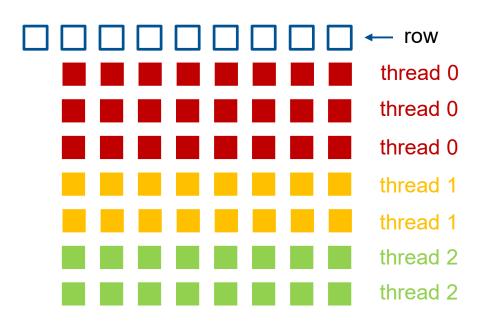
```
void ForwardElimination(double *matrix, double *rhs, int rows, int columns){
   for(int row = 0; row < rows; row++){</pre>
       // Extract Diagonal element
       int diag idx = row*rows + row;
       double diag elem = matrix[diag idx];
        #pragma omp parallel for
       for (int lower_rows=row+1; lower_rows<rows; lower_rows++){</pre>
            int below_diag_idx = lower_rows*rows + row;
            assert(diag elem!=0);
           double elimination_factor = matrix[below_diag_idx]/diag_elem;
            int element idx;
            for (int column=row+1; column<columns; column++){</pre>
                // set the column index of the entry to be operated
                element idx = lower rows*rows + column;
               // subtract the row
                matrix[element_idx] -= elimination_factor*matrix[row*rows+column];
            rhs[lower_rows] -= elimination_factor*rhs[row];
            // set below diagonal elements to 0
            // matrix[below_diag_idx] = 0.;
```

#### **Last Version**

-Schedule:default

-More server friendly

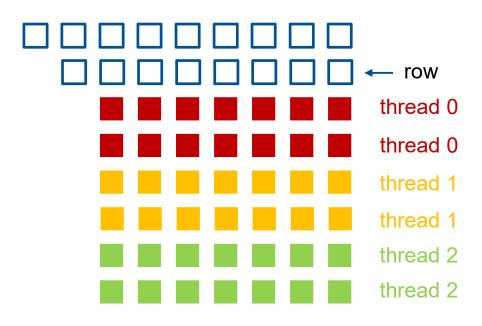




#### **Last version**

-The thread will access similar address in the next round





#### **Last version**

-The thread will access similar address in the next round

# OpenMP - Final Speed-up results, profiling



	Cache misses			
size	dynamic	default		
2048	57%	57%		
4096	84%	83%		
8192	90%	86%		

## OpenMP - Final Speed-up results



size	Overhead of ForwardElimination	Overhead of vmovapd	Speedup (on my PC)	Speedup (on server)
2048	80.99%	69.15%	1.39	≈10(dynamic) ≈ <mark>14</mark> (default)
4096	90.44%	73.9%	1.23	
8192	95.06%	75.82%	1.14	

#### **Result of Final**

-Speedup: Similar as the previous version on my PC, but better on server

# OpenMP - Final Implementation



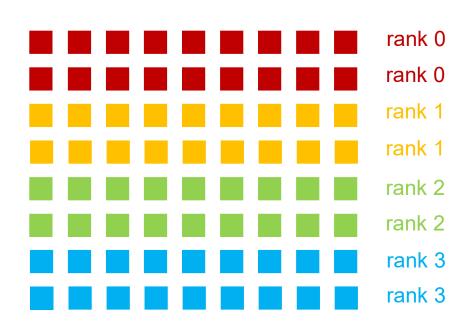
Theoretical speedup: 3~4(with 4 threads) according to Amdahl's law

- Almost all of ForwardElimination is parallized

Bottleneck: Large cost from data transfer

- The bigger the matrix is, the larger the cost from data transfer would be.





#### First try

-Data are equally distributed to each progress in block

-Addresses of data are continuous in each progress



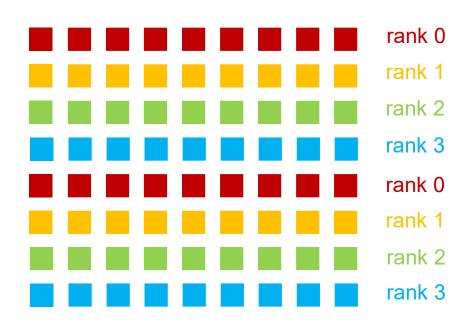






#### MPI - Final Implementation improvements and new speed-up





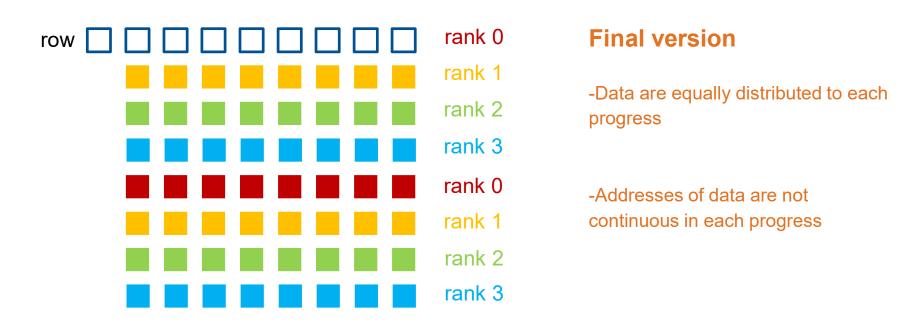
#### **Final version**

-Data are equally distributed to each progress

-Addresses of data are not continuous in each progress

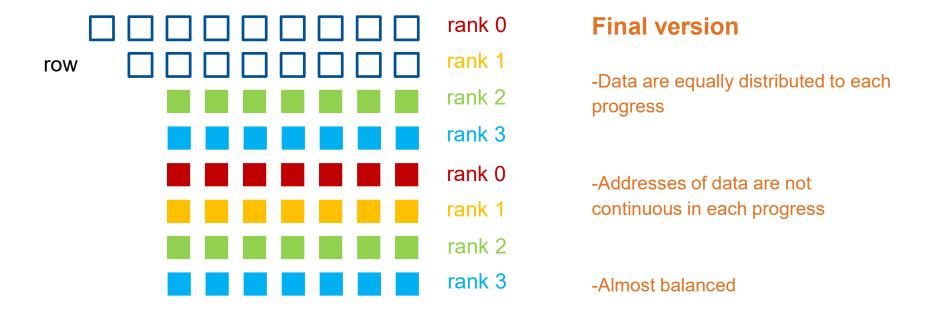
#### MPI - Final Implementation improvements and new speed-up





#### MPI - Final Implementation improvements and new speed-up





## MPI – Speedup



	Speedup	on my PC	Speedup Server	
size	First try	Final version	First try	Final version
2048	1.39	1.39	10.54	10.89
4096	1.21	1.18		
8192	1.13	1.12		

#### Comparison

-Speedup: Very close both on PC and server

-The impact of imbalance is not so obvious



```
Total Lost Samples: 0
# Samples: 2M of event 'cycles'
 Event count (approx.): 1016468430965
 Overhead
               Pid: Command
    25.05%
            7329: mpige
    25.04%
            7330: mpi ge
    25.03%
            7332: mpige
             7331: mpige
    24.88%
    0.00%
              7328: hydra_pmi _proxy
              7326: mpirun
    0.00%
```

#### Comparison

- The impact of imbalance is not so obvious
- The load seems balanced because there is a barrier at each end of the first "for loop" to update "row+1"

```
MPI_Bcast(matrix+(row+1)*rows, columns, MPI_DOUBLE, r, MPI_COMM_WORLD);
MPI_Bcast(rhs+row+1, 1, MPI_DOUBLE, r, MPI_COMM_WORLD);
```

# MPI - Profiling



0170	First tr	-y	Final version		
size	ForwardElimination	vmovapd	ForwardElimination	vmovapd	
2048	45.13%	52.3%	53.37%	67.27%	
4096	57.31%	56.98%	73.14%	72.68%	
0400	04.040/	F7 7F0/	05.000/	74 770/	
8192	64.01%	57.75%	85.03%	74.77%	

#### Comparison

- The data address in the first try are continuous, so that it would be easier to access

# MPI - Final Speed-up results, profiling



64.01%	mpi ge	mpi ge	[.]	MPI::Forw
3.80%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
1.12%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.99%	mpige	[kernel.kallsyms]	[ k]	do_syscal 🛚
0.95%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 \( \)
0.82%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 \( \)
0.81%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.81%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.81%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.80%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.80%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8
0.80%	mpige	libmpich.so.0.0.0	[.]	0 x 0 0 0 0 0 0 0 8

#### **Cost from MPI**

-Using MPI might cause some communication cost

#### MPI - Final Implementation



Theoretical speedup: 3~4(with 4 threads) according to Amdahl's law

- Almost all of ForwardElimination is parallized

Bottleneck: Large cost from data transfer, MPI communication

- The bigger the matrix is, the larger the cost from data transfer would be
- MPI communication would cost some time



#### Combine of OMP and MPI approach

- Two threads for MPI(distributed nodes) and two for OMP(multiple cores)
- The structure is almost the same as MPI approach
- OMP is nested within the MPI process

```
void ForwardElimination(double *matrix, double *rhs, int rows, int columns, int rank, int size){
   for(int row = 0; row < rows-1; row++){</pre>
      int block_idx = row / size;
      int l_rows_start;
                                                          Master thread
      if((row%size)<rank){
         l_rows_start = block_idx*size + rank;
      elsel
          l_rows_start = (block_idx+1)*size + rank;
                                                           Worker threads of the
      #pragma omp parallel for num threads(2)
                                                           master processor
      for (){
                                                          End OpenMP
      int r = (row+1)%size;
      MPI_Bcast(matrix+(row+1)*rows, columns, MPI_DOUBLE, r, MPI_COMM_WORLD);
      MPI_Bcast(rhs+row+1, 1, MPI_DOUBLE, r, MPI_COMM_WORLD);
```

### Hybrid - Final Performance Results



size	Speedup	*with 4 threads on own PC
2048	1.39	
4096	1.20	
8192	1.13	

```
# Overhead Pid: Command

# ......

# 27.61% 4122: hybridge
26.99% 4123: hybridge
18.20% 4125: hybridge
18.00% 4124: hybridge
```

The two threads responsible for the MPI process share more load

Theoretical speedup: 3~4 (with 4 threads) according to Amdahl's law

- Almost all of ForwardElimination is parallized

Bottleneck: Large cost from data transfer, MPI communication

- The bigger the matrix is, the larger the cost from data transfer would be
- MPI communication would cost some time
- Load balance

-But this is a simulation on a single machine, it should perform better on real distribution nodes

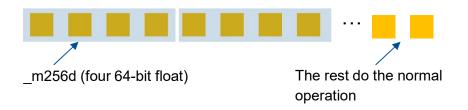
#### Bonus (SIMD) - Parallelized implementation and approach



```
for (int column=row+1; column<columns; column++){</pre>
   // Set the column index of the entry to be operated
   element idx = lower rows*rows + column;
   matrix[element idx] -= elimination factor*matrix[row*rows+column];
```

This is operation on the row and can be vectorized.

#### **AVX Intrinsics**



```
int vectorLen = 4;
int end = columns - columns % vectorLen;
 m256d elimination factor vec= mm256 set1 pd(elimination factor);
 m256d element row, select row;
for(int column = row + 1;column < end;column += vectorLen){</pre>
   element row = mm256 loadu pd(matrix + lower rows*rows + column);
   select row = mm256 loadu pd(matrix + row*rows + column);
     m256d temp row = mm256 mul pd(elimination factor vec, select row);
   element row = mm256 sub pd(element row, temp row);
    mm256 storeu pd(matrix + lower rows*rows + column, element row);
for(int column = end; column < columns; column++){</pre>
   matrix[lower rows*rows + column] -= elimination factor*matrix[row*rows+column];
rhs[lower rows] -= elimination factor*rhs[row];
```

#### Bonus (SIMD) - Final Performance Results



size	Hybrid Speedup	Bonus *with 4 thread Speedup <sup>on own PC</sup>
2048	1.39	1.40
4096	1.20	1.20
8192	1.13	1.13

No significant improvement on speedup.

#### Let's look at assembly of hybridge.cpp

```
.L8:

vmovupd (%rax,%r12), %xmm6

vinsertf128 $0x1, 16(%rax,%r12), %ymm6, %ymm1

vmovupd 0(%r13,%r12), %xmm5

vinsertf128 $0x1, 16(%r13,%r12), %ymm5, %ymm0

vmulpd %ymm3, %ymm1, %ymm1

vsubpd %ymm1, %ymm0, %ymm0

vmovups %xmm0, 0(%r13,%r12)

vextractf128 $0x1, %ymm0, 16(%r13,%r12)

addq $32, %r12

cmpq %rsi, %r12

jne .L8
```

Many vectorized AVX/SSE instructions have been generated after compilation. It's the characteristic of gcc -o3

#### Conclusion



\*Speedup with 4 threads on own PC

			I	I
size	OpenMP	MPI	OMP+MPI	OMP+MPI+SIMD
2048	1.39	1.39	1.39	1.40
4096	1.23	1.18	1.20	1.20
8192	1.14	1.12	1.13	1.13
	I	I		

- Speedup are close among the four approaches.
- Speedup doesn't reach the theoretical speedup mostly due to cost of data transfer
- It would be more appropriate to use MPI and hybrid method for distributed system