

Recursion vs. DFS.

- # DFS is more general scope. it's one kind of search algorithm.
- # can be implemented in either **recursive** or **iterative** ways.

DFS:

recall "pre-order" to traverse a tree.

Easy to use recursion.

DFS Common Questions:

- # 1. Print all subsets of a set.
- # 2. Print all valid permutations of $(1, (1, 1))$.
- # 3. Combinations of Coins
- # 4. **Will be covered later.**

Basic Idea: ☆☆☆

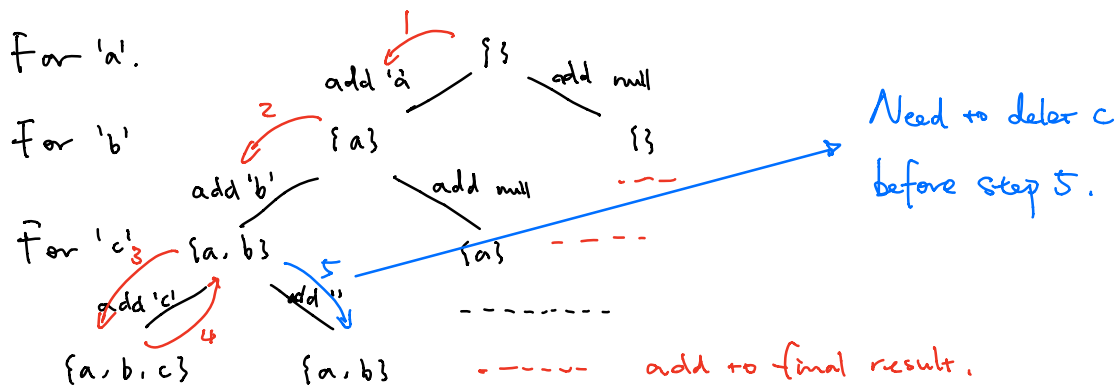
1. How many levels in the recursion tree?
What does it store on each level?
2. How many diff states should we try to put on each level.

Q1. Print all subsets of a set S.

eg: {a, b, c}.

Basic Idea:

1. How many levels in the recursion tree?
3 levels. For each level, add or not add to final set
2. How many diff states should we try to put on each level?
2 states. add new elements / not add.



inputs: (char[] input, int index, StringBuilder sb).
 (current level).

Base Case: if index hits length of input,
 add to final result.

Recursion:

sb.append(input[index]);
 recursion on (index + 1); } add.

★ sb.deleteCharAt(sb.length() - 1);
 recursion on (index + 1); } Not add.

Q2. Valid Permutation using the parenthesis provided. eg. ()()()

Basic Idea:

1. How many levels?
 6. each level represents one position.
2. How many states?
 2. either add "(" or add ")"

Restriction: whenever we add another ")", we make

since # of "(" added so far > # of ")" so far

inputs:

int n — # of pairs of ()

int l — # of left parentheses

int r — # of right parentheses.

StringBuilder.

if (l == n && r == n). add to rst.

if (l < n)

append "(".

recursion.

deleteCharAt

if (r < l).

append ")"

recursion

deleteCharAt.

Time: $O(2^n \cdot n)$.

Space: $O(2n) \rightarrow O(n)$

Q3. All Combinations of Coins. that can sum up to n.

e.g: n = 99. [25, 10, 5, 1]

Soln #1. #1. How many levels?

99: $99/1 = 99$.

#2. How many diff states?

4 states: 25, 10, 5, 1.

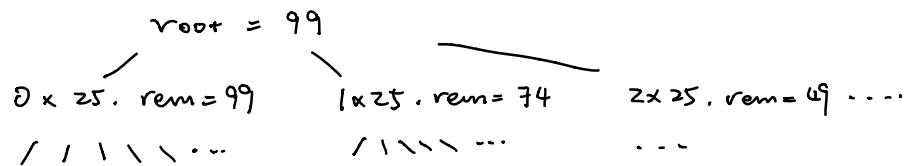
Low efficiency: 4^{99} .

Soln #2 #1. How many levels?

4: 25, 10, 5, 1.

#2. How many states?

Dynamically changing.



Time: $O(99^4)$.

inputs: money left, index level, int[] sol.

Base Case: if index = coins.length.

sol[index] = moneyleft (last coin is 1).

Recursion: for $i = 0 : i < \text{moneyleft} / \text{coins}[\text{index}]$.

sol[index] = i.

recursion(moneyleft - i * coins[index], level+1, sol);

Q4. Permutations of a string. (no duplicates).

Basic idea:

#1. How many levels?

Length of string. Each level represents one position.

#2. How many states?

$n \rightarrow n-1 \rightarrow n-2 \rightarrow \dots \rightarrow 2 \rightarrow 1$.

Decrease by 1 for each level.

input: char[] s, index.

[0... index]. positions that are confirmed.

Index --- $n-1$. positions not sure.

Base case: if index hits length of string,
add to res & return

Recursion: for $i = \text{index}; i < \text{length}; i++$.

Swap (i , index)

recursion (s , index+1);

swap (i , index).

Time: $O(n! \cdot n)$

Space: $O(n)$

Notes: Whenever every permutation contains all elements
in the initial input, then consider **swapping**.



#1. Subset: order NOT matter. States constant

#2. () : order NOT matter. States constant, cut subtree

#3. 99 coins: order NOT matter. States dynamically change

#4. Permutation: order MATTER. **Swap & swap**.