# A. Mex in the Grid

Input file: standard input
Output file: standard output

Time limit: 2 seconds
Memory limit: 256 megabytes

You are given  $n^2$  cards with values from 0 to  $n^2-1$ . You are to arrange them in a n by n grid such that there is **exactly** one card in each cell.

The MEX (minimum excluded value) of a subgrid\* is defined as the smallest non-negative integer that does not appear in the subgrid.

Your task is to arrange the cards such that the sum of MEX values over all  $\left(\frac{n(n+1)}{2}\right)^2$  subgrids is maximized.

#### Input

Each test contains multiple test cases. The first line contains the number of test cases t ( $1 \le t \le 100$ ). The description of the test cases follows.

The first line of each test case contains a single integer n ( $1 \le n \le 500$ ) — the side length of the grid.

It is guaranteed that the sum of n over all test cases does not exceed 1000.

## Output

For each test case, output n lines, each containing n integers representing the elements of the grid.

If there are multiple answers, you can output any of them.

Standard Input	Standard Output
2	0 1
2	2 3
3	8 4 5
	6 0 1
	7 2 3

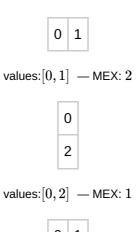
### **Note**

In the first test case, one valid arrangement is:

There are 9 subgrids in total, and the 4 of them with non-zero MEX are shown below:

0

<sup>\*</sup>A subgrid of a n by n grid is specified by four numbers  $l_1, r_1, l_2, r_2$  satisfying  $1 \le l_1 \le r_1 \le n$  and  $1 \le l_2 \le r_2 \le n$ . The element in the i-th row and the j-th column of the grid is part of the subgrid if and only if  $l_1 \le i \le r_1$  and  $l_2 \le j \le r_2$ .



 $\text{values:} [0,1,2,3] \, - \text{MEX:} \, 4$ 

The sum of MEX over all subgrids would be 1+2+1+4=8. It can be proven that no other arrangements have a larger sum of MEX values.

# **B.** Quartet Swapping

Input file: standard input
Output file: standard output

Time limit: 2 seconds
Memory limit: 256 megabytes

You are given a permutation a of length  $n^*$ . You are allowed to do the following operation any number of times (possibly zero):

• Choose an index  $1 \leq i \leq n-3$ . Then, swap  $a_i$  with  $a_{i+2}$ , and  $a_{i+1}$  with  $a_{i+3}$  simultaneously. In other words, permutation a will be transformed from  $[\ldots,a_i,a_{i+1},a_{i+2},a_{i+3},\ldots]$  to  $[\ldots,a_{i+2},a_{i+3},a_i,a_{i+1},\ldots]$ .

Determine the lexicographically smallest permutation<sup>†</sup> that can be obtained by applying the above operation any number of times.

• in the first position where x and y differ, the array x has a smaller element than the corresponding element in y.

## Input

Each test contains multiple test cases. The first line contains the number of test cases t ( $1 \le t \le 1000$ ). The description of the test cases follows.

The first line of each test case contains a single integer n ( $4 \le n \le 2 \cdot 10^5$ ) — the length of permutation a.

The second line contains n integers  $a_1, a_2, \ldots, a_n$  ( $1 \le a_i \le n$ ) — the elements of permutation a.

It is guaranteed that the sum of n over all test cases does not exceed  $2 \cdot 10^5$ .

### Output

For each test case, output the lexicographically smallest permutation that can be obtained by applying the above operation any number of times.

Standard Input	Standard Output
3	1 2 3 4
4	2 1 3 4 5
3 4 1 2	2 1 4 3 6 5 8 7 10 9
5	
5 4 3 1 2	
10	
10 9 8 7 6 5 4 3 2 1	

#### **Note**

In the first test case, an operation can be done on index i=1, and the permutation will become [1,2,3,4], which is the lexicographically smallest permutation achievable.

<sup>\*</sup>A permutation of length n is an array consisting of n distinct integers from 1 to n in arbitrary order. For example, [2,3,1,5,4] is a permutation, but [1,2,2] is not a permutation (2 appears twice in the array), and [1,3,4] is also not a permutation (n=3 but there is 4 in the array).

<sup>&</sup>lt;sup>†</sup> An array x is lexicographically smaller than an array y of the same size if and only if the following holds:

In the second test case, we can do the following sequence of operations:

- Do an operation on index i=2. The permutation becomes [5,1,2,4,3].
- Do an operation on index i=1. The permutation becomes [2,4,5,1,3].
- Do an operation on index i=2. The permutation becomes [2,1,3,4,5].

# C. 23 Kingdom

Input file: standard input
Output file: standard output

Time limit: 4 seconds
Memory limit: 256 megabytes

The *distance* of a value x in an array c, denoted as  $d_x(c)$ , is defined as the largest gap between any two occurrences of x in c.

Formally,  $d_x(c) = \max(j-i)$  over all pairs i < j where  $c_i = c_j = x$ . If x appears only once or not at all in c, then  $d_x(c) = 0$ .

The *beauty* of an array is the sum of the distances of each distinct value in the array. Formally, the beauty of an array c is equal to  $\sum_{1 \le x \le n} d_x(c)$ .

Given an array a of length n, an array b is *nice* if it also has length n and its elements satisfy  $1 \le b_i \le a_i$  for all  $1 \le i \le n$ . Your task is to find the maximum possible beauty of any nice array.

## Input

Each test contains multiple test cases. The first line contains the number of test cases t ( $1 \le t \le 10^4$ ). The description of the test cases follows.

The first line of each test case contains a single integer n ( $1 \le n \le 2 \cdot 10^5$ ) — the length of array a.

The second line of each test case contains n integers  $a_1, a_2, \ldots, a_n$  ( $1 \le a_i \le n$ ) — the elements of array a.

It is guaranteed that the sum of n over all test cases does not exceed  $2 \cdot 10^5$ .

#### Output

For each test case, output a single integer representing the maximum possible beauty among all nice arrays.

Standard Input	Standard Output
4	4
4	1
1 2 1 2	16
2	16
2 2	
10	
1 2 1 5 1 2 2 1 1 2	
8	
1 5 2 8 4 1 4 2	

#### Note

In the first test case, if b = [1, 2, 1, 2], then  $d_1(b) = 3 - 1 = 2$  and  $d_2(b) = 4 - 2 = 2$ , resulting in a beauty of 2 + 2 = 4. It can be proven that there are no nice arrays with a beauty greater than 4.

In the second test case, both b=[1,1] and b=[2,2] are valid solutions with a beauty of 1.

In the third test case, if b=[1,2,1,4,1,2,1,1,1,2] with  $d_1(b)=9-1=8$ ,  $d_2(b)=10-2=8$ , and  $d_4(b)=0$ , resulting in a beauty of 16.

# D. Mani and Segments

Input file: standard input
Output file: standard output

Time limit: 3 seconds
Memory limit: 256 megabytes

An array b of length |b| is *cute* if the sum of the length of its Longest Increasing Subsequence (LIS) and the length of its Longest Decreasing Subsequence (LDS)\* is **exactly** one more than the length of the array. More formally, the array b is cute if LIS(b) + LDS(b) = |b| + 1.

You are given a permutation a of length  $n^{\dagger}$ . Your task is to count the number of non-empty subarrays<sup>‡</sup> of permutation a that are cute.

The longest increasing (decreasing) subsequence of an array is the longest subsequence such that its elements are ordered in strictly increasing (decreasing) order.

# Input

Each test contains multiple test cases. The first line contains the number of test cases t ( $1 \le t \le 10^4$ ). The description of the test cases follows.

The first line of each test case contains a single integer n ( $1 \le n \le 2 \cdot 10^5$ ) — the length of permutation a.

The second line of each test case contains n integers  $a_1, a_2, \ldots, a_n$  ( $1 \le a_i \le n$ ) — the elements of permutation a.

It is guaranteed that the sum of n over all test cases does not exceed  $2\cdot 10^5$ .

#### Output

For each test case, output the number of cute non-empty subarrays of permutation a.

Standard Input	Standard Output
5	6
3	15
3 1 2	9
5	28
2 3 4 5 1	36
4	
3 4 1 2	
7	
1 2 3 4 5 6 7	
10	
7 8 2 4 5 10 1 3 6 9	

<sup>\*</sup>A sequence x is a subsequence of a sequence y if x can be obtained from y by the deletion of several (possibly, zero or all) element from arbitrary positions.

<sup>&</sup>lt;sup>†</sup> A permutation of length n is an array consisting of n distinct integers from 1 to n in arbitrary order. For example, [2,3,1,5,4] is a permutation, but [1,2,2] is not a permutation (2 appears twice in the array), and [1,3,4] is also not a permutation (n=3 but there is 4 in the array).

<sup>&</sup>lt;sup>‡</sup> An array x is a subarray of an array y if x can be obtained from y by the deletion of several (possibly, zero or all) elements from the beginning and several (possibly, zero or all) elements from the end.

# **Note**

In the first test case, all of the 6 non-empty subarrays are cute:

- [3]: LIS([3]) + LDS([3]) = 1 + 1 = 2.
- [1]: LIS([1]) + LDS([1]) = 1 + 1 = 2.
- [2]: LIS([2]) + LDS([2]) = 1 + 1 = 2.
- [3,1]: LIS([3,1]) + LDS([3,1]) = 1 + 2 = 3.
- [1,2]: LIS([1,2]) + LDS([1,2]) = 2 + 1 = 3.
- [3,1,2]: LIS([3,1,2]) + LDS([3,1,2]) = 2 + 2 = 4.

In the second test case, one of the cute subarrays is [2,3,4,5,1] as LIS([2,3,4,5,1])=4 and LDS([2,3,4,5,1])=2, which satisfies 4+2=5+1.

# E. Kia Bakes a Cake

Input file: standard input
Output file: standard output

Time limit: 6 seconds
Memory limit: 512 megabytes

You are given a binary string s of length n and a tree T with n vertices. Let k be the number of 1s in s. We will construct a complete undirected weighted graph with k vertices as follows:

- For each  $1 \le i \le n$  with  $s_i = 1$ , create a vertex labeled i.
- For any two vertices labeled u and v that are created in the above step, define the edge weight between them w(u, v) as the distance\* between vertex u and vertex v in the tree T.

A **simple** path<sup>†</sup> that visits vertices labeled  $v_1, v_2, \ldots, v_m$  in this order is *nice* if for all  $1 \le i \le m-2$ , the condition  $2 \cdot w(v_i, v_{i+1}) \le w(v_{i+1}, v_{i+2})$  holds. In other words, the weight of each edge in the path must be at least twice the weight of the previous edge. Note that  $s_{v_i} = 1$  has to be satisfied for all  $1 \le i \le m$ , as otherwise, there would be no vertex with the corresponding label.

For each vertex labeled i ( $1 \le i \le n$  and  $s_i = 1$ ) in the complete undirected weighted graph, determine the maximum number of vertices in any nice simple path starting from the vertex labeled i.

#### Input

Each test contains multiple test cases. The first line contains the number of test cases t ( $1 \le t \le 10^4$ ). The description of the test cases follows.

The first line of each test case contains a single integer n ( $1 \le n \le 7 \cdot 10^4$ ) — the length of the binary string s and the number of vertices in the tree T.

The second line of each test case contains a binary string with n characters  $s_1 s_2 \dots s_n$  ( $s_i \in \{0, 1\}$ ) — the string representing the vertices to be constructed in the complete undirected weighted graph.

Each of the next n-1 lines contains two integers u and v ( $1 \le u, v \le n$ ) — the endpoints of the edges of the tree T.

It is guaranteed that the given edges form a tree.

It is guaranteed that the sum of n over all test cases does not exceed  $7 \cdot 10^4$ .

## Output

For each test case, output n integers, the i-th integer representing the maximum number of vertices in any nice simple path starting from the vertex labeled i. If there is no vertex labeled i, i.e.,  $s_i = 0$ , output -1 instead.

Standard Output
-1 3 3 3 3
-1 5 4 -1 4 -1 5 5 5 5 -1 4 -1 5 5 -1 3 -1 1
-

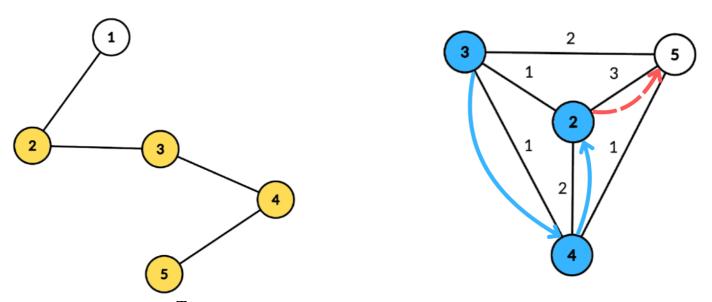
<sup>\*</sup>The distance between two vertices a and b in a tree is equal to the number of edges on the unique simple path between vertex a and vertex b.

 $<sup>^\</sup>dagger$  A path is a sequence of vertices  $v_1,v_2,\ldots,v_m$  such that there is an edge between  $v_i$  and  $v_{i+1}$  for all  $1\leq i\leq m-1$ . A simple path is a path with no repeated vertices, i.e.,  $v_i\neq v_j$  for all  $1\leq i< j\leq m$ .

```
1 2
2 3
3 4
4 5
17
01101011110101101
1 2
2 3
3 4
4 5
6 7
7 8
8 9
9 10
10 11
11 12
12 13
13 14
14 15
15 16
16 17
2
01
1 2
```

### **Note**

In the first test case, the tree T and the constructed graph are as follows:



Left side is the tree T with selected nodes colored yellow. The right side is the constructed complete graph. The nice path shown in the diagram is  $3 \to 4 \to 2$ . The path is nice as w(4,2)=2 is at least twice of w(3,4)=1. Extending the path using  $2 \to 5$  is not possible as w(2,5)=3 is less than twice of w(4,2)=2.

In the second test case, the tree T is a simple path of length 17. An example of a nice path starting from the vertex labeled 2 is  $2 \to 3 \to 5 \to 9 \to 17$ , which has edge weights of 1, 2, 4, 8 doubling each time.

# F. Shoo Shatters the Sunshine

Input file: standard input
Output file: standard output

Time limit: 7 seconds
Memory limit: 256 megabytes

You are given a tree with n vertices, where each vertex can be colored red, blue, or white. The *coolness* of a coloring is defined as the maximum distance\* between a red and a blue vertex.

Formally, if we denote the color of the i-th vertex as  $c_i$ , the coolness of a coloring is  $\max d(u,v)$  over all pairs of vertices  $1 \le u,v \le n$  where  $c_u$  is red and  $c_v$  is blue. If there are no red or no blue vertices, the coolness is zero.

Your task is to calculate the sum of coolness over all  $3^n$  possible colorings of the tree, modulo  $998\ 244\ 353$ .

#### Input

Each test contains multiple test cases. The first line contains the number of test cases t (1  $\leq t \leq$  50). The description of the test cases follows.

The first line of each test case contains a single integer n ( $2 \le n \le 3000$ ) — the number of vertices in the tree

Each of the next n-1 lines contains two integers u and v ( $1 \le u, v \le n$ ) — the endpoints of the edges of the tree

It is guaranteed that the given edges form a tree.

It is guaranteed that the sum of n over all test cases does not exceed 3000.

#### Output

For each test case, output the sum of coolness over all  $3^n$  possible colorings of the tree, modulo  $998\ 244\ 353$ .

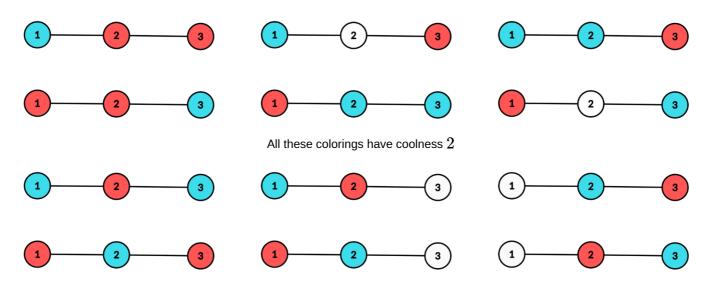
Standard Input	Standard Output
3	18
3	1920
1 2	78555509
2 3	
6	
1 2	
1 3	
1 4	
3 5	
5 6	
17	
1 2	
1 3	
1 4	
1 5	

<sup>\*</sup>The distance between two vertices a and b in a tree is equal to the number of edges on the unique simple path between vertex a and vertex b.

2 6		
2 7		
2 8		
3 9		
3 10		
7 11		
7 12		
11 13		
13 14		
14 15		
10 16		
16 17	,	

# **Note**

In the first test case, there are 12 colorings that have at least one blue and one red node. The following pictures show their coloring and their coolness:



All these colorings have coolness  $\boldsymbol{1}$ 

Therefore, the sum of coolness over all possible colorings is  $6 \cdot 2 + 6 \cdot 1 = 18$ .

In the second test case, the following are some examples of colorings with a coolness of 3:

