

Symbolic Execution

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Lecture #8 out of 10

80 minutes

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In Theory

In Practice

Test Case Generation

Concolic Testing

Further Reading/Watching

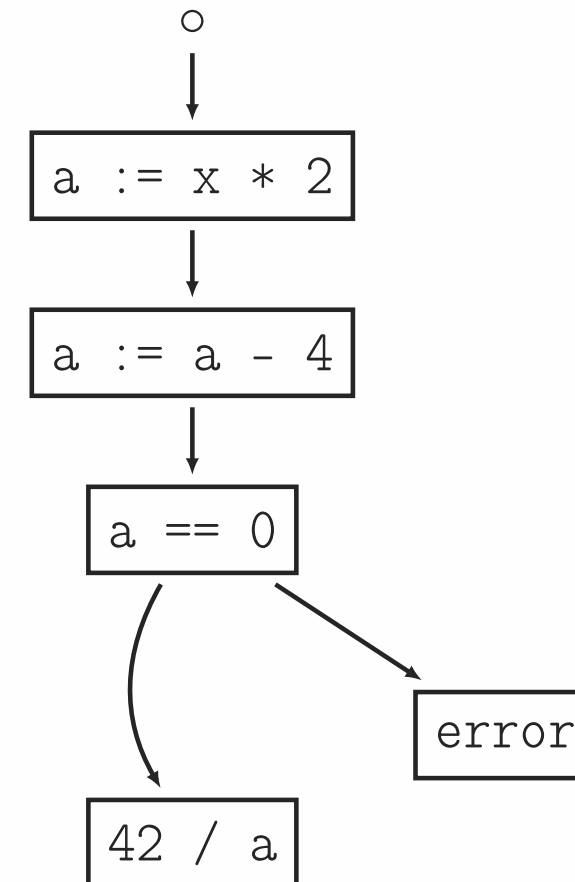
Chapter #1:

In Theory

[[CFG](#) Feasibility Infeasible Symbols PC Solver]

Control Flow Graph

```
int f(int x) {  
    int a = x * 2;  
    a = a - 4;  
    if (a == 0)  
        error("Div by zero!");  
    return 42 / a;  
}
```

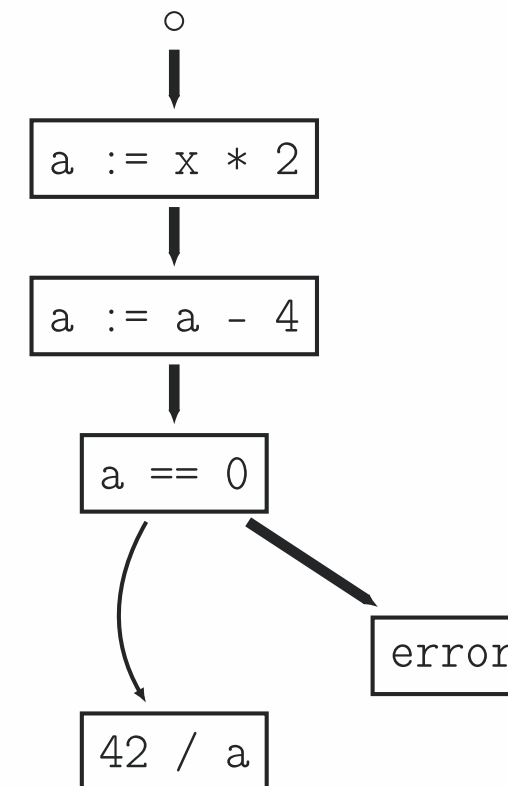


[CFG [Feasibility](#) Infeasible Symbols PC Solver]

Path Feasibility

A path is *feasible* if there exists an input \mathcal{I} to the program that covers the path; i.e., when program is executed with \mathcal{I} as input, the path is taken.

```
int f(int x) {  
    int a = x * 2;  
    a = a - 4;  
    if (a == 0)  
        error("Div by zero!");  
    return 42 / a;  
}
```



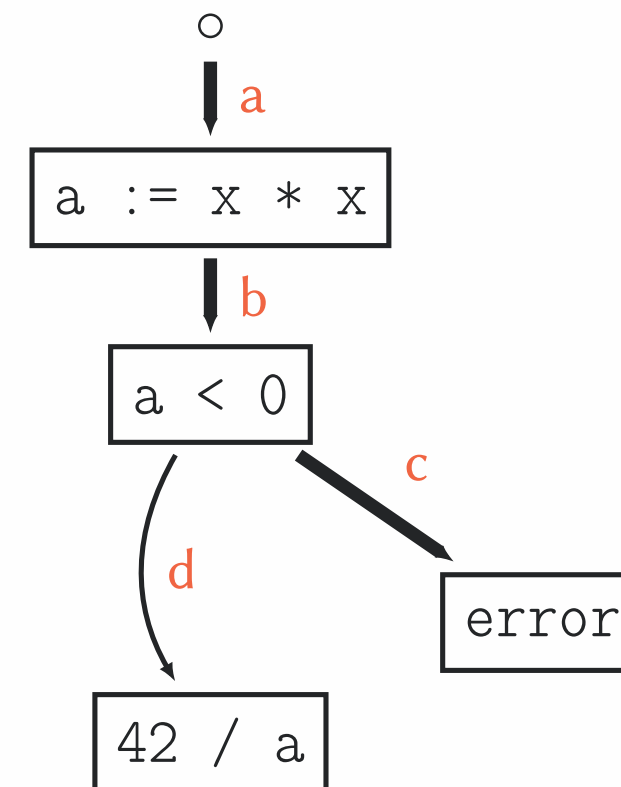
[CFG [Feasibility](#) Infeasible Symbols PC Solver]

[CFG Feasibility [Infeasible](#) Symbols PC Solver]

Infeasible Path

A path is *infeasible* if there exists no input \mathcal{I} that covers the path.

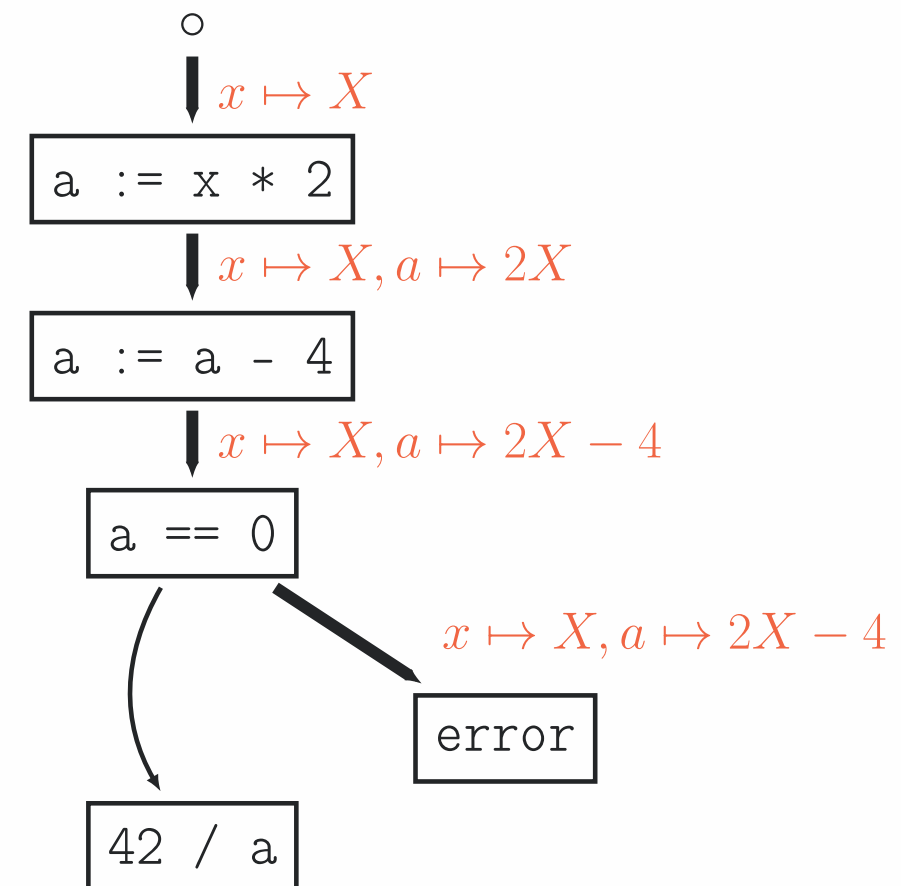
```
int f(int x) {  
  int a = x * x;  
  if (a < 0)  
    error("Too small!");  
  return 42 / a;  
}
```



[CFG Feasibility Infeasible Symbols PC Solver]

Symbols

```
int f(int x) {  
  int a = x * 2;  
  a = a - 4;  
  if (a == 0)  
    error("Div by zero!");  
  return 42 / a;  
}
```

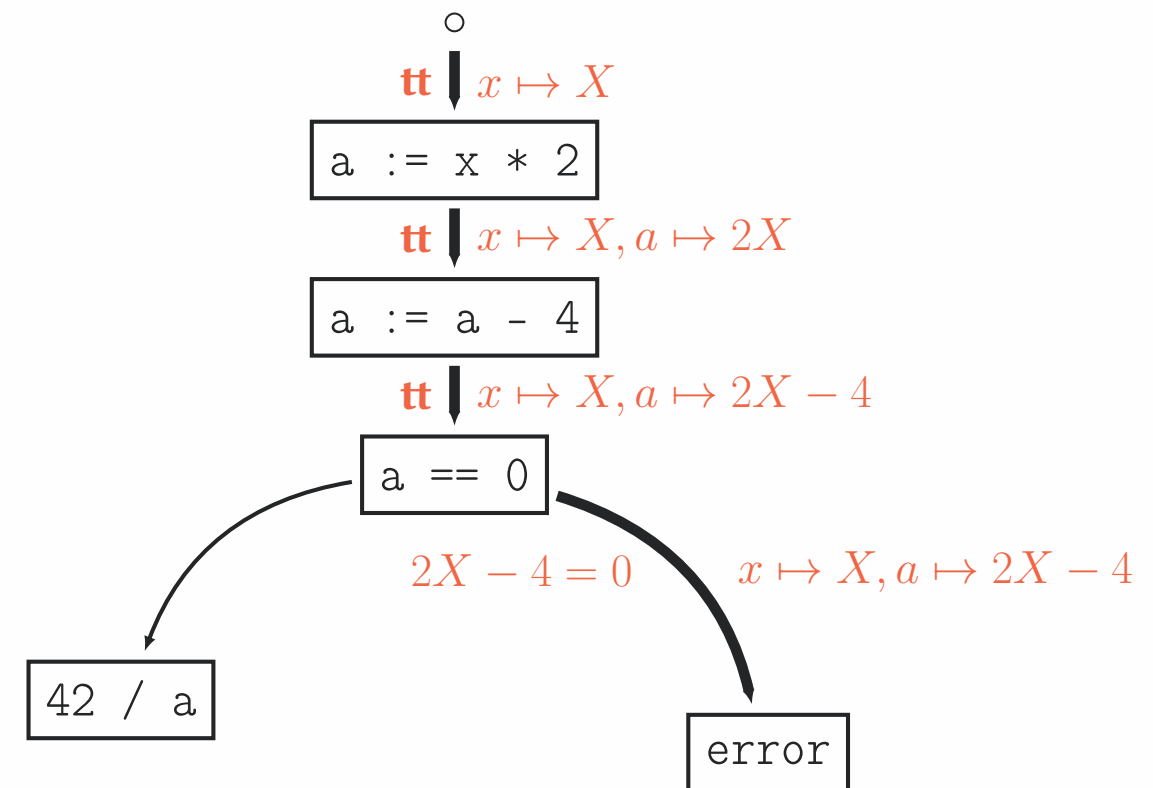


[CFG Feasibility Infeasible Symbols [PC](#) Solver]

Path Conditions

Path condition is a condition on the input symbols such that if a path is feasible its path-condition is satisfiable.

```
int f(int x) {  
  int a = x * 2;  
  a = a - 4;  
  if (a == 0)  
    error("Div by zero!");  
  return 42 / a;  
}
```



[CFG Feasibility Infeasible Symbols [PC](#) Solver]

[CFG Feasibility Infeasible Symbols PC [Solver](#)]

Constraint Solver

A *constraint solver* is a tool that finds satisfying assignments for a *constraint*, if it is satisfiable.

A *solution* of the constraint is a set of assignments, one for each free variable that makes the constraint satisfiable.

Constraint:

$$x \mapsto X, a \mapsto 2X - 4$$
$$2X - 4 = 0$$

Solution:

$$X = 2$$

Chapter #2:

In Practice

[[SAT](#) SMT Unsolvable Explosion ChatGPT CSA]

SAT Solvers

SAT solver is a computer program which aims to solve the *Boolean satisfiability problem*: whether the variables of a given Boolean formula can be consistently replaced by the values TRUE or FALSE in such a way that the formula evaluates to TRUE.

Examples:

$$a \wedge b \rightarrow \dots$$

$$a \wedge b \wedge \neg a \rightarrow \dots$$

$$a \vee b \vee \neg a \rightarrow \dots$$

$$a \wedge (\mathbf{ff} \vee \mathbf{tt}) \rightarrow \dots$$

All expressions are in Boolean logic.

SMT Solvers

SMT solver is a computer program which aims to solve the *satisfiability modulo theories*: determine whether a mathematical formula is satisfiable.

Examples:

$$a < 5 \wedge a > 3 \rightarrow \dots$$

$$a < 5 \wedge f(a) > 42 \rightarrow \dots$$

$$a < 5 \vee a > 10 \vee \neg a \rightarrow \dots$$

$$a \wedge \mathbf{ff} \wedge x = 7 \rightarrow \dots$$

SMT solvers: Z3, cvc5, Yices, and many more...

Unsolvable Constraints

Symbolic execution cannot handle *unsolvable* or almost unsolvable constraints.

```
void enter(String p) {  
    int h = sha256(p);  
    if (!h.endsWith("68f728")) {  
        error("Access denied!");  
    }  
    // You are welcome!  
}
```

Path constraint:

$$p \mapsto P$$
$$H \mapsto \text{sha256}(P)$$
$$\text{endsWith}(H) = \mathbf{tt}$$

[SAT SMT Unsolvable [Explosion](#) ChatGPT CSA]

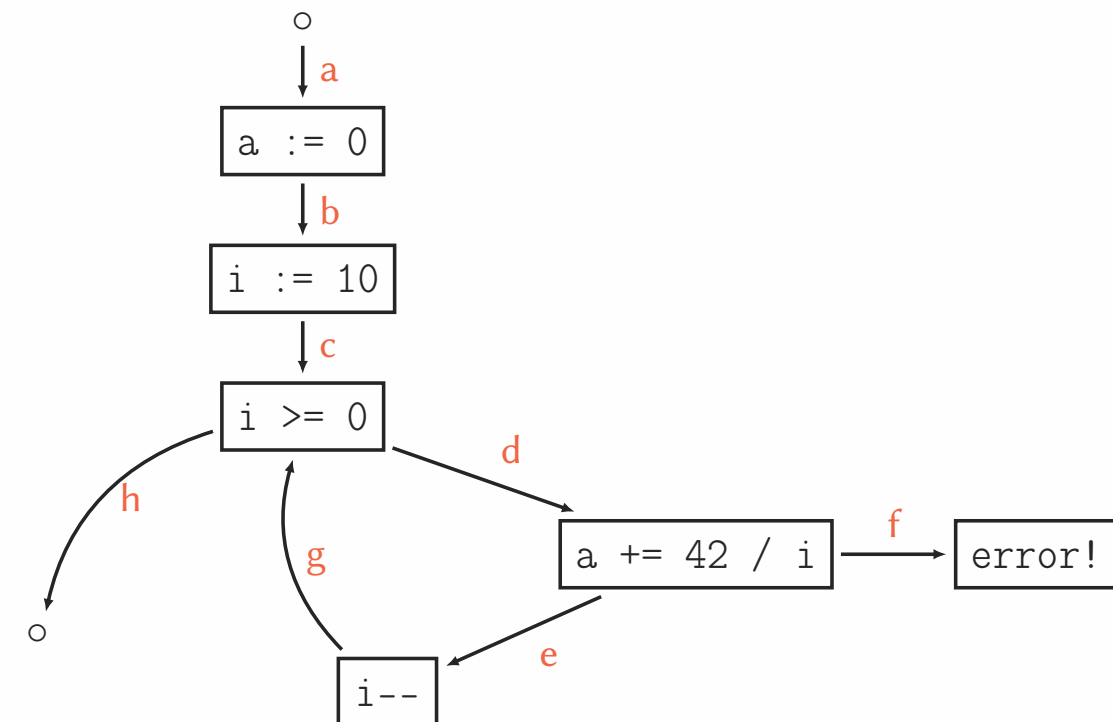
Path Explosion

Path explosion refers to the fact that the number of control-flow paths in a program grows exponentially with an increase in program size and can even be infinite in the case of programs with unbounded loop iterations.

```
int a = 0;  
for (int i = 10; i >= 0; i--) {  
    a += 42 / i;  
}
```

Path:

$(a, b, c, d, e, g, d, e, g, d, e, g, d, \dots, e, g, d, f)$




[SAT SMT Unsolvable [Explosion](#) ChatGPT CSA]


[SAT SMT Unsolvable Explosion [ChatGPT](#) CSA]

Clang Tidy vs. ChatGPT

```
/code/tmp/cpp$ cat a.cpp
int f() {
    int a = 0;
    for (int i = 3; i >= 0; i--) {
        a += 42 / i;
    }
    return a;
}
/code/tmp/cpp$ clang-tidy a.cpp --
1 warning generated.
/code/tmp/cpp/a.cpp:4:13: warning: Division by zero [clang-analyzer-core.DivideZero]
    a += 42 / i;
               ^
/code/tmp/cpp/a.cpp:3:3: note: Loop condition is true.
Entering loop body
    for (int i = 3; i >= 0; i--) {
    ^
/code/tmp/cpp/a.cpp:3:3: note: Loop condition is true.
Entering loop body
/code/tmp/cpp/a.cpp:3:3: note: Loop condition is true.
Entering loop body
/code/tmp/cpp/a.cpp:3:27: note: The value 0 is assigned
to 'i'
    for (int i = 3; i >= 0; i--) {
                          ^
/code/tmp/cpp/a.cpp:3:3: note: Loop condition is true.
Entering loop body
    for (int i = 3; i >= 0; i--) {
    ^
/code/tmp/cpp/a.cpp:4:13: note: Division by zero
    a += 42 / i;
               ^
```

 Do you see any problems with this C++ code:

```
int a = 0;
for (int i = 10; i >= 0; i--) {
    a += 42 / i;
}
```

 There is a problem with this code.

The loop condition is `i >= 0`, which means the loop will run 11 times (i.e., for `i` values of 10, 9, 8, ..., 0`). When `i` reaches 0, the expression 42 / i` will cause a division by zero error, which is undefined behavior in C++.`

To fix this issue, you can change the loop condition to `i > 0` instead of i >= 0`. This will ensure that the loop runs only 10 times, for i` values of 10, 9, 8, ..., 1.`

Here's the corrected code:

css

```
int a = 0;
for (int i = 10; i > 0; i--) {
    a += 42 / i;
}
```

Copy code

This code will execute without any issues and will correctly calculate the sum of the expression `42 / i` for i` values of 10, 9, 8, ..., 1.`

[SAT SMT Unsolvable Explosion ChatGPT [CSA](#)]

Clang Static Analyzer

```
$ cat a.cpp
int main() {
    int a = 0;
    for (int i = 10; i >= 0; i--) {
        a += 42 / i;
    }
    return a;
}
$ clang-tidy a.cpp --
$ clang --analyze -Xclang -analyzer-constraints=z3 \
-Xclang -analyzer-max-loop -Xclang 5 a.cpp
$ clang --analyze -Xclang -analyzer-constraints=z3 \
-Xclang -analyzer-max-loop -Xclang 15 a.cpp
a.cpp:4:13: warning: Division by zero [core.DivideZero]
    a += 42 / i;
           ~~~~
1 warning generated.
```

Chapter #3:

Test Case Generation

[[Input](#) Bitcode TCs Replay]

Symbolic Input

```
#include <climits>
#include "stdlib.h"
int f(int x) {
    int a = x * 2;
    a = a - 4;
    if (a == 0)
        exit(-1);
    return 42 / a;
}
int main(int argc, char** argv) {
    int x = atoi(argv[1]);
    return f(x);
}
```

```
#include <climits>
#include "stdlib.h"
#include "klee/klee.h"
int f(int x) {
    int a = x * 2;
    a = a - 4;
    if (a == 0)
        exit(-1);
    return 42 / a;
}
int main(int argc, char** argv) {
    int x;
    klee_make_symbolic(&x, sizeof(x), "x");
    return f(x);
}
```

[Input [Bitcode](#) TCs Replay]

Compile to LLVM Bitcode

```
$ clang -I /opt/homebrew/Cellar/klee/2.3\_4/include -c -g \  
-emit-llvm -O0 -Xclang -disable-O0-optnone a.cpp
```

```
$ klee a.bc
```

```
KLEE: output directory is "/code/tmp/cpp/klee-out-2"
```

```
KLEE: Using STP solver backend
```

```
KLEE: done: total instructions = 38
```

```
KLEE: done: completed paths = 2
```

```
KLEE: done: partially completed paths = 0
```

```
KLEE: done: generated tests = 2
```

```
$ ls -al klee-out-0/*.ktest
```

```
-rw-r--r--  1 yb  staff   46 Apr  7 17:30 test000001.ktest
```

```
-rw-r--r--  1 yb  staff   46 Apr  7 17:30 test000002.ktest
```

```
$ llvm-bcanalyzer --dump a.bc
```

```
...
```

[Input Bitcode [TCs](#) Replay]

Test Cases

```
#include <climits>
#include "stdlib.h"
#include "klee/klee.h"
int f(int x) {
    int a = x * 2;
    a = a - 4;
    if (a == 0)
        exit(-1);
    return 42 / a;
}
int main(int argc, char** argv) {
    int x;
    klee_make_symbolic(&x, sizeof(x), "x");
    return f(x);
}
```

```
$ ktest-tool klee-last/test000001.ktest
ktest file : 'klee-last/test000001.ktest'
args       : ['a.bc']
num objects: 1
object 0: name: 'x'
object 0: size: 4
object 0: data: b'\x02\x00\x00\x00'
object 0: hex : 0x02000000
object 0: int : 2
object 0: uint: 2
object 0: text: ....
```

```
$ ktest-tool klee-last/test000002.ktest
ktest file : 'klee-last/test000002.ktest'
args       : ['a.bc']
num objects: 1
object 0: name: 'x'
object 0: size: 4
object 0: data: b'\x00\x00\x00\x00'
object 0: hex : 0x00000000
object 0: int : 0
object 0: uint: 0
object 0: text: ....
```

[Input Bitcode TCs [Replay](#)]

Replaying Test Cases

```
$ export LD_LIBRARY_PATH=/opt/homebrew/Cellar/keel/2.3_4/lib:$LD_LIBRARY_PATH
```

```
$ clang -I /opt/homebrew/Cellar/keel/2.3_4/include -L/opt/homebrew/Cellar/keel/2.3_4/lib \
-lkeelRuntest -Xclang -disable-00-optnone a.cpp
```

```
$ KTEST_FILE=keel-last/test000001.ktest ./a.out ; echo $?
255
```

```
$ KTEST_FILE=keel-last/test000002.ktest ./a.out ; echo $?
246
```


Chapter #4:

Concolic Testing

[[Example](#) Steps]

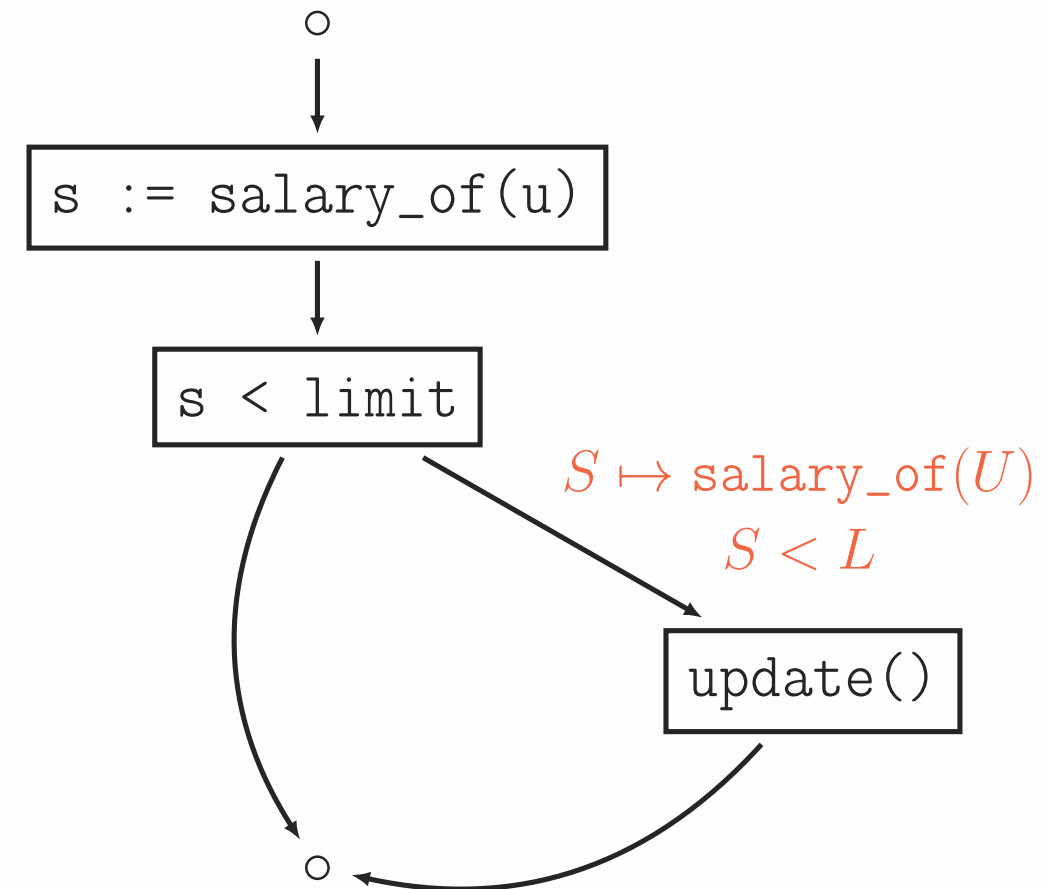
Motivating Example

```
enum user { Viki, Peter, Jeff, Sarah };
```

```
int salary_of(user u) { ... }
```

```
void raise(user u, int limit) {  
    int s = salary_of(u);  
    if (s < limit)  
        update(u, limit);  
}
```

```
// Viki    120  
// Peter   180  
// Jeff    50  
// Sarah   70
```



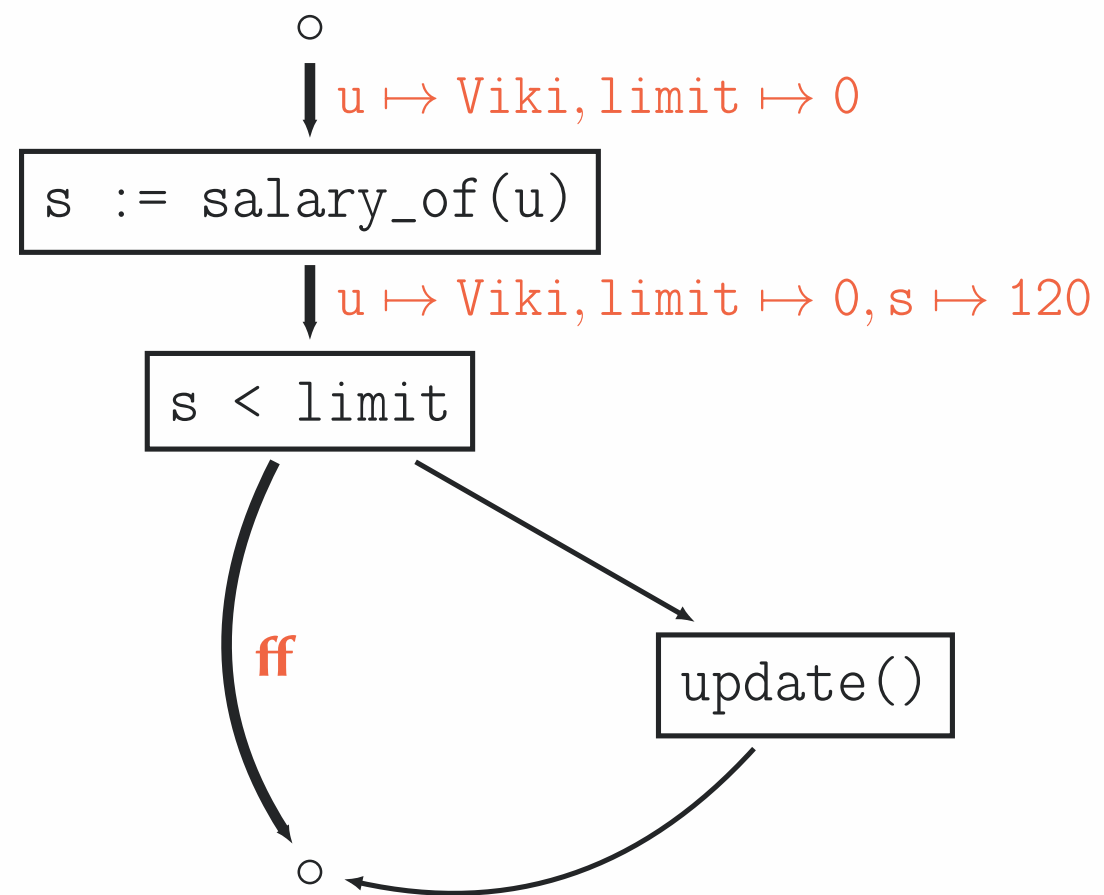
[[Example](#) Steps]

How to find test values of `u` and `limit` for `raise()`? It's impossible :(

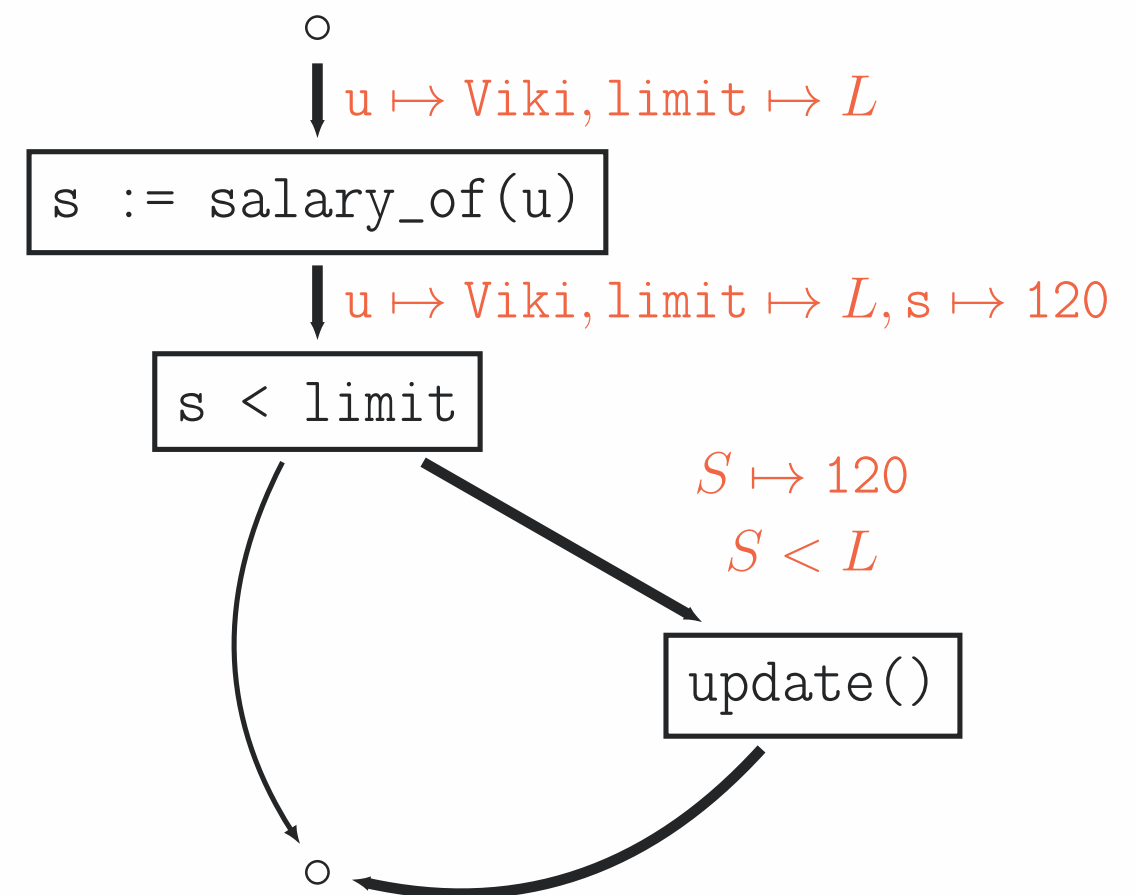
[Example [Steps](#)]

Two Steps

1. Concrete (w/random input):



2. Symbolic (w/neglected condition):



[Example [Steps](#)]

Chapter #5:

Further Reading/Watching

Check this GitHub repo: [ksluckow/awesome-symbolic-execution](https://github.com/ksluckow/awesome-symbolic-execution)

References