

Model Checking

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Lecture #9 out of 10

90 minutes

All videos are in [this YouTube playlist](#).

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Motivating Example

Model-less Checking

Chapter #1:

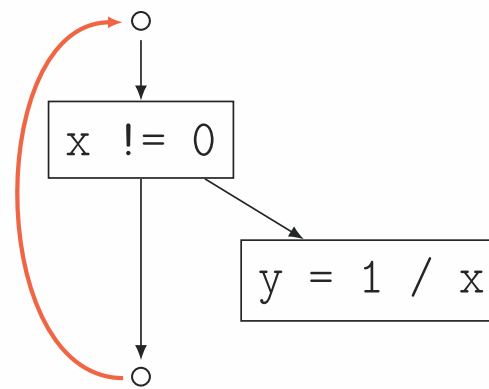
Motivating Example

[[Div by Zero](#) ProMeLa SPIN Monitor Assertion]

Div by Zero

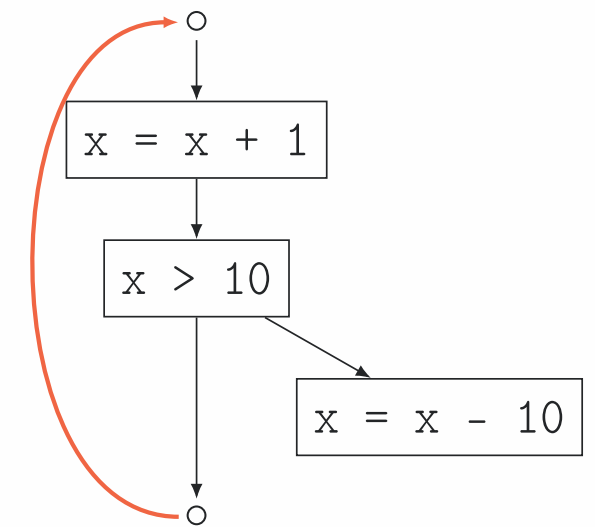
// Process no. 1:

```
extern int x;
extern double y;
int measure() {
    if (x != 0) {
        y = 1.0 / x;
    }
}
```



// Process no. 2:

```
extern int x;
void roll() {
    x += 1;
    if (x > 10) {
        x -= 10;
    }
}
```



Can we detect “division by zero” using symbolic execution? Is “division by zero” the only error here?

ProMeLa (Process Meta Language)

```
extern int x;
extern double y;
int measure() {
    if (x != 0) {
        y = 1.0 / x;
    }
}
void roll() {
    x += 1;
    if (x > 10) {
        x -= 10;
    }
}
```

```
int x; bool dbz;
active proctype measure() {
    do :: true ->
        if
            :: (x != 0) -> dbz = (x == 0)
            :: skip
        fi
    od
}
active proctype roll() {
    do :: true ->
        x = x + 1;
        if
            :: x > 10 -> x = x - 10
            :: skip
        fi
    od
}
```

SPIN (Simple ProMeLa Interpreter)

```
int x; bool dbz;
active proctype measure() {
  do :: true ->
    if
      :: (x != 0) -> dbz = (x == 0)
      :: skip
    fi
  od
}
active proctype roll() {
  do :: true ->
    x = x + 1;
    if
      :: x > 10 -> x = x - 10
      :: skip
    fi;
    printf("x = %d\n", x);
  od
}
```

```
$ spin main.pml | head
      x = 1
      x = 2
      x = 3
      x = 4
      x = 5
      x = 6
      x = 7
      x = 8
      x = 9
      x = 10
$ spin main.pml | tail
...
```

Just checkout [this repo](#) and run make, the spin binary will be compiled.

Monitoring Process

```
int x; bool dbz;
active proctype measure() {
  do :: true ->
    if
      :: (x != 0) -> dbz = (x == 0)
      :: skip
    fi
  od
}
active[2] proctype roll() {
  do :: true ->
    x = x + 1;
    if
      :: x > 10 -> x = x - 10
      :: skip
    fi
  od
}
```

```
active proctype monitor() {
  do :: true ->
    assert(!dbz);
    assert(x >= 0);
  od
}
```

Pay attention to the [2] suffix after the active keyword. It tells SPIN to start two instances of the roll process.

Fail on Assertion

```
int x; bool dbz;
active proctype measure() {
  do :: true ->
    if
      :: (x != 0) -> dbz = (x == 0)
      :: skip
    fi
  od }
active[2] proctype roll() {
  do :: true ->
    x = x + 1;
    if
      :: x > 10 -> x = x - 10
      :: skip
    fi
  od }
active proctype monitor() {
  do :: true -> assert(!dbz); assert(x >= 0); od
}
```

```
$ spin main.pml
spin: main.pml:22, Error: assertion violated
spin: text of failed assertion: assert((x>=0))
#processes: 4
    x = -9
    dbz = 0
584:  proc  3 (monitor:1) main.pml:22 (state 3)
584:  proc  2 (roll:1) main.pml:17 (state 7)
584:  proc  1 (roll:1) main.pml:18 (state 9)
584:  proc  0 (measure:1) main.pml:9 (state 8)
4 processes created
```


Chapter #2:

Model-less Checking

Java Pathfinder

Further Reading/Watching

Introduction lecture by [Joost-Pieter Katoen](#)

[A Primer on Model Checking](#) by Mordechai Ben-Ari