

Program Analysis

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YEGOR BUGAYENKO

Lecture #6 out of 10

90 minutes

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Basics

Quality of Analysis

Abstract Interpretation

Approximation

Chapter #1: Basics

Syntactic & Semantic Property

Semantic property can be completely defined with respect to the set of executions of a program, while a syntactic property can be decided directly based on the program text.

Rice's Theorem

Rice's theorem states that all non-trivial semantic properties of programs are undecidable.

A property is non-trivial if it is neither true for every partial computable function, nor false for every partial computable function.

Halting problem is the problem of determining, from 1) a description of an arbitrary computer program and 2) an input, whether the program will finish running, or continue to run forever. A general algorithm to solve the halting problem for all possible program–input pairs **cannot exist**.

Non-trivial Properties

Examples of a non-trivial properties:

- A program exits
- A program prints “Hello”
- A program dies with “Segmentation Fault”
- A program prints user password to the console

Static Analysis

Consider two C++ programs given to a static analyzer (e.g. Clang Tidy):

```
int f() {  
    int x = 0;  
    return 42 / x;  
}
```

```
int f(int x) {  
    return 42 / x;  
}
```

Expected answers from Clang Tidy:

Yes! :)

No :(

Style Checking

Consider two C++ programs given to a style checker (e.g. cpplint):

```
int f (int x)
{
    return 42 / x;
}
```

```
int f(int x) {
    return 42 / x;
}
```

Expected answers from cpplint:

Extra space before (in
function call ; { should
almost always be at the end
of the previous line

No :(

Dynamic Analysis

Consider this C++ programs given to a dynamic analyzer:

```
int* foo() {  
    return (int*) 42;  
}  
  
int main() {  
    int* p = foo();  
    return *p;  
}
```

```
$ gcc main.cpp  
$ ./a.out  
Segmentation fault: 11
```

Chapter #2:

Quality of Analysis

Sound & Complete



Precision & Recall

Precision is the fraction of relevant instances among the retrieved instances (100% precision means soundness).

Recall is the fraction of relevant instances that were retrieved (100% recall means completeness).

$$\text{Precision} = \frac{TP}{TP + FP} \quad \text{Recall} = \frac{TP}{TP + FN} \quad \text{Accuracy} = \frac{TP + TN}{TP + TN + FP + FN}$$

Experiment

Say, we give a few programs to a static analyzer:

a	b	c	d	e	f	g	h
Yes	Yes	No	Yes	No	Yes	Yes	No

$TP = \underline{\hspace{1cm}}$ $FP = \underline{\hspace{1cm}}$ $TN = \underline{\hspace{1cm}}$ $FN = \underline{\hspace{1cm}}$
 Precision = $\frac{TP}{TP + FP} = \underline{\hspace{2cm}}$ Recall = $\frac{TP}{TP + FN} = \underline{\hspace{2cm}}$
 Accuracy = $\frac{TP + TN}{TP + TN + FP + FN} = \underline{\hspace{2cm}}$

Chapter #3:

Abstract Interpretation

Chapter #4:

Approximation