Program Analysis

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Lecture #6 out of 10
90 minutes

Chapter #1:

Basics

Syntactic & Semantic Properties

Semantic property can be completely defined with respect to the set of executions of a program, while a syntactic property can be decided directly based on the program text.

```
if (x) { printf("大家好"); }
```

Which properties are dynamic?

- A program may print a text to the console
- A program may call `printf()` C library function
- A program prints to the console
- A program consists of one line of code

Rice's Theorem

Rice's theorem states that all non-trivial semantic properties of programs are undecidable.

A property is non-trivial if it is neither true for every partial computable function, nor false for every partial computable function.

Halting problem is the problem of determining, from 1) a description of an arbitrary computer program and 2) an input, whether the program will finish running, or continue to run forever. A general algorithm to solve the halting problem for all possible program–input pairs **cannot exist**.

Non-trivial Properties

Examples of a non-trivial properties:

- A program exits
- A program prints "Hello"
- A program finishes in less than 5 seconds
- A program dies with "Segmentation Fault"
- A program prints user password to the console

Suggest a few properties.

Static Analysis

Consider two C++ programs given to a static analyzer (e.g. Clang Tidy):

```
int f() {
  int x = 0;
  return 42 / x;
}
```

```
int f(int x) {
  return 42 / x;
}
```

Expected answers from Clang Tidy:

Yes! :)

No :(

Style Checking

Consider two C++ programs given to a style checker (e.g. cpplint):

```
int f (int x)
{
  return 42 / x;
}
```

```
int foo(int x) {
  return 42 / x;
}
```

Expected answers from cpplint:

Extra space before (in function call ; { should almost always be at the end of the previous line

No :(

Dynamic Analysis

Consider this C++ programs given to a dynamic analyzer (AddressSanitizer):

```
int foo(int i) {
  int a[5];
  return a[i];
}
int main() {
  return foo(6);
}
```

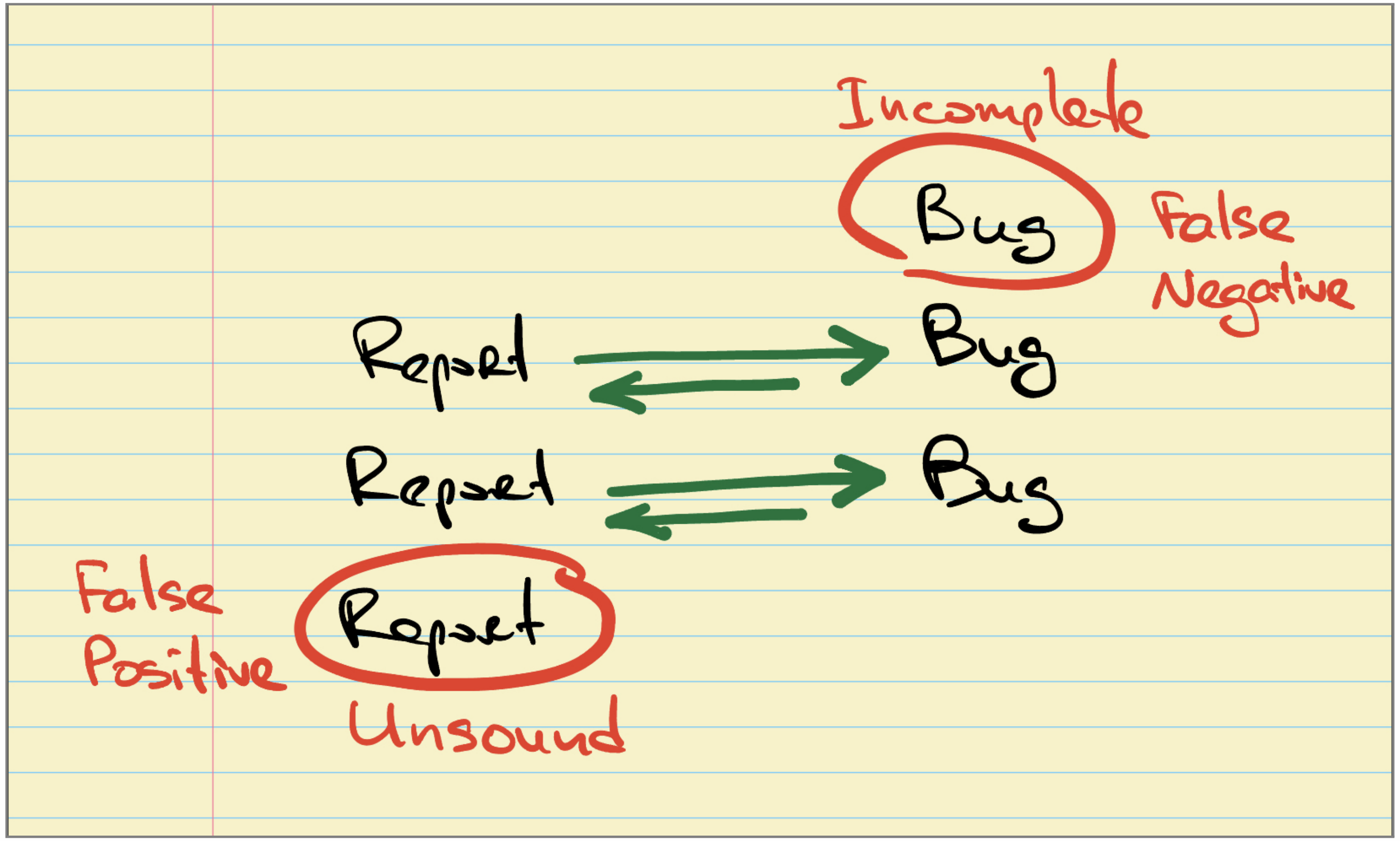
```
$ gcc -fsanitize=address -g a.cpp
$ ./a.out
```

```
=================================================================
==76375==ERROR: AddressSanitizer: stack-buffer-overflow on address 0x00016babf0d8
READ of size 4 at 0x00016babf0d8 thread T0
    #0 0x104343e54 in foo(int) a.cpp:9
    #1 0x104343f38 in main a.cpp:12
    #2 0x1a07c7e4c  (<unknown module>)

Address 0x00016babf0d8 is located in stack of thread T0 at offset 56 in frame
    #0 0x104343cf0 in foo(int) a.cpp:7

  This frame has 1 object(s):
    [32, 52) 'a' (line 8) <== Memory access at offset 56 overflows this variable
```

Chapter #2:

Quality of Analysis

Sound & Complete



Precision & Recall

Precision is the fraction of relevant instances among the retrieved instances (100% precision means soundness).

Recall is the fraction of relevant instances that were retrieved (100% recall means completeness).

$$\text{Precision} = \frac{TP}{TP + FP} \qquad \text{Recall} = \frac{TP}{TP + FN} \qquad \text{Accuracy} = \frac{TP + TN}{TP + TN + FP + FN}$$

Experiment

Say, we give a few programs to a static analyzer:

a	b	c	d	e	f	g	h
Yes	Yes	No	Yes	No	Yes	Yes	No

$$TP = \_\_ \qquad FP = \_\_ \qquad TN = \_\_ \qquad FN = \_\_$$

$$\text{Precision} = \frac{TP}{TP + FP} = \_\_\_\_\_\_\_\_ \qquad \text{Recall} = \frac{TP}{TP + FN} = \_\_\_\_\_\_\_\_$$

$$\text{Accuracy} = \frac{TP + TN}{TP + TN + FP + FN} = \_\_\_\_\_\_\_\_$$

Flip of Terminology

Soundness and Completeness: With Precision by Prof. Bertrand Meyer, in Blog@CACM: "It is very easy to obtain soundness if we forsake completeness: reject every case."

Chapter #3: Lattice

Total Order

Total order is a binary relation $\leq$ (strict total order is $<$).

Lineary ordered set (loset) is a set equipped with a total order.

Which of them are losets:

$\{1, -5, 2, 0, 42\}$

$\{3, 5, -9, 5, 12\}$

$\{3, 5, \text{"Hello"}, 12\}$

$\{x, y, z\}$

$\varnothing$

Partially Ordered Set

Partial order is total order but only between some elements.

Partially ordered set (poset) is a set equipped with a partial order.

Which of them are posets:

$$\{1, \text{"apple"}, 2, -7, \text{"orange"}\}$$

$$\{3, 5, -9, 5, 12\}$$

$$\{3, 5, 42, 12\}$$

$$\{x, y, z\}$$

$$\varnothing$$

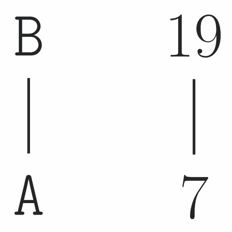
Lattice

Lattice is a poset where each two elements (x, y) have least upper bound $(x \sqcup y)$ and greatest lower bound $(x \sqcap y)$.

$\{42, 2, 13\}$

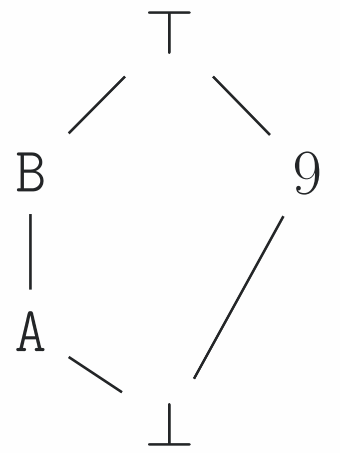
$42 \sqcup 2 = 42$

$\{\mathtt{A}, 7, 19, \mathtt{B}\}$



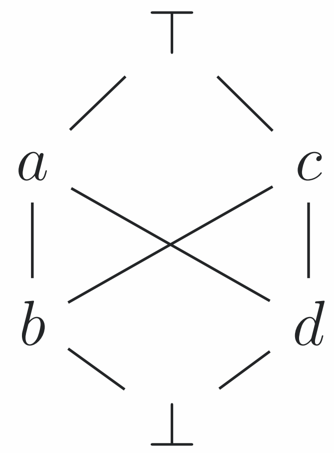
$\mathtt{A} \sqcup 7 = \varnothing$

$\{\mathtt{A}, \top, 9, \mathtt{B}, \bot\}$



$\mathtt{A} \sqcup 9 = \top$ $\mathtt{B} \sqcap 9 = \bot$

$\{\top, \bot, a, b, c, d\}$



$b \sqcup d = \varnothing$ $a \sqcap c = \varnothing$

Chapter #4:

Abstract Interpretation

There is a compromise to be made between the precision of the analysis and its decidability (computability), or tractability (computational cost).

Chapter #5:

Approximation

Further Reading/Watching

Lecture by Patrick Cousot, on YouTube