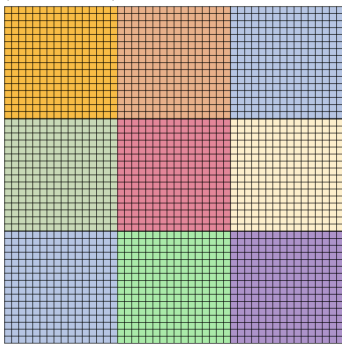


Parallel Programming hw4-1 ----Blocked Folyd-Warshall algo.

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Implementation:

block size: K

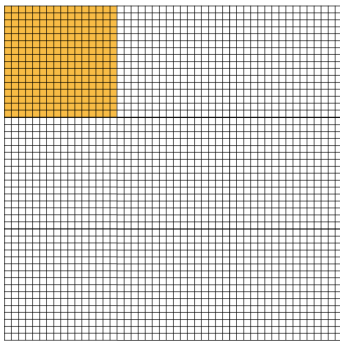


$K:32$

將 distance matrix(data)依照每個 block size ($K*K$) 來分割 data。並依照演算法 3 個 kernel function(phase)用到的 data dependency block 來宣告 block 數以及 share memory 數。

共需執行 $\text{Round} = \text{ceil}(V/k)$ 次， $\text{Round}*\text{Round}$ 個 block

Phase1



phase1 每個 round 只需運算一個 block(黃色)的大小 data ($\text{block}_{(k,k)}$)

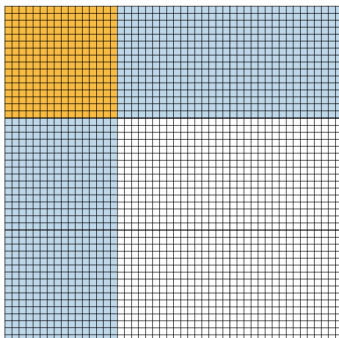
故 phase1 所需 thread 及 block :

$\text{dim3 ThreadPerBlock}(k,k)$

$\text{dim3 BlockPerGrid}(1,1)$

並且每個 block 所使用到的 share memory size = $K*K$

Phase2



Phase2 每個 round 需運算一個 $2*\text{Round}$ 個 block(藍色+黃色)

故 phase2 所需 thread 及 block :

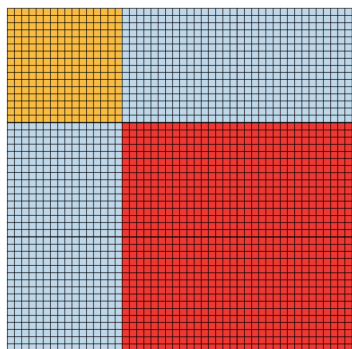
$\text{dim3 ThreadPerBlock}(k,k)$

$\text{dim3 BlockPerGrid}(\text{round},2)$

而每個 block 會需要用到自己及 phase1 的 data

故每個 block 所使用到的 share memory size = $2*K*K$

Phase3



Phase3 每個 round 需運算一個 $(\text{round}-1)*(\text{round}-1)$ 個 block(藍色+紅色)
故 phase3 所需 thread 及 block :

$\text{dim3 ThreadPerBlock}(k,k)$

$\text{dim3 BlockPerGrid}(\text{round},\text{round})$

而每個 block 會需要用到自己及 phase2 的 row 及 col data
故每個 block 所使用到的 share memory size = $3*K*K$

Profiling Results

Occupancy

kernel	Achieved occupancy(avg.)
Phase1	0.496326
Phase2	0.942633
Phase3	0.896393

sm efficiency

kernel	sm efficiency
Phase1	3.66%
Phase2	98.39%
Phase3	99.97%

Share memory load/store throughput

kernel	Load throughput	Store throughput
Phase1	57.631GB/s	19.755GB/s
Phase2	2429.2GB/s	1249.9GB/s
Phase3	2970.5GB/s	92.716GB/s

Global load/store throughput

kernel	Load throughput	Store throughput
Phase1	364.46MB/s	741.18MB/s
Phase2	16.554GB/s	44.787GB/s
Phase3	19.902GB/s	56.468GB/s

Experiment and Analysis

System Spec: hades server

Time distribution:

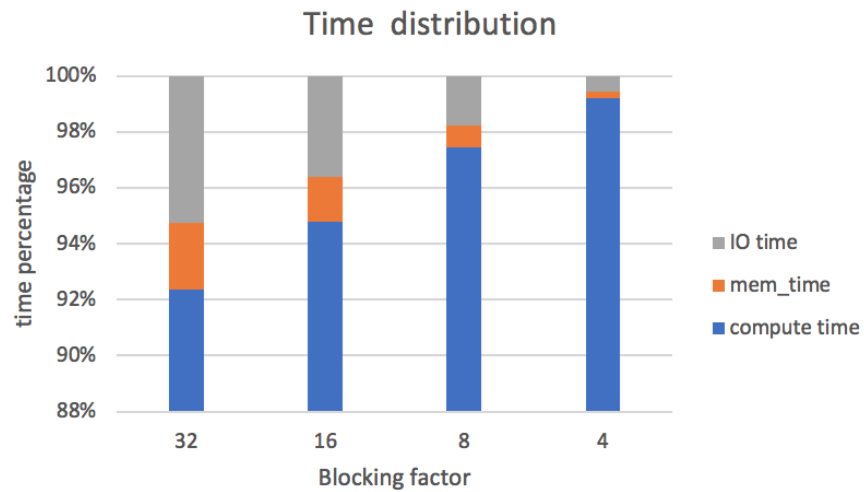
採用 nvprof 來查詢 kernel compute 以及 memry htd,dth time

使用 timespec 來紀錄 I/O time

(a.)

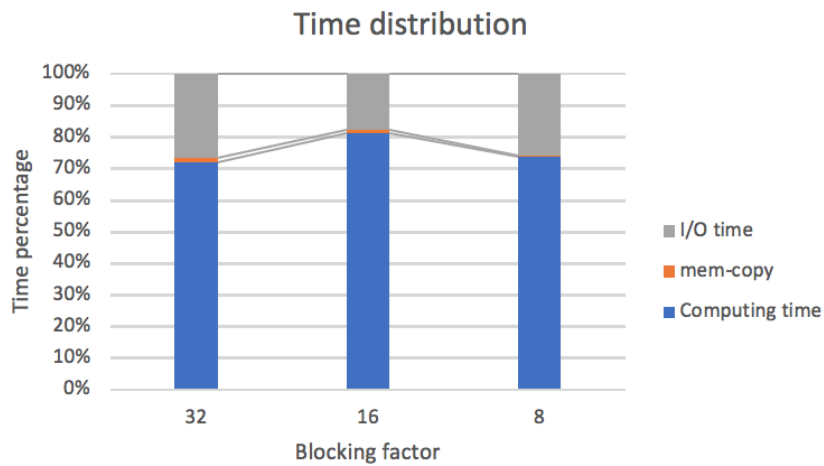
Test case:p21k

Block-factor	Computing time	copy time	I/O time
Block =32	20.813 s	0.53733 s	1.185324 s
Block =16	33.0087 s	0.54889 s	1.199476 s
sBlock =8	68.9008 s	0.55201 s	1.256658 s
Block=4	227.112 s	0.54833 s	1.29825 s



(b.)

Test case:p31k



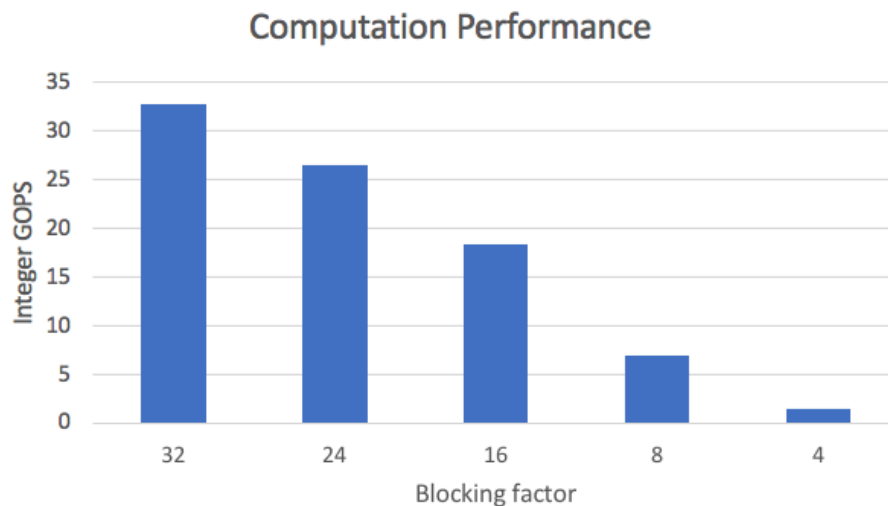
Block-factor	Computing time	copy time	I/O time
Block =32	64.9908	1.21186	23.973353
Block =16	100.281	1.20125	22.019499
Block =8	203.924	1.20078	71.302255

Blocking factor:

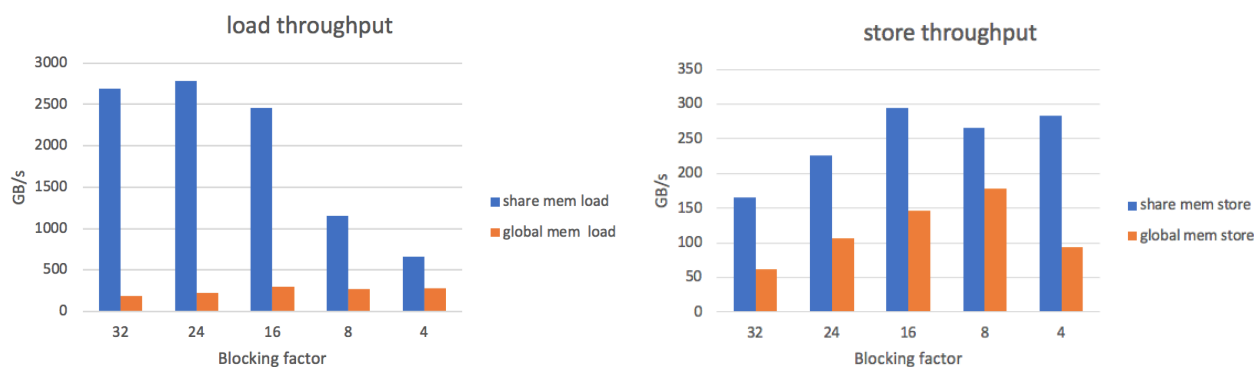
採用 phase3 作為 compute time 標準

Computation Performance:

$$\text{GOPS} = \frac{\#Instruction \times 4 \text{ byte}}{\text{compute time} \times 10^9}$$

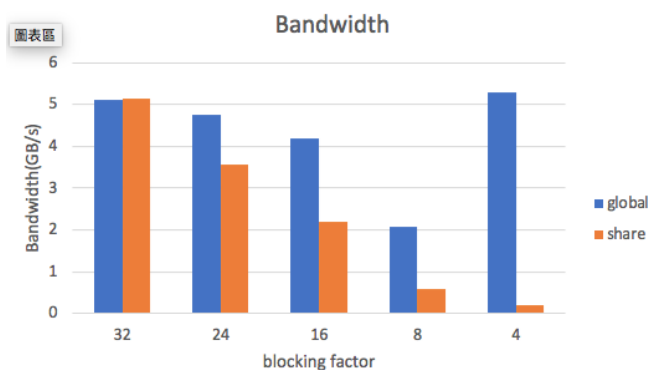


Global memory memory / shared memory load/store throughput :



Global memory memory / shared memory load/store bandwidth :

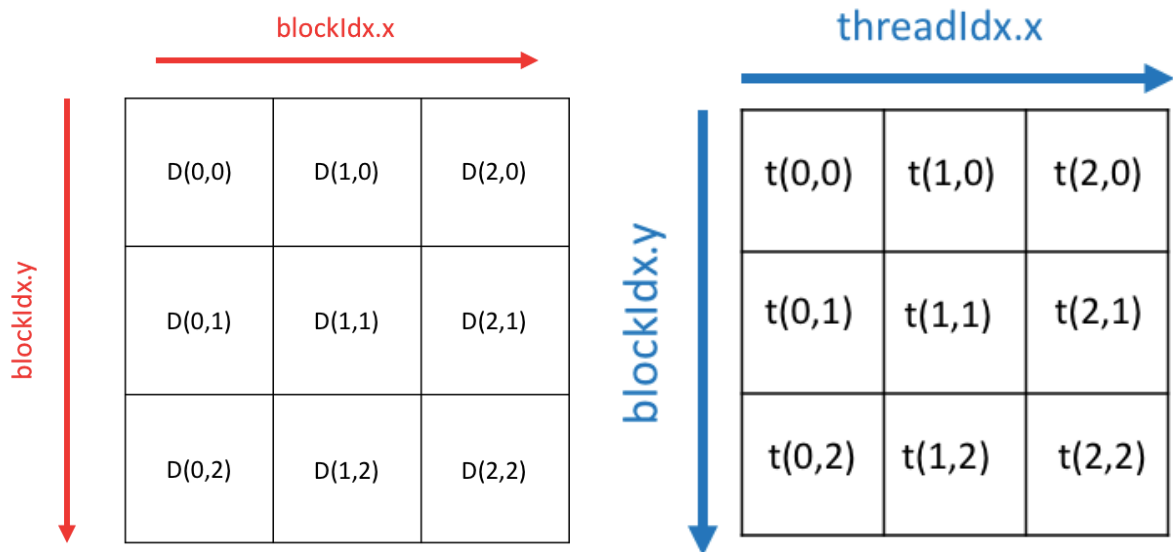
$$\text{Bandwidth} = \frac{(\text{Load transction} + \text{load transction}) \times 32(\text{bytes})}{\text{compute time}}$$



Optimization :

1. using 2 dim block
2. memory coalesced

memory access 為 column major,而 threadIdx 編號是以 x 為 major,若直接以 threadIdx.x,threadIdx.y 轉換為 x,y 座標,則會發生不連續 memory access



3. shared memory

使用 share memory 來提升存取速度

- 4.avoid bank conflict

用 memory padding(share memory:block_size*(block_size+1)) 方式來避免

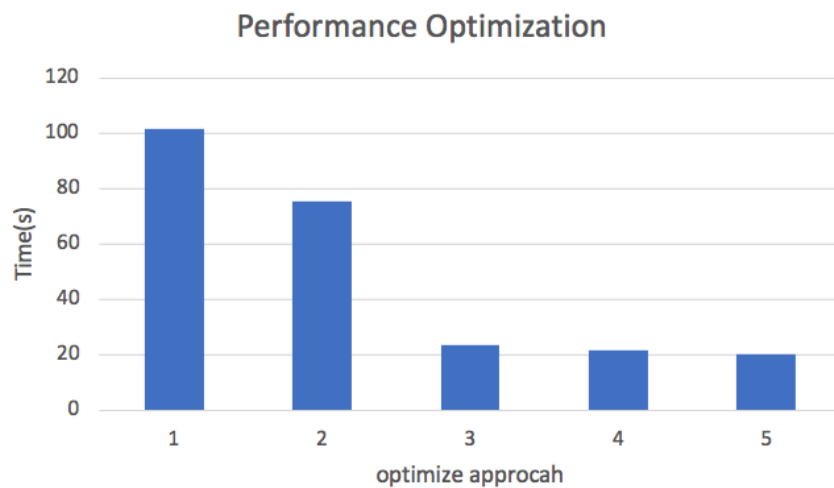
Invocations	Avg	Total	Event Name	Min
Max				
Device "GeForce GTX 1080 (0)"				
Kernel: phase1(int*, int, int)				
655	0	0	shared_ld_bank_conflict	0
0	655	0	shared_st_bank_conflict	0
0	0	0		
Kernel: phase2(int*, int, int)				
655	0	0	shared_ld_bank_conflict	0
0	655	0	shared_st_bank_conflict	0
0	0	0		
Kernel: phase3(int*, int, int)				
655	0	0	shared_ld_bank_conflict	0
0	655	0	shared_st_bank_conflict	0
0	0	0		

- 5.unrolling

使用 pragma unroll 來展開迴圈

1	2D	101.599
2	2D + share_memory	75.4432
3	2D + share_memory +memory coalsed	23.3685

4	2D + share_memory + bankconflict+memory coalsed	21.477
5	2D + share_memory + bankconflict+memory coalsed +unroll	20.3403



Experiment and Analysis

從這次作業我學到如何運用 **cuda** 來做 **gpu** 的運算，來加速程式運算。透過存取 **Share memory** 方式凸顯 **global memory** 存取速度上的差異性，並且可以透過 **coalesced memory** ,**memory padding** 等等各式各樣的優化方式讓效能最佳化。本次作業困難點就是 **block** 與 **block** 彼此之間資料的傳輸，**data** 該如何分割，如何透過 **blockIdx**,以及 **thredIdx mapping** 到 **data distance** 上，並存入 **shared memory**。另外發現若使用過多的 **share memory** 效能不一定會越好，過多的 **thread** 在 **syncthreads** 會嚴重影響效能。因此我覺得還可以改進的部分就是減少 **syncthreads** 時間。