

# **Set Theory I**

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# 1 Permutations

Permutation  $\sigma$  is a bijection from a set  $S$  onto itself.

Every permutation can be decomposed into one or more disjoint cycles (or orbits), thus, they can also be defined by them like this:

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 5 & 4 & 3 & 1 \end{pmatrix} = (3\ 4)(1\ 2\ 5)$$

1. A cycle of one element is called a **fixed point**
2. A permutation without fixed points is called a **derangement**
3. A permutation that's an orbit of 2 elements is called a **transposition**

## 1.1 The symmetric group

A symmetric group defined over a set is the group whose elements are all the permutations over the set, and whose group operation is the composition of functions.

Reminder: A group is an algebraic structure with the following characteristics:

- Closure
- Associativity
- An identity permutation exists
- $\forall \sigma \in S_n (\exists \pi : \sigma \circ \pi = Id)$

## 2 Hall's theorem

**Hall's theorem -** *In a finite bipartite graph  $G(X, Y, E)$*   
 $\forall W \subseteq X (|W| \leq |N_g(W)|) \iff \text{An } X\text{-perfect matching exists.}$

( $\Leftarrow$ ) Suppose we have an  $X$  perfect matching  $M$ , since for any given  $W$  all vertices in  $W$  have a distinct matching vertex in  $Y$  by  $M$  we get

$$\forall W \subseteq X (|W| \leq |N_g(W)|)$$

( $\Rightarrow$ ) We'll prove by contradiction. Assume an  $X$ -perfect matching doesn't exist, we'll denote the maximal matching  $M$ , and the sets of vertices in  $X, Y$  that appear in  $M$  as  $S, T$ . An  $X$ -perfect matching doesn't exist  $\Rightarrow X \setminus S \neq \emptyset$ , so we can choose a vertex  $u_0 \in X \setminus S$  and consider all alternating paths of the form  $P = (u_0, v_1, v_2, \dots)$  such that odd edges are not in  $M$  and even edges are in  $M$ . Denote:

$$A = \{u : u \in P \wedge u \in X\}, \quad B = \{v : v \in P \wedge v \in Y\}$$

We know every vertex in  $B$  is matched by  $M$  to a vertex in  $A$  because otherwise we could create a bigger matching by toggling whether each of the edges belong to  $M$  or not.

$$\Rightarrow |B| \leq |A \setminus \{u_0\}| \Rightarrow |B| < |A|$$

but also for any vertex in  $a \in A$  any of its neighbors  $b$  are in  $B$ . We can show that an alternating path to  $b$  exists either by removing the matched edge  $ab$  from the alternating path to  $a$ , or by adding the unmatched edge  $ab$  to the alternating path to  $a$ .

$$\Rightarrow B = N_g(A)$$

$$\Rightarrow |N_g(A)| < |A|$$

That's a contradiction so an  $X$ -perfect matching must exist.

### 3 Cantor's theorem

$$|A| < |P(A)|$$

We can define  $f : A \rightarrow P(A)$  as such

$$f(a) = a$$

$$\Rightarrow |A| \leq |P(A)|$$

Assume  $|A| = |P(A)|$ . That means there's a bijection  $g : A \rightarrow P(A)$ .  
consider the following set:

$$D = \{a : a \notin g(a)\}$$

Since  $g$  is a bijection  $\exists b \in A : f(b) = D$ . Now look at the different cases:

$$\begin{cases} b \in D, & b \notin g(b) = D \Rightarrow \text{contradiction} \\ b \notin D = g(b), & b \in D \Rightarrow \text{contradiction} \end{cases}$$

Therefore  $|A| \neq |P(A)| \Rightarrow |A| < |P(A)|$

## 4 Equivalence Relations

An equivalence relation is a binary relation (a set of ordered pairs) that is

- Reflexive
- Symmetric
- Transitive

### 4.1 Some Terminology

Suppose we have an equivalence relation  $R$  on a set  $X$ :

**Equivalence Class:**  $[a]_R = \{b \in X : bRa = 1\}$

**Quotient Set:**  $X/R = \{[a]_R : a \in X\}$

**Projection:** The projection of  $R$  is  $\pi : X \rightarrow X/R$  such that  $\pi(x) = [x]_R$

**A Cut:** A cut of  $X$  is a set with only one element of each Equivalence class.

Equivalence relations can be defined by their Quotient set. Thus they can also be defined by a function or a partition. The numbers of partitions of a set  $|X| = n$  are known as Bell's numbers and can be calculated recursively as such:

$$B_{n+1} = \sum_{k=0}^n \binom{n}{k} B_k$$

Think why.

## 5 König's Theorem

**König's Theorem** - For an index set  $I$ , suppose  $\forall i \in I$  and  $\kappa_i, \lambda_i$  we know  $\kappa_i < \lambda_i$  then  $\sum_{i \in I} \kappa_i < \prod_{i \in I} \lambda_i$   
We'll show this by proving for any

$$f : \sum_{i \in I} B_i \rightarrow \prod_{i \in I} C_i \quad |B_i| = \kappa_i, |C_i| = \lambda_i$$

That  $f$  is not surjective. Let's define the function  $f_i$  as such:

$$f_i : B_i \rightarrow C_i$$

$$f_i(x) = f(x)_i$$

$$\forall i \in I (|B_i| < |C_i|) \Rightarrow \forall i \in I (f_i \text{ is not surjective}) \Rightarrow \exists c_i \in C_i \setminus \text{Im} f_i$$

Consider the vector:

$$\hat{c} = \langle c_i : i \in I \rangle$$

$$\text{If } \hat{c} \in \text{Im} f \Rightarrow \exists i \in I, b \in B_i : f(b) = \hat{c}$$

$$\Rightarrow f(b)_i = c_i \Rightarrow f_i(b) = c_i \text{ but } c_i \in C_i \setminus \text{Im} f_i$$

That's a contradiction so we got

$$|\sum_{i \in I} B_i| < |\prod_{i \in I} C_i| \Rightarrow \sum_{i \in I} \kappa_i < \prod_{i \in I} \lambda_i$$



## 6 Partial Orders

A **Weak/Non-Strict** Partial Order is a homogeneous relation  $\leq$  on a set  $P$  that is:

- Reflexive
- Antisymmetric<sup>1</sup>
- Transitive

A **Strong/Strict Partial** Order is a homogeneous relation  $<$  on a set  $P$  that is:

- Irreflexive
- Asymmetric<sup>2</sup>
- Transitive

note:  $< \cup \leq_{Id} = \leq$

## 7 Partially Ordered Sets

A Partially Ordered Set (aka a poset) is a set on which a partial order is defined  $(A, \leq)$ . We say two elements  $a, b \in A$  are comparable  $\iff a \leq b \vee b \leq a$ . If two elements are incomparable they're linearly independent. A **linear/total order** is a partial order under which every pair of elements is comparable. All ordered subsets(chains) are linearly independent of each other.

### 7.1 Extrema

**Greatest Element** - *an element that's comparable and greater than all other elements*

**Maximal Element** - *an element that doesn't have a greater element than him*

**Upper/Lower Bounds in sets** - *a is a bound in A of  $B \subseteq A$  if  $a \in A \wedge \forall b \in B(b \leq a)$*

#### 7.1.1 About Lattices

Let A be a partially ordered set:

A is a lattice  $\iff \forall S \subseteq A (|S| = 2 \Rightarrow SupS, InfS \text{ exist})$

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<sup>1</sup> $a \leq b \wedge b \leq a \Rightarrow a = b$   
<sup>2</sup> $a < b \Rightarrow \neg b < a$

## 8 Cardinals

Cardinal numbers are the "numbers" we use to represent the cardinality of sets. Their "size". All cardinal numbers are based on the size of  $\mathbb{N}$  that is  $\aleph_0$ . This subject is rather simple, yet hard to start from scratch. Thus I encourage you to try to prove the following:

1.  $|\mathbb{N}| < |\mathbb{R}|$
2.  $\aleph_0 = \aleph_0 + n = \aleph_0 \times n = \aleph_0 \times \aleph_0 = |\mathbb{Z}| = |\mathbb{Q}|$
3.  $\aleph = 2^{\aleph_0} = |(0, 1)^{\aleph_0}| = \aleph \times \aleph_0 = \aleph \times \aleph = |(0, 1)| = |[0, 1]|$
4. A plane can't be covered by  $\aleph_0$  lines.
5. Let  $A$  be an infinite set,  $\exists S \subseteq A : |S| = \aleph_0$
6.  $\aleph = |P(\mathbb{N})| = |P(\mathbb{Q})|$
7. let  $A = \{\text{The set of all finite subsets of } \mathbb{N}\}$  prove  $|A| = \aleph_0$
8.  $\aleph_0^{\aleph} = \aleph$
9.  $|\mathbb{R}^{\mathbb{R}}| = |P(\mathbb{R})| = 2^{\aleph}$
10. |The disjoint union of  $\mathbb{N}$  sets of size  $\mathbb{N}$ | =  $\aleph_0$
11.  $\aleph_0^{\aleph_0} = \aleph$
12. |The set of all invertible functions  $\mathbb{R} \rightarrow \mathbb{R}$ | =  $2^{\aleph}$
13.  $|A| = |\text{The set of all algebraic numbers}| = \aleph_0$
14.  $|B| = |\mathbb{R} \setminus A| = |\text{The set of all transcendental number}| = \aleph$
15. |All subsets of  $\mathbb{R}$  with cardinality  $\aleph/\aleph_0$ |
16. Let  $\aleph_0$  people with a natural number of hats on their head guess how many hats they have. How many options are there, given only a finite number of people guessed right/wrong?

## 9 Schröder–Bernstein Theorem

**Schröder–Bernstein Theorem** -  $|A| \leq |B| \wedge |B| \leq |A| \iff |A| = |B|$  There are more proofs that rely on similar ideas. Here's one:  
We're given two injective functions

$$f : A \rightarrow B$$

$$g : B \rightarrow A$$

Without loss of generality assume  $A, B$  are disjoint (Why can we do this?)  
Considering the partially defined functions  $f^{-1}, g^{-1}$  we can create a sequence for every element of  $A \cup B$  in the following way:

$$\dots f^{-1}g^{-1}(a) \rightarrow g^{-1}(a) \rightarrow a \rightarrow f(a) \rightarrow g(f(a)) \dots$$

The sequence can keep going forever to the right, but to the left it may stop eventually since the inverse functions are partial<sup>1</sup>. We can see that every element in  $A \cup B$  has a sequence and that if an element appears in two sequences they'll be identical since they're injective and by our construction. Thus those sequences form a partition of  $A \cup B$  so it's sufficient to create bijections for all partitions, and we're finished. Our bijection will be:

$$h(x) = \begin{cases} f(x), & \text{for } x \in A \text{ in an A-stop} \\ g^{-1}(x), & \text{for } x \in A \text{ in a B-stop} \\ f^{-1}(x), & \text{for } x \in B \text{ in an A-stop} \\ g(x), & \text{for } x \in B \text{ in a B-stop} \end{cases}$$

And of course any element of an A-stop can go one step right with  $f$  and we can get any element's left neighbor by applying  $f^{-1}$ , and so we get a bijection for any A-stop sequence. The proof for B-stops is similar. Finally, we get that  $h$  is the bijection we looked for and that  $|A| = |B|$ .

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<sup>1</sup>We'll call those who stop from the left on an element of  $A$  A-stops and the rest B-stops - even though they may not always stop!

## 10 Homomorphism and Isomorphism of Ordered Sets

### 10.1 Homomorphisms

Let  $(X, \leq_1), (Y, \leq_2)$  be partially ordered sets.

$$F \text{ is a Homomorphism} \iff \forall x_1, x_2 \in X (x_1 \leq_1 x_2 \Rightarrow F(x_1) \leq_2 F(x_2))$$

### 10.2 Isomorphisms

Let  $(X, \leq_1), (Y, \leq_2)$  be partially ordered sets.

$$F \text{ is an Isomorphism} \iff F : X \rightarrow Y \text{ is a bijection} \wedge \forall x_1, x_2 \in X (x_1 \leq_1 x_2 \Rightarrow F(x_1) \leq_2 F(x_2))$$

An isomorphism is reflexive, symmetric and transitive so it's an equivalence relation.

If  $F$  is an Isomorphism and the orders are total orders,  $F^{-1}$  is also an isomorphism.

### 10.3 Lexicographic Order

Also known as a Dictionary Order, is an order That's defined as such:

$$(x_1, y_1) \leq (x_2, y_2) \iff x_1 <_x x_2 \vee (x_1 = x_2 \wedge y_1 \leq_y y_2)$$

This is a partial order on  $X \times Y$

## 11 Zorn's Lemma

**Zorn's lemma** - Let  $F$  be a non-empty poset. If for every chain<sup>1</sup> in  $F$  exists an upper bound in  $F$ , then  $F$  has at least one maximal element.

### 11.1 Proof All Vector Spaces Have a Base

Let  $V$  be a vector space: If  $V = \{0\}$  then  $\emptyset$  is its basis. If  $V$  is finitely generated then we can add vectors from  $V$  to  $\emptyset$  until it's spanning  $V$ . Suppose  $V$  is not finitely generated, let's Define  $F$  as the set of all linearly independent sets of vectors.  $F$  is partially ordered by the order of inclusion of sets. Let  $C = (A_i)_{i \in I}$  be a chain in  $F$ ,  $A = \bigcup_{i \in I} A_i$ .  $A$  is clearly a maximal element of the chain. Let's prove it's in  $F$ . Assume  $A$  isn't in  $F \Rightarrow$  there exists a finite series of linearly dependent vectors, each is an element of a finite series of elements of  $C$ . Since that series is finite, and linearly ordered as a subset of  $C$ , There exists a maximal element that must contain all the vectors in the linearly independent vector series, but that element is in  $F$  so it's both linearly dependent and independent at the same time! contradiction! We get that  $A \in F$  so by Zorn's lemma  $F$  has a maximal element  $T$ . That  $T$  is our basis.

### 11.2 Comparing Cardinals

We'll show that for every two cardinals  $\alpha, \beta$  other than 0 we get  $\alpha \leq \beta \vee \beta \leq \alpha$ . Let  $A, B$  be two sets of cardinality  $\alpha, \beta$ . Define  $F$  to be the set of all ordered pairs  $(X, f)$  such that  $f : X \rightarrow B$  is an injective function ( $X \subseteq A$ ). Now we'll define an order in the following way:

$$(X_1, f_1) \leq (X_2, f_2) \iff X_1 \subseteq X_2 \wedge f_2|_{X_1} = f_1$$

Let  $C = ((X_1, f_1), (X_2, f_2), \dots)$  be a chain in  $F$ ,  $(X, g) = (\bigcup A_i, \bigcup f_i)$

$$\Rightarrow \forall i ((X_i, f_i) \leq (X, g)).$$

Assume  $g$  isn't a function, we get  $(x, y), (x, z) \in G$

$$\Rightarrow \exists i, j \text{ such that: } f_i(x) = y, f_j(x) = z$$

$C$  is a chain so we without lose of generailty we get:

$$f_i \leq f_j$$

$$\Rightarrow f_j|_{X_i} = f_i$$

$$\Rightarrow f_i(x) = f_j(x)$$

$$\Rightarrow y = z$$

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<sup>1</sup>a totally ordered subset

That means  $g$  is a function, and since it's a union of injective functions, it must also be injective. That means it's in  $F$  and using Zorn's lemma we get a maximal element in  $F$ , which we'll denote  $(D, h)$ . If  $D = X$  then  $h$  is injective and we get  $A \leq B$ . If it's not, it must be surjective or we get a contradiction to  $(D, h)$ 's maximality and thus  $B \leq A$

We can also prove  $\alpha + \alpha = \alpha$ . We know that  $\alpha + \alpha = 2\alpha$  so we'll just prove  $\alpha = 2\alpha$ . We'll build  $F$  using bijections this time. Denote the maximal element  $M = (X, g)$ . If  $|X| = 2\alpha$  We finished, else we get that there's a set of size  $\aleph_0$  that can be mapped "bijectively" to the set of  $2\alpha$  contradicting  $M$ 's maximality.

### 11.3 Corollaries

- $\alpha + \beta = \max\{\alpha, \beta\}$
- $|A \setminus B| = |A| \iff |B| \leq |A|$
- $\alpha * \alpha = \alpha$  (not a direct corollary)
- $\alpha^\alpha = 2^\alpha$

## 12 Axiom of Choice

First let's define what is a choice function. A Choice Function - is a function from an indexed family of sets  $(S_i)_{i \in I}$  to their union such that  $\forall i \in I (f(S_i) \in S_i)$ . Now for the axiom itself, The Axiom of Choice:

$$\forall X [\emptyset \notin X \Rightarrow \exists f : X \rightarrow \bigcup X \forall A \in X (f(A) \in A)]$$

### 12.1 Nomenclature

**AC** - *Axiom of Choice*

**ZF** - *Zermelo-Fraenkel set theory omitting AC*

**ZFC** - *ZF extended to include AC*

## 13 More Axioms In ZF

Axiom of extensionality

$$\forall A \forall B (\forall x (x \in A \iff x \in B) \Rightarrow A = B)$$

Axiom of union

$$\forall A \exists B \forall c (c \in B \iff \exists D (c \in D \wedge D \in A))$$

Axiom of infinity

$$\exists I (\emptyset \in I \wedge \forall x \in I (x \cup \{x\} \in I))$$

Axiom of power set

$$\forall x \exists y \forall z [z \in y \iff \forall w (w \in z \Rightarrow w \in x)]$$

Axiom of regularity

$$\forall x (x \neq \emptyset \rightarrow \exists y (y \in x \wedge y \cap x = \emptyset))$$

Axiom of pairing

$$\forall A \forall B \exists C \forall D [D \in C \iff (D = A \wedge D = B)]$$

**Axiom schema of specification** - any definable subclass of a set is a set.

**Axiom schema of replacement** - the image of any set under any definable mapping is also a set

## 14 Measure

Measure theory is complex and goes well beyond what I can show in this section but let's talk about it anyway. A measure is a way to generalize the length, volume, and such for sets. Let  $X$  be a set and  $\Sigma$  a  $\sigma$ -algebra over  $X$ . A set function  $\mu$  from  $\Sigma$  to the extended real number line is called a measure if

- $\forall E \in \Sigma : \mu(E) \geq 0$
- $\mu(\emptyset) = 0$
- $\sigma$ -additivity: For all countable collections  $\{E_k\}_{k=1}^{\infty}$  of pairwise disjoint sets in  $\Sigma$ ,

$$\mu\left(\bigcup_{k=1}^{\infty} E_k\right) = \sum_{k=1}^{\infty} \mu(E_k)$$

If the condition of non-negativity is dropped then  $\mu$  is called a signed measure. The pair  $(X, \Sigma)$  is called a measurable space, and the members of  $\Sigma$  are called measurable sets. A triple  $(X, \Sigma, \mu)$  is called a measure space. A probability measure is a measure with total measure  $\mu(X) = 1$ . A probability space is a measure space with a probability measure.

### 14.1 Lebesgue Measure

Now we have to simplify things so we'll consider only Lebesgue measure of bounded sets on the real number line. First, if the set is of the form  $X = (a, b)$  or  $X = [a, b]$  and such the measure must satisfy  $\mu(X) = b - a$ , if  $Y = f(X)$  and  $f$  is an isometric function then  $\mu(X) = \mu(Y)$ . Denote  $\mathcal{Y} = (Y_i)_{i \in I}$  such that  $|I| < \aleph_0$  and  $X \subseteq \bigcup_{i \in I} Y_i$  and each  $Y_i$  is an interval on  $\mathbb{R}$ . Denote  $s(\mathcal{Y})$  The sum of lengths of intervals in  $\mathcal{Y}$ . The outer measure of  $X$  is

$$\mu^*(X) = \inf_{\mathcal{Y}} s(\mathcal{Y})$$

And the inner measure is defined with an interval  $X \subseteq [a, b]$  and  $X' = [a, b] \setminus X$

$$\mu_*(X) = (b - a) - \mu^*(X')$$

It's easy to show that the outer measure is always greater than the inner measure. The Lebesgue measure is defined if they are equal and is equal to

$$\mu_*(X) = \mu^*(X) = \mu(X)$$



## 15 Well Order

A partially ordered set  $(X, \leq)$  is well ordered

$$\iff$$

$$\forall S \subseteq X (S \neq \emptyset \rightarrow \exists b \in S (b \text{ is a minimal element of } S))$$

Think about the following theorems:

1. Every finite totally ordered set is well ordered.
2. If  $\leq$  is a well order then it's a linear order as well.
3. Let  $(X, \leq)$  be a linearly ordered set. It's well ordered  $\iff$  it doesn't include an infinite decreasing series.

We'll proceed to define two very similar terms.

**Risha** - If  $X$  is well ordered  $A \subseteq X$  (usually we mean  $A \subset X$ ) is a Risha if  $x \in A \wedge y < x \rightarrow y \in A$

**Initial segment** -  $I_x(a) = \{x \in X : x < a\}$  aka initial segment of  $a$  in  $X$

note:  $[0, 0.5]$  in  $[0, 1] \in \mathbb{R}$  is a Risha but not an initial segment. Prove a Risha and an initial segment are the same in wosets.

### 15.1 Some Lemmas

1. let  $X$  be a woset,  $f : X \rightarrow X$  a one-to-one homomorphism  $\rightarrow \forall x \in X (x \leq f(x))$
2. let  $(X, \leq_x) \cong (Y, \leq_y)$  be isomorphic wosets, there's only one unique isomorphism between them (proof using previous theorem)
3. in a woset  $X$  a risha can't be have an isomorphism with  $X$
4. in wosets  $I_x(a) \cong I_x(b) \Rightarrow a = b$
5. let  $f : X \rightarrow Y$  be an isomorphism between wosets s.t.  $y_0 = f(x_0) \Rightarrow I_x(x_0) = I_y(y_0)$

### 15.2 A lemma about partial orders

If  $(X, \leq_x), (Y, \leq_y)$  are partial orders, and  $\leq_x$  is a total order, then if  $f$  is an invertible homomorphism then it's an isomorphism, and  $\leq_y$  is a total order.

## 16 Comparison of Well Ordered Sets

If  $X, Y$  are wosets then exactly one of the following is true

1.  $(X, \leq_x) \cong (Y, \leq_y)$
2.  $\exists y_0 \in Y : (X, \leq_x) \cong (I_y(y_0), \leq_y)$
3.  $\exists x_0 \in X : (Y, \leq_y) \cong (I_x(x_0), \leq_x)$

If  $X = \emptyset \vee Y = \emptyset$  the proof is trivial. Assuming they're not empty we'll define:

$$A = \{x \in X : \exists y \in Y (I_X(x) \cong I_Y(y))\}$$

$$B = \{y \in Y : \exists x \in X (I_X(x) \cong I_Y(y))\}$$

$$\phi : A \rightarrow B$$

$$\phi(x) = y : I_X(x) \cong I_Y(y)$$

It's clear why  $\phi$  is a bijection, we will show it's an isomorphism. Consider  $a_1 < a_2 \in A$  and  $\phi(a_1) = b_1, \phi(a_2) = b_2$ . Since  $I_X(a_2) \cong I_Y(b_2)$  we'll denote their isomorphism  $\alpha$ .  $a_1 < a_2 \Rightarrow a_1 \in \text{Dom} \alpha \Rightarrow \alpha(a_1) \in \text{Im} \alpha = I_Y(b_2) \Rightarrow \alpha(a_1) < b_2$ . By one of our previous lemmas<sup>0</sup>  $I_X(a_1) \cong I_Y(\alpha(a_1))$  and we know  $I_X(a_1) \cong I_Y(b_1) \Rightarrow b_1 = \alpha(a_1)$ . Recall that  $\alpha(a_1) < b_2 \Rightarrow b_1 < b_2$ . Since  $\phi$  is a bijection and a homomorphism it's an isomorphism  $\Rightarrow A \cong B$ .

By cases we'll get:

1. If  $A = X, B = Y \Rightarrow (1)$ .
2. If  $B = Y \wedge A \subset X \neq \emptyset$  denote  $A \setminus X$ 's minimal element  $c$  and then  $\Rightarrow I_X(c) = A^1 \Rightarrow (3)^2$
3. If  $A = X \wedge B \subset Y \neq \emptyset$  denote  $Y \setminus B$  minimal element  $d$  and then  $\Rightarrow I_X(d) = B \Rightarrow (2)$
4. If  $A \subset X \wedge B \subset Y \Rightarrow I_X(c) \cong A \wedge I_Y(d) \cong B. A \cong B \Rightarrow I_X(c) \cong I_Y(d) \Rightarrow c \in A$  but  $c \notin A$  by our construction  $\Rightarrow$  contradiction.

Now we'll show only one of (1), (2), (3) can be true for any  $X, Y$ :

(2) + (3)  $\Rightarrow \exists \delta : X \rightarrow I_Y(d)$  isomorphism  $\Rightarrow^0 I_X(c) \cong I_Y(\delta(c))$  and since we know  $Y \cong I_X(c)$  we get that  $Y \cong I_Y(\delta(c))$  which we know can't be.

(1) + (3)/(1) + (2)  $\Rightarrow$  an initial segment of  $X/Y$  is isomorphic to  $X/Y$  and that can't be!

---

<sup>0</sup>Refer to 15.1.5

<sup>1</sup>Think why(two sided inclusion).

<sup>2</sup>since  $A \cong B$

## 17 Ordinals

Ordinals are the generalization of ordinal numerals aimed to extend enumeration to infinite sets. The finite ordinals will be defined as such:

$$k = \text{ord}(\{0, 1, \dots, k-1\}) = \text{ord}(I_{\mathbb{N}}(k))$$

$$\text{ord}(\emptyset) = 0$$

$$\text{ord}(\mathbb{N}) = \omega$$

By the comparability of wosets we can define an order on the ordinals as such:

$$A \cong B \iff \text{ord}(A) = \text{ord}(B)$$

$$A < B \iff \text{ord}(A) < \text{ord}(B)$$

$$A > B \iff \text{ord}(A) > \text{ord}(B)$$

Now we'll define a new set function on ordinals  $W(\alpha)$

$$W(\alpha) = \{\beta : \beta < \alpha\}$$

$W(\alpha)$  is a woset and  $\text{ord}(W(\alpha)) = \alpha$

proof by constructing the isomorphism:  $\phi : A^1 \rightarrow W(\alpha)$

$$\phi(a) = W(\text{ord}(I_A(a)))$$

So by our construction, every set of ordinals  $A$  is a woset. A proof can be made by showing that every  $A' \subseteq A$  has an element  $a$ , which if not already minimal, has a  $W(a) \cup A'$  that contains the minimal element since it's a woset as a subset of a woset.

### 17.1 Cesare Burali-Forti Paradox

The set of all ordinals can't be well defined. Suppose it were a set, it'll be a woset, we'll denote it  $O \Rightarrow \text{ord}(O) \in O \Rightarrow W(O) \subseteq O$  but we know  $\text{ord}(W(O)) = O$  and that means  $I_O(\text{ord}(O)) \cong O$  Thus an initial segment of the set is isomorphic to it which is a contradiction.

### 17.2 Russell's Paradox

**Russell's Paradox** - *Let  $R$  be the set than contains all the sets that don't contain themselves.*

If  $R$  contains itself, it must not contain itself.

If  $R$  doesn't contain itself, then it must contain itself.

Paradox.

---

<sup>1</sup> $|A| = \alpha$  and  $A$  is a woset

### 17.3 Kinds of Ordinals

There are two kinds of ordinals:

**Successor Ordinals** - *ordinals that immediatly success another ordinal*

**Limit Ordinals** - *the rest.*

### 17.4 Ordinal Arthimetic

#### 17.4.1 addition

Let  $(X, \leq_x), (Y, \leq_y)$  be two **disjoint** wosets such that  $(ord(A), ord(B) = (\alpha, \beta))$ .

We denote:  $(X \cup Y, \leq)$ :

$$a \leq b \iff \begin{cases} a, b \in X & a \leq_x b \\ a, b \in Y & a \leq_y b \\ a \in X & b \in Y \end{cases}$$

As  $X \oplus Y$ , And by definition  $\alpha + \beta = ord(X \oplus Y)$

Ordinals are associative but not commutative with addition

- $n + \omega = \omega$
- $\alpha + 0 = \alpha$
- $\omega < \omega + 1 < \omega + 2 < \dots < \omega + k < \dots < 2\omega$

#### 17.4.2 multiplication

Let  $(X, \leq_x), (Y, \leq_y)$  be two **disjoint** wosets such that  $(ord(A), ord(B) = (\alpha, \beta))$ .

We denote:  $(X \times Y, \leq_{dictionary})$  As  $X \odot Y$ , and by definition  $\alpha * \beta = ord(X \odot Y)$

It's possible to show  $\omega = k\omega$  by constructing an isomorphism

$$\phi : \mathbb{N} \rightarrow \{0, 1, \dots, k-1\} \times \mathbb{N}$$

$$\phi(n) = (\lfloor n/k \rfloor, n \bmod k)$$

- $\omega * 0 = 0$
- $\alpha * 1 = \alpha$
- $2\omega = \omega < \omega 2 = \omega + \omega < \omega 3 < \dots < \omega k < \dots < \omega^2$

Ordinals are left distributive but not right distributive. Why?

#### 17.4.3 Powers

$$\alpha^\gamma = \begin{cases} 1 & \gamma = 0 \\ \alpha^{\gamma-1} \alpha & \gamma \text{ is a succesor ordinal} \\ \min_{\delta < \gamma} \{\mu : \alpha^\delta < \mu\} & \gamma \text{ is a limit ordinal} \end{cases}$$

From that we infer the biggest ordinal so far is  $\omega^\omega$

Ordinals are usually expressed as polynomials of powers of  $\omega$  with natural coefficients. And I lied earlier...

$$\omega^\omega < \omega^\omega + 1 < \omega^\omega + 2 < \dots < \omega^{\omega^\omega} < \dots$$

#### 17.4.4 $2^\omega$ and $\omega^\omega$

By our previous definition we can conclude that  $2^\omega$  is

$$\begin{aligned} & \min_{\delta < \omega} \{\mu : 2^\delta < \mu\} \\ &= \min\{2^1, 2^2, \dots, 2^k, \dots\} \end{aligned}$$

Since this series doesn't have an upper bound the result is the smallest infinite ordinal or  $2^\omega = \omega$

By our previous definition we can conclude that  $\omega^\omega$  is

$$\begin{aligned} & \min_{\delta < \omega} \{\mu : \omega^\delta < \mu\} \\ &= \min\{\omega^1, \omega^2, \dots, \omega^k, \dots\} \end{aligned}$$

Let's consider

$$X = X_1 \oplus X_2 \oplus X_3 \dots (\forall n \in \mathbb{N}, \text{ord}(X_n) = \omega^n)$$

We see that this is what we looked for<sup>1</sup> thus  $\text{ord}(X) = \omega^\omega$  but also this is a sum of a countable amount of groups of a countable size so surprisingly  $|X| = \aleph_0$  and this is the case for sets of ordinals  $\omega^{\omega^\omega}$  and so on...

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<sup>1</sup>Verify this is indeed what we looked for

## 18 The Well Ordering Theorem

The Well Ordering Theorem(WOT) states that any set can be well ordered and is equivalent to Zorn's lemma and AC.

### 18.1 Proof Form AC

Let the set we're trying to well-order be  $A$  and let  $f$  be a choice function for the family of non-empty subsets of  $A$ . Now for every ordinal  $\alpha$  define:

$$\begin{cases} a_\alpha = f(A \setminus \{a_\xi \mid \xi < \alpha\}), & A \setminus \{a_\xi \mid \xi < \alpha\} \neq \emptyset \\ a_\alpha = \text{UNDEFINED}, & \text{otherwise} \end{cases}$$

Then

$$\langle a_\alpha \mid a_\alpha \text{ is defined} \rangle$$

Is a well order on  $A$ .

### 18.2 Proof of AC using WOT

To make a choice function for a collection of non-empty sets,  $E$ , take the union of the sets in  $E$  and call it  $X$ . There exists a well-ordering of  $X$ ; let  $R$  be such an ordering. The function that to each set  $S$  of  $E$  associates the smallest element of  $S$ , as ordered by (the restriction to  $S$  of)  $R$ , is a choice function for the collection  $E$ .<sup>2</sup>

It's worth noting the difference between choosing this one choice function here ( $R$ ), and applying the WOT to all the sets  $S \in E$  separately, and choosing the minimal element in each set separately. While the first is allowed under ZF since we're only making a single choice, the latter is not allowed when there are infinitely many elements in  $E$  without assuming the axiom of choice itself, and thus is not a valid way to prove AC.

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<sup>2</sup>This proof was taken straight from wikipedia.

## 19 The Continuum Conjecture

Suppose  $X$  is a set of cardinality  $\aleph$ , and consider the set of all Risha's with a cardinality greater than  $\aleph_0$ <sup>1</sup>, since  $X$  can be a woset, that set has a minimal element  $m$  and  $|m| = \aleph_1$ . The conjecture is that  $\aleph = \aleph_1$ . This was proven to be unsolvable under ZFC. We can also define an  $\aleph$  greater than all  $\aleph$  of the form  $\aleph_n$  where  $n \in \mathbb{N}$  by looking at sets  $|A_n| = \aleph_n$  and at  $B = \bigcup_{i \in \mathbb{N}} \aleph_i$ .  $|B| > \aleph_n (\forall n \in \mathbb{N}) \Rightarrow \exists x \in B : |I_B(x)| > \aleph_n (\forall n \in \mathbb{N})$ . We denote the minimal element of the set of all such  $x$ 's  $M$ , and  $|M| = \aleph_\omega = \sum_{i \in \mathbb{N}} \aleph_i$  and after all countable ordinals we'll reach  $\aleph_\Omega$ , the  $\aleph_1$ -th ordinal, and the first uncountable one. The generalized Continuum Conjecture is:

$$2^{\aleph_\alpha} = \aleph_{\alpha+1}$$

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<sup>1</sup>If it's empty define any one element to be maximal

## 20 Transfinite Induction

Transfinite induction is an extension of mathematical induction to well-ordered sets, for example to sets of ordinal numbers or cardinal numbers. Its correctness is a theorem of ZFC. When given  $X$  a set, and  $P$  the set of elements who have a certain property, the principle that's derived by the WOT can be written like this:

$$\forall x \in X (\forall y < x (y \in P)) \Rightarrow x \in P$$

The steps of transfinite induction:

1. the 0 case ( $0 \in P$ )
2. The successor ordinal case ( $\alpha \in P \Rightarrow \alpha + 1 \in P$ )
3. The limit ordinal case case ( $\forall \beta < \gamma (\beta \in P) \Rightarrow \gamma \in P$ )

### 20.1 Proof That The Only Isomorphism from a Well-Ordered Set to Itself is the Identity Isomorphism

Consider the property  $P$  that "*this element is transformed to itself under all iso- morphisms*". Now consider an element  $a$  such that all elements that are lesser than  $a$  are in  $P$ . This can always be done by choosing the minimal element by WOT.  $a$  can't be transformed to an element lesser than  $a$  because then the isomorphism won't be injective, and also not to an element greater than it, because  $a$  must also have a source, since the isomorphism is surjective, but then we get a contradiction to the the fact the isomorphism is a homomorphism.

**note:** transfinite induction works because of WOT but there are of course sets like  $\mathbb{R}$  with normal ordering that isn't a woset so we can't use transfinite induction on it. A counter example for our proof may be  $f(x) = x + 1$



## 21 Extras

### 21.1 A Bit About Constructions

Constructions of sets are the way to formally define sets like  $\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}$

#### 21.1.1 Construction of $\mathbb{N}$

There are multiple ways<sup>3</sup> to define  $\mathbb{N}$  one in ZF is recursively defining the natural numbers as such:

$$0 = \{\} = \emptyset$$

$$1 = \{0\} = \{\{\}\}$$

$$2 = \{0, 1\} = \{\emptyset, \{\emptyset\}\}$$

And so on defining numbers using the successor function  $S(n) = n \cup \{n\}$ .  $\mathbb{N}$  is the smallest set containing 0 and closed under  $S(n)$

#### 21.1.2 Construction of $\mathbb{Z}$

$\mathbb{Z}$  was constructed as  $\mathbb{N} \times \mathbb{N}/R$

$$\langle a, b \rangle R \langle c, d \rangle \iff c - d = a - b$$

For example  $-1 = \{\langle 2, 3 \rangle, \langle 5, 6 \rangle, \dots, \langle n, n+1 \rangle\}$   
Define  $\mathbb{Z}_+, \mathbb{Z}_*$

#### 21.1.3 Construction of $\mathbb{Q}$

$\mathbb{Z}$  was constructed as  $\mathbb{Z} \times \mathbb{Z}'/R$

$$\langle a, b \rangle R \langle c, d \rangle \iff ad = cb$$

For example  $\frac{1}{2} = \{\langle 1, 2 \rangle, \langle -2, -4 \rangle, \dots, \langle n, 2n \rangle\}$   
Define  $\mathbb{Q}_+, \mathbb{Q}_*$

#### 21.1.4 About the Construction of $\mathbb{R}$

The construction of  $\mathbb{R}$  is more difficult than you may expect. It should be studied in a number theory course, and is usually very unrigorously defined as all numbers in the interval  $(-\infty, \infty)$

### 21.2 Discrete Functions

**Discrete Function** - A function that is defined only for a set of numbers that can be listed, such as the set of whole numbers or the set of integers.

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<sup>3</sup>One of them is by isomorphism classes of finite sets

<sup>4</sup> $\mathbb{Z}' = \mathbb{Z} \setminus \{0\}$

## 21.3 More Definitions

### 21.3.1 Saturated Functions

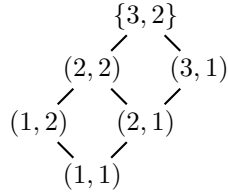
For a function  $f : X \rightarrow Y$

$$\forall A \subseteq X, f_*(A) = \{f(x) : x \in A (A \in P(X))\}$$

$$\forall B \subseteq Y, f^*(B) = \{x : f(x) \in B (B \in P(Y))\}$$

### 21.3.2 Hasse Diagrams

Hasse diagrams represent posets. For example the Hasse Diagram of the the set  $\{1, 2, 3\} \times \{1, 2\}$  with the standard order.



### 21.3.3 Some Denotions

- A Singleton is a set containing only one element.
- $P(A) = \{B : B \subseteq A\}$
- $A \triangle B = \{A \cup B\} \setminus \{A \cap B\}$
- $|\mathbb{R}| = c = \beth_1 = \aleph$
- $A^c = \{b : b \notin A\}$
- $\prod_{i \in I} X_i = \{f : I \rightarrow \bigcup_{i \in I} X_i \mid \forall i \in I (f(i) \in X_i)\}$
- A pairing function is a bijection  $\pi : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$
- The indicator function of  $A \subseteq X$  is  $1_A(x) = I_A(x) = \chi_A(x) = 1 \iff x$  is in  $A$  and equals 0 otherwise