

COMP 321: Introduction to Computer Systems

Project 4: Unix Shell

Assigned: 2/18/16, **Due: 3/10/16, 11:55 PM**

Important: This project may be done individually or in pairs. Be sure to carefully read the course policies for assignments (including the honor code policy) on the assignments page of the course web site:

<http://www.clear.rice.edu/comp321/html/assignments.html>

Overview

You will write a simple Unix shell program that supports job control. The goals of this project are as follows:

- Work with a more realistic system application.
- Become familiar with the concepts of process control.
- Become familiar with the concepts of signalling.

General Overview of Unix Shells

A *shell* is an interactive command-line interpreter that runs programs on behalf of the user. A shell repeatedly prints a prompt, waits for a *command line* on `stdin`, and then carries out some action, as directed by the contents of the command line.

The command line is a sequence of ASCII text words delimited by whitespace. The first word in the command line is either the name of a built-in command or the name of an executable file. The remaining words are command-line arguments. If the first word is a built-in command, the shell immediately executes the command in the current process. Otherwise, the shell assumes that the first word is the name of a program to execute. In this case, the shell forks a child process, then attempts to load and run the program in the context of the child. The child processes created as a result of interpreting a single command line are known collectively as a *job*. In general, a job can consist of multiple child processes connected by Unix pipes.

If the first word in the command line starts with a directory, such as “.”, “/”, or the name of a subdirectory in the current directory, the word is assumed to be the path name of an executable program, for example, “./tsh”, “/usr/bin/ls”, or “comp321/count/count”. Otherwise, the first word is assumed to be the name of an executable that is contained in one of the directories in the shell’s search path, which is an ordered list of directories that the shell searches for executables.

If the command line ends with an ampersand (“&”), then the job runs in the *background*, which means that the shell does not wait for the job to terminate before printing the prompt and awaiting the next command line. Otherwise, the job runs in the *foreground*, which means that the shell waits for the job to terminate before awaiting the next command line. Thus, at any point in time, at most one job can be running in the foreground. However, an arbitrary number of jobs can run in the background.

For example, typing the command line

```
tsh> jobs
```

causes the shell to execute the built-in `jobs` command.

In contrast, typing the command line

```
tsh> /bin/ls -l -d
```

runs the `ls` program in the foreground. By convention, the shell ensures that when the program begins executing its main routine

```
int main(int argc, char *argv[])
```

the `argc` and `argv` arguments have the following values:

- `argc == 3`,
- `argv[0] == "/bin/ls"`,
- `argv[1] == "-l"`,
- `argv[2] == "-d"`.

Alternatively, typing the command line

```
tsh> /bin/ls -l -d &
```

runs the `ls` program in the background.

Unix shells support the notion of *job control*, which allows users to move jobs back and forth between the background and foreground, and to change the process state (running, stopped, or terminated) of the processes in a job. Typing `ctrl-c` causes a `SIGINT` signal to be delivered to each process in the foreground job. The default action for `SIGINT` is to terminate the process. Similarly, typing `ctrl-z` causes a `SIGTSTP` signal to be delivered to each process in the foreground job. The default action for `SIGTSTP` is to place a process in the stopped state, where it remains until it is awakened by the receipt of a `SIGCONT` signal. Unix shells also provide various built-in commands that support job control. For example:

- `jobs`: List the running and stopped background jobs.
- `bg <job>`: Change a stopped job to a running background job.
- `fg <job>`: Change a stopped or running background job to a running job in the foreground.

The Tiny Shell (`tsh`) Specification

Your `tsh` shell should have the following features:

- The prompt should be the string `"tsh> "`.
- The search path should be the string contained by the `PATH` environmental variable. The search path is a string that contains a list of directories, where each directory is separated by the `":"` character. If the search path contains an empty string, this should be interpreted as the current directory being located in the search path. The search path can contain an empty string if the entire path is an empty string, if there are two `":"` characters in a row, or if the path starts or ends with the `":"` character. The search path should be loaded when the program is initialized.

- The command line typed by the user should consist of a `name` and zero or more arguments, all separated by one or more spaces. If `name` is a built-in command, then `tsh` should handle it immediately and wait for the next command line. Otherwise, `tsh` should assume that `name` is the name of an executable file, which it loads and runs in the context of an initial child process. This initial child process is called a *job*.

If `name` does not start with a directory and the search path is not `NULL`, then `tsh` should assume that `name` is the name of an executable file contained in one of the directories in the search path. In this case, `tsh` should search the directories in the search path, in order, for an executable named `name`. The first executable named `name` found in the search path that `tsh` can execute is the program that `tsh` loads and runs. If `name` starts with a directory or the search path is `NULL`, `tsh` should assume that `name` is the path name of an executable file.

- `tsh` must use the `execve` function. In other words, `tsh` may not use any other member of the `exec` family, such as `execvp`.
- `tsh` need not support pipes (`|`) or I/O redirection (`<` and `>`).
- Typing `ctrl-c` (`ctrl-z`) should cause a `SIGINT` (`SIGTSTP`) signal to be sent to the current foreground job, as well as any descendents of that job (e.g., any child processes that it forked). If there is no foreground job, then the signal should have no effect.
- If the command line ends with an ampersand (`&`), then `tsh` should run the job in the background. Otherwise, it should run the job in the foreground.
- Each job can be identified by either a process ID (PID) or a job ID (JID), which is a positive integer assigned by `tsh`. JIDs should be denoted on the command line by the prefix `%`. For example, `%5` denotes JID 5, and `5` denotes PID 5. (We have provided you with all of the routines you need for manipulating the job list.)
- `tsh` should support the following built-in commands:
 - The `quit` command terminates the shell.
 - The `jobs` command lists all background jobs.
 - The `bg <job>` command restarts `<job>` by sending it a `SIGCONT` signal, and then runs it in the background. The `<job>` argument can be either a PID or a JID.
 - The `fg <job>` command restarts `<job>` by sending it a `SIGCONT` signal, and then runs it in the foreground. The `<job>` argument can be either a PID or a JID.
- `tsh` should reap all of its zombie children. If any job terminates because it receives a signal that it didn't catch, then `tsh` should recognize this event and print a message with the job's PID and a description of the offending signal of the form

```
Job [1] (11639) terminated by signal SIGINT
```

Where the job ID is 1 and the process id is 11639 and the job terminated because of a `SIGINT` signal. The provided `signame` array is useful for converting signal numbers into names. Specifically, use the signal number as the index for selecting an element of the array. The element will be the address of a string containing the symbol for that signal, without the “SIG” prefix.

Provided Files

The provided files are all available in

```
/clear/www/htdocs/comp321/assignments/shell/
```

To begin working on the project, do the following:

- Copy all of the files in this directory to the directory in which you plan to do your work.
- Type the command `make` to compile and link a skeleton shell and some simple programs you can run from within your shell to test it.
- Type your name(s) and NetID(s) in the header comment at the top of `tsh.c`.

If you choose to work in a pair, you may find it helpful to share files via Subversion. Instructions for doing so can be found in the subversion lab on the course web site.

Looking at the `tsh.c` (*tiny shell*) file, you will see that it contains a functional skeleton of a simple Unix shell. To help you get started, we have already implemented the less interesting functions. Your assignment is to complete the remaining empty functions listed below. As a sanity check for you, we've listed the approximate number of lines of code for each of these functions in our reference solution (which includes lots of comments).

- `initpath`: Performs any necessary initialization of the search path. The complexity of this function affects the complexity of the `eval` function. Specifically, a more complicated `initpath` function may allow for a simpler `eval` function. [1–20 lines]
- `eval`: Main routine that parses and interprets the command line. [90–110 lines]
- `builtin_cmd`: Recognizes and interprets the built-in commands: `quit`, `fg`, `bg`, and `jobs`. [25 lines]
- `do_bgfg`: Implements the `bg` and `fg` built-in commands. [50 lines]
- `waitfg`: Waits for a foreground job to complete. [20 lines]
- `sigchld_handler`: Catches `SIGCHLD` signals. [80 lines]
- `sigint_handler`: Catches `SIGINT` (`ctrl-c`) signals. [15 lines]
- `sigtstp_handler`: Catches `SIGTSTP` (`ctrl-z`) signals. [15 lines]

Each time you modify your `tsh.c` file, type `make` to recompile it. To run your shell, type `tsh` to the command line:

```
UNIX% ./tsh
tsh> [type commands to your shell here]
```

Checking Your Work

We have provided some tools to help you check your work.

You may assume that `MAXLINE`, `MAXARGS`, `MAXJOBS`, and `MAXJID` are the maximum number of characters in any command line, the maximum number of arguments on any command line, the maximum number of jobs at any time, and the maximum job ID, respectively. You do not need to check to ensure that these maximums will be met, rather you may assume that no one will type any input into the shell that would cause them to be exceeded.

Reference solution. The executable `tshref` is the reference solution for the shell. Run this program to resolve any questions you have about how your shell should behave. *Your shell should emit output that is identical to the reference solution* (except for PIDs, of course, which change from run to run).

Shell driver. The `sdriver.pl` program executes a shell as a child process, sends it commands and signals as directed by a *trace file*, and captures and displays the output from the shell.

Use the `-h` argument to find out the usage of `sdriver.pl`:

```
UNIX% ./sdriver.pl -h
Usage: sdriver.pl [-hv] -t <trace> -s <shellprog> -a <args>
Options:
  -h                Print this message
  -v                Be more verbose
  -t <trace>        Trace file
  -s <shell>        Shell program to test
  -a <args>         Shell arguments
  -g                Generate output for autograder
```

We have also provided sample trace files (`trace*.txt`) that you can use in conjunction with the shell driver to test the correctness of your shell. The lower-numbered trace files do very simple tests, and the higher-numbered tests do more complicated tests.

You can run the shell driver on your shell using trace file `trace01.txt` (for instance) by typing:

```
UNIX% ./sdriver.pl -t trace01.txt -s ./tsh -a "-p"
```

(the `-a "-p"` argument tells your shell not to emit a prompt), or

```
UNIX% make test01
```

Similarly, to compare your result with the reference shell, you can run the trace driver on the reference shell by typing:

```
UNIX% ./sdriver.pl -t trace01.txt -s ./tshref -a "-p"
```

or

```
UNIX% make rtest01
```

For your reference, `tshref.out` gives the output of the reference solution for the provided traces. This might be more convenient for you than manually running the shell driver on all trace files.

The neat thing about the trace files is that they generate the same output you would have gotten had you run your shell interactively (except for an initial comment that identifies the trace). For example:

```
UNIX% make test07
./sdriver.pl -t trace07.txt -s ./tsh -a "-p"
#
# trace07.txt - Forward SIGINT only to foreground job.
#
```

```
tsh> ./myspin 4 &
[1] (27778) ./myspin 4 &
tsh> ./myspin 5
Job [2] (27780) terminated by signal SIGINT
tsh> jobs
[1] (27778) Running ./myspin 4 &
UNIX%
```

Hints

- Read every word of Chapter 8 (Exceptional Control Flow) in your textbook. **Significant changes were made to this chapter in the third edition of the textbook.** In particular, the material in Section 8.5.5 on safe signal handling did not exist in the previous edition.
- Read every word of this project handout.
- Use the trace files to guide the development of your shell. Starting with `trace01.txt`, make sure that your shell produces the *identical* output as the reference shell. Then move on to trace file `trace02.txt`, and so on.
- The functions `waitpid`, `kill`, `fork`, `execve`, `setpgid`, and `sigprocmask` will be the fundamental building blocks of your shell. Make sure that you understand what these functions do and how to use them. The `WUNTRACED` and `WNOHANG` options to `waitpid` will also be useful.
- When you implement your signal handlers, be sure to send `SIGINT` and `SIGTSTP` signals to the entire foreground process group, using “-pid” instead of “pid” in the argument to the `kill` function. The `sdriver.pl` program tests for this error.
- One of the tricky parts of the assignment is deciding on the allocation of work between the `waitfg` and `sigchld_handler` functions. We recommend the following approach:
 - In `waitfg`, use a loop around the `sleep` function that tests whether or not the specified process is still running in the foreground.
 - In `sigchld_handler`, use `waitpid` to reap zombie children.

While other solutions are possible, such as calling `waitpid` in both `waitfg` and `sigchld_handler`, these can be very confusing. It is simpler to do all reaping in the handler.

- When you implement your signal handlers, be sure to follow the rules described in Section 8.5.5. In particular, per rule *GI*, you must not call functions such as `printf` or `fprintf` to print messages.
- In `eval`, the parent must use `sigprocmask` to block `SIGCHLD` signals before it forks the child, and then unblock these signals, again using `sigprocmask` after it adds the child to the job list by calling `addjob`. Since children inherit the blocked vectors of their parents, the child must be sure to then unblock `SIGCHLD` signals before it execs the new program.

The parent needs to block the `SIGCHLD` signals in this way in order to avoid the race condition where the child is reaped by `sigchld_handler` (and thus removed from the job list) *before* the parent calls `addjob`.

- Programs such as `more`, `less`, `vi`, and `emacs` do strange things with the terminal settings. Don’t run these programs from your shell. Stick with simple text-based programs such as `/bin/ls`, `/bin/ps`, `/bin/echo`, and the included test programs (such as `myspin`).

- When you run your shell from the standard Unix shell, your shell is running in the foreground process group. If your shell then creates a child process, by default that child will also be a member of the foreground process group. Since typing `ctrl-c` sends a `SIGINT` to every process in the foreground group, typing `ctrl-c` will send a `SIGINT` to your shell, as well as to every process that your shell created, which obviously isn't correct.

Here is the workaround: After the `fork`, but before the call to `execve`, the child process should call `setpgid(0, 0)`, which puts the child in a new process group whose group ID is identical to the child's PID. This ensures that there will be only one process, your shell, in the foreground process group. When you type `ctrl-c`, the shell should catch the resulting `SIGINT` and then forward it to the appropriate foreground job (or more precisely, the process group that contains the foreground job).

- As described, your shell will not be able to handle programs that read from the standard input, and you are not required to do so. Note that this means you will not be able to run your shell from within your shell. Handling programs that read from the standard input requires your shell to give control of the terminal to its foreground child, which is beyond the scope of this project. We suggest you stick to the simple text based programs mentioned above and to the provided test programs (such as `mypin`).
- A string is guaranteed to start with a directory if the string contains a `"/`.

Evaluation

To turn in your solution, you should use the `turnin` process on CLEAR (<https://docs.rice.edu/confluence/x/qAiUAQ>). First, create a `shell` directory within your `comp321` directory. Then, be sure that your solution is contained in a file named `tsh.c` and your writeup is contained in a file named `writeup.txt` inside this `shell` directory. Finally, submit your solution by running the following command from inside your `comp321` directory:

```
UNIX% turnin comp321-S16:shell
```

Assignments can be submitted multiple times. Grading will always be done on the most recent submission before the deadline, unless slip days are used. If you would like, you can also turn in additional files with testing code. **However, remember that we will *only* compile your `tsh.c` file and nothing else!**

The project will be graded as follows:

- Writeup: 15%
- Style: 15%
- Correctness: 70%

Your solution will be tested for correctness on a CLEAR machine, using the same shell driver and trace files that were provided (along with additional traces). Your shell should produce **identical** output on these traces, with only two exceptions:

- The PIDs can (and will) be different.
- The output of the `/bin/ps` commands will be different from run to run. However, the running states of any processes run from within your shell in the output of the `/bin/ps` command should be identical.

For each test, you will only receive credit for that test if the output of your shell is identical to the output of the reference shell (other than the exceptions listed above).