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## Lecture Notes 5: Sequences, Strings and Relations

(2024A, Week 5, Monday, April 1, 16:30-18:00)<sup>5</sup>

A **sequence** s is a function whose domain D is a subset of integers. The notation  $s_n$  is typically used instead of the more general function notation s(n). The term n is the **index** of the sequence. If D is a finite set, we call s a **finite** sequence; otherwise, s is an **infinite** sequence.

- (a) A sequence s is **increasing** (respectively, **decreasing**) if for all i and j in the domain of s, if i < j, then  $s_i < s_j$  (respectively,  $s_i > s_j$ ).
- (b) A sequence s is **nondecreasing** (respectively, **nonincreasing**) if for all i and j in the domain of s, if i < j, then  $s_i \le s_j$  (respectively,  $s_i \ge s_j$ ).
- (c) A **subsequence** of a sequence s is a sequence obtained from s by choosing certain terms of s in the same order in which they appear in s.
- (d) If  $\{a_i\}_{i=m}^n$  is a sequence, we define

$$\sum_{i=m}^{n} a_i = a_m + a_{m+1} + \dots + a_n, \quad \prod_{i=m}^{n} a_i = a_m a_{m+1} + \dots + a_n. \tag{2}$$

The formalism  $\sum_{i=m}^{n}$  is the **sum** (or **sigma**) notation and  $\prod_{i=m}^{n}$  is the **product** notation. In (2), i is the **index**, m is the **lower limit** and n is the **upper limit**.

A **string** over X, where X is a finite set, is a finite sequence of (repeatable) elements from X. Since a string is a sequence, order is taken into account.

- (a) Repetitions in a string can be specified by **superscripts**. For example, the string '110001' over the binary set  $\{0,1\}$  can be written as '1<sup>2</sup>0<sup>3</sup>1'.
- (b) The string with no element is the **null string** and is denoted by  $\lambda$ .
- (c)  $X^*$  denotes the set of all strings over X including the null string, and  $X^+$  denotes the set of all nonnull strings over X.
- (d) The **length** of a string  $\alpha$ , denoted by  $|\alpha|$ , is the number of elements in  $\alpha$ .
- (e) If  $\alpha$  and  $\beta$  are two strings, the string consisting of  $\alpha$  followed by  $\beta$ , written  $\alpha\beta$  (or  $\alpha + \beta$ ), is the **concatenation** of  $\alpha$  and  $\beta$ .
- (f) A string  $\beta$  is a **substring** of a string  $\alpha$  if there are strings  $\gamma$  and  $\delta$  such that  $\alpha = \gamma \beta \delta$ .

<sup>&</sup>lt;sup>5</sup>Most of the content of this document is taken from the book [1].

A (binary) **relation** R from a set X to a set Y is a subset of the Cartesian product  $X \times Y$ . If  $(x,y) \in R$ , we write xRy and say that x is related to y. If X = Y, R is a relation on X.

An informative way to picture a relation on a set is to draw its **digraph**. To draw the digraph of a relation on a set X, we first draw dots or vertices to represent the elements of X. Next, if the element (x, y) is in the relation, we draw an arrow (a **directed edge**) from x to y. An element of the form (x, x) in a relation corresponds to a directed edge from x to x. Such an edge is a **loop**.

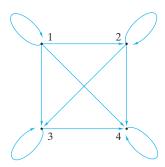


Figure 7: Digraph of the relation xRy if  $x \le y$  on  $\{1, 2, 3, 4\}$ .

Let R be a relation on a set X.

- (a) R is **reflexive** if  $(x, x) \in R$  for every  $x \in X$ .
- (b) R is **symmetric** if for all  $x, y \in X$ , if  $(x, y) \in R$ , then  $(y, x) \in R$ .
- (c) R is **antisymmetric** if for all  $x, y \in X$ , if  $(x, y) \in R$  and  $(y, x) \in R$ , then x = y.
- (d) R is **transitive** if for all  $x, y, z \in X$ , if (x, y) and  $(y, z) \in R$ , then  $(x, z) \in R$ .
- (e) R is a **partial order** if R is reflexive, antisymmetric and transitive. If R is a partial order on a set X, the notation  $x \leq y$  is used to indicate that  $(x,y) \in R$ . This notation suggests that we are interpreting the relation as an ordering of the elements in X. Let R be a partial order on X. If  $x,y \in X$  and either  $x \leq y$  or  $y \leq x$ , we say that x and y are **comparable**. If  $x,y \in X$  and  $x \nleq y$  and  $y \nleq x$ , we say that x and y are **incomparable**. If every pair of elements in X is comparable, R is a **total order**.
- (f) R is an **equivalence** relation if R is reflexive, symmetric, and transitive. Let R be an equivalence relation on a set X. For each  $a \in X$ , let  $[a] = \{x \in X \mid xRa\}$ . Then  $S = \{[a] \mid a \in X\}$  is a partition of X. The set [a] is the **equivalence class** of a in X given by the relation R.

Let R be a relation from X to Y. The **inverse** of R, denoted by  $R^{-1}$ , is the relation from Y to X defined by

$$R^{-1} = \{(y, x) \mid (x, y) \in R\}.$$

Let  $R_1$  be a relation from X to Y and  $R_2$  be a relation from Y to Z. The **composition** of  $R_1$  and  $R_2$ , denoted by  $R_2 \circ R_1$ , is the relation from X to Z defined by

$$R_2 \circ R_1 = \{(x, z) \mid (x, y) \in R_1 \text{ and } (y, z) \in R_2 \text{ for some } y \in Y\}.$$

A matrix is a convenient way to represent a relation R from X to Y. We label the rows with the elements of X (in some order) and the columns with the elements of Y (again, in some order). We then set the entry in row x and column y to 1 if xRy and to 0 otherwise. This matrix is the **matrix of the relation** R relative to the chosen orderings of X and Y.

## References

1. Johnsonbaugh, R.: Discrete Mathematics - Eighth Edition. *Pearson Education*, New York (2018).