

# 430.211 Programming Methodology (프로그래밍 방법론)

## Sorting Algorithm part I

Lab2 #02

Fall 2025

# Attendance check

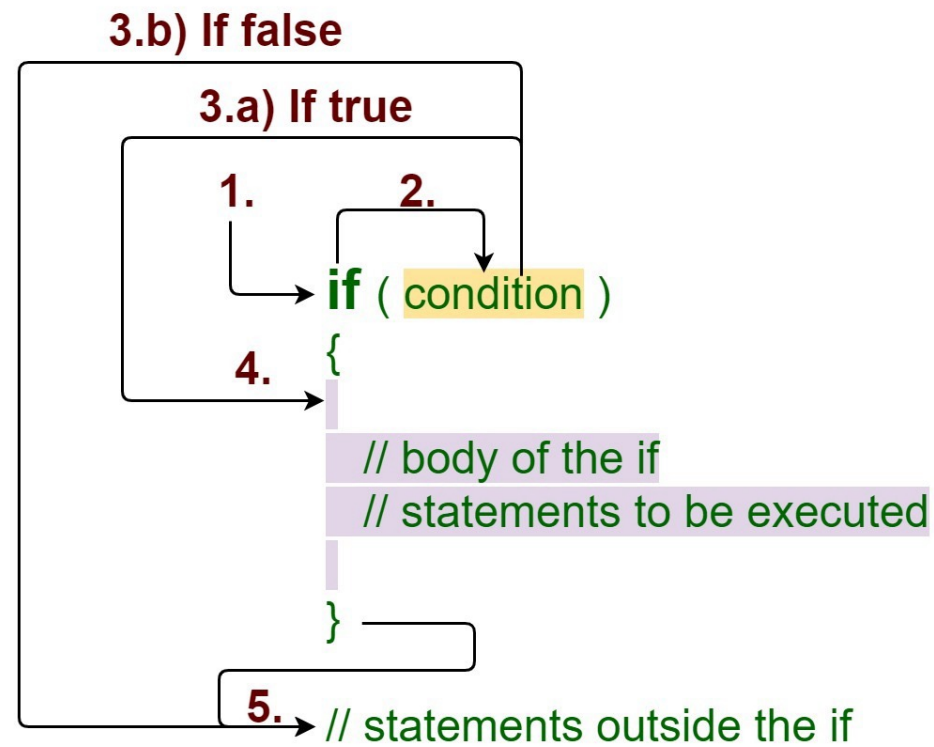
- Attendance will be taken after the class ends
  - With Quiz system in ELICE platform

# Outline

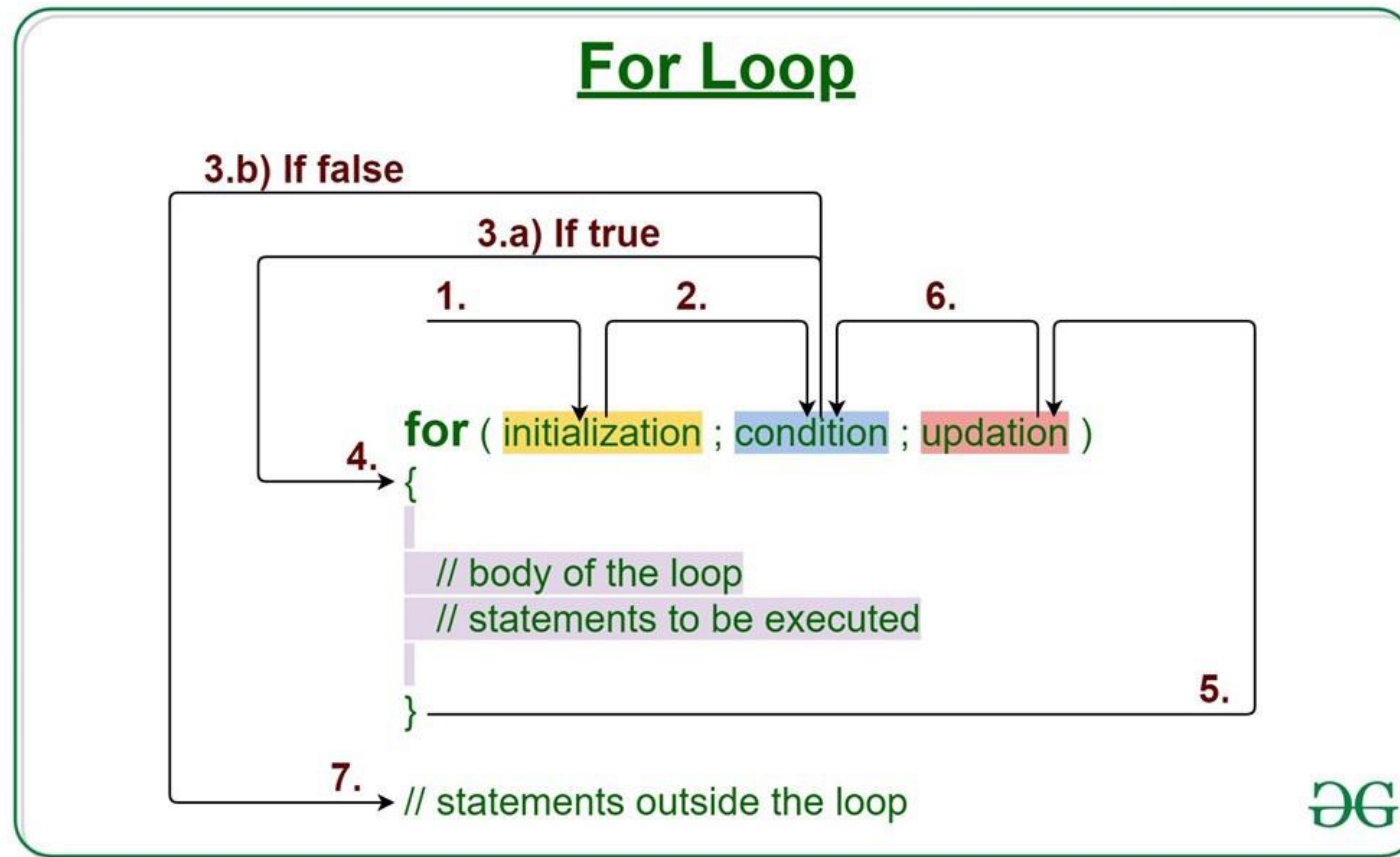
- **If / For loop / While loop**
  - Since we are ahead of the **Wednesday** class schedule
- Array
- What is sorting?
- Incremental Algorithm
  - Bubble sort
  - Insertion sort
  - Selection sort
  - Heap sort

# If

## If statement



# For loop



- **Initialization** is executed (one time) before the execution of the code block.
- **Condition** defines the condition for executing the code block.
- **Update** is executed (every time) after the code block has been executed.

# For loop

```
#include <bits/stdc++.h>
using namespace std;

int main() {

    // for loop to print "Hi" 5 times
    for (int i = 5; i < 10; i++) {
        cout << "Hi" << endl;
    }

    return 0;
}
```

## Output

```
Hi
Hi
Hi
Hi
Hi
```

# For loop

```
#include <iostream>
using namespace std;

int main() {

    // Initial value of number
    int n = 5;

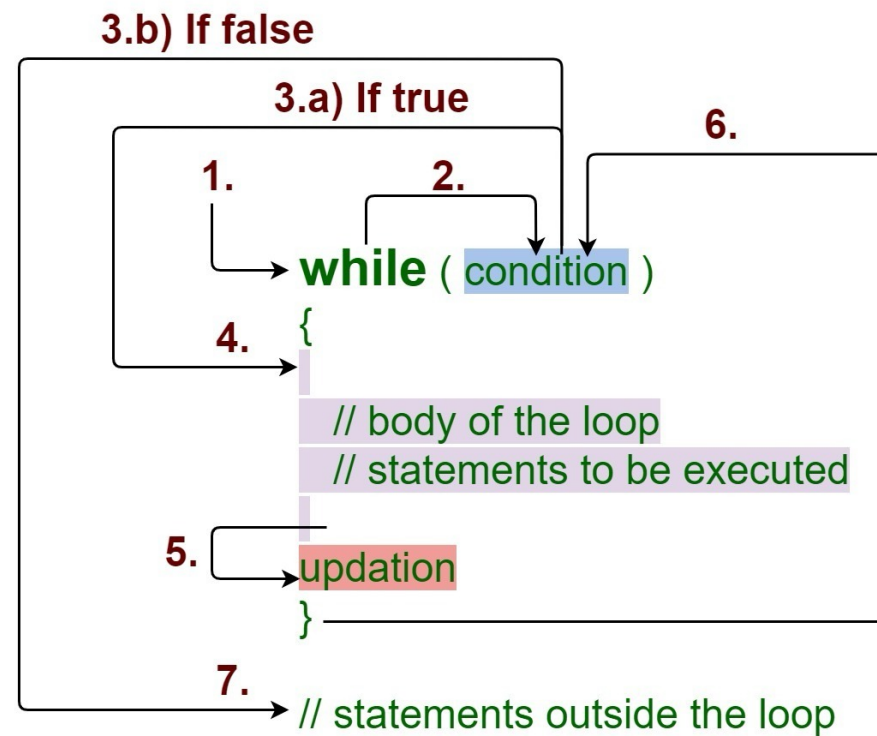
    // Initialization of loop variable
    int i;
    for (i = n; i >= 1; i--)
        cout << i << " ";
    return 0;
}
```

## Output

5 4 3 2 1

# While loop

## While Loop





# While loop

```
#include <bits/stdc++.h>
using namespace std;

int main() {

    // Declaration and Initialization of loop variable
    int i = 1;

    // while loop to print numbers from 1 to 5
    while (i <= 5) {
        cout << i << " ";

        // Updating loop varialbe
        i++;
    }

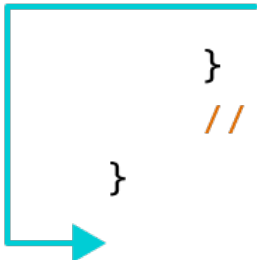
    return 0;
}
```

## Output

1 2 3 4 5

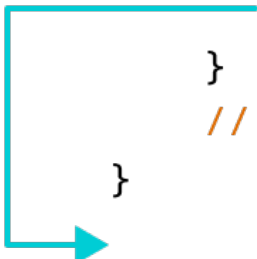
# Break statement

```
for (init; condition; update) {  
    // code  
    if (condition to break) {  
        break;  
    }  
    // code  
}
```



---

```
while (condition) {  
    // code  
    if (condition to break) {  
        break;  
    }  
    // code  
}
```



# Outline

- If / For loop / While loop
- **Array**
- What is sorting?
- Incremental Algorithm
  - Bubble sort
  - Insertion sort
  - Selection sort
  - Heap sort

# Array

- Array definition:
  - A collection of data of a same type
  - `int, float, double, char` are simple data types
- Array indexes always start with zero!
  - To computer scientists, “zero” is the "first" number

# Array

- Initialization:

```
int arr[6] = {2, 12, 1, 7, 3, 4}; // declaration and initialization
```

- Equivalent to the following:

```
int arr[6]; // declaration
```

```
arr[0] = 2; // assignment
```

```
arr[1] = 12;
```

```
arr[2] = 1;
```

```
arr[3] = 7;
```

```
arr[4] = 3;
```

```
arr[5] = 4;
```

<b>2</b>	<b>12</b>	<b>1</b>	<b>7</b>	<b>3</b>	<b>4</b>
----------	-----------	----------	----------	----------	----------

arr[0]   arr[1]   arr[2]   arr[3]   arr[4]   arr[5]

- Value in brackets called index or subscript

- Numbered from **0 to (size - 1)**

- Elements are stored sequentially in a contiguous block of memory.

# Outline

- For loop and while loop
- Array
- **What is sorting?**
- Incremental Algorithm
  - Bubble sort
  - Insertion sort
  - Selection sort
  - Heap sort

# What is Sorting?

- Arranging items in ascending or descending order

Unsorted array

12	4	18	1	16	7	10	15	2
----	---	----	---	----	---	----	----	---

Sorted array in ascending order.

1	2	4	7	10	12	15	16	18
---	---	---	---	----	----	----	----	----

Sorted array in descending order.

18	16	15	12	10	7	4	2	1
----	----	----	----	----	---	---	---	---

# What is Sorting?

- There are various sorting algorithms

- Bubble sorting
- Insertion sorting
- Selection sorting
- Heap sorting
- Merge sorting
- Quick sorting
- Counting sorting..., and so on

Today lecture



# What is Sorting?

- Which sorting algorithm is appropriate ?
  - Time complexity / Space complexity (Big O notation)
  - But, we will mainly focus on implementation in this course.

Formal definition of Big-Oh:

$f(N) = O(g(N))$ , if there exists positive constants  $c$ ,  $N_0$  such that

$$f(N) \leq c \cdot g(N) \text{ for all } N \geq N_0.$$

If  $c = 3, N_0 = 10$

Regular	Big-O	
2	$O(1)$	--> It's just a constant number
$2n + 10$	$O(n)$	--> $n$ has the largest effect
$5n^2$	$O(n^2)$	--> $n^2$ has the largest effect
$f(N)$	$g(N)$	

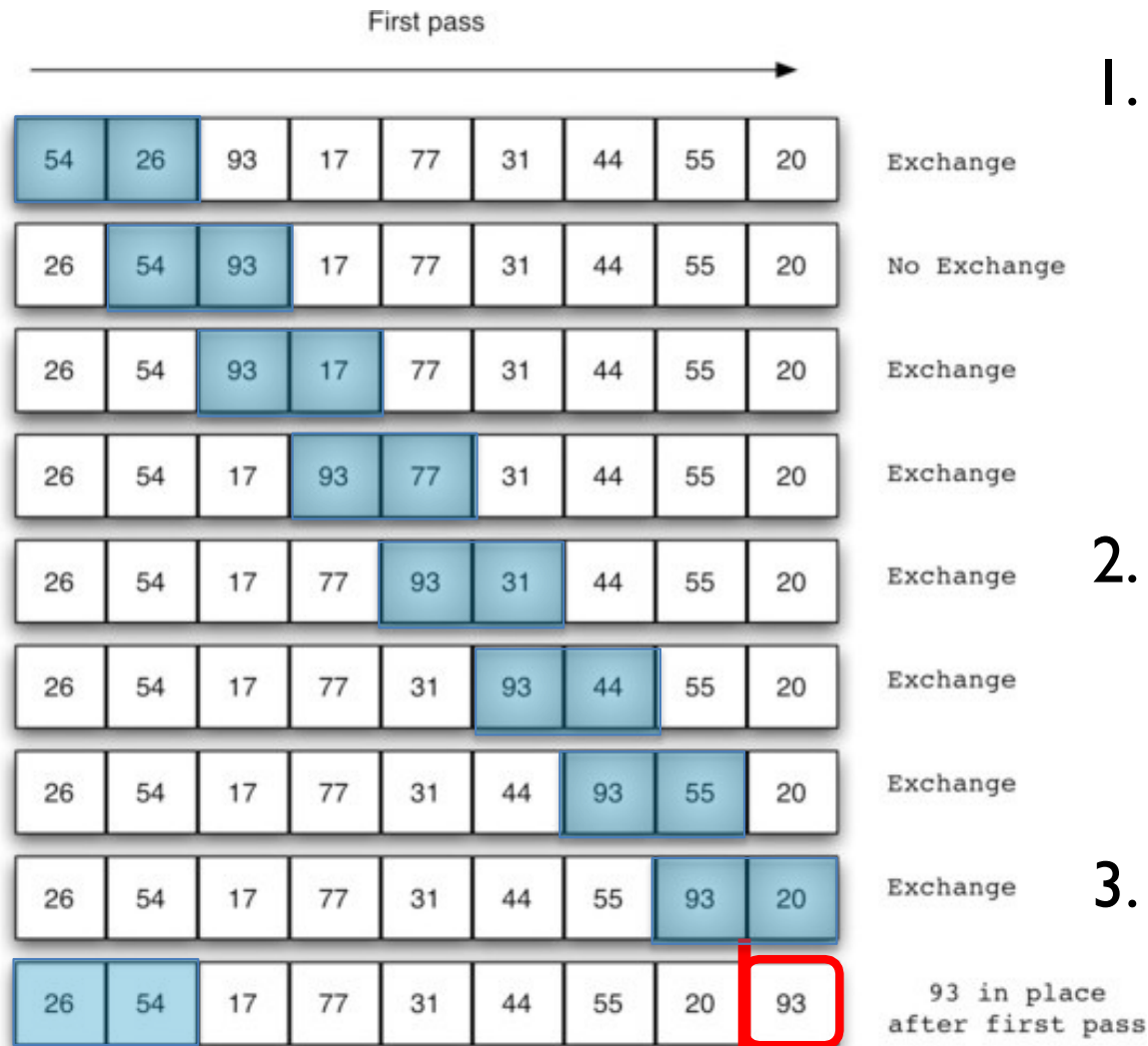
# What is Sorting?

- Which sorting algorithm is appropriate ?
  - **Time** complexity / **Space** (memory) complexity (Big O notation)
    - Let # of items to sort =  $n$
    - Count the number of operations in the program logic
      - Add / divide / multiply two items
      - Compare two items
      - ...
  - (ex) if algorithm A requires a total of  $2n^2 + 4n$  operations, then algorithm A has a time complexity of  $O(n^2)$

# Outline

- If / For loop / While loop
- Array
- What is sorting?
- **Incremental Algorithm**
  - **Bubble sort**
  - Insertion sort
  - Selection sort
  - Heap sort

# Bubble Sort



## 1. First pass

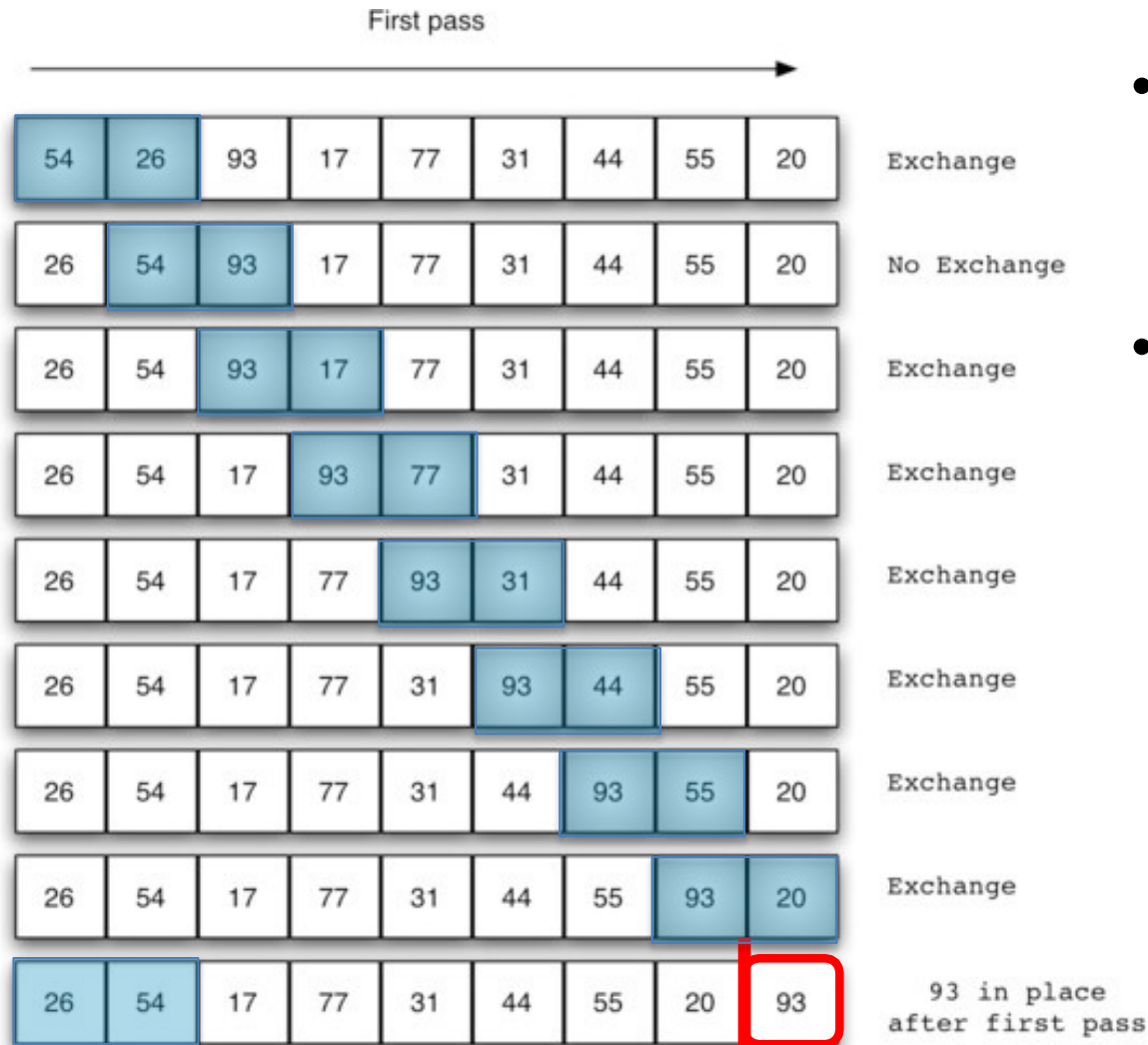
- Compare **two elements**
- Swap in the desirable order until the last element
- **The largest value has bubbled up to the rightmost position**

## 2. Second Pass

- Compare **two elements**
- Swap in the right order until the second-to-last element

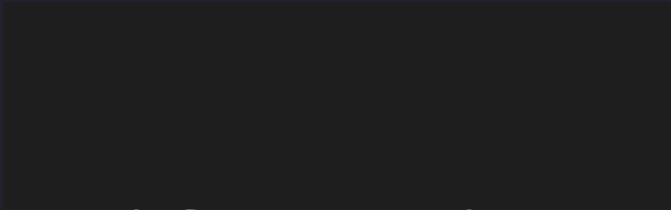
## 3. Continue...

# Bubble Sort



- Pros :
  - Easy to implement
- Cons :
  - As the length of input increases, time increases exponentially.
  - Time complexity is  $O(n^2)$

# Bubble Sort

```
for (unsigned int i = 0; i < size - 1; i++)
    for (unsigned int j = 0; j < size - 1 - i; j++)
    {
        if (array[j] > array[j + 1])
        {
            //#####implement here#####
            
            //#####
        }
    }
```

- Hint
  - Index “i” : Used for passes
  - Index “j” : Used for comparing two indexes (Blue box)

# Bubble Sort

```
for (unsigned int i = 0; i < size - 1; i++)
    for (unsigned int j = 0; j < size - 1 - i; j++)
    {
        if (array[j] > array[j + 1])
        {
            //#####implement here#####
            int temp = array[j];
            array[j] = array[j+1];
            array[j+1] = temp;
            //#####
        }
    }
```

- Hint
  - Index “i” : Used for passes
  - Index “j” : Used for comparing two indexes  
(Blue box)

# Bubble Sort

```
for (unsigned int i = 0; i < size - 1; i++)  
    for (unsigned int j = 0; j < size - 1 - i; j++)  
    {  
        if (array[j] > array[j + 1])  
        {  
            //#####implement here#####  
            int temp = array[j];  
            array[j] = array[j+1];  
            array[j+1] = temp;  
            //#####  
        }  
    }
```

$\approx N^2/2$  iterations

3 operations

$= O(N^2)$  time complexity

- Hint
  - Index “i” : Used for passes
  - Index “j” : Used for comparing two indexes  
(Blue box)



# Outline

- For loop and while loop
- Array
- What is sorting?
- **Incremental Algorithm**
  - Bubble sort
  - **Insertion sort**
  - Selection sort
  - Heap sort

# Insertion Sort



- Start from **index=1**
  - Compare with **the indexes** in red blocks
  - Insert the **target** in an **appropriate** location
- Start from **index=2**
  - ... continue

# Insertion Sort

0	1	2	3	4	5	6	7
5	6	7	9	15	17	10	11

target = 10  
Target Index = 6

0	1	2	3	4	5	6	7
5	6	7	9	15	17	10	11
5	6	7	9	15	17	17	11

Compare index = 5  
( 17 > target )

No  
Swap!

5	6	7	9	15	17	17	11
5	6	7	9	15	15	17	11

Compare index = 4  
15 > target

No  
Swap!

5	6	7	9	15	15	17	11
5	6	7	9	15	15	17	11

Compare index = 3  
9 < target

5	6	7	9	10	15	17	11
---	---	---	---	----	----	----	----

array [ Compare index + 1 ] = target

# Insertion Sort



- Pros :
  - Could swap less times than bubble sort
  - Insertion Sort is very fast when the data is nearly sorted
- Cons :
  - Worst case have time complexity  $O(n^2)$

# Insertion Sort

```
for (unsigned int i = 1; i < size; i++) {  
    int target = array[i];  
    int j = i - 1;  
    while (j >= 0) {  
        #####implement here#####  
  
        #####  
    }  
    array[j + 1] = target;  
}
```

- Hint
  - Index “i” :Target (Green box)
  - Index “j” : Compare (Red box)

# Insertion Sort

```
for (unsigned int i = 1; i < size; i++) {
    int target = array[i];
    int j = i - 1;
    while (j >= 0) {
        //#####implement here#####
        if (array[j] <= target) break;
        else {
            array[j+1] = array[j];
            j--; // decrease j by 1 (the same as "j = j-1;")
        }
        //#####
    }
    array[j + 1] = target;
}
```

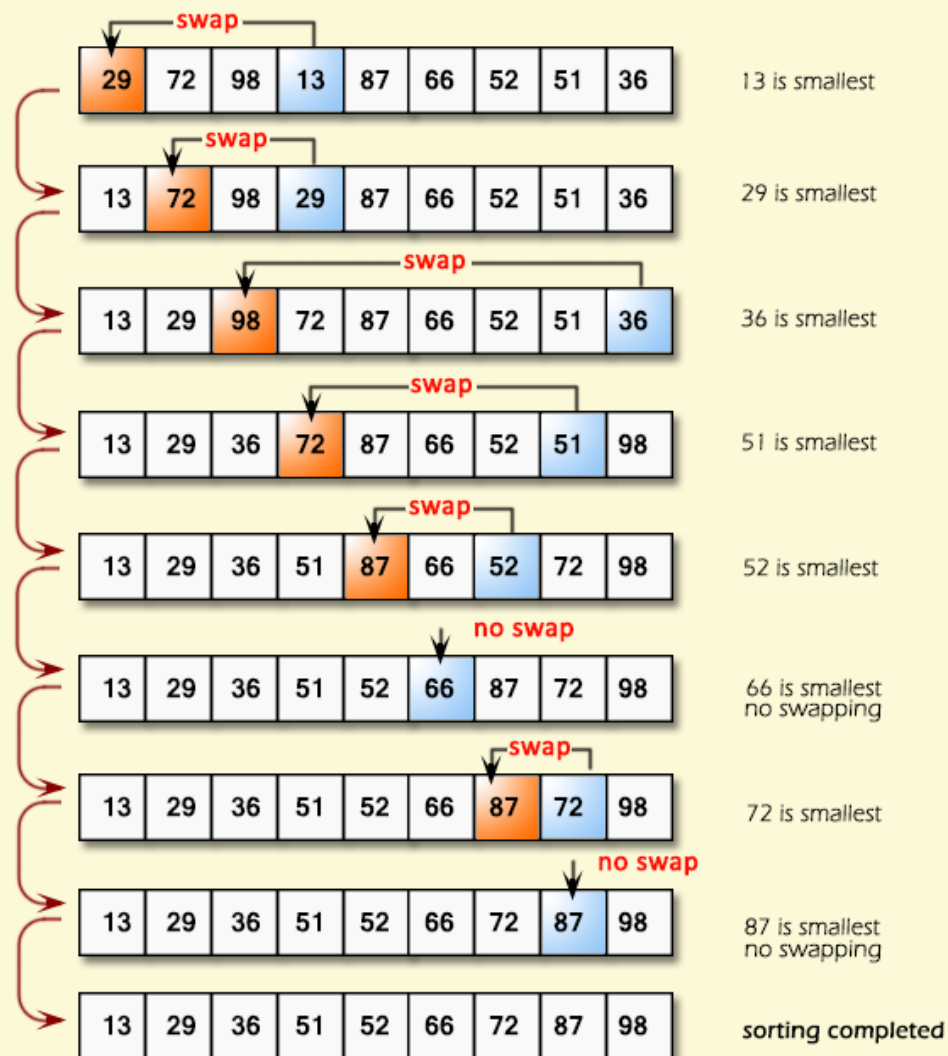
- Hint
  - Index “i” :Target (Green box)
  - Index “j” : Compare (Red box)

# Outline

- If / For loop / While loop
- Array
- What is sorting?
- **Incremental Algorithm**
  - Bubble sort
  - Insertion sort
  - **Selection sort**
  - Heap sort

# Selection Sort

## Selection Sort



1. Start from the index 0

Swap with the minimum value in the remaining list

2. Start from the index 1

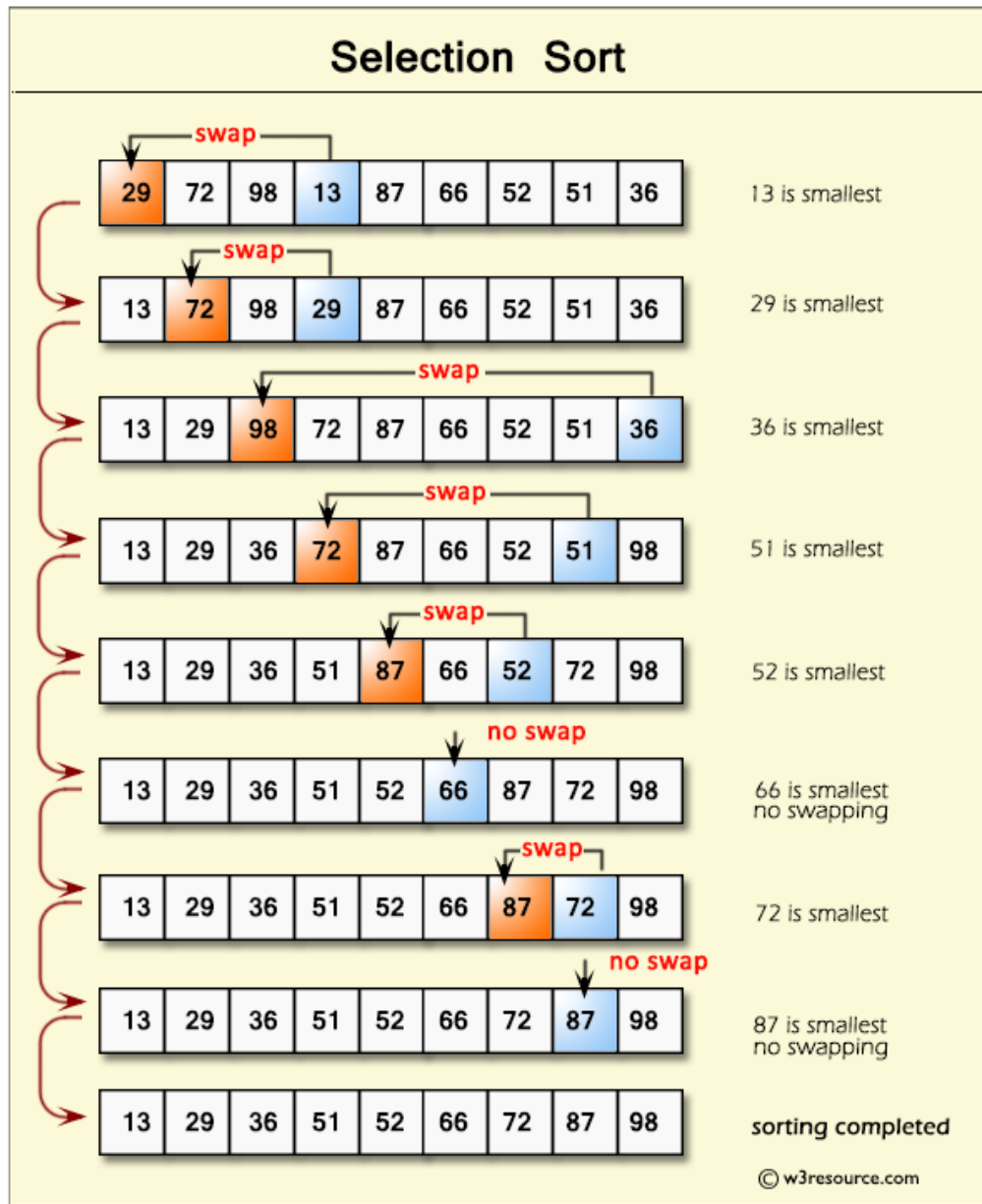
Swap with the minimum value in the remaining list

.. continue




# Selection Sort

- Pros
  - A few swap
- Cons
  - A lot of compare
  - Time complexity  $O(n^2)$



# Selection Sort

```
for (unsigned int i = 0; i < size - 1; i++)
{
    int min = array[i];
    int min_idx = i;
    for (unsigned int j = i + 1; j < size; j++)
    {
        #####implement here#####
        
        #####
    }
    int temp = array[i];
    array[i] = array[min_idx];
    array[min_idx] = temp;
}
```

- **Hint**

- Index “i” : Red box in the “figure”
- Index “j” : Blue box in the “figure”

# Selection Sort

```
for (unsigned int i = 0; i < size - 1; i++)
{
    int min = array[i];
    int min_idx = i;
    for (unsigned int j = i + 1; j < size; j++)
    {
        #####implement here#####
        if (array[j] < min) {
            min = array[j];
            min_idx = j;
        }
        #####
    }
    int temp = array[i];
    array[i] = array[min_idx];
    array[min_idx] = temp;
}
```

- Hint

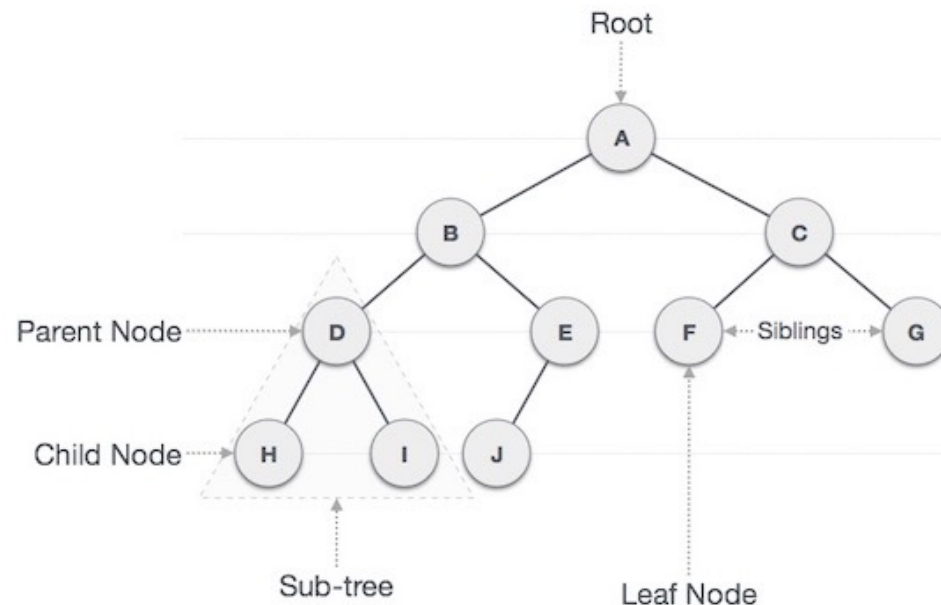
- Index “i” : Red box in the “figure”
- Index “j” : Blue box in the “figure”

# Outline

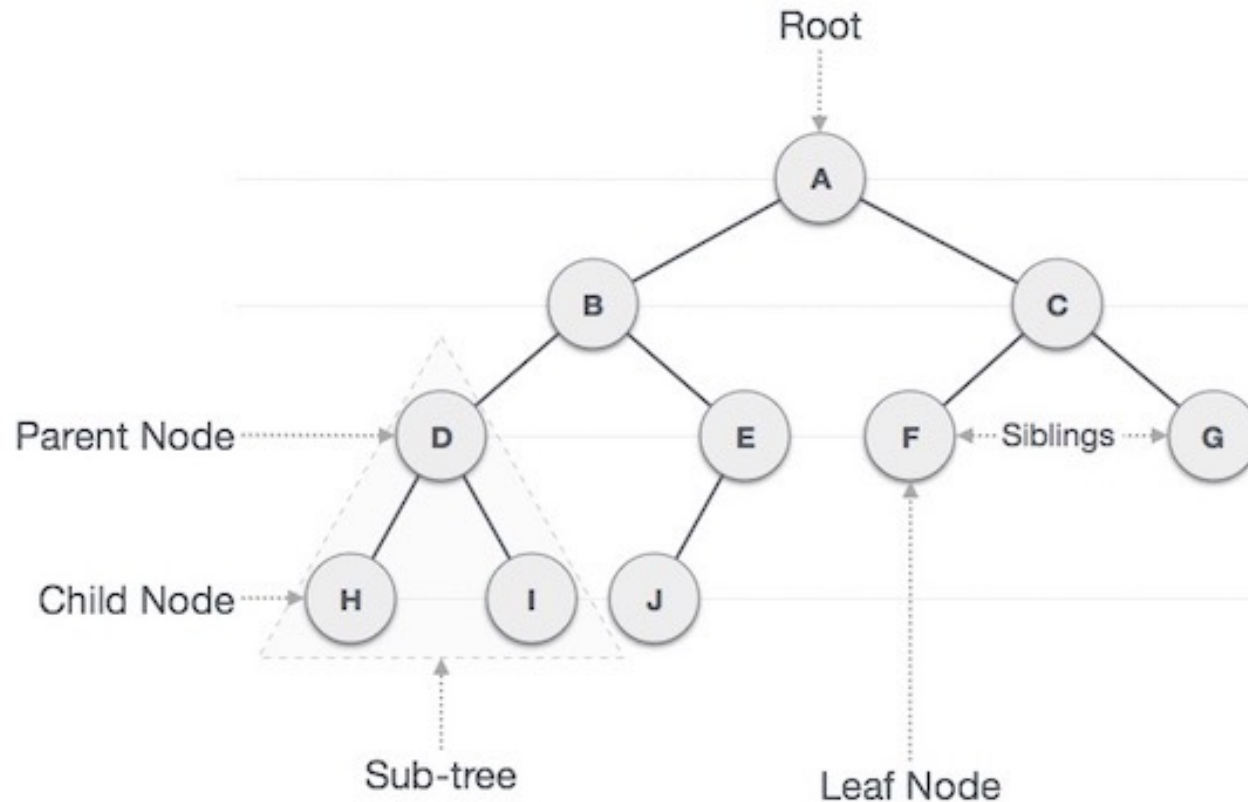
- If / For loop / While loop
- Array
- What is sorting?
- **Incremental Algorithm**
  - Bubble sort
  - Insertion sort
  - Selection sort
  - **Heap sort**

# Heap

- Heap
  - Complete binary tree-based structure that satisfies Heap property
- Tree
  - Every node (except root node) has only one parent node.
  - No cycle.

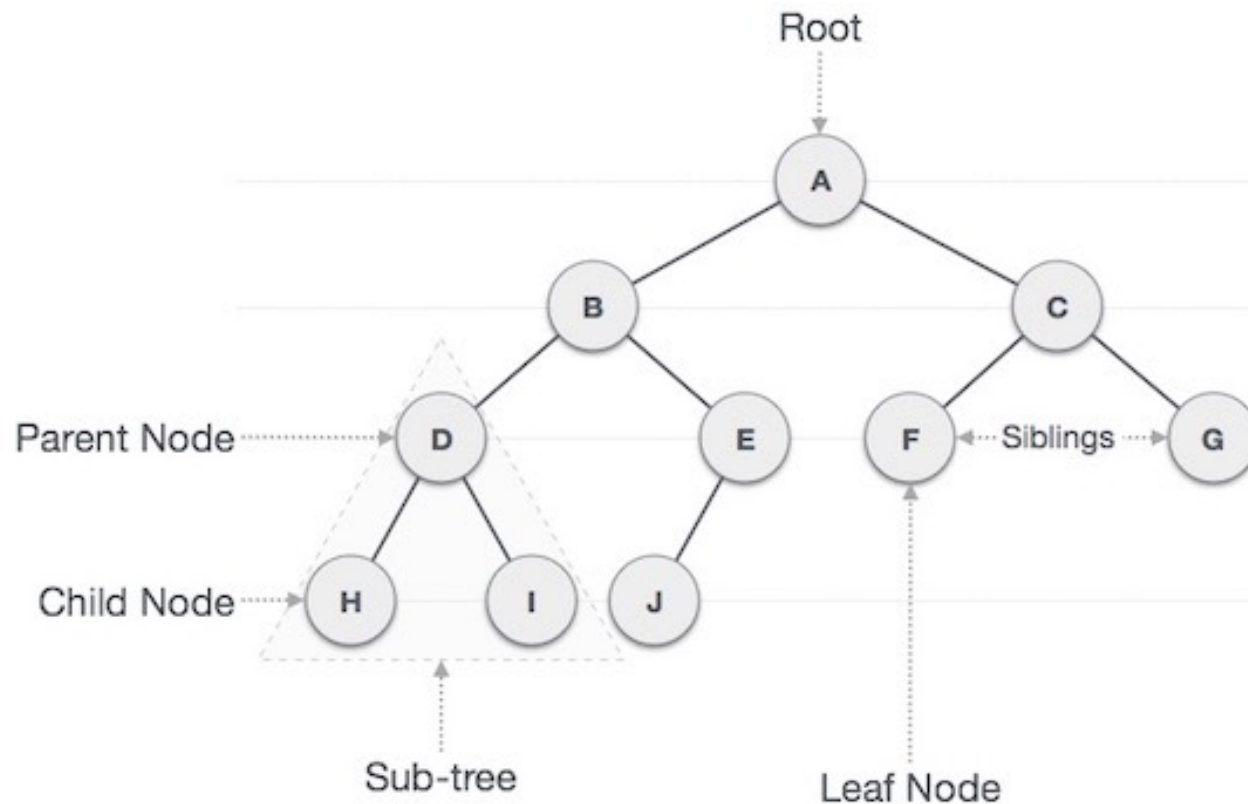


# Heap



- **Node** : Elements composing a tree structure
- **Root node/Root** : Starting point of a tree. The highest node that has no parent node
- **Parent node** : A upper node connected towards the root node

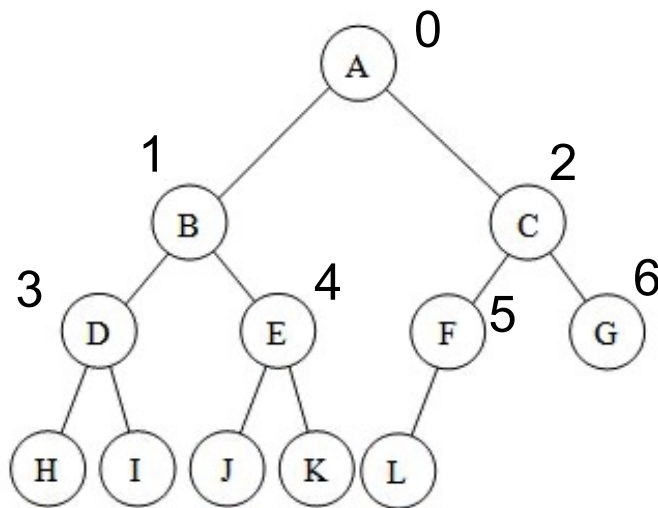
# Heap



- **Child node** :A lower node connected towards the opposite of the root node
- **Siblings node** :The nodes with the same parent node
- **Leaf node** :The node that has no child node

# Heap

- Complete binary tree
  - Binary tree: each parent node has at most two children
  - Complete binary tree:
    - 1) All levels except the last are completely filled
    - 2) Nodes in the last level are filled from left to right



Left child index =  $2 * \text{index} + 1$

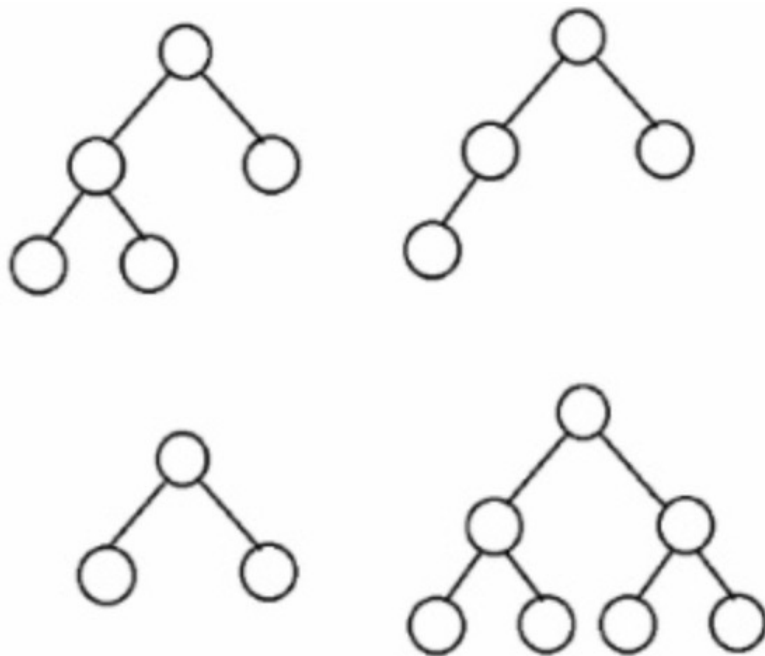
Right child index =  $2 * \text{index} + 2$

[A, B, C, D, E, F, G, H, I, J, K, L]

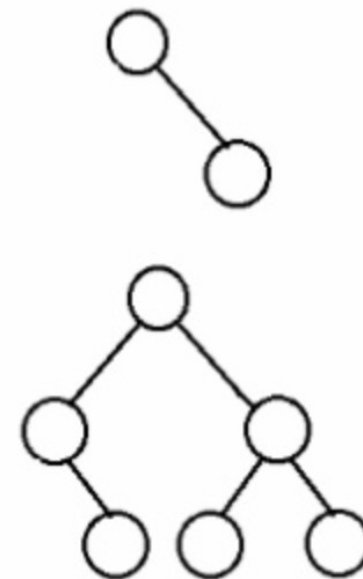


# Heap

- Complete binary tree



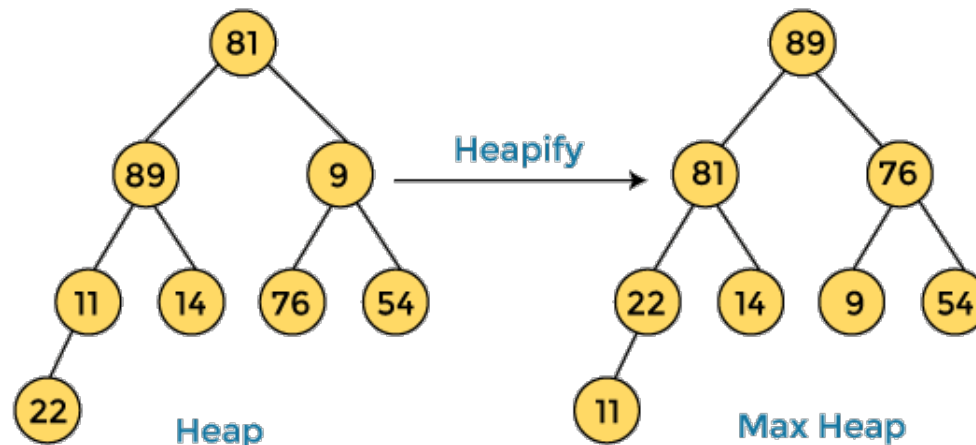
(Complete binary tree)



(**Not** Complete binary tree)

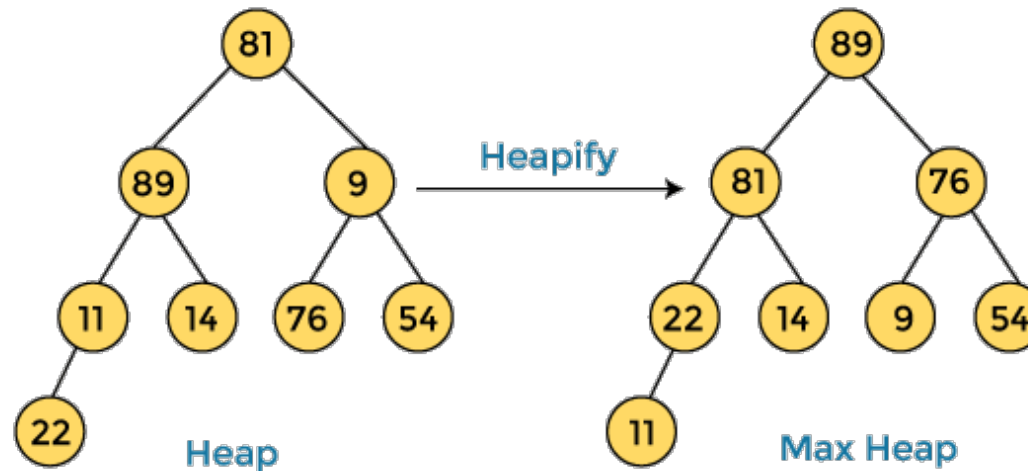
# Heap

- Heap
  - Complete binary tree-based structure with Heap property
  - Heap property
    - **Max heap property:**
      - Parent node's value  $\geq$  children's values (root is largest)
    - **Min heap property:**
      - Parent node's value  $\leq$  children's values (root is smallest)



# Heap

- Max-Heapify



81	89	9	11	14	76	54	22
----	----	---	----	----	----	----	----

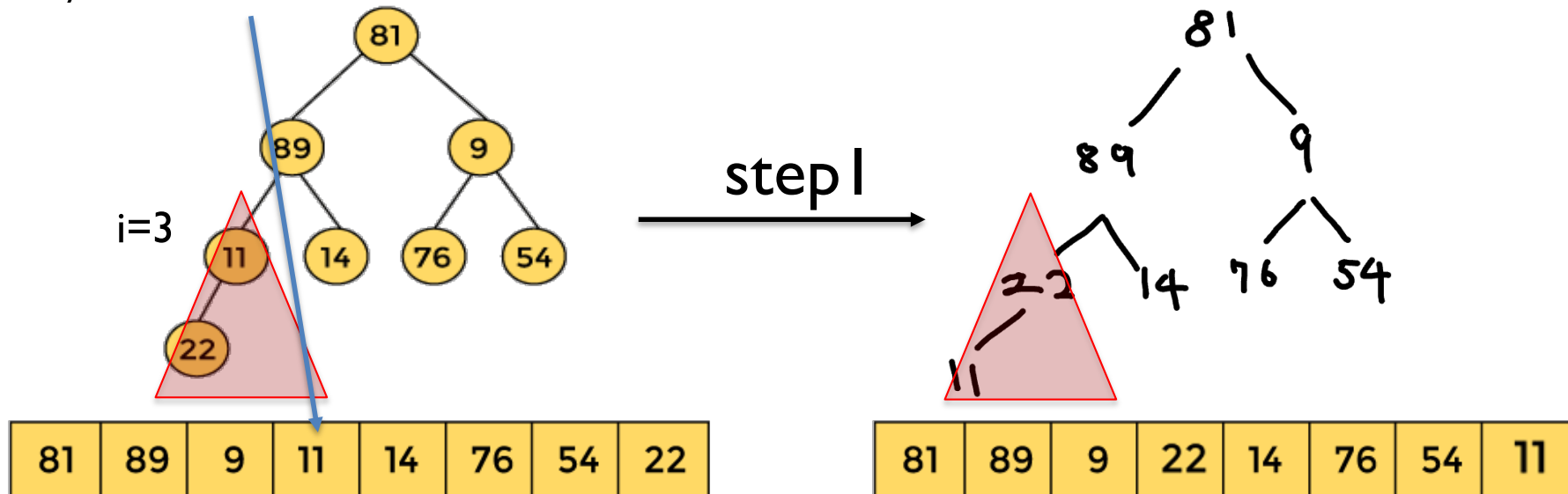
89	81	76	22	14	9	54	11
----	----	----	----	----	---	----	----

# Heap

## How to make a max heap

initial index = (tree size) / 2 - 1

$$8/2 - 1 = 3$$



### Build-Max-Heap(A)

heap\_size = length(A)

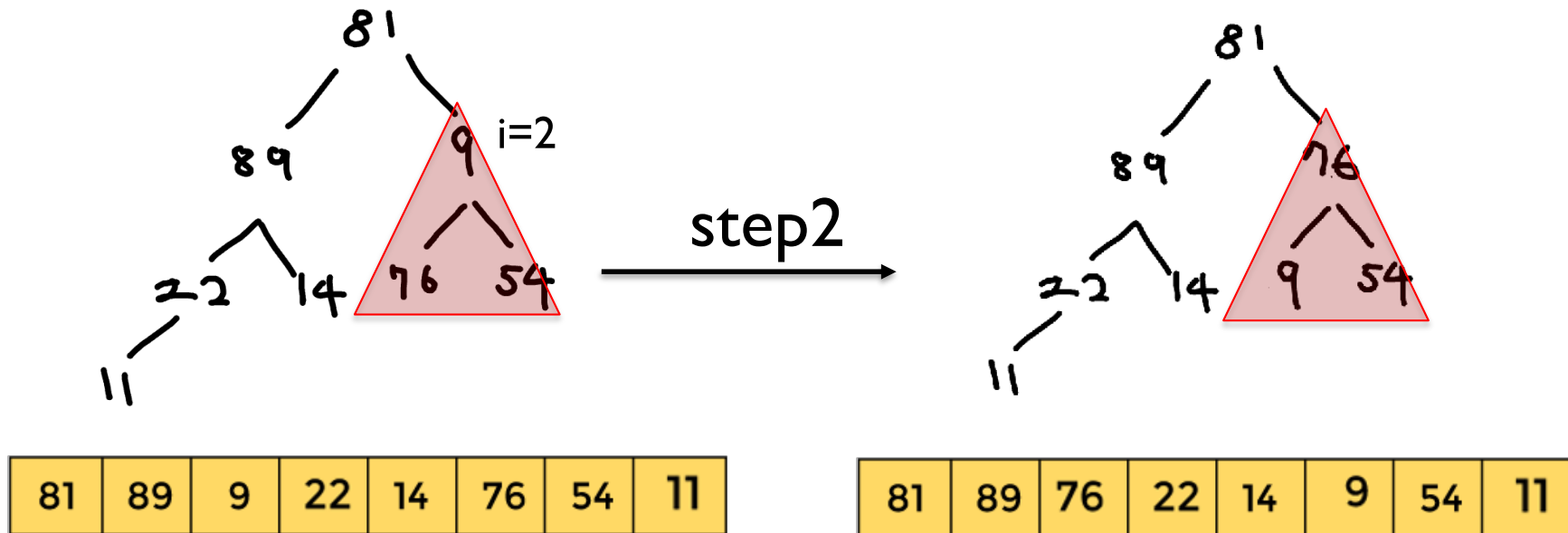
for  $i \leftarrow \text{heap\_size}/2 - 1$  **down to** 0

do **Max-Heapify**(A,i)

# Heap

## How to make a max heap

Swap the parent's value with the largest of the three numbers



### Build-Max-Heap(A)

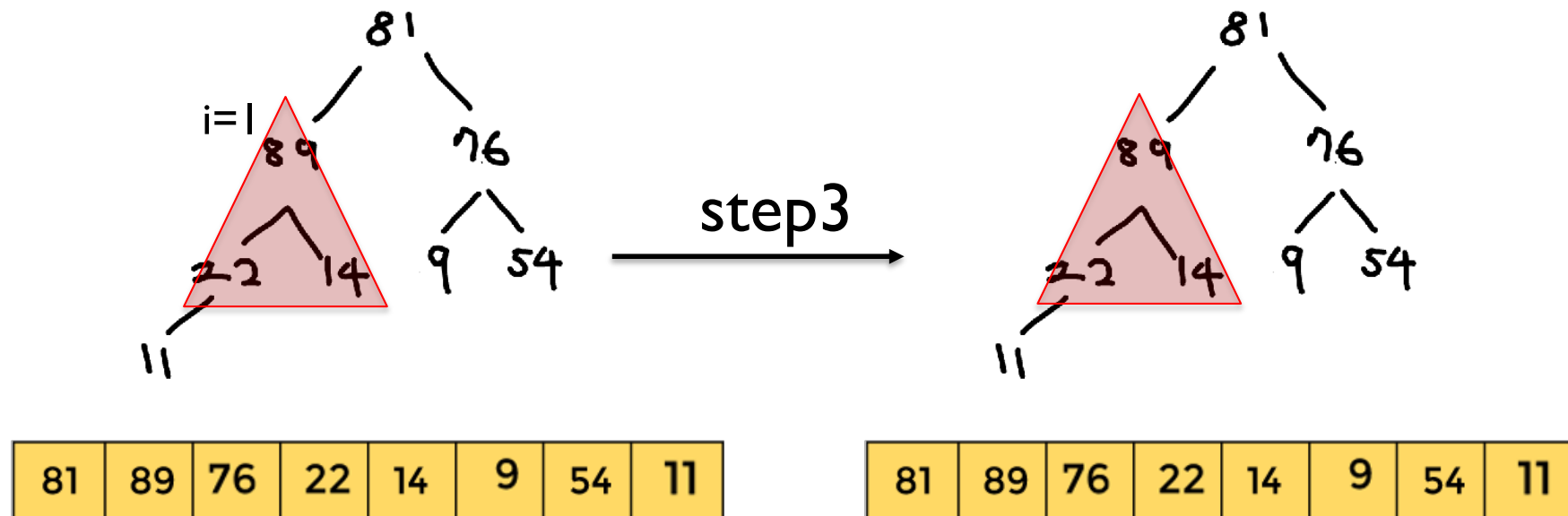
heap\_size = length(A)

for  $i \leftarrow \text{heap\_size}/2 - 1$  **down to** 0

do **Max-Heapify**(A,i)

# Heap

## How to make a max heap



### Build-Max-Heap(A)

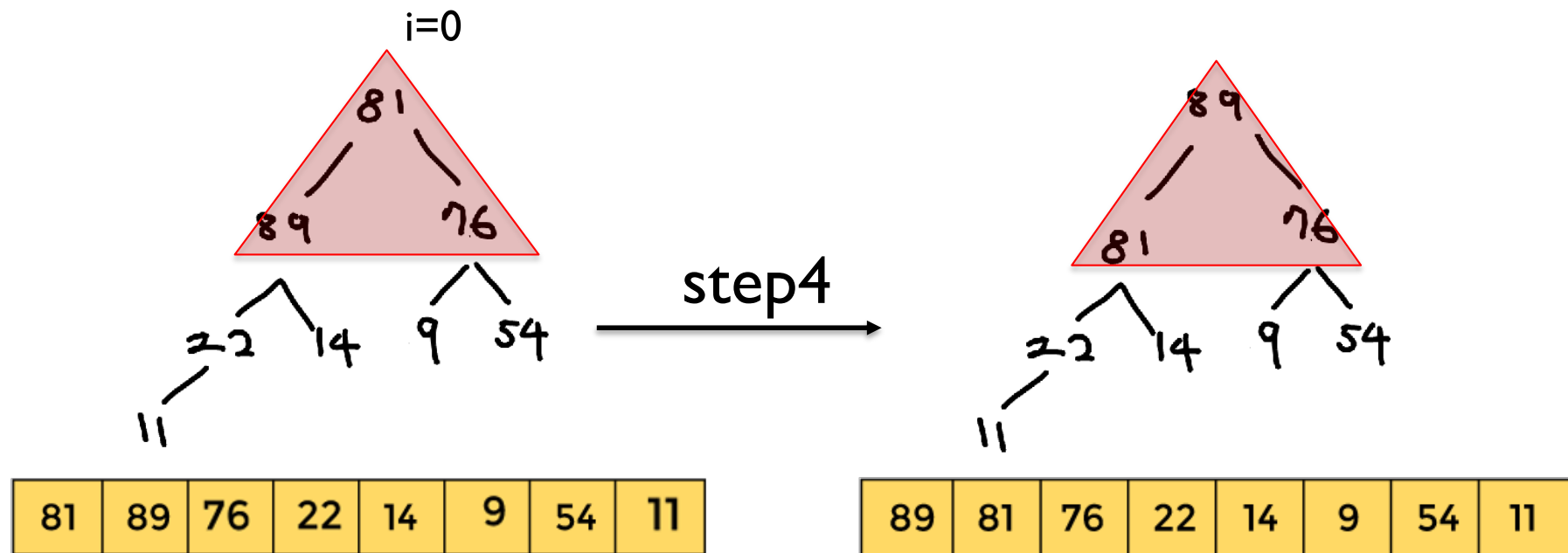
heap\_size = length(A)

for  $i \leftarrow \text{heap\_size}/2 - 1$  **down to** 0

do **Max-Heapify**(A,i)

# Heap

## How to make a max heap



### Build-Max-Heap(A)

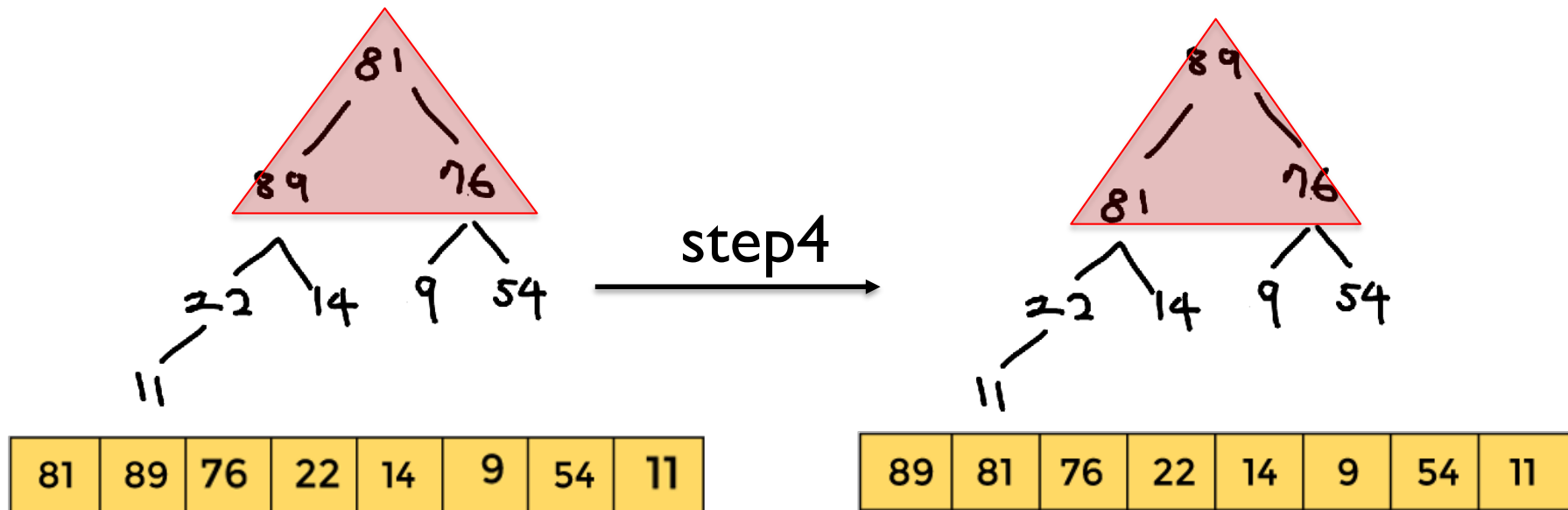
heap\_size = length(A)

for  $i \leftarrow \text{heap\_size}/2 - 1$  **down to** 0

do **Max-Heapify**(A,i)

# Heap

## How to make a max heap



### Build-Max-Heap(A)

heap\_size = length(A)

for  $i \leftarrow \text{heap\_size}/2 - 1$  **down to** 0  
do Max-Heapify(A,i)

MAX-HEAPIFY(A, i)

```

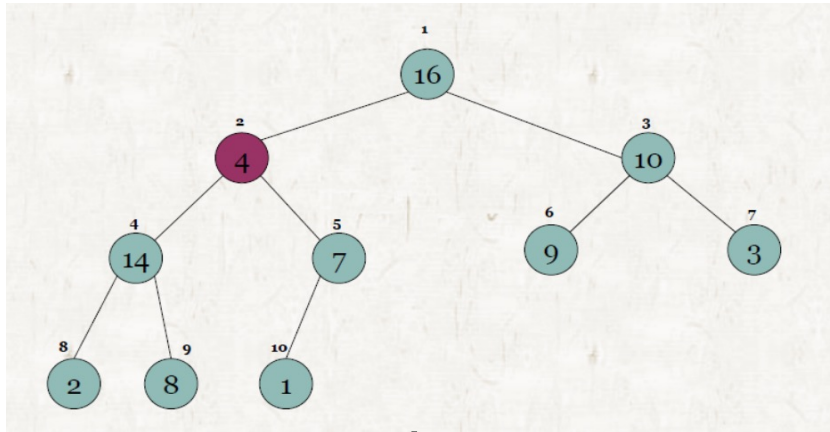
1  l ← LEFT(i)
2  r ← RIGHT(i)
3  if l < heap-size[A] and A[l] > A[i]
4    then largest ← l
5  else largest ← i
6  if r < heap-size[A] and A[r] > A[largest]
7    then largest ← r
8  if largest ≠ i
9    then exchange A[i] ↔ A[largest]
10   MAX-HEAPIFY(A, largest)
    
```

**Recursive call  
Why?**



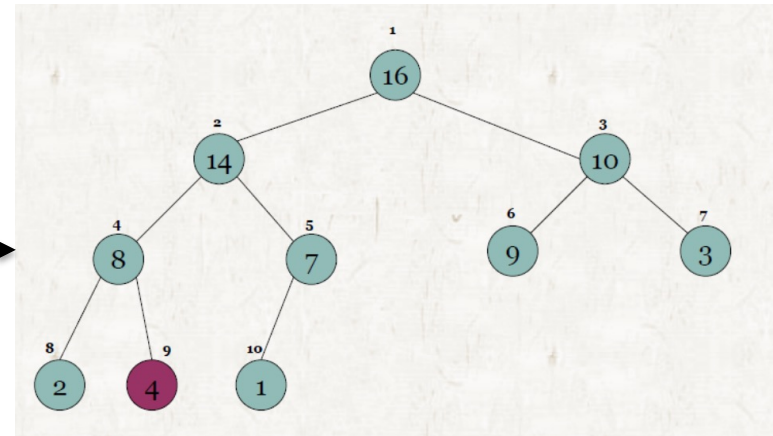
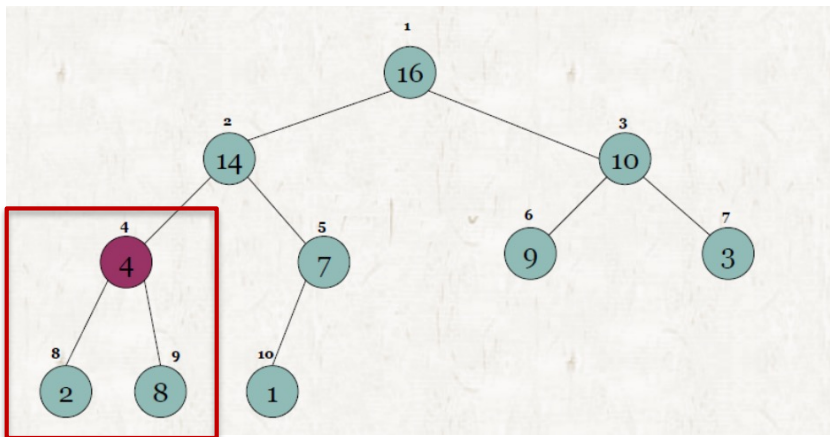
# Heap

## How to make a max heap



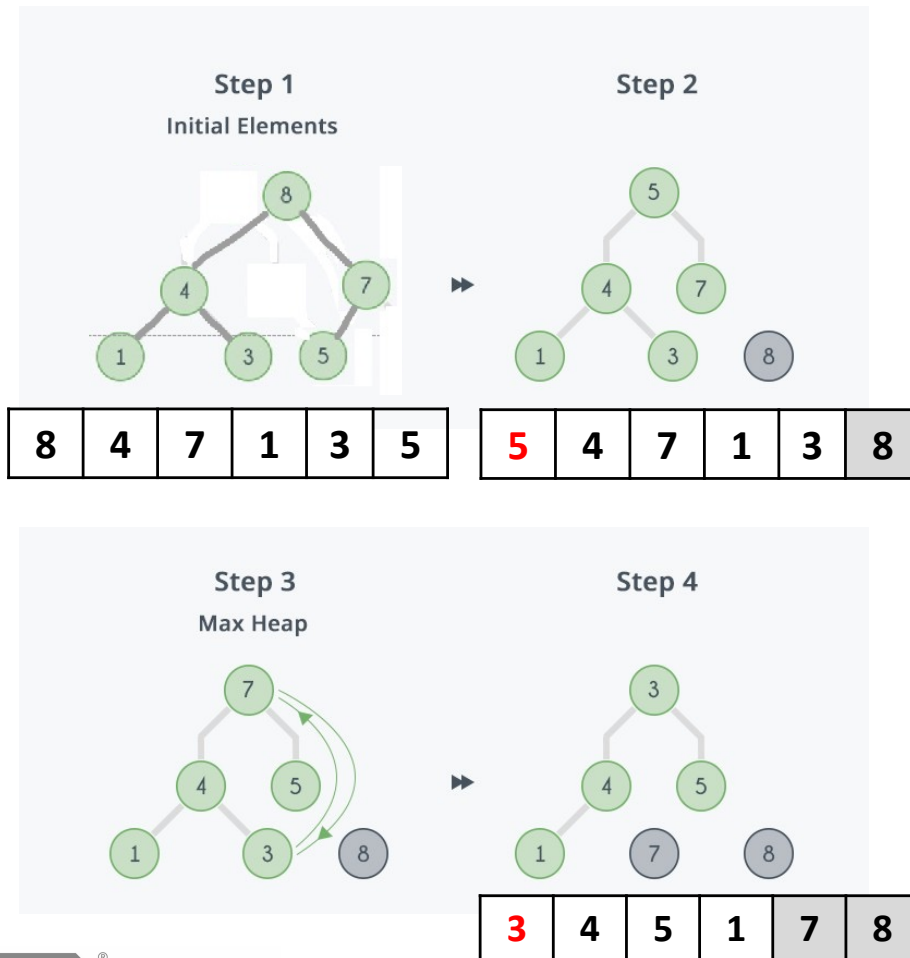
**Assume the scenario where we are heapifying at index 2**

- During Max-Heapify, 4 moves down the tree, and the Max-Heap property of the lower subtree is violated
- Therefore, we need to heapify that subtree again



# Heap Sort

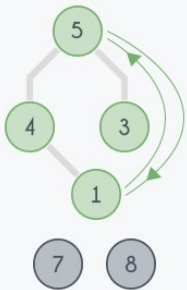
- From Max Heap



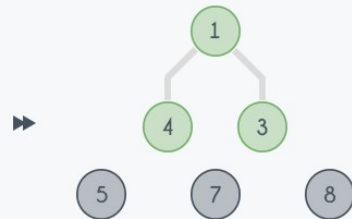
1. Build a max heap from **Arr**
2. The root element,  $\text{Arr}[0]$  is the maximum element  
→ Swap with the last element of **Arr**
3. Max-Heapify from **root node** because all node satisfies the max heap property, except for the root node
4. Repeat 2-3 step

# Heap Sort

Step 5  
Max Heap

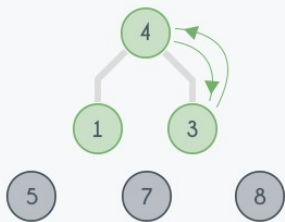


Step 6



1	4	3	5	7	8
---	---	---	---	---	---

Step 7  
Max Heap



Step 8



3	1	4	5	7	8
---	---	---	---	---	---

Step 9  
Max Heap



Step 10



1	3	4	5	7	8
---	---	---	---	---	---

# Heap Sort

- Pros
  - Nice time complexity
  - The most efficient among the sorting algorithms with  $O(n \log n)$ , ideally
- Cons
  - According to the state of the input, slower than some algorithms (quick sort, merge sort) in practice.

## Build-Max-Heap(A)

heap\_size = length(A)

$O(n)$  for  $i \leftarrow \text{heap\_size}/2 - 1$  down to 0  
do Max-Heapify(A, i)

MAX-HEAPIFY(A, i)

```
1 l ← LEFT(i)
2 r ← RIGHT(i)
3 if l < heap-size[A] and A[l] > A[i]
4   then largest ← l
5   else largest ← i
6 if r < heap-size[A] and A[r] > A[largest]
7   then largest ← r
8 if largest ≠ i
9   then exchange A[i] ↔ A[largest]
10  MAX-HEAPIFY(A, largest)
```

$O(\log n)$

( $\because$  recursive call and  
tree depth is  $\log n$ )

# Heap Sort

```
void max_heapify(int* array, int index, int heap_size) {  
    int L = 2 * index + 1;  
    int R = 2 * index + 2; Array argument  
    int largest;  
  
    //#####implement here#####  
    largest = index;  
    if (L < heap_size && array[L] > array[index])  
        largest = L;  
    if (R < heap_size && array[R] > array[largest])  
        largest = R;  
    if (largest != index) {  
        //#####  
        array[index] = array[largest];  
        array[largest] = array[index];  
        max_heapify(array, largest, heap_size);  
    }  
    //#####  
}
```

```
MAX-HEAPIFY(A, i)  
1  l ← LEFT(i)  
2  r ← RIGHT(i)  
3  if l < heap-size[A] and A[l] > A[i]  
4      then largest ← l  
5      else largest ← i  
6  if r < heap-size[A] and A[r] > A[largest]  
7      then largest ← r  
8  if largest ≠ i  
9      then exchange A[i] ↔ A[largest]  
10     MAX-HEAPIFY(A, largest)
```

# Heap Sort

## Build-Max-Heap(A)

```
heap_size = length(A)
```

```
for i ← heap_size/2 - 1 down to 0
```

do **Max-Heapify**(A,i)

```
int heap_size = size;
for (int i = size/2 - 1; i >= 0; i--)
    max_heapify(array, i, heap_size);
```

## Build-max-heap

# Sorting

# Heap Sort

```
void max_heapify(int* array, int index, int heap_size) {  
    int L = 2 * index + 1;  
    int R = 2 * index + 2; Array argument  
    int largest;  
  
    //#####implement here#####  
    largest = index;  
    if (L < heap_size && array[L] > array[index])  
        largest = L;  
    if (R < heap_size && array[R] > array[largest])  
        largest = R;  
    if (largest != index) {  
        int temp = array[index];  
        array[index] = array[largest];  
        array[largest] = temp;  
        max_heapify(array, largest, heap_size);  
    }  
    //#####  
}
```

```
MAX-HEAPIFY(A, i)  
1  l ← LEFT(i)  
2  r ← RIGHT(i)  
3  if l < heap-size[A] and A[l] > A[i]  
4      then largest ← l  
5      else largest ← i  
6  if r < heap-size[A] and A[r] > A[largest]  
7      then largest ← r  
8  if largest ≠ i  
9      then exchange A[i] ↔ A[largest]  
10     MAX-HEAPIFY(A, largest)
```

# Heap Sort

- Time complexity  $O(n \log n)$

```
int heap_size = size;
for (int i = size/2 - 1; i >= 0; i--)
    max_heapify(array, i, heap_size); // O(logn)

for (int i = size - 1; i > 0; i--) {
    //#####implement here#####
    int temp = array[0];
    array[0] = array[i];
    array[i] = temp;
    max_heapify(array, 0, --heap_size); // O(logn)
    //#####
}
```

$O(n \log n)$

$O(n \log n)$

$O(n \log n)$

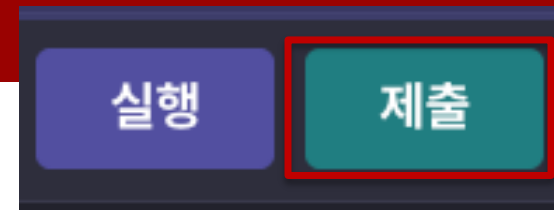


# Assignment

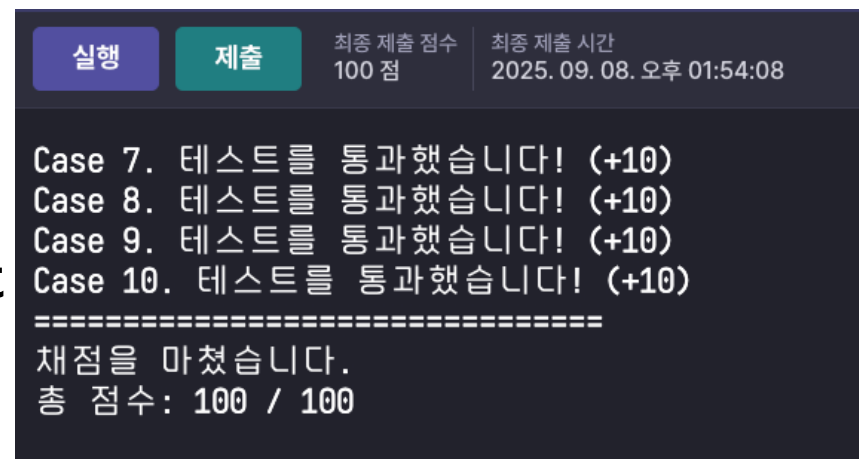
- Problem
  - Implement Bubble sort / Insertion sort / Selection sort / Heap sort (**Descending order**)
- Example
  - **Input:** array size, sorting type, and array with random order.
  - **Output:** sorted array with descending order.
  - Ex) Size 5, bubble sort, input array= [12, 8, 11, 2, 19]  
Output array= [19, 12, 11, 8, 2]

```
Give me the size : 5
Give me the type of algorithm (0: Bubble, 1: Insertion, 2: Selection, 3: Heap): 0
12 8 38 3 1
38 12 8 3 1
Reset
```





# Assignment



- Due date : 9/15(Mon) 14:30 PM
- Push “제출” button on Elice
- There is no hidden case
  - A result like the one on the right means you got 100 points.



- Once the due date has passed, the assignment will be hidden and can no longer be submitted

- ☐ 학습 수업 자료
- ☐  practice 
- ☐  assignment 

# Lab 2 week 1 assignment

1. Environment setting assignment

2. Code of Ethics

Due date : 9/15 (Mon) 14:30 PM

As some students have joined during the  
“수강신청변경기간” (course add/drop) period,  
the deadline will be extended by one week

# Attendance Check

- Quiz in Elice